



Article

## Near Hexagons of Order $(3, t)$ having a Big Quad

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**Abstract:** We classify all near hexagons of order  $(3, t)$  that contain a big quad. We show that, up to isomorphism, there are ten such near hexagons.

**Keywords:** Near hexagon, Big quad

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### 1. Introduction

A *generalised quadrangle* is a point-line geometry that satisfies the following properties:

(GQ1) There exist two disjoint lines.

(GQ2) For every non-incident point-line pair  $(x, L)$ , there exists a unique point on  $L$  collinear with  $x$ .

An *ovoid* of a generalised quadrangle is a set of points that meets each line in a singleton. A *rosette of ovoids* is a set of ovoids through a given point  $x$  that partitions the set of points noncollinear with  $x$ . The point  $x$  is then called the *centre* of the rosette.

A *near hexagon* is a point-line geometry that satisfies the following properties:

(NH1) Every two distinct points are incident with at most one line.

(NH2) For every point  $x$  and every line  $L$ , there exists a unique point  $\pi_L(x)$  on  $L$  which is nearest to  $x$  with respect to the distance function  $d(\cdot, \cdot)$  of the collinearity graph  $\Gamma$ .

(NH3) The diameter of  $\Gamma$  is equal to 3.

A *quad* of a near hexagon  $\mathcal{S}$  is a set  $Q$  of points satisfying the following properties:

(Q1) The diameter of  $Q$  is equal to 2.

(Q2) If  $x$  and  $y$  are two distinct points of  $Q$ , then any point collinear with  $x$  and  $y$  also belongs to  $Q$ .

(Q3) The point-line geometry formed by the points of  $Q$  and those lines of  $\mathcal{S}$  that have all their points in  $Q$  is a generalised quadrangle.

Such a quad is called *big* if every point  $x$  of  $\mathcal{S}$  outside  $Q$  is collinear with a (necessarily unique) point  $\pi_Q(x)$  of  $Q$ .

A point-line geometry is said to have *order*  $(s, t)$  if every line is incident with precisely  $s + 1$  points and if every point is incident with exactly  $t + 1$  lines. A generalised quadrangle  $Q$  with at least three points on each line has a constant number  $t_Q + 1$  of lines through each point.

A near hexagon is called *dense* if every line is incident with at least three points and if every two points at distance 2 have at least two common neighbours. In [1], we classified finite dense near hexagons with four points per line. There are ten known examples of such near hexagons, namely the near hexagons\*  $\mathbb{L}_4 \times \mathbb{L}_4 \times \mathbb{L}_4$ ,  $W(3) \times \mathbb{L}_4$ ,  $Q(4, 3) \times \mathbb{L}_4$ ,  $GQ(3, 5) \times \mathbb{L}_4$ ,  $Q^-(5, 3) \times \mathbb{L}_4$ ,  $DW(5, 3)$ ,  $DQ(6, 3)$ ,  $DH(5, 9)$ ,  $GQ(3, 5) \otimes GQ(3, 5)$  and  $Q^-(5, 3) \otimes Q^-(5, 3)$ . Each of these has a big quad. In [1], we showed that the other (unknown) finite dense near hexagons with four points per line fall into four classes and that none of them has a big quad. This note arose from investigations whether these ten known examples are the only near hexagons of order  $(3, t)$  with a big quad. Previously, also all near hexagons of order  $(2, t)$  with a big quad had been classified and also there all examples turned out to be dense, see [2, §6] and [3, §10.7].

