# Realizability of p-Point, q-Line Graphs with Prescribed Maximum Degree, and Point Connectivity

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### Abstract

It is well known that some graph-theoretic extremal questions play a significant role in the investigation of communication network vulnerability. Answering questions concerning the realizability of graph invariants also solves several of these extremal problems. We define a  $(p, q, \kappa, \Delta)$  graph as a graph having p points, q lines, point connectivity  $\kappa$  and maximum degree  $\Delta$ . An arbitrary quadruple of integers (a, b, c, d) is called  $(p, q, \kappa, \Delta)$  realizable if there is a  $(p, q, \kappa, \Delta)$  graph with  $p = a, q = b, \kappa = c$  and  $\Delta = d$ . Necessary and sufficient conditions for a quadruple to be  $(p, q, \kappa, \Delta)$  realizable are derived. In earlier papers, Boesch and Suffel gave necessary and sufficient conditions for  $(p, q, \kappa)$ ,  $(p, q, \lambda)$ ,  $(p, q, \delta)$ ,  $(p, \Delta, \delta, \lambda)$  and  $(p, \Delta, \delta, \kappa)$  realizability, where  $\lambda$  denotes the line connectivity of a graph and  $\delta$  denotes the minimum degree for all points in a graph.

### Introduction

Here we consider an undirected graph G = (V, X) with a finite point set V and a set X whose elements, called lines, are two point subsets of V. The number of points is denoted by p, and the number of lines |X| is denoted by q(G) or q. This paper uses the notation and terminology of Harary [14]; however a few basic concepts are now reproduced.

The line connectivity of a graph G (denoted by  $\lambda(G)$  or  $\lambda$ ) is the minimum number of lines whose removal results in a disconnected graph. A graph is called trivial if it has just one point. The point connectivity (denoted by  $\kappa(G)$  or  $\kappa$ ) is the minimum number of points whose removal results in a disconnected or trivial graph. The number of lines connected to a point v of G is the degree of that point, denoted by  $d_v(G)$  or  $d_v$ . The minimum degree is denoted by  $\delta$  or  $\delta(G)$  and the maximum degree is denoted by  $\Delta$  or  $\Delta(G)$ . If  $\delta = \Delta$ , the graph is called regular. A p point graph with  $\delta = p-1$  is called complete and is denoted by  $K_p$ . A set of  $\kappa$  points whose removal disconnects G, or makes G trivial, is called a minimum point disconnecting set. The graph obtained from  $C_p$  ( the cycle on p points) by adding lines between all pairs of points that are distance at least two but not greater than A apart is denoted by  $C_p^A$ .

It is well known that some graph-theoretic extremal questions play a significant role in the investigation of communication network vulnerability [1-13]. Harary [15] found the maximum point connectivity among all graphs with a given number of points and a given number of lines. Answering questions concerning the realizability of graph invariants also solves several of these extremal problems. We define a  $(p, q, \kappa, \Delta)$  graph as a graph having p points, q lines, point connectivity  $\kappa$  and maximum degree  $\Delta$ . An arbitrary quadruple of integers (a, b, c, d) is called  $(p, q, \kappa, \Delta)$  realizable if there is a  $(p, q, \kappa, \Delta)$  graph with p = a, q = b,  $\kappa = c$  and  $\Delta = d$ . Necessary and sufficient conditions for a quadruple to be  $(p, q, \kappa, \Delta)$  realizable (or, more briefly, realizable) are derived. Boesch and Suffel derived necessary and sufficient conditions for  $(p, q, \kappa)$ ,  $(p, q, \lambda)$ ,  $(p, q, \delta)$ ,  $(p, \Delta, \delta, \lambda)$  and  $(p, \Delta, \delta, \kappa)$  realizability in earlier papers [6-8].

### **Preliminaries**

First we recall a result given by Harary [15].

**Lemma 1:** If  $2 \le \delta \le p-1$ , then there is a graph on p points with  $q = \lceil \frac{1}{2} p \delta \rceil$  and  $\lambda = \delta = \kappa$ . (This graph is a power of cycle and is usually called the Harary graph on p points).

