

Generalized Index of Boolean Matrices

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ABSTRACT. In this paper we introduce a new parameter related to the index of convergence of Boolean matrix – the generalized index. The parameter is motivated by memoryless communication system. We obtain the values of this parameter for reducible, irreducible and symmetric matrices.

1 Introduction

The set, B_n of $n \times n$ Boolean matrices forms a finite multiplicative semigroup of order 2^{n^2} .

Note that we use Boolean arithmetic when calculating the powers of a matrix. Let $A \in B_n$. Since the sequence of powers $A^0 = I, A, A^2, \dots$ forms a finite subsemigroup $\langle A \rangle$ of B_n , then there exists a least nonnegative integer $k = k(A)$ such that $A^k = A^{k+t}$ for some $t > 0$, and there exists a

*This research was partially supported by National Natural Science Foundation of China and Guangdong Provincial Natural Science Foundation of China.

least positive integer $p = p(A)$ such that $A^k = A^{k+p}$. We call the integer $k = k(A)$ the index of convergence of A and the integer $p = p(A)$ the period of A .

A matrix $A \in B_n$ is reducible if there is an $n \times n$ permutation matrix P such that

$$PAP^{-1} = \begin{bmatrix} B & 0 \\ C & D \end{bmatrix},$$

where B and C are square (nonvacuous) matrices. The matrix A is irreducible if it is not reducible. Let IB_n and RB_n denote the set of irreducible and reducible matrices of order n , respectively.

If $A \in B_n$ ($n > 1$) is irreducible, then $p(A) = 1$ if and only iff A is primitive and in this case $k(A)$ is just the primitive exponent of A , i.e., the least positive integer k such that A^k is the matrix of all 1's. Let P_n denote the set of primitive matrices in B_n .

There is a natural 1 – 1 correspondence between the set of $n \times n$ Boolean matrices and the set of labeled digraphs of order n . We associate with the matrix $A = (a_{ij})$ the digraph $D = D(A)$ with vertex $V(D) = \{1, 2, \dots, n\}$ and arc set $\{(i, j) : a_{ij} \neq 0\}$. Thus the study of the properties of Boolean matrices can be turned into that of the corresponding associated digraphs. The index of convergence, period of A are called the index of convergence, period of $D(A)$, and is denoted by $k(D(A))$ and $p(D(A))$ equivalently. It is well known that if A is irreducible, then $p(A)$ is the greatest common divisor of the distinct lengths of the circuits of $D(A)$, and that if A is reducible, then $p(A)$ is the least multiple of $p(A_1), \dots, p(A_m)$, where A_1, \dots, A_m are the irreducible constituents of A . It is also well known that Boolean matrix A is irreducible if and only if $D(A)$ is strongly connected.

Let $A = (a_{ij}) \in B_n$. The (i, j) -entry of A^m is denoted by $a_{ij}^{(m)}$. Then $a_{ij}^{(m)} \neq 0$ if and only if there is a walk of length m from i to j in $D(A)$.

Recently, R.A. Brualdi and Bolian Liu [1] introduced generalized exponents of primitive matrices from memoryless communication system. Note that an adjacency matrix of a network need not be primitive. In this paper we introduce a new parameter that is a generalization of the exponent for the sequence of powers of Boolean matrices in [1].

Let D be a digraph with period $p = p(D)$ where $D = D(A)$ for some $A = (a_{ij}) \in B_n$. Define $k_D(i, j) := \min\{k : \text{for any integer } m \geq k, \text{ there exists a walk of length } m \text{ from } i \text{ to } j \text{ in } D \text{ if and only if there exists a walk of length } m + p \text{ from } i \text{ to } j \text{ in } D\}$.

Clearly, if A is primitive, then $k_D(i, j)$ is the local exponent in [1].

$k_D(i, j)$ is called the local index of convergence (or simply local index) from i to j . Clearly, $k_D(i, j)$ is the integer k such that the sequence $\{a_{ij}^{(m)} : m = 0, 1, 2, \dots\}$ is with period p from a beginning term $a_{ij}^{(k)}$, i.e.,

$k_D(i, j)$ is the least nonnegative integer k such that $a_{ij}^{(i)} = a_{ij}^{(t+p)}$ for all $t \geq k$. Clearly $k(D) = \max_{1 \leq i, j \leq n} k_D(i, j)$ (see [2]).