After successfully classifying all near hexagons of order  $(3, t)$  containing a big quad, we noticed that the used arguments could be generalised to cover a more general result. We will therefore prove our results in this more general setting. Specifically, we show the following;

**Theorem 1.** *Suppose  $\mathcal{S}$  is a finite near hexagon of order  $(s, t)$ , where  $s \geq 2$ . Let  $\mathcal{Q}$  denote the set of all quads of  $\mathcal{S}$  and denote by  $(s, t_Q)$  the order of a quad  $Q \in \mathcal{Q}$ . Suppose that at least one quad of  $\mathcal{S}$  is big and that all numbers  $t_Q$ ,  $Q \in \mathcal{Q}$ , are odd. Suppose also that every quad  $Q \in \mathcal{Q}$  that is not a grid has no partition in ovoids or no rosette of ovoids. Then  $\mathcal{S}$  must be dense.*

Suppose now that  $\mathcal{S}$  is a near hexagon of order  $(3, t)$  with a big quad. We will show in Lemma 4 that  $\mathcal{S}$  is finite. Every quad of  $\mathcal{S}$  only has lines of size four. By [4] (see also [5, §6.2]), there are only five generalised quadrangles with four points per line, namely the  $(4 \times 4)$ -grid of order  $(3, 1)$ , the symplectic generalised quadrangle  $W(3)$  of order  $(3, 3)$ , the parabolic generalised quadrangle  $Q(4, 3)$  of order  $(3, 3)$ , the unique generalised quadrangle  $GQ(3, 5)$  of order  $(3, 5)$  and the elliptic generalised quadrangle  $Q^-(5, 3)$  of order  $(3, 9)$ . Only one of these five generalised quadrangles has both partitions in ovoids and rosettes of ovoids, namely the  $(4 \times 4)$ -grid, see Lemma 3. Theorem 1 then implies.

**Corollary 1.** *Suppose  $\mathcal{S}$  is a (possibly infinite) near hexagon of order  $(3, t)$  having a big quad. Then  $\mathcal{S}$  is finite and dense, and therefore isomorphic to either  $\mathbb{L}_4 \times \mathbb{L}_4 \times \mathbb{L}_4$ ,  $W(3) \times \mathbb{L}_4$ ,  $Q(4, 3) \times \mathbb{L}_4$ ,  $GQ(3, 5) \times \mathbb{L}_4$ ,  $Q^-(5, 3) \times \mathbb{L}_4$ ,  $DW(5, 3)$ ,  $DQ(6, 3)$ ,  $DH(5, 9)$ ,  $GQ(3, 5) \otimes GQ(3, 5)$  or  $Q^-(5, 3) \otimes Q^-(5, 3)$ .*

In a general near hexagon with four points per line, we show in Corollary 3 that any quad isomorphic to  $W(3)$ ,  $GQ(3, 5)$  or  $Q^-(5, 3)$  has to be big. Combining this with Corollary 1, we thus find the following;

**Corollary 2.** *Suppose  $\mathcal{S}$  is a (possibly infinite) near hexagon of order  $(3, t)$  that is not isomorphic to one of the ten examples mentioned in Corollary 1. Then every quad of  $\mathcal{S}$  is isomorphic to the  $(4 \times 4)$ -grid or to the generalised quadrangle  $Q(4, 3)$ .*

## 2. Preliminaries

Suppose  $\mathcal{S}$  is a near hexagon. For every point  $x$ ,  $\Gamma_i(x)$  with  $i \in \mathbb{N}$  denotes the set of points at distance  $i$  from  $x$ . For every nonempty set  $X$  of points,  $\Gamma_i(X)$  with  $i \in \mathbb{N}$  denotes the set of points at distance  $i$  from  $X$ , i.e. the set of points  $x$  for which the minimal distance from  $x$  to a point of  $X$  is equal to  $i$ . The following result will be useful in our discussion;

**Proposition 1** ([6, Proposition 2.5]). *Suppose  $x$  and  $y$  are two points of a near hexagon  $\mathcal{S}$  at distance 2 from each other such that  $x$  and  $y$  have two common neighbours  $a$  and  $b$  such that at least one of the*

\*For a definition of these near hexagons, see e.g. [1].

four lines  $xa, xb, ay, by$  contains at least three points. Then  $x$  and  $y$  are contained in a unique quad  $Q(x, y)$ .

