

Spanning Trees Orthogonal to One-Factorizations of K_{2n}

John Krussel*
Department of Mathematical Sciences
Lewis and Clark College
Portland, Oregon 97219

Susan Marshall
Equipe Combinatoire,
Université de Paris VI
4, Place Jussieu
75252 Paris Cedex

Helen Verrall†
Department of Mathematics and Statistics
Simon Fraser University
Burnaby, British Columbia V5A 1S6

Abstract: In [3] Brualdi and Hollingsworth conjectured that for any one-factorization \mathcal{F} of K_{2n} there exists a decomposition of K_{2n} into spanning trees orthogonal to \mathcal{F} . They also showed that two such spanning trees always existed. We construct three such trees and exhibit an infinite class of complete graphs with an orthogonal decomposition into spanning trees with respect to the one-factorization GK_{2n} .

A *one-factorization* \mathcal{F} of the complete graph K_{2n} is a partition of $E(K_{2n})$ into spanning subgraphs $F_1, F_2, \dots, F_{2n-1}$ called *one-factors* so that each F_i is regular of degree one. See [4] for a survey on one-factorizations of the complete graph. A subgraph H of K_{2n} is said to be *orthogonal* to \mathcal{F} if $|E(H) \cap E(F_i)| = 1, 1 \leq i \leq 2n-1$. (If, instead of a one-factorization we consider an edge-coloring of K_{2n} so that each color class is a one-factor, then H is said to be *multicolored* as in [3].) A survey of results on orthogonal subgraphs can be found in [1].

Following the notation of [2] we will denote by GK_{2n} the “standard” one-factorization of K_{2n} constructed in the following manner. Label one vertex with ∞ and the rest of the vertices with $0, 1, \dots, 2n-2$. A one-factor consists of the edge $\{\infty, a\}$ (which we call the *infinite edge*) along with all (finite) edges of the form $\{x, y\}$ where $x + y \equiv 2a \pmod{2n-1}$. Call a the *center* of each of these edges

* This work was completed while on sabbatical at Simon Fraser University.

† This work was completed during a post-doctoral appointment supported by Simon Fraser University.

and define the *length* of an edge to be $\ell(x, y) = \min\{x - y \pmod{2n-1}, y - x \pmod{2n-1}\}$, where x and y are the finite labels on its endvertices.

In [3] it was shown that for any one-factorization \mathcal{F} of K_{2n} there exist two edge-disjoint spanning trees orthogonal to \mathcal{F} . We improve this result to three edge-disjoint spanning trees orthogonal to \mathcal{F} and construct them in the process.

THEOREM 1. *If \mathcal{F} is a one-factorization of K_{2n} ($n > 2$), then there exist three edge-disjoint spanning trees orthogonal to \mathcal{F} .*

Proof. First let $n = 3$. It is well-known that there is a unique one-factorization of K_6 . Without loss of generality, we can consider this factorization to be GK_6 . Then the three desired spanning trees are

$$\begin{aligned} & \{ \{\infty, 2\}, \{\infty, 3\}, \{\infty, 4\}, \{0, 2\}, \{1, 4\} \}; \\ & \{ \{\infty, 0\}, \{0, 1\}, \{0, 4\}, \{1, 2\}, \{3, 4\} \}; \\ & \{ \{\infty, 1\}, \{0, 3\}, \{1, 3\}, \{2, 3\}, \{2, 4\} \}. \end{aligned}$$

Now we assume that $n \geq 4$. Construct T as described in [3], i.e. consider the spanning star S_0 on vertex v_0 and delete edges $\{v_0, v_1\} \in F_1$ and $\{v_0, v_2\} \in F_2$, forming S'_0 . Then there is some edge $\{v_1, u_1\} \in F_2$ and an edge $\{v_2, u_2\} \in F_1$. (It is possible that $u_1 = u_2$.) Join these two to S'_0 to form S''_0 , still an orthogonal spanning tree. Now consider the spanning star S_1 on v_1 and note that $S_1 \cap S''_0 = \{v_1, u_1\}$. To form S'_1 we delete from S_1 the edge $\{v_1, u_1\}$ and some other edge $\{v_1, v_3\} \in F_3$, $v_3 \neq v_0, v_2, u_2$, chosen so that $\{v_0, u_1\} \notin F_3$. (Say $\{v_0, u_1\} \in F_6$.) This is always possible since there are $2n - 1$ one-factors and $n \geq 4$. Then we join to S'_1 the edge $\{u_1, w_1\} \in F_3$ and the edge $\{v_3, u_3\} \in F_2$ to form S''_1 , a second orthogonal spanning tree. Note that it is possible that $w_1 = u_3$ or one or both of w_1 and u_3 is the same as u_2 or $w_1 = v_2$ if $u_2 \neq u_1$.