Let $D = D(A)$ where $A = (a_{ij}) \in B_n$. The least positive period of the sequence $\{a_{ij}^{(m)} : m = 0, 1, 2, \dots\}$ is called the local period from i to j of digraph D , denoted by p_{ij} . Clearly $p_{ij} | p(A)$.

For $i = 1, 2, \dots, n$, define

$$k_D(i) := \max_{j \in D} k_D(i, j);$$

$p_D(i) := \min\{p : \text{there exists a walk of length } m \text{ for every integer } m \geq k_D(i, j) \text{ from } i \text{ to } j \text{ in } D \text{ if and only if there exists a walk of length } m + p \text{ from } i \text{ to } j, j \in \{1, 2, \dots, n\} \text{ in } D\}$.

$k_D(i)$ and $p_D(i)$ are called the index and period of vertex i of D respectively.

Lemma 1. ([2] *Let*

$$a_0, a_1, a_2, \dots, a_m, \dots \tag{1}$$

be a sequence with period d (> 0) from a beginning term a_{m_d} , i.e., m_d is the least nonnegative integer such that $a_{t+d} = a_t$ for all $t \geq m_d$. If $d | p$, and (1) is also with period p from a beginning term a_{m_p} , then $m_d = m_p$.

Since $p_{ij} | p(A)$, $p_A(i) | p(A)$ holds. It follows from Lemma 1 that the sequence $\{a_{ij}^{(m)} : m = 0, 1, \dots\}$ is with period $p_A(i)$ and $p(A)$ from the same beginning term. Hence we use period $p(A)$ to define $k_A(i, j)$ and $k_A(i)$.

The numbers $k_D(i)$ have an interpretation in terms of a memoryless communication system associated with D . Suppose that at time $t = 1$ each vertex of D with some information passes the information to each of its neighbours (those vertices reachable by a walk of length 1) and then forgets its information. But it may receive information from another vertex. The system continues in this way. Vertex i always passes the information to the same vertices, the set of which is denoted by $N(i)$, every period of time after $t = t_0$. This time t_0 is $k_D(i)$. The period of time is $p_D(i)$. Clearly if $N(i) = V(D)$, then $k_D(i)$ is the generalized exponent in [1].

We choose to order the vertices of D in such a way that

$$k_D(1) \leq k_D(2) \leq \dots \leq k_D(n),$$

and call $k_D(i)$ the i th generalized index of D , denoted by $k(D, i)$. We write $k(A, i) = k(D(A), i)$ for any $A \in B_n$ with $1 \leq i \leq n$ and call $k(A, i)$ the i th generalized index of A . Clearly $k(A, i)$ is the smallest nonnegative integer k such that i rows of A^k and A^{k+p} are equal.

Thus for all $n \times n$ Boolean matrices A ,

$$k(A, n) = k(A) = k(D(A)) \leq (n - 1)^2 + 1. \quad (\text{see [1], [5]})$$

Define $k(n, i)$, $k^R(n, i)$, and $k^I(n, i)$, $k^S(n, i)$ to be the maximum of $k(A, i)$ where the maximum is taken over all Boolean matrices, all reducible Boolean matrices, all irreducible Boolean matrices and all symmetric Boolean matrices of order n , respectively.

The numbers $k(n, i)$, $k^R(n, i)$, $k^I(n, i)$, $k^S(n, i)$ are called generalized index of the corresponding classes of Boolean matrices of order n .

2 $k^I(n, i)$

We first establish the following.

Lemma 2. *There exists a matrix $\Gamma_n \in IB_n$ such that $k(\Gamma_n, i) = n^2 - 3n + 2 + i$ for $1 \leq i \leq n$.*

Proof: By Theorem 2.3 of [1], there exists $\Gamma_n \in P_n \subseteq IB_n$ such that $k(\Gamma_n, i) = k_{\Gamma_n}(i) = \exp_{\Gamma_n}(i) = n^2 - 3n + 2 + i$. Thus Lemma 2 follows \square

Lemma 3. ([2]) *Suppose that A is an irreducible Boolean matrices with period p , and the length of a shortest circuit of $D(A)$ is s , then*

$$k(A) \leq n + s\left(\frac{n}{p} - 2\right).$$

Theorem 1. $k^I(n, i) = n^2 - 3n + i + 2$, $1 \leq i \leq n$.