For every point  $x$  of  $\mathcal{S}$ , the *local space*  $\mathcal{L}_x$  at  $x$  is the point-line geometry whose points and lines are the lines and quads through  $x$ , respectively, with incidence being containment. This local space is always a *partial linear space*. Proposition 1 implies the following.

- In a dense near hexagon, every local space is linear.
- In a nondense near hexagon with more than two points on each line, at least one local space is not linear, i.e. there exist two intersecting lines that are not contained in a quad, or equivalently, there exist two points at distance 2 that have a unique common neighbour.

We now collect a number of known properties of quads in a near hexagon  $\mathcal{S}$ . We refer to the literature (e.g. [3, 7]) for proofs.

If  $(x, Q)$  is a point-quad pair of  $\mathcal{S}$ , then  $d(x, Q) \leq 2$ . If  $d(x, Q) = 1$ , then  $Q$  contains a unique point  $\pi_Q(x)$  collinear with  $x$  and  $d(x, y) = 1 + d(\pi_Q(x), y)$  for every point  $y$  of  $Q$ . If  $d(x, Q) = 2$ , then the set  $O_x := \Gamma_2(x) \cap Q$  is an ovoid of  $Q$ .

If  $(x, Q)$  is a point-quad pair with  $d(x, Q) = 2$ , then there are two possibilities for a line  $L$  through  $x$ :

- We have  $L \subseteq \Gamma_2(Q)$ . In this case, the ovoids  $O_y, y \in L$ , form a partition of ovoids of  $Q$ .
- The line  $L$  contains a unique point  $u \in \Gamma_1(Q)$  and  $L \setminus \{u\} \subseteq \Gamma_2(Q)$ . In this case, the ovoids  $O_y, y \in L \setminus \{u\}$ , form a rosette of ovoids of  $Q$  with centre  $\pi_Q(u)$ .

If  $Q$  is a quad and  $L$  a line contained in  $\Gamma_1(Q)$ , then  $\pi_Q(L) := \{\pi_Q(x) \mid x \in L\}$  is a line of  $Q$ , and the map  $x \mapsto \pi_Q(x)$  defines a bijection between  $L$  and  $\pi_Q(L)$ . From this it can easily be deduced that if  $Q_1$  and  $Q_2$  are two disjoint quads with  $Q_1$  big, then  $\pi_{Q_1}(Q_2) := \{\pi_{Q_1}(x) \mid x \in Q_2\}$  is a subquadrangle of  $Q_1$  isomorphic to  $Q_2$ . On the other hand, if  $Q_1$  and  $Q_2$  are two distinct intersecting quads of  $\mathcal{S}$  with  $Q_1$  big, then  $Q_1 \cap Q_2$  must be a line.

Also the following lemmas will be useful in our treatment;

**Lemma 1.** *Suppose  $\mathcal{S}$  is a finite near hexagon of order  $(s, t)$  having  $v$  points. Then  $|\Gamma_2(x)| = n_2 := \frac{v}{s+1} - 1 + s^2t - st$  and  $|\Gamma_3(x)| = n_3 := \frac{sv}{s+1} - s - s^2t$  for every point  $x$  of  $\mathcal{S}$ .*

*Proof.* For every point  $x$  of  $\mathcal{S}$ , we have  $|\Gamma_0(x)| + |\Gamma_1(x)| + |\Gamma_2(x)| + |\Gamma_3(x)| = v$ , where  $|\Gamma_0(x)| = 1$  and  $|\Gamma_1(x)| = s(t+1)$ . By [3, Theorem 1.2], we also know that  $|\Gamma_0(x)| - \frac{1}{s}|\Gamma_1(x)| + \frac{1}{s^2}|\Gamma_2(x)| - \frac{1}{s^3}|\Gamma_3(x)| = 0$ . The values of  $|\Gamma_2(x)|$  and  $|\Gamma_3(x)|$  are now easily computed. They are respectively equal to  $n_2$  and  $n_3$ .  $\square$

A generalized quadrangle  $Q$  of order  $(s, t)$  contains  $(s+1)(st+1)$  points, and for any point  $x$  of  $Q$  there are  $(s+1)(st+1) - 1 - |\Gamma_1(x)| = (s+1)(st+1) - 1 - s(t+1) = s^2t$  points in  $Q$  noncollinear with  $x$ .