The next lemma contains a complete list of the relevant lower bounds for q.

Lemma 2: The following holds;

- (1)  $\Delta \leq q$ ,
- (2) if  $\kappa = 1$ , then  $p 1 \le q$ ,
- (3) if  $\kappa = 2$ , then  $p + \Delta 2 \le q$ , and
- $(4) \lceil \frac{1}{2} ((p-1) \kappa + \Delta) \rceil \leq q.$

**Proof:** We pause to note that (1) and (4) are pertinent when  $\kappa=0$  and  $\kappa\geq 3$ , respectively. As (1) is obvious and (2) is well known we will now prove (3). Let G be a graph with  $\kappa=2$  and b be a point with  $d_b(G)=\Delta$ . Since  $G-\{b\}$  is connected; it follows that  $q(G-\{b\})\geq p-2$  and  $q(G)\geq p+\Delta-2$ , thus proving (3). Noting that (4) follows from  $\kappa\leq \delta$ , concludes this proof.

The remaining three lemmas find relevant upper bounds for q.

**Lemma 3:** (1) If 
$$\kappa = 1$$
 and  $p = 2 \Delta + 2$ , then  $q < \frac{1}{2} p \Delta$ .  
(2) For all graphs  $q \le \kappa + \frac{1}{2} (p - 1)(p - 2)$ .

**Proof:** First we will prove (1). Suppose G is a graph with  $\kappa=1$ , p=2  $\Delta+2$  and  $q=\frac{1}{2}$  p  $\Delta$ . As G is not complete there is a point c whose removal disconnects G into at least two components. Thus the point set of  $G-\{c\}$  may be partitioned into two sets T and U such that no lines of  $G-\{c\}$  join T and U. Since G is regular of degree  $\Delta$ , we have  $|T \cup \{c\}| \geq \Delta+1$  and  $|T| \geq \Delta$ . It is also true that  $|U| \geq \Delta$ . Because p=2  $\Delta+2$  either  $|T|=\Delta$  or  $|U|=\Delta$ . Suppose without loss of generality that  $|T|=\Delta$ . Therefore c is adjacent to every point in T and  $d_c(G) \geq \Delta+1$ . As this is impossible, no such G exists. Obviously  $q \leq \frac{1}{2}$  p  $\Delta$  and the result follows. The inequality in (2) was established for all graphs in [8].

**Lemma 4:** If 
$$p < 2 \Delta + 2 - \kappa$$
, then  $q \le \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor$ .

**Proof:** Let G be a graph with  $p < 2 \Delta + 2 - \kappa$  and S be a minimum point disconnecting set of G. The point set of G - S may be partitioned into two sets T and U such that no lines of G - S join T and U. Noting that either  $|T| \le \Delta - \kappa$  or  $|U| \le \Delta - \kappa$  we assume without loss of generality that  $|T| \le \Delta - \kappa$ . If |T| = N then

$$q \le \frac{1}{2} N (N-1) + \frac{1}{2} (p-N) \Delta + \frac{1}{2} N \kappa$$
.

Our goal is to maximize the right side of the inequality on  $1 \le N \le \Delta - \kappa$ . The fact that this quantity is a quadratic in N with a leading term of 1/2 N<sup>2</sup> tells us the maximum must take place at one of the bounds of the interval. The value of the

right side of the inequality at either bound is  $\frac{1}{2}(p-1)\Delta + \frac{1}{2}\kappa$  and the result follows.

**Lemma 5:** For all graphs either  $q \le \lfloor \frac{1}{2} (p-1) \Delta + \frac{1}{2} \kappa \rfloor$  or

$$q \leq \lfloor \frac{1}{2} p \Delta - \frac{1}{2} \max[(p - \Delta - 1)(2 \Delta + 2 - p) - \kappa(\kappa - 1), 0] \rfloor.$$