Now we can construct a third orthogonal spanning tree in the same way. Notice that the only vertices which may have degree greater than 2 in $S''_0 \cup S''_1$ are $v_0, v_1, u_1, u_2, u_3, w_1$. Now v_3 can be none of these: $v_3 \neq v_0, u_2$ since it is chosen to be distinct from v_0, v_2, u_2 ; $v_3 \neq v_1, u_3$ since v_3 is adjacent to both; $v_3 \neq u_1$ since $\{u_1, v_1\}$ and $\{v_3, v_1\}$ are edges in distinct one-factors; and $v_3 \neq w_1$ since v_3 is incident with exactly one edge in F_3 but $\{u_1, w_1\}$ and $\{v_1, v_3\}$ are both in F_3 and $u_1 \neq v_1$. Therefore the degree of v_3 is 2. So now consider the spanning star S_3 on v_3 . Delete from S_3 the edge $\{v_3, u_3\} \in S''_1 \cap F_2$ and the edge $\{v_0, v_3\} \in S''_0$ which is also in, say, F_4 . We now have S'_3 . Now join to S'_3 the edges $\{v_0, v_2\} \in F_2$ and $\{u_3, v_4\} \in F_4$, forming S''_3 , a third orthogonal spanning tree.

Notice, however, that nothing precludes the edge $\{u_3, v_1\} \in S''_1$ from also being in F_4 (i.e. $v_1 = v_4$) in which case the construction of S''_3 is prevented. One possible way around this is to let u_3 play the role of v_3 and construct S'''_1

instead of S_1'' . Then $S_1''' = S_1'' \setminus \{\{v_1, u_3\}, \{u_1, w_1\}\} \cup \{\{v_1, v_3\}, \{u_1, v_6\}\}$, where $\{u_1, v_6\} \in F_4$. We know that v_6 can't be either v_0 or v_1 . Let S_4 be the star centered at u_3 . Then, $S_3''' = S_4 \setminus \{\{u_3, v_0\}, \{u_3, v_3\}\} \cup \{\{v_0, v_2\}, \{v_3, v_5\}\}$, where $\{u_3, v_0\}, \{v_3, v_5\} \in F_7$.

This gives the same problem as before if $v_5 = v_1$ and hence $F_7 = F_3$. However, in this case we must have the following sets of edges in the given one-factors: $\{u_3, v_3\} \in F_2, \{v_1, v_3\}, \{v_0, u_3\} \in F_3, \{v_1, u_3\}, \{v_0, v_3\} \in F_4$. Then we can modify the above construction, taking advantage of this symmetry. In constructing S_0'' we eliminate from S_0 not only $\{v_0, v_1\} \in F_1$ and $\{v_0, v_2\} \in F_2$ but also $\{v_0, v_3\} \in F_4$ and $\{v_0, u_3\} \in F_3$. These are replaced by $\{v_1, u_1\} \in F_2, \{v_2, u_2\} \in F_1, \{v_1, v_3\} \in F_3$ and $\{v_1, u_3\} \in F_4$.

Then the construction of S_1'' is modified by deleting from S_1 the edges $\{v_1, u_1\} \in F_2, \{v_1, v_3\} \in F_3$, and $\{v_1, u_3\} \in F_4$ as well as $\{v_1, u_4\} \in F_5$, where F_5 is chosen arbitrarily as long as it is distinct from F_1, \dots, F_4, F_6 . We replace these edges with $\{v_0, v_3\} \in F_4, \{v_3, u_3\} \in F_2, \{u_4, x_1\} \in F_3$, and $\{u_1, x_2\} \in F_5$. Finally, for S_3'' we delete from S_3 the edges $\{v_3, v_1\} \in F_3, \{v_3, u_3\} \in F_2, \{v_3, v_0\} \in F_4$, and $\{v_3, w_3\} \in F_5$. The construction is completed by replacing these with $\{v_0, v_2\} \in F_2, \{v_0, u_3\} \in F_3, \{v_1, u_4\} \in F_5$, and $\{w_3, y_3\} \in F_4$. ■

Note that the given spanning trees of K_6 can be constructed with the algorithm of the proof if $v_0 = \infty, v_1 = 0, v_2 = 1, u_1 = 2, u_2 = 4, v_3 = 3$, but it is easier to display them than to prove that the construction produces three trees for any v_0 .