**Lemma 2.** *Suppose  $\mathcal{S}$  is a finite near hexagon of order  $(s, t)$ ,  $s \geq 2$ , having  $v$  points. Let  $x$  be a point of  $\mathcal{S}$  and denote by  $\mathcal{R}$  the set of all quads through  $x$ . For every  $R \in \mathcal{R}$ , denote by  $(s, t_R)$  the order of the quad  $R$ . Then*

$$\sum_{R \in \mathcal{R}} t_R^2 = t(t+1) - \frac{1}{s^2} \left( \frac{v}{s+1} - 1 + s^2t - st \right).$$

*Proof.* A quad  $R \in \mathcal{R}$  contains  $s^2t_R$  points of  $\Gamma_2(x)$ . So, the number of points of  $\Gamma_2(x)$  not contained in a quad together with  $x$  equals  $n_2 - \sum_{R \in \mathcal{R}} s^2t_R$ , where  $n_2$  is the number defined in Lemma 1. Now, counting pairs  $(u, v)$  of collinear points in  $\Gamma_1(x) \times \Gamma_2(x)$  gives  $s(t+1) \cdot st = \sum_{R \in \mathcal{R}} s^2t_R \cdot (t_R+1) + (n_2 - \sum_{R \in \mathcal{R}} s^2t_R) \cdot 1$ , from which the mentioned equality readily follows.  $\square$

Regarding ovoids of generalized quadrangles of order  $(3, t)$ , the following can be said.

**Lemma 3.** (a) *The generalized quadrangles  $W(3)$  and  $Q^-(5, 3)$  do not have ovoids.*

(b) *The generalized quadrangle  $GQ(3, 5)$  does not have rosettes of ovoids.*

(c) *The generalized quadrangle  $Q(4, 3)$  does not have partitions in ovoids.*

*Proof.* This is precisely Lemma 1 of [1]. Proofs for the nonexistence of ovoids in the generalized quadrangles  $W(3)$  and  $Q^-(5, 3)$  can be found in [5, 3.4.1]. Payne [8, VI.1] classified all ovoids of  $GQ(3, 5)$ , and from this classification it readily follows that there cannot be rosettes of ovoids. The classification of the ovoids of  $Q(4, 3)$  can be found in [7, p. 160], and from this classification, it follows that there cannot be partitions in ovoids.  $\square$

**Corollary 3.** *In a near hexagon with four points per line, every quad isomorphic to either  $W(3)$ ,  $GQ(3, 5)$  or  $Q^-(5, 3)$  is big.*

*Proof.* We must show that  $\Gamma_2(Q) = \emptyset$  for every such quad  $Q$ . Suppose to the contrary that  $\Gamma_2(Q) \neq \emptyset$  and let  $x \in \Gamma_2(Q)$ . Then  $\Gamma_2(x) \cap Q$  is an ovoid of  $Q$ . As  $x \in \Gamma_2(Q)$ , there exists a line through  $x$  meeting  $\Gamma_1(Q)$ . This line would induce a rosette of ovoids of  $Q$ . A contradiction follows from Lemma 3.  $\square$

Regarding near hexagons of order  $(3, t)$ , the following can be said.