**Proof:** Since the result is obvious if  $(p - \Delta - 1)(2 \Delta + 2 - p) - \kappa (\kappa - 1) \le 0$ , henceforth we assume  $(p - \Delta - 1)(2 \Delta + 2 - p) - \kappa (\kappa - 1) > 0$ . We pause to note that  $p - \Delta - 1 \ge 0$  and  $\kappa(\kappa - 1) \ge 0$  implies  $p \le 2 \Delta + 1$ . It also follows that  $p \ge 3$  and  $0 < \Delta < p - 1$ . Our goal is to show that if  $q > \lfloor \frac{1}{2} (p - 1) \Delta + \frac{1}{2} \kappa \rfloor$  then

$$q \le \lfloor \frac{1}{2} p \Delta - \frac{1}{2} (p - \Delta - 1)(2 \Delta + 2 - p) + \frac{1}{2} \kappa (\kappa - 1) \rfloor$$
.

Let G be a graph with  $(p - \Delta - 1)(2 \Delta + 2 - p) - \kappa (\kappa - 1) > 0$  and  $q > \lfloor \frac{1}{2}(p-1)\Delta + \frac{1}{2}\kappa \rfloor$ . Let S, T and U be defined as they were in the proof of lemma 4. We wish to show that

$$2q = \sum_{i \in G} d_i \le p \Delta - (p - \Delta - 1) (2 \Delta + 2 - p) + \kappa(\kappa - 1).$$

Let V denote the union of T and U. As  $\sum_{i \in G} d_i = \sum_{i \in S} d_i + \sum_{i \in V} d_i$  we seek upper bounds for  $\sum_{i \in S} d_i$  and  $\sum_{i \in V} d_i$ . Note that  $\sum_{i \in S} d_i \le \kappa \Delta$ . To find an upper bound for  $\sum_{i \in V} d_i$  first find bounds for  $\|T\|$  and  $\|U\|$ . Suppose without loss of generality that  $\|T\| \le \|U\|$ . We also assume that  $\|T\| \le \Delta - \kappa$ . An argument similar to the one used to prove lemma 4 infers  $q \le \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor$ , contradicting our assumption that  $q > \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor$ . Therefore  $\|T\| \ge \Delta + 1 - \kappa$ ,  $\|U\| \ge \Delta + 1 - \kappa$  and  $\|T\| \le p - \Delta - 1$ . Letting  $\|T\| = N$  gives us

$$\sum_{i \in V} d_i \leq N(N-1) + (p-N-\kappa)(p-N-\kappa-1) + \kappa \Delta.$$

We want to maximize the right side of this inequality on  $\Delta + 1 - \kappa \le N \le p - \Delta - 1$ . This quantity is a quadratic in N with a leading term of  $2N^2$ , and thus is maximum at one of the bounds of the interval. Substituting for N shows this quadratic has the same value at each bound of the interval, giving us

$$\sum_{i \in V} d_i \leq \kappa \Delta + (\Delta + 1 - \kappa)(\Delta - \kappa) + (p - \Delta - 1)(p - \Delta - 2).$$

Consequently

$$\textstyle \sum_{i \in G} d_i = \sum_{i \in S} d_i + \sum_{i \in V} d_i \leq 2 \ \kappa \ \Delta + (\Delta + 1 - \kappa)(\Delta - \kappa) + (p - \Delta - 1)(p - \Delta - 2).$$

# Showing that

$$2\kappa\Delta + (\Delta + 1 - \kappa)(\Delta - \kappa) + (p - \Delta - 1)(p - \Delta - 2) \le p\Delta - (p - \Delta - 1)(2\Delta + 2 - p) + \kappa(\kappa - 1)$$

will complete this proof. In fact, simplifying both sides shows they are equal and we are done.

At this point we contemplate whether or not lemma 4 can be used to simplify the statement of lemma 5, and vice versa. To see this is not possible consider the following four cases, each of which satisfies

 $(p-\Delta-1)(2\ \Delta+2-p)$  -  $\kappa(\kappa-1)>0$ . The first two cases are  $p=11, \Delta=8, \kappa=4$  and  $p=6, \Delta=3, \kappa=1$ , each of which satisfies  $p<2\ \Delta+2-\kappa$ . The remaining two cases are  $p=8, \Delta=4, \kappa=2$  and  $p=16, \Delta=8, \kappa=4$ , each of which satisfies  $p\geq 2\ \Delta+2-\kappa$ . Substituting each of these four cases into the upper bounds for q given in lemmas 4 and 5 shows the two lemmas cannot be simplified.

