



Article

## Rainbow Vertex Connection Numbers and Total Rainbow Connection Numbers of Middle and Total Graphs

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**Abstract:** A vertex-colouring of a graph  $\Gamma$  is rainbow vertex connected if every pair of vertices  $(u, v)$  in  $\Gamma$  there is a  $u-v$  path whose internal vertices have different colours. The rainbow vertex connection number of a graph  $\Gamma$ , is the minimum number of colours needed to make  $\Gamma$  rainbow vertex connected, denoted by  $rvc(\Gamma)$ . Here, we study the rainbow vertex connection numbers of middle and total graphs. A total-colouring of a graph  $\Gamma$  is total rainbow connected if every pair of vertices  $(u, v)$  in  $\Gamma$  there is a  $u-v$  path whose edges and internal vertices have different colours. The total rainbow connection number of  $\Gamma$ , is the minimum number of colours required to colour the edges and vertices of  $\Gamma$  in order to make  $\Gamma$  total rainbow connected, denoted by  $trc(\Gamma)$ . In this paper, we also research the total rainbow connection numbers of middle and total graphs.

**Keywords:** Rainbow vertex connection number; total rainbow connection number; middle graph; total graph

**Mathematics Subject Classification:** 05C15, 05C35, 05C40.

### 1. Introduction

We consider finite and simple graphs only. That is, we do not allow the existence of loops and multiple edges. We follow the terminology and notation of [1] for those not described here. A graph  $\Gamma$  is an ordered pair  $(V(\Gamma), E(\Gamma))$  consisting of a set  $V(\Gamma)$  of vertices and a set  $E(\Gamma)$  of edges. Let  $P_s, C_s, K_s$  and  $K_{1,s-1}$ , denote the path, cycle, complete graph, and star graph with  $s$  vertices, respectively.

An *edge-colouring* of a connected graph  $\Gamma=(V(\Gamma), E(\Gamma))$  is a mapping  $c : E \rightarrow S$ , where  $S$  is a set of colours. Usually, the set  $S$  of colours is taken to be  $\{1, 2, \dots, t\}$ ,  $t \in \mathbb{N}$ . A  $u-v$  path  $P$  in an edge-coloured graph is defined as a *rainbow path* if it does not exist two edges on this path coloured the same. An edge-coloured graph  $\Gamma$  is *rainbow connected* if any two vertices of  $\Gamma$  are connected by a rainbow path. The *rainbow connection number* of a connected graph, denoted by  $rc(\Gamma)$ , is the minimum  $t$  for which are needed to make the graph rainbow connected. The concept of rainbow connection number was first introduced and researched by Chartrand et al. in [2].

Likewise the concept of rainbow connection number, Krivelevich and Yuster [3] put forward the concept of rainbow vertex connection number. A *vertex-colouring* of a connected graph  $\Gamma = (V(\Gamma), E(\Gamma))$  is a mapping  $c : V \rightarrow S$ , where  $S$  is a set of colours. Usually, the set  $S$  of colours is taken to be  $\{1, 2, \dots, t\}$ ,  $t \in \mathbb{N}$ . A path in a vertex-coloured graph is defined as a *vertex rainbow*

path if its internal vertices are assigned distinct colours. A vertex-coloured graph  $\Gamma$  is *rainbow vertex connected* if any two vertices of  $\Gamma$  are connected by a vertex rainbow path, while the colouring is defined as *rainbow vertex-colouring*. The *rainbow vertex connection number* of  $\Gamma$ , is the minimum  $t$  for which are needed to make  $\Gamma$  rainbow vertex connected, denoted by  $\text{rvc}(\Gamma)$ . Let  $\Gamma$  be a connected graph. The graph  $\Gamma$  is complete, we have  $\text{rvc}(\Gamma) = 0$ , and  $\text{rvc}(\Gamma) \geq \text{diam}(\Gamma) - 1$  with equality if and only if the diameter of a graph is 1 or 2. Moreover, if  $\Sigma$  is a connected spanning subgraph of  $\Gamma$ (that is  $V(\Gamma) = V(\Sigma)$ ), then  $\text{rvc}(\Gamma) \leq \text{rvc}(\Sigma)$ .

The research of rainbow connection has attracted tremendous attention in the literature, and many conclusions have been published, see [3–11] for example. For more results, The reader can refer to the survey [12] and a new monograph [13].