Also in [3] it was conjectured that for any one-factorization \mathcal{F} of K_{2n} there exists a decomposition of the complete graph into spanning trees each orthogonal to \mathcal{F} . While this appears to be a difficult problem in general, we have found an infinite class of complete graphs with such an orthogonal decomposition with respect to GK_{2n} .

LEMMA 2: If $2n - 1$ is prime and there exists a spanning tree of K_{2n} containing exactly one edge of each length $1, 2, \dots, n - 1$ which is orthogonal to the one-factorization GK_{2n} , then there exists a decomposition of K_{2n} into spanning trees orthogonal to GK_{2n} .

Proof. Let E_l be the set of all edges in K_{2n} of length l and let $e_l \in E_l$. Since each one-factor of GK_{2n} contains exactly one edge of each length, a spanning tree consisting of $E_l - \{e_l\}$ along with the infinite edge to the center of e_l is orthogonal to GK_{2n} . Thus the spanning tree described in the statement of the lemma defines which edges e_l will be deleted from E_l for $l = 1, 2, \dots, n - 1$ to form the other spanning trees. ■

Without the spanning tree described in the statement of the lemma, if $2n - 1 = p$ is prime the proof of the lemma describes the construction of $n - 1$ spanning trees orthogonal to GK_{2n} , namely, $E_l - f_l$ together with the infinite edge centred at f_l , $1 \leq l \leq n - 1$, where the f_l are chosen so that no two have the same centre and f_l has length l . In general, the construction provides $\frac{\phi(2n-1)}{2}$ spanning trees orthogonal to GK_{2n} because E_l is a tree if and only if l and $2n - 1$ are relatively prime.

THEOREM 3: If $2n - 1 = p$ is a prime of the form $8m + 7$ then there exists an orthogonal decomposition of K_{2n} into spanning trees with respect to the one-factorization GK_{2n} .

Proof: We will construct a spanning tree of the type described in the lemma. To that end consider the group G whose elements are $\{1, 2, \dots, n - 1\}$ and whose operation is defined as $a \cdot b = \min\{ab, p - ab\} \pmod{p}$. If $a \cdot b = p - ab$ we will say that the result has been reduced by absolute value since the group operation is the same as reducing ab modulo p to one of the values $\{-(n - 1), \dots, 0, 1, \dots, n - 1\}$ and then taking the absolute value of the result. Now consider the set $H = \{2^j | j = 0, 1, \dots, k\}$, where the multiplication performed as in G , and k is the smallest value so that $2^{k+1} = 1$. Thus $2^k \equiv \pm \frac{p-1}{2} \pmod{p}$. For the following we calculate 2^j recursively, for $j \neq 0$, $2^j = 2 \cdot 2^{j-1}$. Then 2^j , $j \neq 0$, is an odd number in H iff 2^j has been reduced by absolute value. Thus if we were not taking absolute value, but leaving the numbers as positive or negative, each odd number in H would correspond to a change of sign. Hence $2^k \equiv -\frac{p-1}{2} \equiv \frac{p+1}{2}$ iff H contains an even number of odd numbers. Let $K = \{2^j | j = 0, 1, \dots, k\}$ where 2^j is reduced modulo p but not by absolute value.

We now construct the spanning tree needed for the application of Lemma 2. We begin with the spanning star centered at ∞ . From now on, whenever we add an edge of finite length, we also remove the infinite edge containing its center. So add all the edges with one endpoint 0 and the other endpoint in K , and remove the appropriate infinite edges. The center of each edge $\{0, 2^j\}$ is 2^{j-1} and is thus adjacent to 0 if $j \neq 0$. The center of the edge $\{0, 1\}$ is $\frac{p+1}{2}$ which is adjacent to 0 iff H contains an even number of odd numbers. Thus we still have a connected spanning graph iff H contains an even number of odd numbers.