# The $(p, q, \kappa, \Delta)$ realizability theorem

**Theorem.** A quadruple of non-negative integers  $(p, q, \kappa, \Delta)$  is realizable if and only if exactly one of the following conditions holds:

(I) 
$$0 = \kappa \le \Delta < p-1, \Delta \le q \le \lfloor \frac{1}{2} p \Delta \rfloor$$
 and if  $p \le 2 \Delta + 1$ , then  $q \le \lfloor \frac{1}{2} (p-1) \Delta \rfloor$ ;

- (II)  $1 = \kappa \le \Delta$ ,  $p 1 \le q$ , and if  $\Delta = 1$ , then p = 2:
  - (A)  $\Delta + 2 \le p \le 2 \Delta + 1$  and  $q \le \lfloor \frac{1}{2} ((p-1) \Delta + 1) \rfloor$ ;
  - (B)  $p = \Delta + 1$  and  $q \le 1 + \frac{1}{2}(p-1)(p-2)$ ;
  - (C)  $p = 2 \Delta + 2 \text{ and } q < \frac{1}{2} p \Delta;$
  - (D)  $p \ge 2 \Delta + 3$ ,  $q \le \lfloor \frac{1}{2} p \Delta \rfloor$  and if  $\Delta = 2$ , then q < p;

(III) 
$$2 \le \kappa = \Delta \le p - 1$$
,  $p \cdot \Delta$  is even and  $q = \frac{1}{2} p \Delta$ ;

- (IV)  $2 \le \kappa < \Delta$ ,  $\lceil \frac{1}{2} ((p-1)\kappa + \Delta) \rceil \le q$  and if  $\kappa = 2$ , then  $p + \Delta 2 \le q$ :
  - (A)  $\Delta + 1 and <math>q \le \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor$ ;
  - (B)  $p = \Delta + 1$  and  $q \le \kappa + \frac{1}{2}(p-1)(p-2)$ ;
  - (C)  $p \ge 2 \Delta + 2 \kappa$  and either  $q \le \lfloor \frac{1}{2} ((p-1)\Delta + \kappa) \rfloor$  or  $q \le \lfloor \frac{1}{2} p \Delta \frac{1}{2} \max [(p-\Delta-1)(2\Delta+2-p) \kappa(\kappa-1), 0] \rfloor;$

(V) 
$$p = 1$$
 and  $q = \kappa = \Delta = 0$ .

**Proof:** The necessity of  $\Delta < p-1$  in (I) follows from the fact that  $\Delta = p-1 \ge 1$  implies  $\kappa > 0$ . The other conditions in (I) are a consequence of lemmas 2, 4 and 5, and some obvious facts about graphs. We now consider (II). If a connected graph is regular of degree one, then it has exactly two points. Thus  $\kappa = \Delta = 1$  implies p = 2. Substituting  $\kappa = 1$  and p = 2  $\Delta + 1$  into lemma 5 gives us  $q \le \lfloor \frac{1}{2} ((p-1) \Delta + 1) \rfloor$ . The only connected regular graph of degree two is a cycle. Therefore  $\kappa = 1$  and  $\Delta = 2$  implies q < p. The other conditions in (II) follow from lemmas 2, 3, 4 and 5 together with well known facts concerning graphs. The conditions in (III) are a result of properties of a regular graph. The conditions in (IV) are a consequence of lemmas 2, 3, 4 and 5, and basic graph theory. The conditions in (V) are trivial.

We now provide constructions to prove sufficiency.