As a natural generalization, Uchizawa et al. [14] and Liu et al. [15] introduced the concept of total rainbow connection number, respectively. A *total-colouring* of a graph  $\Gamma$  is a mapping  $c : V \cup E \rightarrow S$ , where  $S$  is a set of colours. Usually, the set  $S$  of colours is taken to be  $\{1, 2, \dots, t\}$ ,  $t \in \mathbb{N}$ . A  $u - v$  path  $P$  in a graph is defined as a *total rainbow path* if its internal vertices and edges are assigned distinct colours. A total-coloured graph  $\Gamma$  is *total rainbow connected* if any two vertices in  $\Gamma$  are connected by a total rainbow path, while the colouring is defined as *total rainbow colouring*. The *total rainbow connection number* of  $\Gamma$ , is the minimum  $t$  for which are needed to make  $\Gamma$  total rainbow connected, denoted by  $\text{trc}(\Gamma)$ . A simple observation is that the graph  $\Gamma$  is complete if and only if  $\text{trc}(\Gamma) = 1$ , otherwise  $\text{trc}(\Gamma) \geq 3$ . Moreover,  $\text{trc}(\Gamma) = 3$  if  $\text{diam}(\Gamma) = 2$  and  $\text{rc}(\Gamma) = 2$ . We also noticed some trivial fact that  $\text{trc}(\Gamma) \leq \text{trc}(\Sigma)$ , where  $\Sigma$  is a connected spanning subgraph of  $\Gamma$ . For more results about the function  $\text{trc}(\Gamma)$ , the reader can see [16–22] for details.

**Definition 1** ([23]). *The middle graph  $M(\Gamma)$  of a connected graph  $\Gamma$  is defined as follows,*

- (i)  $V(M(\Gamma)) = V(\Gamma) \cup E(\Gamma)$ .
- (ii) *Joining edges with those pairs of these vertices  $(x \in E(\Gamma), y \in E(\Gamma) ; x \in E(\Gamma), y \in V(\Gamma))$  which is adjacent (incident) in  $\Gamma$ .*

In [24], Li investigated the rainbow connection numbers of middle graphs of  $\Gamma$ , where  $\Gamma \cong P_s, C_s, K_{1,s}$  or  $K_s$ . Here, we research the total rainbow connection numbers (rainbow vertex connection numbers) for the above middle graphs. Our first main result is stated as follows.

**Theorem 1.** (i) *Let  $M(P_s)$  be the middle graph of  $P_s$ . Then  $\text{rvc}(M(P_s)) = s - 1$  and  $\text{trc}(M(P_s)) = 2s - 1$ .*  
 (ii) *Let  $M(C_s)$  be the middle graph of  $C_s$ . Then*

$$\text{rvc}(M(C_s)) = \begin{cases} \frac{s}{2} & \text{if } s \text{ is even;} \\ \frac{s+1}{2} & \text{if } s \text{ is odd.} \end{cases}$$

and

$$\text{trc}(M(C_s)) = \begin{cases} s + 1 & \text{if } s \text{ is even;} \\ s \text{ or } s + 1 & \text{if } s \text{ is odd.} \end{cases}$$

- (iii) *Let  $M(K_{1,s})$  be the middle graph of  $K_{1,s}$ . Then  $\text{rvc}(M(K_{1,s})) = s$  and  $\text{trc}(M(K_{1,s})) = 2s$ .*
- (iv) *Let  $M(K_s)$  be the middle graph of  $K_s$ . Then  $\text{rvc}(M(K_s)) = 1$  and  $3 \leq \text{trc}(M(K_s)) \leq s + 1$ .*

**Definition 2** ([25]). *The total graph  $T(\Gamma)$  of a graph  $\Gamma$  is defined as follows.*

- (i)  $V(T(\Gamma)) = V(\Gamma) \cup E(\Gamma)$ .
- (ii) *Joining edges with those pairs of these vertices  $(x \in E(\Gamma), y \in E(\Gamma) ; x \in E(\Gamma), y \in V(\Gamma) ; x \in V(\Gamma), y \in V(\Gamma))$  which is adjacent (incident) in  $\Gamma$ .*

Li [24] considered the rainbow connection numbers of total graphs of  $\Gamma$ , where  $\Gamma \cong P_s, C_s, K_{1,s}$  or  $K_s$ . Here, we also consider the total rainbow connection numbers (rainbow vertex connection numbers) for the above total graphs. Our second main result is stated as follows.

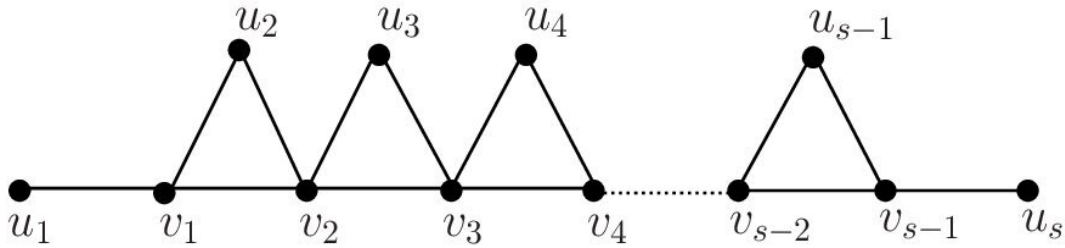


Figure 1. The middle graph of  $P_s$ .