Proof: For any $A \in IB_n$, let $D(A)$ be the associated digraph of A . We consider the following two cases.

Case 1. $A \in P_n$. By Theorem 3.4 of [1],

$$k(A, i) = k(D(A), i) \leq n^2 - 3n + i + 2.$$

Case 2. $A \in IB_n \setminus P_n$. Then $p = p(A) \geq 2$ and $D(A)$ is strongly connected. Let s be the length of a shortest circuit of $D(A)$. By Lemma 3,

$$\begin{aligned} k(A, i) &\leq k(A, n) = k(A) \\ &\leq n + s\left(\frac{n}{p} - 2\right) \\ &\leq n + n\left(\frac{n}{2} - 2\right) \\ &\leq n^2 - 3n + 3 \\ &\leq n^2 - 3n + i + 2. \end{aligned}$$

Summarizing the above conclusions, we have $k(A, i) \leq n^2 - 3n + i + 2$ for any $A \in IB_n$, and by Lemma 2 there exists a $\Gamma_n \in IB_n$ such that $k(\Gamma_n, i) = n - 3n + i + 2$. Hence

$$k^I(n, i) = \max_{A \in IB_n} k(A, i) = n - 3n + i + 2.$$

\square

3 $k^R(n, i)$

Lemma 4. ([2]) *If $X \in B_n$ has the following form*

$$X = \begin{bmatrix} B & 0 \\ \alpha & a \end{bmatrix},$$

where B is an $(n-1) \times (n-1)$ Boolean matrix. Then

$$\begin{cases} k(B) \leq k(X) \leq k(B) + 1 & \text{if } a = 0, \\ k(B) \leq k(X) \leq \max\{k(B), n-1\} & \text{if } a = 1. \end{cases}$$

We now show the following.

Theorem 2. $k^R(n, i) = (n-3)(n-2) + i, 1 \leq i \leq n, n \geq 2$.

Proof: By Lemma 2, there exists a strongly connected digraph $G = (V, E)$ (of order $n-1$) with vertex set $V = \{1, 2, \dots, n-1\}$ such that $k(G, i) = (n-2)(n-3) + i$ for all $1 \leq i \leq n-1$. Let G' be (of order n) obtained from G by adding a new vertex named n and an arc $(n, n-1)$ to G . Then G' is not strongly connected, and it is easy to verify that

$$k(G', i) = \begin{cases} k(G, i) = (n-2)(n-3) + i & \text{if } 1 \leq i \leq n-1, \\ k(G, n-1) + 1 = (n-2)(n-3) + n & \text{if } i = n, \end{cases}$$

That is $k(G', i) = (n-2)(n-3) + i$ for all $1 \leq i \leq n$. Thus Theorem 2 follows immediately from the following

Theorem 2'. *If A is a reducible Boolean matrix of order n , then $k(A, i) \leq (n-2)(n-3) + i$ for all $1 \leq i \leq n$.*

Proof: We use induction on n to prove the theorem. Theorem 2' is true for all reducible Boolean matrices of order less than n and $p = p(A)$.

Claim 1. $k(B, i) \leq (|B|-1)(|B|-2) + i, 1 \leq i \leq |B|$, for any Boolean matrix B of order less than n . (Here we use $|B|$ to denote the order of B for convenience). To justify this claim, if B is irreducible, then Claim 1 follows from Theorem 1; otherwise if B is reducible, then Claim 1 follows from the induction hypothesis. (Theorem 2' holds for $|B| < n$).

The proof is now divided into the following three cases.

Case 1. There exists an $n \times n$ permutation matrix P such that $PAP^{-1} = \begin{bmatrix} B & 0 \\ \alpha & a \end{bmatrix}$, where B is an $(n-1) \times (n-1)$ matrix. Then, for all $t \geq 1$,

$$PA^tP^{-1} = \begin{bmatrix} B^t & 0 \\ \alpha B^{t-1} + a\alpha \sum_{i=0}^{t-2} B^i & a \end{bmatrix}.$$

Since $k(A, i)$ is the smallest nonnegative integer k such that i rows of A^k and A^{k+p} are equal, $k(A, i) \leq k(B, i) \leq (n-2)(n-3) + i$ for all $1 \leq i \leq n-1$. (The last inequality follows from Claim 1). Also by Lemma 4 and Claim 1, $k(A, n) \leq \max\{k(B, n-1) + 1, n-1\} \leq (n-2)(n-3) + n$.