Case 1. Suppose that  $\lceil \frac{1}{2} ((p-1) \kappa + \Delta) \rceil \le q \le \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor$ ,  $3 \le \kappa < \Delta \le p-1$  and if  $\Delta = p-1$ , then  $q \le \kappa + \frac{1}{2} (p-1)(p-2)$ . Let  $H_1$  denote the Harary graph on p-1 points with  $\lceil \frac{1}{2} (p-1)(\kappa - 1) \rceil$  lines and  $\kappa(H_1) = \kappa - 1$ . Form a new graph by taking the union of a single point (which is denoted by b) and  $H_1$ . Note that  $H_1$  is not complete. If  $\kappa$  is odd we proceed as follows. Observing that  $H_1$  is a power of cycle which does not include diameters, we now add  $\lceil \frac{1}{2} (p-\Delta-1) \rceil$  independent diameters to  $H_1$ . At this point in our construction there are either  $\Delta$  or  $\Delta$  - 1 points in  $H_1$  that have degree  $\kappa$  - 1. Make b adjacent to all of the points of degree  $\kappa$  - 1 in  $H_1$ . If  $d_b = \Delta$  - 1, then make b adjacent to another point in  $H_1$ . Presently our graph contains  $\lceil \frac{1}{2} ((p-1) \kappa + \Delta) \rceil$  lines. If  $\Delta = p-2$  let c denote the point in  $H_1$  which is not adjacent to b, otherwise let c denote any point in  $H_1$ . To complete our construction we add  $q - \lceil \frac{1}{2} ((p-1) \kappa + \Delta) \rceil$  lines, none of which are incident to c, in such a way that no point has degree exceeding  $\Delta$ .

A minimum point disconnecting set of a power of cycle with no diameters is composed of two separate consecutive strings of points. As a result, it can easily be shown that our graph has the desired point connectivity. The graph also fulfills the other requirements of this case.

We now consider the case when  $\kappa$  is even. The construction is essentially the same as when  $\kappa$  was odd, except for the following differences. Here  $H_1$  includes diameters, however if p-1 is odd we may again add independent diameters to  $H_1$ . If p-1 is even we may add lines of the form ( i, i+p/2-1) to  $H_1$ , where i is an arbitrary point in  $H_1$ . Since  $\Delta \geq 5$  a sufficient number of lines can be added to  $H_1$ . It is easily verified that our graph has the desired properties. Therefore any quadruple satisfying this case is realizable.

Case 2. Suppose that  $\lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor + 1 \leq q \leq \lfloor \frac{1}{2} p \Delta \rfloor$ ,  $2 \leq \kappa < \Delta$  and  $p \geq 2 \Delta + 2$ . Let  $H_2$  denote the Harary graph on  $p - \Delta - 1$  points with  $\lceil \frac{1}{2} (p - \Delta - 1)(\Delta - 1) \rceil$  lines and  $\kappa(H_2) = \Delta - 1$ . Take the union of  $H_2$  and  $K_{\Delta + 1}$  to form a single graph. Observe that  $H_2$  is not complete. If  $\kappa$  is even we proceed

as follows. If  $p-\Delta-1$  is even and  $\Delta-1$  is odd add  $\lfloor \frac{1}{2}(p-\Delta-1-\kappa) \rfloor$  lines of the form (i,i+p/2-1) to  $H_2$ , where i is an arbitrary point in  $H_2$ . Otherwise add  $\lfloor \frac{1}{2}(p-\Delta-1-\kappa) \rfloor$  independent diameters to  $H_2$ . We pause to note that  $H_2$  has either  $\kappa$  or  $\kappa+1$  points of degree  $\Delta-1$ . Add  $\kappa$  independent lines joining points in  $K_{\Delta+1}$  to points of degree  $\Delta-1$  in  $H_2$ . Next, partition the points of degree  $\Delta+1$  in  $K_{\Delta+1}$  into pairs. For each of these pairs, delete the line which joins the points in that pair. Presently our graph contains  $\lfloor \frac{1}{2}p\Delta \rfloor$  lines. We finish our construction by deleting  $\lfloor \frac{1}{2}p\Delta \rfloor-q$  independent lines in  $K_{\Delta+1}$ , none of which are adjacent to lines previously deleted in  $K_{\Delta+1}$ . Since an independent set of lines was deleted from  $K_{\Delta+1}$  and  $\kappa < \Delta$ , our graph has the desired point connectivity. It can easily be proven that our graph has the other desired properties.