**Theorem 2.** (i) Let  $T(P_s)$  be the total graph of  $P_s$ . Then  $rvc(T(P_s)) = s - 2$  and  $trc(T(P_s)) = 2s - 3$ .  
 (ii) Let  $T(C_s)$  be the total graph of  $C_s$ . Then

$$rvc(T(C_s)) = \begin{cases} \frac{s-2}{2} \text{ or } \frac{s}{2} & \text{if } s \text{ is even;} \\ \frac{s-1}{2} \text{ or } \frac{s+1}{2} & \text{if } s \text{ is odd.} \end{cases}$$

and

$$trc(T(C_s)) = \begin{cases} s - 1, s \text{ or } s + 1 & \text{if } s \text{ is even;} \\ s \text{ or } s + 1 & \text{if } s \text{ is odd.} \end{cases}$$

(iii) Let  $T(K_{1,s})$  be the total graph of  $K_{1,s}$ . Then  $rvc(T(K_{1,s})) = 1, trc(T(K_{1,2})) = 3, trc(T(K_{1,3})) = 4$  and  $trc(T(K_{1,s})) = 5$  with  $s \geq 4$ .

(iv) Let  $T(K_s)$  be the total graph of  $K_s$ . Then  $rvc(T(K_s)) = 1$  and  $3 \leq trc(T(K_s)) \leq s + 1$ .

## 2. Proof of Theorem 1

**Proposition 1.** Let  $M(P_s)$  be the middle graph of  $P_s$ . Then  $rvc(M(P_s)) = s - 1$  and  $trc(M(P_s)) = 2s - 1$ .

*Proof.* The graph  $M(P_s)$  is depicted in Figure 1. First we prove that  $rvc(M(P_s)) = s - 1$ . Let  $c$  be a vertex-colouring of  $M(P_s)$  defined as follows:  $c(u_1) = 1, c(v_i) = c(u_{i+1}) = i$  for  $1 \leq i \leq s - 2, c(v_{s-1}) = c(u_s) = s - 1$ . We will see that  $M(P_s)$  is rainbow vertex connected, and so  $rvc(M(P_s)) \leq s - 1$ . On the other hand, since  $\text{diam}(M(P_s)) = s$ , this implies  $rvc(M(P_s)) \geq s - 1$ , and so  $rvc(M(P_s)) = s - 1$ .

Now we prove that  $trc(M(P_s)) = 2s - 1$ . Let  $c$  be a total-colouring of  $M(P_s)$  defined as follows:  $c(u_1v_1) = 1, c(v_iu_{i+1}) = c(v_iv_{i+1}) = c(u_{i+1}v_{i+1}) = i + 1$  for  $1 \leq i \leq s - 2, c(v_{s-1}u_s) = s, c(u_1) = s + 1, c(v_i) = c(u_{i+1}) = s + i$  for  $1 \leq i \leq s - 2, c(v_{s-1}) = c(u_s) = 2s - 1$ . We will see that  $M(P_s)$  is total rainbow connected. Thus  $trc(M(P_s)) \leq 2s - 1$ . On the other hand, since  $\text{diam}(M(P_s)) = s$ , this implies  $trc(M(P_s)) \geq 2s - 1$ , which follows that  $trc(M(P_s)) = 2s - 1$ .  $\square$

**Proposition 2.** Let  $M(C_s)$  be the middle graph of  $C_s$ . Then

$$rvc(M(C_s)) = \begin{cases} \frac{s}{2} & \text{if } s \text{ is even;} \\ \frac{s+1}{2} & \text{if } s \text{ is odd.} \end{cases}$$

and

$$trc(M(C_s)) = \begin{cases} s + 1 & \text{if } s \text{ is even;} \\ s \text{ or } s + 1 & \text{if } s \text{ is odd.} \end{cases}$$

*Proof.* The graph  $M(C_s)$  is depicted in Figure 2. First we prove that

$$rvc(M(C_s)) = \begin{cases} \frac{s}{2} & \text{if } s \text{ is even;} \\ \frac{s+1}{2} & \text{if } s \text{ is odd.} \end{cases}$$

Suppose  $s$  is even with  $s = 2t$ . Since  $\text{diam}(M(C_s)) = t + 1$ , we have  $rvc(M(C_s)) \geq t$ . Let  $c$  be a vertex-colouring of  $M(C_s)$  defined as follows:  $c(u_i) = c(v_i) = i$  for  $1 \leq i \leq t, c(u_i) = c(v_i) = i - t$