Case 2. There exists an $n \times n$ permutation matrix P such that $PAP^{-1} = \begin{bmatrix} B & \beta \\ 0 & a \end{bmatrix}$, where B is an $(n-1) \times (n-1)$ matrix. Then, for all $t \geq 1$,

$$PA^tP^{-1} = \begin{bmatrix} B^t & (1-a)B^{t-1}\beta a(\sum_{i=0}^{t-1} B^i)\beta \\ 0 & a \end{bmatrix}.$$

Since the n th row of PA^tP^{-1} is independent of $t \geq 1$, we have $k(A, 1) \leq 1$. Recall that B is an $(n-1) \times (n-1)$ matrix. Thus, for all $2 \leq i \leq n$ and $t \geq \max\{k(B, i-1) + 1, n-1\}$, $\sum_{i=0}^{t-1} B^i = \sum_{i=0}^{n-2} B^i$ and so

$$PA^tP^{-1} = \begin{bmatrix} B^t & (1-a)B^{t-1}\beta + a(\sum_{i=0}^{n-2} B^i)\beta \\ 0 & a \end{bmatrix}.$$

Since $k(A, i)$ is the smallest nonnegative integer k such that i rows of A^k and A^{k+p} are equal, $k(A, i) \leq \max\{k(B, i-1) + 1, n-1\} \leq (n-2)(n-3) + i$ for all $2 \leq i \leq n$. (The last inequality follows from Claim 1).

Case 3. A does not satisfy the conditions in Cases 1 or 2. By [2, Lemma 7], $k(A, i) \leq k(A) \leq n^2 - 5n + 9 \leq (n-2)(n-3) + i$ for all $3 \leq i \leq n$. Now suppose $1 \leq i \leq 2$. Since A is reducible, there exists an $n \times n$ permutation matrix P such that $PAP^{-1} = \begin{bmatrix} C & 0 \\ E & D \end{bmatrix}$, where C and D are

square matrices with orders at most $n-2$. Then $n \geq 4$. By Claim 1, $k(C, i) \leq (|C| - 1)(|C| - 2) + i \leq (n-3)(n-4) + 2 \leq (n-2)(n-3) + i$. Again since $k(A, i)$ is the smallest nonnegative integer k such that i rows of A^k and A^{k+p} are equal, $k(A, i) \leq k(C, i) \leq (n-2)(n-3) + i$ for $1 \leq i \leq 2$.

Combining the above three cases, we complete the proof of Theorem 2' \square

It follows from Theorems 1 and 2 that

$$k(n, i) = \max\{k^I(n, i), k^R(n, i)\} = n^2 - 3n + 2 + i$$

for all $1 \leq i \leq n$.

4 $k^S(n, s)$

We turn to discuss symmetric irreducible matrices. Let ISB_n be the set of symmetric irreducible matrices and let

$$k^{IS}(n, i) := \max\{k(A, i) : A \in ISB_n\}.$$

Lemma 5. ([1]) *If A is an $n \times n$ symmetric primitive matrix, then $k(A, i) \leq n - 2 + i$ and there is an $n \times n$ symmetric primitive matrix A such that $k(A, i) = n - 2 + i$.*

By the way, we want to point out a minor error in Figure 2 of [1]. Its labeling of vertices $(n, n - 1, \dots, 1)$ should be $(1, 2, \dots, n)$.

Note that if $A \in ISB_n \setminus P_n$, then $p(A) = 2$, and the length of any circuit of $D(A)$ is even.

Lemma 6. ([4]) *Let $A \in ISB_n \setminus P_n$. Then $k(A) \leq n - 2$.*

By Lemmas 5 and 6, for any $A \in ISB_n$, we have $k(A, i) \leq n - 2 + i$ if A is primitive and $k(A, i) \leq k(A, n) = k(A) \leq n - 2 < n - 2 + i$ otherwise. And by Lemma 5 again, there is a matrix $A_0 \in ISB_n \cap P_n$ such that $k(A_0, i) = n - 2 + i$. Hence we have

Theorem 3. $k^{IS}(n, i) = n - 2 + i$.