On the other hand, if  $\kappa$  is odd we alter the previous construction as follows. Instead of adding  $\kappa$  lines joining points in  $K_{\Delta+1}$  to points in  $H_2$ , add  $\kappa+1$  such lines in the following manner. Two lines are made incident to the same point in  $K_{\Delta+1}$  (denote this point by e) while the remaining  $\kappa-1$  lines are independent. Since one more line joins points in  $K_{\Delta+1}$  to points in  $H_2$ , we change the number of diameters we add to  $H_2$  accordingly. Next, delete lines in  $K_{\Delta+1}$  in such a way that all points in  $K_{\Delta+1}$  have degree  $\Delta$ . Finally, complete this construction by deleting  $\lfloor \frac{1}{2} p \Delta \rfloor - q$  lines in the manner of the previous construction. Note that the set of lines removed from  $K_{\Delta+1}$  is not independent. Let F denote one of the deleted lines which was incident to e. As F can be replaced by a path containing points in  $H_2$  and the remaining deleted lines form an independent set, our graph has point connectivity  $\kappa$ . The graph also fulfills the other requirements of this case.

Case 3. Suppose that  $\lfloor \frac{1}{2}((p-1)\Delta + \kappa) \rfloor + 1 \leq q \leq \lfloor \frac{1}{2}p\Delta - \frac{1}{2}\max[(p-\Delta-1)(2\Delta+2-p)-\kappa(\kappa-1),0] \rfloor$ ,  $2 \leq \kappa < \Delta$  and  $2\Delta+2-\kappa \leq p \leq 2\Delta+1$ . Since  $\lfloor \frac{1}{2}((p-1)\Delta+\kappa) \rfloor + 1 \leq q$  lemma 4 results in the above lower bound of p and lemma 5 yields the above upper bound of q. First, we consider the possibility that  $(p-\Delta-1)(2\Delta+2-p) \geq \kappa(\kappa-1)$ . If this is the case, let C denote the graph consisting of  $\kappa$  isolated points and form the union of C and  $K_{\Delta+1-\kappa}$ . Next, join every point of C to every point in  $K_{\Delta+1-\kappa}$ . We now form the union of this graph and  $K_{p-\Delta-1}$ . Denote the points in  $K_{\Delta+1-\kappa}$  by A and the points in  $K_{p-\Delta-1}$  by B. At this juncture all of the points of A have degree  $\Delta$ , all of the points of B have degree  $p-\Delta-2$  and all of the points of C have degree  $\Delta+1-\kappa$ .

We wish to prove that  $p \ge \kappa + \Delta$ . Assuming this inequality is false gives us  $p - \Delta - 1 < \kappa - 1$ . As a result of  $(p - \Delta - 1)$   $(2 \Delta + 2 - p) \ge \kappa(\kappa - 1)$  it follows that  $2 \Delta + 2 - p > \kappa$ , contradicting  $2 \Delta + 2 - \kappa \le p$ . Therefore  $p \ge \kappa + \Delta$  and B contains at least  $\kappa - 1$  points. If  $|B| > \kappa - 1$  join  $\kappa$  points of B in a one-to-one

fashion to the points of C. On the other hand, if  $|B| = \kappa - 1$  then join the points of B in a one-to-one fashion to  $\kappa - 1$  points of C. Note that  $|B| = \kappa - 1 > \kappa - 1 > \kappa - 1$  and  $|C| = \kappa$  Before lines ioining points of B to points

 $|B| = p - \Delta - 1 \ge \kappa - 1$  and  $|C| = \kappa$ . Before lines joining points of B to points of C were added, we had

$$\Delta(p-\Delta-1)-\sum_{i\in B}d_i=(p-\Delta-1)(\Delta-\Delta-1)$$

and  $\Delta \kappa - \sum_{i \in C} d_i = \kappa(\kappa - 1)$ . Consequently, we can now add lines joining points of B to points of C so that every point of C will have degree  $\Delta$  and no point of B will have degree exceeding  $\Delta$ . This is done in such a way that the degrees of the points of B differ by at most one. At this point, our graph contains

 $\left[ \begin{array}{cc} \sqrt{k} \ p \ \triangle - \sqrt{k} \ (p - \triangle - 1)(2 \ \triangle + 2 - p) + \sqrt{k} \ \kappa(\kappa - 1) \end{array} \right]$ 

independent lines in  $K_{\Delta+1-\kappa}$  finishes the construction. As a result of the lower bound of q, every point has degree of at least  $\kappa+1$ . It can easily be shown that the final graph has the desired properties.