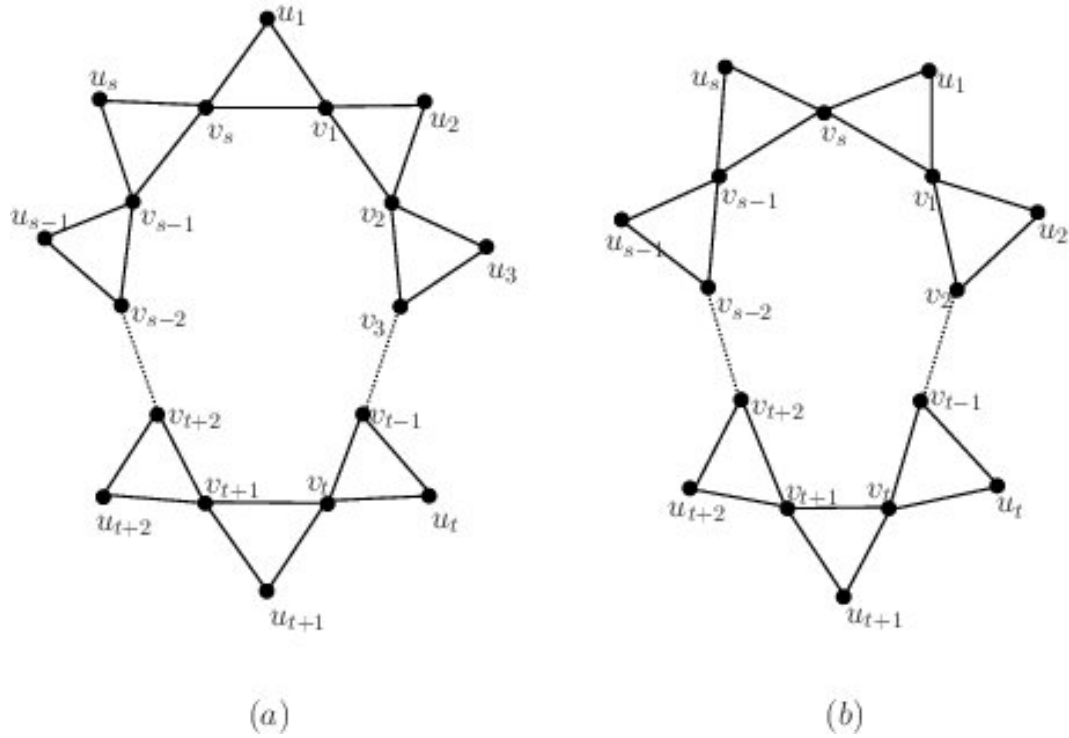


Figure 2. The middle graph of  $C_s$ .

for  $t + 1 \leq i \leq s$ . We know that  $M(C_s)$  is rainbow vertex connected, and so  $\text{rvc}(M(C_s)) \leq t$ . Thus  $\text{rvc}(M(C_s)) = t$ .

Suppose  $s$  is odd with  $s = 2t + 1$ . Since  $\text{diam}(M(C_s)) = t + 1$ , it follows that  $\text{rvc}(M(C_s)) \geq t$ . Now we show that  $\text{rvc}(M(C_s)) \neq t$ . To the contrary, assume that  $M(C_s)$  exists a rainbow vertex-colouring  $c$  with  $t$  colours. Considering  $u_1$  and  $u_{t+1}$ ,  $u_1v_1v_2 \cdots v_{t-1}v_tu_{t+1}$  must be a vertex rainbow  $u_1 - u_{t+1}$  path. Without loss of generality, let  $c(v_i) = i$  for  $1 \leq i \leq t$ . Considering  $u_2$  and  $u_{t+2}$ ,  $u_2v_2v_3 \cdots v_tv_{t+1}u_{t+2}$  must be a vertex rainbow  $u_2 - u_{t+2}$  path. Thus  $c(v_{t+1}) = 1$ . By the same steps, we know that  $c(v_{t+i}) = i$  for  $2 \leq i \leq t$ . Considering  $u_{t+2}$  and  $u_1$ ,  $u_{t+2}v_{t+2}v_{t+3} \cdots v_{s-1}v_su_1$  must be a vertex rainbow  $u_{t+2} - u_1$  path, and so  $c(v_s) = 1$ . But then, there does not exist a vertex rainbow path connecting  $u_{t+3}$  and  $u_2$ , a contradiction. Thus  $\text{rvc}(M(C_s)) \neq t$ . Let  $c$  be a vertex-colouring of  $M(C_s)$  defined as follows:  $c(v_i) = c(u_i) = i$  for  $1 \leq i \leq t$ ,  $c(v_i) = c(u_i) = i - t$  for  $t + 1 \leq i \leq 2t$ ,  $c(v_s) = c(u_s) = t + 1$ . We will see that  $M(C_s)$  is rainbow vertex connected, and so  $\text{rvc}(M(C_s)) = t + 1$ .