Theorem 4. *For any $A \in ISB_n \setminus P_n$, $k(A, i) \leq \lceil \frac{n-1}{2} \rceil + i - 2$, where $\lceil x \rceil$ denotes the smallest integer greater than or equal to x .*

Proof: Let $A \in ISB_n \setminus P_n$. Then $p(A) = 2$. And let T be a spanning tree of $D(A)$. Then T has one or two centres (see [6] page 27), one of which is labeled u . Let $d = \max_{v \in V(D)} d(u, v)$, where $d(u, v)$ denotes the distance from vertex u to v in T , and $D = D(A)$.

Now we consider $k_D(u)$. Take any vertex $v \in V(D)$.

If there is a walk W of length $d - 1$ from u to v in D , then there is a walk of length $d - 1 + 2 = d + 1$ from u to v by attaching a circuit of length 2 to W .

On the other hand, if there is a walk of length $d + 1$ from u to v in D , since $d(u, v)$ and $d + 1$ have the same parity and $d(u, v) \leq d$, we have $d(u, v) \leq d - 1$. Assume $d - 1 = d(u, v) + 2b$ for some nonnegative integer b . By attaching circuits of length 2 to a path of length $d(u, v)$ from u to v , we obtain a walk of length $d(u, v) + 2b = d - 1$ from u to v .

Hence there is a walk of length $d - 1$ from u to v in D if and only if there is a walk of length $d - 1 + 2 = d + 1$ from u to v . By the arbitrary of v , we have $k(A, 1) \leq k_D(u) \leq d - 1 \leq \lceil \frac{1}{2}d(A) \rceil - 1$, where $d(A)$ denotes the diameter of $D(A)$.

And we show that $k(A, i) \leq k(A, i - 1) + 1$ ($2 \leq i \leq n$) as follows.

Suppose $k(A, j) = k_D(v_j)$ for $j = 1, 2, \dots, i - 1$. Since $A \in ISB_n$, D is strongly connected. There is an arc from some vertex $v \in V(D) \setminus \{v_1, v_2, \dots, v_{i-1}\}$ to some vertex $v_j \in \{v_1, v_2, \dots, v_{i-1}\}$. So $k_D(v) \leq k(A, i - 1) + 1$, which implies that $k(A, i) \leq k(A, i - 1) + 1$ for $2 \leq i \leq n$.

This follows that

$$\begin{aligned} k(A, i) &\leq k(A, 1) + i - 1 \\ &\leq \lceil \frac{1}{2}d(A) \rceil - 1 + i - 1 \\ &\leq \lceil \frac{1}{2}(n-1) \rceil + i - 2. \end{aligned}$$

□

Remark: We consider the digraph D_1 obtained from an indirected path of order n with any edge uv replaced by arcs (u, v) and (v, u) . The matrix A with associated digraph D_1 satisfies $d(A) = n - 1$, $\lceil \frac{1}{2}d(A) \rceil = \lceil \frac{n-1}{2} \rceil$, $k(A, 1) = \lceil \frac{1}{2}d(A) \rceil - 1 = \lceil \frac{n-1}{2} \rceil - 1$.

Finally, we show our main result in this section.

Theorem 5. $k^S(n, i) = n - 2 + i$.

Proof: Let $A \in B_n$ is symmetric. If $A \in ISB_n$, then by Theorem 3, we have $k(A, i) \leq n - 2 + i$. Suppose A is reducible. Denote by D_1, D_2, \dots, D_t ($t \geq 2$) all strong components of $D(A)$, the number of vertices are n_1, n_2, \dots, n_t , respectively. Then for every D_j ($1 \leq j \leq t$), there is no walk from any vertex of D_j to any vertex of another component. We choose to order the vertices of D_j ($1 \leq j \leq t$) such that

$$k_{D_j}(1^{(j)}) \leq k_{D_j}(2^{(j)}) \leq \dots \leq k_{D_j}(n_j^{(j)}).$$

By Theorem 3, we have $k_{D_j}(r^{(j)}) \leq n_j - 2 + r < n - 2 + r$ for $1 \leq j \leq t$, $1 \leq r \leq n_j$ with $r^{(j)} \in V(D_j)$, which implies that $k(A, i) = k(D(A), i) < n - 2 + i$ for $1 \leq i \leq n$.

Hence we have $k(A, i) \leq n - 2 + i$ no matter A is irreducible or not. By Lemma 5, there is a symmetric matrix A_0 in B_n such that $k(A_0, i) = n - 2 + i$. It follows that $k^S(n, i) = n - 2 + i$.

We complete the proof. □

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