We now consider the case when  $(p - \Delta - 1)(2 \ \Delta + 2 - p) < \kappa(\kappa - 1)$  and

 $p>\kappa+\Delta$ . First, repeat the previous construction excluding the step where lines were deleted. However, in this case, every point of B has degree  $\Delta$  and C

Were desired. However, in this case, every point of D has degree at least one point which has degree less than  $\Delta$ . Add the appropriate Harary graph to the points of C and/or join pairs of points of C, so that our graph has  $\lceil \sqrt{x} p \Delta \rceil$  lines. Deleting  $\lceil \sqrt{x} p \Delta \rceil$  - q independent lines in  $K_{\Delta+1-\kappa}$ 

completes this construction.

If  $(p - \Delta - 1)(2 \Delta + 2 - p) < \kappa(\kappa - 1)$  and  $p \le \kappa + \Delta$  then form the union of  $K_{\Delta - \kappa + 1}$  and the graph composed of  $p - \Delta - 1 + \kappa$  isolated points. Denote  $\kappa$  of the isolated points by C, the remaining  $p - \Delta - 1$  isolated points by B and the points of  $K_{\Delta - \kappa + 1}$  by A. Join each point of C to each point of A and B. We pause to note that every point of A has degree  $\kappa$  and every

that every point of A has degree  $\triangle$ , every point of B has degree A and every point of C has degree p -  $\kappa \le \Delta$ . Mext, add the appropriate Harary graph to the points of B and/or join pairs of points of B and, if necessary, do the same to the points of C. It may also be necessary to delete a line which joins a point of B to a point of C. Presently our graph contains  $\lceil \log p \Delta \rceil$  lines. Delete lines as in the previous construction and we are finished with case 3.

previous construction and we are finished with case 3. Case 4. Suppose that  $p + \Delta - 2 \le q \le \lfloor \sqrt{((p-1))\Delta + \kappa)} \rfloor$ ,  $2 = \kappa < \Delta \le p - 1$  and if  $\Delta = p - 1$ , then  $q \le \kappa + \sqrt{(p-1)(p-2)}$ . Consider the cycle on p points and let f and g denote two adjacent points of the cycle. Add  $q - p \ge \Delta - 2$  lines so that the following holds. Exactly  $\Delta - 2$  of these lines are incident to g, none of these lines are incident to f and no point has degree exceeding  $\Delta$ . This graph fulfills

the requirements of this case.

Case 5. Suppose that  $1 = \kappa \le \Delta$ ,  $p-1 \le q \le \lfloor \frac{1}{2} p \Delta \rfloor$ , and if  $\Delta = 1$ , then p=2. We also assume that the following conditions hold: (1) if  $\Delta + 2 \le p \le 2 \Delta + 1$ , then  $q \le \lfloor \frac{1}{2} ((p-1)\Delta + 1) \rfloor$ , (2) if  $p = \Delta + 1$ , then  $q \le 1 + \frac{1}{2} (p-1)(p-2)$ , and (3) if  $\Delta = 2$  or  $p = 2 \Delta + 2$ , then  $q < \frac{1}{2} p \Delta$ .

First let us consider when  $q \le \lfloor \frac{1}{2}((p-1)\Delta + 1) \rfloor$ . Let T be any tree on p points with  $\Delta(T) = \Delta$  and let h be a point of degree one in T. Adding q - p + 1 lines, none of which are incident to h, in such a way that all points in T have degree at most  $\Delta$  yields a graph satisfying our present assumptions.