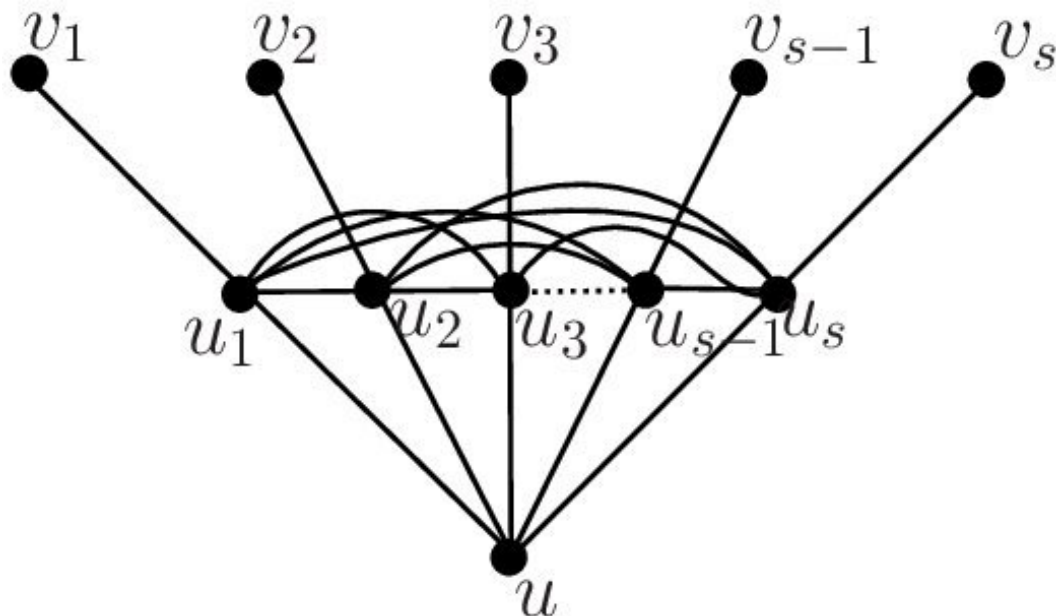
Now we prove that

$$\text{trc}(M(C_s)) = \begin{cases} s + 1 & \text{if } s \text{ is even;} \\ s \text{ or } s + 1 & \text{if } s \text{ is odd.} \end{cases}$$

Suppose  $s$  is even with  $s = 2t$ . Since  $\text{diam}(M(C_s)) = t + 1$ , we have  $\text{trc}(M(C_s)) \geq 2t + 1$ . Let  $c$  be a total-colouring of  $M(C_s)$  defined as follows:  $c(u_1v_1) = c(v_s v_1) = 1$ ,  $c(u_i v_i) = c(v_{i-1} v_i) = i$  for  $2 \leq i \leq t$ ,  $c(u_i v_i) = c(v_{i-1} v_i) = i - t$  for  $t + 1 \leq i \leq s$ , assign all other edges with  $t + 1$ ,  $c(v_i) = c(u_i) = t + i + 1$  for  $1 \leq i \leq t$ ,  $c(v_i) = c(u_i) = i + 1$  for  $t + 1 \leq i \leq s$ . We will see that  $M(C_s)$  is total rainbow connected, and so  $\text{trc}(M(C_s)) \leq s + 1$ . Thus  $\text{trc}(M(C_s)) = s + 1$ .

Suppose  $s$  is odd with  $s = 2t + 1$ . Let  $c$  be a total-colouring of  $M(C_s)$  defined as follows:  $c(u_1v_1) = c(v_s v_1) = 1$ ,  $c(u_1v_s) = t + 1$ ,  $c(v_i u_i) = c(v_{i-1} v_i) = i$  and  $c(v_{i-1} u_i) = t + 1$  for  $2 \leq i \leq t$ ,  $c(v_t v_{t+1}) = c(v_t u_{t+1}) = c(u_{t+1} v_{t+1}) = t + 1$ ,  $c(v_{t+i} u_{t+i+1}) = c(v_{t+i} v_{t+i+1}) = c(v_{t+i+1} u_{t+i+1}) = i$  for  $1 \leq i \leq t$ ,  $c(v_i) = c(u_i) = t + i + 1$  for  $1 \leq i \leq t$ ,  $c(v_i) = c(u_i) = i + 1$  for  $t + 1 \leq i \leq s$ . We obtained that  $M(C_s)$  is total rainbow connected. Since  $\text{diam}(M(C_s)) = t + 1$ , we have  $\text{trc}(M(C_s)) \geq 2t + 1 = s$ , which follows that  $s \leq \text{trc}(M(C_s)) \leq s + 1$ .  $\square$

**Proposition 3.** Let  $M(K_{1,s})$  be the middle graph of  $K_{1,s}$ . Then  $\text{rvc}(M(K_{1,s})) = s$  and  $\text{trc}(M(K_{1,s})) = 2s$ .