**Lemma 4.** *If  $\mathcal{S}$  is a near hexagon of order  $(3, t)$  having a big quad  $Q$ , then  $\mathcal{S}$  is finite.*

*Proof.* It suffices to prove that  $t$  is finite. Any quad  $R$  of  $\mathcal{S}$  also has four points on each line and hence must be finite by the main result of [9]. So,  $t_R \in \mathbb{N}$  if  $(3, t_R)$  denotes the order of  $R$ . In particular,  $t_Q \in \mathbb{N}$ . Let  $x$  be an arbitrary point outside  $Q$  and let  $L$  be the unique line through  $x$  meeting  $Q$  in a point (namely  $\pi_Q(x)$ ). Let  $K$  be a line through  $x$  distinct from  $L$ , let  $y$  be a point of  $K$  distinct from  $x$ . Now, the points  $x$  and  $\pi_Q(y)$  have at least two common neighbours (namely  $\pi_Q(x)$  and  $y$ ) and so are contained in a unique quad by Proposition 1. This quad contains the lines  $K, L$  and meets  $Q$  in a line. This implies that  $t = \sum_{R \in \mathcal{R}'} t_R$ , where  $\mathcal{R}'$  is the set of quads through the line  $L$ . Recall that each number  $t_R, R \in \mathcal{R}'$ , is finite. Since any quad through  $L$  meets  $Q$  in a line, we have  $|\mathcal{R}'| \leq t_Q + 1$  and so also  $\mathcal{R}'$  is finite. We conclude that  $t = \sum_{R \in \mathcal{R}'} t_R$  is finite as well.  $\square$

### 3. Proof of Theorem 1

Suppose  $\mathcal{S}$  is a finite near hexagon of order  $(s, t)$ ,  $s \geq 2$ , containing a big quad of order  $(s, \alpha)$ .

**Lemma 5.** *The total number  $v$  of points of  $\mathcal{S}$  is equal to  $(s + 1)(s\alpha + 1)(s(t - \alpha) + 1) = (s + 1)(s^2\alpha(t - \alpha) + st + 1)$ .*

*Proof.* Let  $Q$  be a big quad of order  $(s, \alpha)$ . The total number of points of  $Q$  is equal to  $(s + 1)(s\alpha + 1)$ . Every point  $x$  of  $Q$  is collinear with  $s(t - \alpha)$  points outside  $Q$ , namely  $s$  on each of the  $t - \alpha$  lines through  $x$  not contained in  $Q$ . On the other hand, each point  $y$  outside  $Q$  is collinear with a unique point of  $Q$ , namely  $\pi_Q(y)$ . So,  $|\mathcal{P} \setminus Q| = |Q| \cdot s(t - \alpha)$  and  $|\mathcal{P}| = |Q| + |Q| \cdot s(t - \alpha) = (s + 1)(s\alpha + 1)(1 + s(t - \alpha))$ .  $\square$

**Lemma 6.** *Every quad of order  $(s, \alpha)$  or  $(s, t - \alpha)$  is big. There are no quads of order  $(s, \beta)$  with  $\min(\alpha, t - \alpha) < \beta < \max(\alpha, t - \alpha)$ .*

*Proof.* Suppose  $R$  is a quad of order  $(s, \beta)$  with  $\min(\alpha, t - \alpha) \leq \beta \leq \max(\alpha, t - \alpha)$ . With a completely similar reasoning as in Lemma 5, we find that the number of points at distance at most 1 from  $R$  is equal to  $(s + 1)(s\beta + 1)(1 + s(t - \beta)) = (s + 1)(s^2\beta(t - \beta) + st + 1)$ . This number is at least  $(s + 1)(s^2\alpha(t - \alpha) + st + 1) = v$  with equality if and only if  $\beta$  is equal to either  $\alpha$  or  $t - \alpha$ . The conclusions of the lemma follow.  $\square$