Next we entertain the possibility that  $\lfloor \frac{1}{2} ((p-1)\Delta + 1) \rfloor + 1 \leq q$ ,  $p \geq 2\Delta + 2$ , and that one of the following holds;  $\Delta$  and  $p - \Delta - 1$  are both odd or  $q < \lfloor \frac{1}{2} p \Delta \rfloor$ . Note that  $\Delta \geq 2$ . Form the union of  $K_{\Delta+1}$  and a path containing  $p - \Delta - 1 \geq \Delta + 1$  points. Denote the points of  $K_{\Delta+1}$  by A and the points of the path by B. Join a point of B with degree one to one of the points of A. Delete a line of  $K_{\Delta+1}$  which is incident to the point of degree  $\Delta + 1$ . Observe that our graph now has  $\frac{1}{2}\Delta(\Delta+1) + p - \Delta - 2$  lines. To finish this construction, add  $q - \frac{1}{2}\Delta(\Delta+1) - p + \Delta + 2$  lines joining points of B, without generating any points of degree greater than  $\Delta$ . Notice that if  $q = \lfloor \frac{1}{2} p \Delta \rfloor$ , then  $\Delta$  and  $p - \Delta - 1$  are both odd and every point of B has degree  $\Delta$ .

Case 6. Suppose that  $0 = \kappa \le \Delta \le p-1$ ,  $\Delta \le q \le \lfloor \frac{1}{2}p\Delta \rfloor$  and if  $p \le 2\Delta + 1$ , then  $q \le \lfloor \frac{1}{2}(p-1)\Delta \rfloor$ . Furthermore, assume  $\Delta = p-1$  if and only if p = 1. If p = 1, then use the graph  $K_1$ . Henceforth, we consider only  $p \ge 2$ . If  $q \le \lfloor \frac{1}{2}(p-1)\Delta \rfloor$ , then we proceed as follows. Take p isolated points and denote one of them by p. Join p to p to p the points and let p denote one of the remaining isolated points. Adding p and p the points and let p denote one of the remaining isolated points. Adding p and p the points and let p denote the points are incident to p, so that no point has degree greater than p completes our construction.

However, if  $q \ge \lfloor \frac{1}{2} (p-1) \Delta \rfloor + 1$  form the union of  $K_{\Delta+1}$  and the graph composed of  $p - \Delta - 1 \ge \Delta + 1$  isolated points. We end this construction by adding  $q - \frac{1}{2} \Delta(\Delta + 1)$  lines, making sure that every point has degree at most  $\Delta$ . Case 7. Suppose that  $2 \le \kappa = \Delta \le p - 1$ ,  $q = \frac{1}{2} p \Delta$  and  $p \cdot \Delta$  is even. In this case the Harary graph on p points with  $\frac{1}{2} p \Delta$  lines and point connectivity  $\Delta$  is sufficient.

We are finished with our constructions and will now show sufficiency. If we assume  $2 \le \kappa < \Delta$  and  $q \le \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor$  then cases 1 and 4 show the sufficiency of the conditions of the theorem. Cases 2 and 3 show the sufficiency of the conditions of the theorem for  $2 \le \kappa < \Delta$  and  $q \ge \lfloor \frac{1}{2} ((p-1) \Delta + \kappa) \rfloor + 1$ . Cases 5,6 and 7 show the conditions of the theorem are sufficient if we also have  $\kappa = \Delta$  or  $\kappa < 2$  and our proof is finished.

## Conclusion

The  $(p, q, \kappa, \Delta)$  realizability theorem in this paper solves several extremal problems. If any three of the parameters p, q,  $\kappa$  and  $\Delta$  are given we can find the range of values for the unknown parameter. Here we will look at the problem of finding the maximum value of  $\kappa$  among all  $(p, q, \Delta)$  graphs, which is denoted by max  $(\kappa \mid p, q, \Delta)$ . The solution is given below (the proof is straightforward).

$$\max \left(\kappa \mid p,\, q,\, \Delta\right) = \begin{cases} 0 & \text{, if } q < p-1 \text{ or } p=1 \\ \\ \min(\left\lfloor (2q-\Delta)/(p-1)\right\rfloor,\, \Delta-1), \text{ if } q \geq p-1 \geq 1 \text{ and } p\Delta \text{ is odd} \end{cases}$$
 
$$\left\lfloor (2q-\Delta)/(p-1)\right\rfloor & \text{, otherwise.} \end{cases}$$

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