**Figure 3.** The middle graph of  $K_{1,s}$ .

*Proof.* The graph  $M(K_{1,s})$  is depicted in Figure 3. First we prove that  $\text{rvc}(M(K_{1,s})) = s$ . Since the cut vertices must be coloured by different colours, we have  $\text{rvc}(M(K_{1,s})) \geq s$ . Assign  $u_i$  with  $i$  for  $1 \leq i \leq s$ , and assign all other vertices with 1. We will see that  $M(K_{1,s})$  is rainbow vertex connected, and so  $\text{rvc}(M(P_s)) \leq s$ . This implies  $\text{rvc}(M(K_{1,s})) = s$ . Now we prove that  $\text{trc}(M(K_{1,s})) = 2s$ . Since the cut edges and cut vertices must be coloured by different colours, we obtain  $\text{trc}(M(K_{1,s})) \geq 2s$ . Let  $c$  be a total-colouring of  $M(K_{1,s})$  defined as follows:  $c(u_i v_i) = i$  for  $1 \leq i \leq s$ ,  $c(u u_1) = s$ ,  $c(u u_i) = i - 1$  for  $2 \leq i \leq s$ ,  $c(u_i u_{i+1}) = i + 2$  for  $1 \leq i \leq s - 2$ ,  $c(u_{s-1} u_s) = 1$ ,  $c(u_i u_j) = j - 1$  for  $1 \leq i, j \leq s$  and  $j - i \geq 2$ ,  $c(u_i) = s + i$  for  $1 \leq i \leq s$ , assign all other vertices with  $s + 1$ . Note that for any two different vertices  $v_i$  and  $v_j$  for  $1 \leq i, j \leq s$ , we know that  $v_i u_i u_j v_j$  is a total rainbow  $v_i - v_j$  path. Thus  $M(K_{1,s})$  is total rainbow connected, and hence  $\text{trc}(M(K_{1,s})) \leq 2s$ . Therefore,  $\text{trc}(M(K_{1,s})) = 2s$ .  $\square$

**Proposition 4.** Let  $M(K_s)$  be the middle graph of  $K_s$ . Then  $\text{rvc}(M(K_s)) = 1$  and  $3 \leq \text{trc}(M(K_s)) \leq s + 1$ .

*Proof.* Note that  $\text{diam}(M(K_s)) = 2$ , we have  $\text{rvc}(M(K_s)) = 1$ . Now we prove that  $3 \leq \text{trc}(M(K_s)) \leq s + 1$ . Obviously,  $\text{trc}(M(K_s)) \geq 3$  since  $\text{diam}(M(K_s)) = 2$ . The structure of  $M(K_s)$  is depicted as follows:  $M(K_s) = H_1 \cup H_2 \cup \dots \cup H_s$ , where  $H_i \cong K_s$  for  $1 \leq i \leq s$ , and for any  $i, j \in \{1, 2, \dots, s\}$ ,  $H_i$  and  $H_j$  only intersect a different vertex. Let  $c$  be a total-colouring of  $M(K_s)$  defined as follows: For any  $e \in H_i$ ,  $c(e) = i$ , assign  $s + 1$  to all vertices. We will see that  $M(K_s)$  is total rainbow connected, and so  $\text{trc}(M(K_s)) \leq s + 1$ . Hence  $3 \leq \text{trc}(M(K_s)) \leq s + 1$ .  $\square$

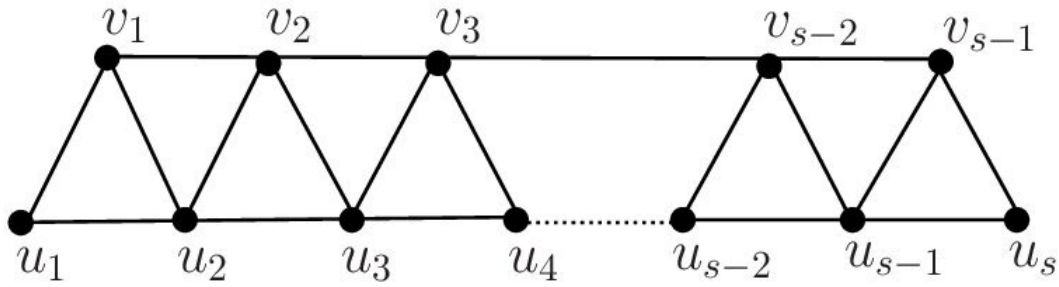
Combining Propositions 1, 2, 3, 4, Theorem 1 is immediate.