**Lemma 7.** *If  $x$  is a point of a big quad  $Q$ , then the local space at the point  $x$  is linear.*

*Proof.* We need to prove that any two distinct lines  $K$  and  $L$  through  $x$  are contained in a (necessarily unique) quad. This is certainly the case if  $K$  and  $L$  are contained in  $Q$ . So, without loss of generality we may suppose that  $L$  is not contained in  $Q$ . Let  $\mathcal{R}'$  denote the set of quads through  $L$  and take a point  $y$  on  $L$  distinct from  $x$ . The number of lines through  $x$  distinct from  $L$  and contained in a quad together with  $L$  is equal to  $\sum_{R \in \mathcal{R}'} t_R$ . This number also equals the number of lines through  $y$  distinct from  $L$  and contained in a quad together with  $L$ . So, it suffices to prove that if  $L'$  is a line through  $y$  distinct from  $L$ , then there is a (necessarily unique) quad through  $L$  and  $L'$ . Put  $L'' := \pi_Q(L')$ . Take  $z \in L' \setminus \{y\}$  and put  $u := \pi_Q(z)$ . The points  $u$  and  $y$  have at least two common neighbours, namely  $z$  and  $x$ . As  $s \geq 2$ , Proposition 1 implies that  $u$  and  $y$  are contained in a unique quad. This unique quad also contains the lines  $L$  and  $L'$ .  $\square$

Using Lemmas 5, 6 and 7, we can now prove Theorem 1. So, suppose that  $\mathcal{S}$  satisfies the conditions of that theorem. By way of contradiction, suppose that  $\mathcal{S}$  is not dense.

Let  $Q$  be a big quad of order  $(s, \alpha)$ . Since  $\mathcal{S}$  is not dense, there exists a point  $x$  whose local space is not linear. By Lemmas 6 and 7,  $x \notin Q$  and  $x$  is not contained in quads of order  $(s, \alpha)$ . Let  $L$  denote the unique line through  $x$  meeting  $Q$  in the point  $y := \pi_Q(x)$ . By Lemma 7, the local space at  $y$  is linear and so there exists a quad through the line  $L$ . Since the local space at  $x \in L$  is not linear, this quad is not big by Lemma 7. If  $(s, \beta)$  is the order of this quad, then from Lemma 6 we know that  $\beta$  cannot be contained in the interval  $[\min(\alpha, t - \alpha), \max(\alpha, t - \alpha)]$ . This implies that  $\alpha \neq 1$ .

Let  $\mathcal{R}$  denote the set of all quads through  $x$  and let  $\mathcal{R}'$  denote the set of all quads through the line  $L$ . Since the local space at  $y$  is linear, every two distinct lines through  $y$  are contained in a unique quad. Since  $Q$  is big, every quad through  $L$  meets  $Q$  in a line and so there are precisely  $t_Q + 1 = \alpha + 1$  of them. We thus have  $\sum_{R \in \mathcal{R}'} t_R = t$  with  $|\mathcal{R}'| = t_Q + 1$ . Since  $t_R, R \in \mathcal{R}'$ , is odd and  $|\mathcal{R}'| = t_Q + 1$  is even, we thus have that  $\sum_{R \in \mathcal{R}'} t_R = t$  is even.

Suppose  $R$  is a quad of order  $(s, \beta)$ ,  $\beta \geq 2$ , through  $x$  that is disjoint from  $Q$ , i.e.  $R$  does not contain the line  $L$ . Since  $x$  is not contained in quads of order  $(s, \alpha)$ , we have  $\beta \neq \alpha$ . Since  $\pi_Q(R)$  is a subquadrangle of  $Q$  isomorphic to  $R$ , we thus have that  $\beta < \alpha$  and  $\pi_Q(R)$  is a proper subquadrangle of  $Q$ . There thus exists a point  $z \in Q \setminus \pi_Q(R)$ . This point lies at distance 2 from  $R$ . Any line through  $z$  containing a point of  $\Gamma_1(R)$  defines a rosette of ovoids of  $R$ . So, we see that  $R$  cannot have a partition in ovoids. This implies that there are no lines through  $z$  contained in  $\Gamma_2(R)$ , i.e. every line through  $z$  contains a unique point of  $\Gamma_1(R)$ . Now,  $\Gamma_2(z) \cap R$  is an ovoid of  $R$  and since the local space in  $z \in Q$  is linear (see Lemma 7), there exists a unique quad  $Q(z, u)$  through  $z$  and any point  $u$  of  $\Gamma_2(z) \cap R$ . We denote by  $\mathcal{U}$  the set of all quads through  $z$  meeting  $R$  in a point of  $\Gamma_2(z) \cap R$ . The quads of  $\mathcal{U}$  determine a partition of the set of all lines through  $z$  and so we have  $t + 1 = \sum_{U \in \mathcal{U}} (t_U + 1)$ . As  $t_U + 1$  is even for every  $U \in \mathcal{U}$ , we have that also  $t + 1$  is even. But this is impossible since we have already shown above that  $t$  is even.