If $H = G$ then the spanning tree is finished since there is an even number of odd numbers less than $\frac{p+1}{2}$ when $p = 8m + 7$. If $H \neq G$ then G can be partitioned by the cosets of H . Thus we continue the construction of the spanning tree by letting the elements of the corresponding cosets of K determine the endpoints of new edges incident with 0. Again, we remove the infinite edges containing the centers of these new edges and the centers determined by cK will also be in cK iff there is an even number of odd numbers in cH .

Thus it remains to show that H and cH contain an even number of odd numbers. We show first that the parity of the number of odds (evens) in H is the same as the parity of the number of odds (evens) in cH . We can assume that c is odd. Otherwise we could write $c = 2^d(2r + 1)$ and since $\pm 2^d \in H$, $cH = (2r + 1)H$. Now suppose $H = \{1, b_1, \dots, b_k\}$, where $b_i = 2^i$ reduced to an element of G , and consider $cH = \{c, cb_1, \dots, cb_k\}$ unreduced modulo p or by absolute value. Since c is odd, the parity of cb_i is the same as that of b_i . And when we reduce cb_i to an element of G , note that the parity of cb_i changes with each reduction by p as well as with reduction by absolute value. Thus we want to show that the total number of reductions in cH is even. We will do this by showing that cb_i is reduced an odd number of times by p iff cb_{i-1} is reduced by absolute value.

Suppose that $lp < cb_i < (l + 1)p$ so that cb_i must be reduced l times by p . Then consider cb_{i-1} and note that $b_i \equiv 2b_{i-1}$ or $b_i \equiv p - 2b_{i-1}$. If $b_i \equiv 2b_{i-1}$ then $cb_{i-1} = \frac{cb_i}{2}$ so $\frac{lp}{2} < cb_{i-1} < \frac{(l+1)p}{2}$. Thus cb_{i-1} is or is not reduced by absolute value as l is odd or even.

If $b_i = p - 2b_{i-1}$ then $lp < cp - 2cb_{i-1} < (l + 1)p$ or $\frac{c-1}{2}p > cb_{i-1} > \frac{c-(l+1)}{2}p$ and again cb_{i-1} is or is not reduced by absolute value as l is odd or even. Taking the same calculations backwards yields the converse. Thus cb_{i-1} is reduced by absolute value iff cb_i is reduced an odd number of times by p . Note that cb_k is never reduced by absolute value since $b_k = \frac{p-1}{2}$ so $cb_k = \frac{cp-c}{2}$ and $(\frac{c-1}{2})p < \frac{cp-c}{2} < \frac{cp}{2}$. Thus if there is an even number of odd values in H then the same will be true for all cosets of H and the spanning tree can be completed.

Now if $p = 8m + 7$ there are $2m + 2$ odd values and $2m + 1$ evens among the $4m + 3$ finite edge lengths. Since the parity of the number of odd values remains the same among the cosets of H , and there is an odd number of cosets of H , each coset must contain an even number of odd values. ■

Since by Dirichlet's Theorem there is an infinite number of primes of the form $8m + 7$ we have an infinite family of K_{2n} with an orthogonal decomposition. It should be noted that while it is true that the construction in the proof of Theorem 3 works for some prime numbers of the form $8m + 1$, 73 and 89 for instance, it does not work in general since it is possible to have an even number of cosets each containing an odd number of odd values. It is not clear whether it works for infinitely many primes of this form. And, of course, if $p = 8m + 3$ or $8m + 5$ there is an odd number of odd values so the construction cannot work.

REFERENCES

1. B. Alspach, K. Heinrich, and G. Liu, Orthogonal Factorizations of Graphs, in *Contemporary Design Theory: A Collections of Surveys*, J. H. Dinitz and D. R. Stinson, eds., John Wiley & Sons, 1992.

2. B. A. Anderson, Finite topologies and Hamiltonian paths, *J. Combin. Theory Ser. B* 14 (1973) 87-93.
3. R. A. Brualdi and S. Hollingsworth, Multicolored trees in complete graphs, *J. Combin. Theory Ser. B* 68 (1996) 310-313.
4. E. Mendelsohn and A. Rosa, One-factorizations of the complete graph—a survey, *J. Graph Theory* 9 (1985) 43-65.