### 3. Proof of Theorem 2

**Proposition 5.** Let  $T(P_s)$  be the total graph of  $P_s$ . Then  $\text{rvc}(T(P_s)) = s - 2$  and  $\text{trc}(T(P_s)) = 2s - 3$ .

*Proof.* The graph  $T(P_s)$  is depicted in Figure 4. First we prove that  $\text{rvc}(T(P_s)) = s - 2$ . Since  $\text{diam}(T(P_s)) = s - 1$ , we have  $\text{rvc}(T(P_s)) \geq s - 2$ . Let  $c$  be a vertex-colouring of  $T(P_s)$  defined as follows:  $c(v_i) = c(u_{i+1}) = i$  for  $1 \leq i \leq s - 2$ , assign all other vertices with 1. We will see that  $T(P_s)$  is rainbow vertex connected, and so  $\text{rvc}(T(P_s)) \leq s - 2$ . Thus  $\text{rvc}(T(P_s)) = s - 2$ .

Now we prove that  $\text{trc}(T(P_s)) = 2s - 3$ . Let  $c$  be a total-colouring of  $T(P_s)$  defined as follows:  $c(u_i v_i) = c(u_i u_{i+1}) = c(v_i u_{i+1}) = i$  for  $1 \leq i \leq s - 1$ , assign all other edges with 1,  $c(v_i) = c(u_{i+1}) =$



**Figure 4.** The total graph of  $P_s$ .

$s + i - 1$  for  $1 \leq i \leq s - 2$ , all other vertices coloured with  $s$ . We will see that  $T(P_s)$  is total rainbow connected with the above total-colouring. Thus  $\text{trc}(T(P_s)) \leq 2s - 3$ . On the other hand, since  $\text{diam}(T(P_s)) = s - 1$ , this implies  $\text{trc}(T(P_s)) \geq 2s - 3$ . This follows that  $\text{trc}(T(P_s)) = 2s - 3$ .  $\square$

**Proposition 6.** Let  $T(C_s)$  be the total graph of  $C_s$ . Then

$$\text{rvc}(T(C_s)) = \begin{cases} \frac{s}{2} - 1 \text{ or } \frac{s}{2} & \text{if } s \text{ is even;} \\ \frac{s-1}{2} \text{ or } \frac{s+1}{2} & \text{if } s \text{ is odd.} \end{cases}$$

and

$$\text{trc}(T(C_s)) = \begin{cases} s - 1, s \text{ or } s + 1 & \text{if } s \text{ is even;} \\ s \text{ or } s + 1 & \text{if } s \text{ is odd.} \end{cases}$$

*Proof.* By [24], we know that

$$\text{diam}(T(C_s)) = \begin{cases} \frac{s}{2} & \text{if } s \text{ is even;} \\ \frac{s+1}{2} & \text{if } s \text{ is odd.} \end{cases}$$

On the other hand, note that  $M(C_s)$  is a connected spanning subgraph of  $T(C_s)$ , this proposition follows from Proposition 2.  $\square$

**Proposition 7.** Let  $T(K_{1,s})$  be the total graph of  $K_{1,s}$ . Then  $\text{rvc}(T(K_{1,s})) = 1, \text{trc}(T(K_{1,2})) = 3, \text{trc}(T(K_{1,3})) = 4$  and  $\text{trc}(T(K_{1,s})) = 5$  with  $s \geq 4$ .

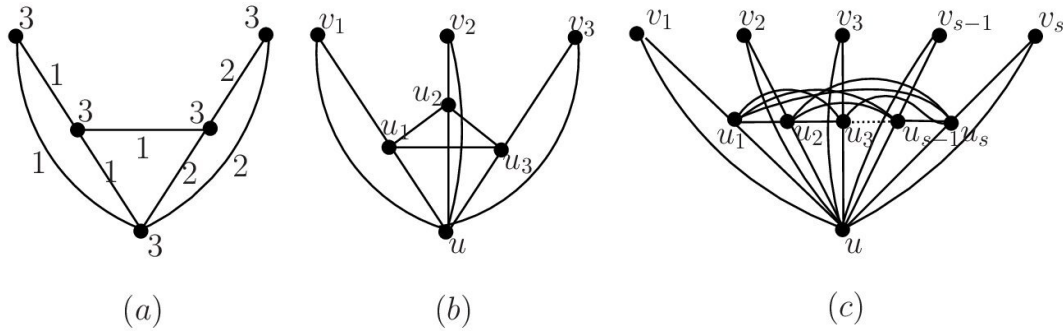
*Proof.* Since  $\text{diam}(T(K_{1,s})) = 2$ , we have  $\text{rvc}(T(K_{1,s})) = 1$  and  $\text{trc}(T(K_{1,s})) \geq 3$  for  $s \geq 2$ .