So, every quad through  $x$  not containing  $L$  is a grid. The total number of such grid-quads is at most

$$\frac{t(t-1)}{2} - \sum_{R \in \mathcal{R}'} \frac{t_R(t_R-1)}{2}.$$

Indeed, there is at most one such quad through every two lines through  $x$  distinct from  $L$ . Moreover, for each  $R \in \mathcal{R}'$ , there are  $\frac{t_R(t_R-1)}{2}$  pairs  $\{L_1, L_2\}$  of distinct lines through  $x$  such that  $L \neq L_1, L_2 \subseteq R$ , and for none of these pairs  $\{L_1, L_2\}$ , the unique quad containing  $L_1$  and  $L_2$  (namely  $R$ ) is a grid. So,

$$\sum_{R \in \mathcal{R}} t_R^2 \leq \frac{t(t-1)}{2} - \sum_{R \in \mathcal{R}'} \frac{t_R(t_R-1)}{2} + \sum_{R \in \mathcal{R}'} t_R^2.$$

Since  $\sum_{R \in \mathcal{R}'} t_R = t$ , we have

$$\sum_{R \in \mathcal{R}} t_R^2 \leq \frac{1}{2} \left( t^2 + \sum_{R \in \mathcal{R}'} t_R^2 \right). \quad (1)$$

Since

- $t_R \in \mathbb{N} \setminus \{0\}$  for every  $R \in \mathcal{R}'$ ,
- $\sum_{R \in \mathcal{R}'} t_R = t$  and  $|\mathcal{R}'| = \alpha + 1$ ,
- $a^2 + b^2 < (a - 1)^2 + (b + 1)^2$  for all  $a, b \in \mathbb{N}$  with  $2 \leq a \leq b$ ,

we have

$$\sum_{R \in \mathcal{R}'} t_R^2 \leq \alpha \cdot 1^2 + 1 \cdot (t - \alpha)^2. \quad (2)$$

By (1), (2) and the fact that  $\alpha > 1$ , we have that

$$\sum_{R \in \mathcal{R}} t_R^2 \leq \frac{\alpha + \alpha^2 - 2\alpha t + t^2 + t^2}{2} < \alpha^2 - \alpha t + t^2. \quad (3)$$

On the other hand, by Lemmas 1 and 2, we know that  $\sum_{R \in \mathcal{R}} t_R^2 = t(t+1) - \frac{n_2}{s^2}$ , where  $n_2 = \frac{v}{s+1} - 1 + s^2 t - st$ . Since  $v = (s+1)(s^2 \alpha(t - \alpha) + st + 1)$ , we have  $n_2 = s^2(t + \alpha t - \alpha^2)$  and

$$\sum_{R \in \mathcal{R}} t_R^2 = t(t+1) - \frac{n_2}{s^2} = t^2 - \alpha t + \alpha^2. \quad (4)$$

A contradiction follows from (3) and (4).

## Acknowledgment

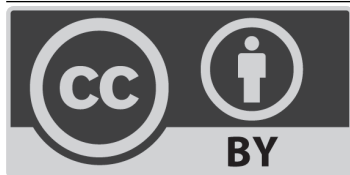
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## Conflict of Interest

The authors declare no conflict of interest.

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