Suppose  $s = 2$ . We can easily verify that the total-colouring shown in Figure 5(a) is total rainbow. Thus  $\text{trc}(T(K_{1,2})) = 3$ .

Suppose  $s = 3$ . Let  $c$  be a total-colouring of  $T(K_{1,3})$  defined as follows:  $c(uv_1) = c(uu_1) = 1, c(uv_2) = c(uu_2) = 2, c(uv_3) = c(uu_3) = 3$ , assign all other edges with 1, and assign all vertices with 4. We can easily verify that  $T(K_{1,3})$  is total rainbow connected with the above total-colouring, and so  $\text{trc}(T(K_{1,3})) \leq 4$ . Now we only need to prove that  $\text{trc}(T(K_{1,3})) \neq 3$ . To the contrary, assume that  $T(K_{1,3})$  exists a total rainbow colouring with 3 colours. Considering  $v_1$  and  $v_2$ , the total rainbow  $v_1 - v_2$  path must be  $v_1uv_2$ . Without loss of generality, assume that  $c(v_1u) = 1, c(u) = 2, c(uv_2) = 3$ . Considering  $v_1$  and  $v_3$ , the total rainbow  $v_1 - v_3$  path must be  $v_1uv_3$ . Thus  $c(uv_3) = 3$ . But then, there is no total rainbow  $v_2 - v_3$  path, a contradiction. Hence  $\text{trc}(T(K_{1,3})) \neq 3$ , and so  $\text{trc}(T(K_{1,3})) = 4$ .

Suppose  $s \geq 4$ . Let  $c$  be a total-colouring of  $T(K_{1,s})$  defined as follows: assign 1 to the edges  $uv_i$  for  $1 \leq i \leq s$ , assign 2 to the edges  $u_i v_i$  for  $1 \leq i \leq s$ , assign 3 to all other edges, assign 4 to  $u$ , and assign 5 to all other vertices. We will see that  $T(K_{1,s})$  is total rainbow connected with the above total-colouring. Thus  $\text{trc}(T(K_{1,s})) \leq 5$ . Now we only need to prove that  $\text{trc}(T(K_{1,s})) \neq 4$ . To the contrary, assume that  $T(K_{1,s})$  exists a total rainbow colouring with 4 colours. Considering  $v_1$  and  $v_2$ ,  $v_1uv_2$  must be a total rainbow  $v_1 - v_2$  path. Without loss of generality, assume that  $c(u) = 1, c(v_1u) = 2, c(uv_2) = 3$ .





**Figure 5.** The total graph of  $K_{1,s}$ .

Considering  $v_1$  and  $v_3$ ,  $v_1uv_3$  must be a total rainbow  $v_1 - v_3$  path. Thus  $c(uv_3) = 4$ , otherwise there is no total rainbow  $v_2 - v_3$  path. Considering  $v_1$  and  $v_4$ ,  $v_1uv_4$  must be a total rainbow  $v_1 - v_4$  path. Hence  $c(uv_4) = 3$  or  $4$ . But then there is no total rainbow  $v_2 - v_4$  path or  $v_3 - v_4$  path. Hence  $\text{trc}(T(K_{1,s})) \neq 4$ , and so  $\text{trc}(T(K_{1,s})) = 5$ .  $\square$

**Proposition 8.** Let  $T(K_s)$  be the total graph of  $K_s$ . Then  $\text{rvc}(T(K_s)) = 1$  and  $3 \leq \text{trc}(T(K_s)) \leq s + 1$ .

*Proof.* Note that  $\text{diam}(T(K_s)) = 2$ . Then  $\text{rvc}(T(K_s)) = 1$ . Obviously,  $M(K_s)$  is a connected spanning subgraph of  $T(K_s)$ . By Proposition 4, we have  $3 \leq \text{trc}(T(K_s)) \leq s + 1$ .  $\square$

Combining Propositions 5, 6, 7, 8, Theorem 2 is immediate.

#### 4. Conclusion

The concept of total rainbow connection number was proposed in recent years. Moreover, Chen et al. [16] proved that the calculating of  $\text{trc}(\Gamma)$  is NP-hard. Subsequently, there is a great interest towards determining the total rainbow connection numbers of some graph classes. In this paper, we mainly consider the total rainbow connection numbers of middle and total graphs.

#### Conflict of Interest

The authors declare no conflict of interests.

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