

Inducing regularization of graphs, multigraphs and pseudographs

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Abstract

For a given structure (graph, multigraph, or pseudograph) G and an integer $r \geq \Delta(G)$, a smallest inducing r -regularization of G (which is an r -regular superstructure of the smallest possible order, with bounded edge multiplicities, and containing G as an induced substructure) is constructed.

1 Introduction

Graphs are simple, multigraphs without loops, and pseudographs may contain both loops and multiple edges. For terminology and notation we refer to Chartrand and Lesniak [2].

For a given structure (graph, multigraph, or pseudograph) G and an integer $r \geq \Delta(G)$, an r -regular superstructure containing G as an induced substructure is called an *inducing r -regularization/regularization* of G . The problem of inducing r -regularization is originated in 1916 by König [5, 6]. Therein an r -regular multigraph containing a given multigraph G and its copy G' (if G is not r -regular), both as disjoint induced submultigraphs, is constructed so that additional (possibly multiple) edges join a vertex x of G to its copy x' in G' , r being any cardinal number, $r \geq \Delta(G)$. Optimal (i.e. the smallest) inducing regularization of a simple graph G within simple graphs with the smallest possible degree $\Delta(G)$ is published in two papers by Erdős and Kelly [3, 4]. Inducing r -regularization of a simple graph G within simple r -regular graphs with $r \geq \Delta(G)$ is achieved in Chartrand and Lesniak [2] by splitting König's method into $r - \delta(G)$ steps. Then only simple xx' edges are added at each step and the order of the resulting r -regular graph is $2^{r-\delta(G)}$ times as large as that of G .

We extend Erdős and Kelly's result both to pseudographs and multi-graphs with any multiplicity bound p , maintaining optimality for any upper bound r on the maximum degree Δ . We present examples to show that the upper bound we give on the number of necessary new vertices is sharp. If G is an n -vertex graph and regularization of G is also a graph then in Erdős and Kelly's case (with $r = \Delta(G)$) the optimal inducing r -regularization of G may require up to n new vertices, which is the case if $n \geq 4$ and $G = K_n - e$, see [3, 4]; whereas, if r is odd and $r > n$ then exactly $r + 1$ new vertices are necessary precisely for all graphs G with $\delta(G) = 0$ and even n , otherwise up to $\max\{r, n\}$ new vertices only.

2 Main result

For $X \in \{P, M\}$, let X -graph stand for P -graph (pseudograph) if $X = P$, otherwise for M -graph (multigraph) if $X = M$. Let G be an X -graph of order n with vertex set $V = \{v_1, v_2, \dots, v_n\}$. Let δ and Δ stand for the minimum and maximum degrees among vertices of G . Let r be any integer such that $r \geq \Delta$. Call $a_i := r - \deg_G(v_i)$ to be the r -deficiency of the vertex v_i . Then $r - \delta$ is the *maximum r -deficiency* among vertices in G . Let $\sigma = \sum_i a_i$ be the sum of the r -deficiencies. The number of edges joining two vertices v, u is called the *multiplicity* of the vertex-pair v, u . The end of a claim's proof will be marked by the symbol \square .

Theorem *Given an X -graph G of order n , with minimum and maximum degrees δ and Δ , let p and r be integers such that $r \geq \Delta$ and p is an upper bound on the maximum multiplicity in G . Let F denote any r -regular X -graph with maximum multiplicity at most p and containing an induced sub- X -graph isomorphic to G . The necessary and sufficient condition that $n + t$ be the minimal order possible for F is that t is the least nonnegative integer such that*

- (i) $tr \geq \sigma$;
- (ii) $pt \geq r - \delta$;
- (iii) $(t + n)r$ is an even integer;
- (iv) either (P) $pt^2 - (r - p)t + \sigma \geq 0$ if $X = P$ or
(M) $pt^2 - (r + p)t + \sigma \geq 0$ if $X = M$.

Moreover, $t \leq t_0$ where

$$t_0 := \begin{cases} \lceil r/p \rceil + 1 & \text{if both } r \text{ and } \lceil r/p \rceil + n \text{ are odd,} \\ & \lceil r/p \rceil > n, \text{ and } \delta < r + p - p\lceil r/p \rceil, \\ \max\{\lceil r/p \rceil, n\} & \text{otherwise.} \end{cases} \quad (1)$$

Proof. Necessity. Suppose the order $n + t$ of the super- X -graph F is minimal. Let G' be the sub- X -graph of F isomorphic to G and let H

be the sub- X -graph induced by the t vertices of F not in G' . Then in F there are σ edges between the sub- X -graphs G' and H . Since each of the t vertices of H is incident with at most r of these edges, (i) follows. The vertex of G' whose r -deficiency is $r - \delta$ has at least $\lceil (r - \delta)/p \rceil$ neighbors in H , whence (ii) holds. Clearly $(t + n)r$ is an even number so (iii) holds. The sum of the degrees in H of the vertices of H is $tr - \sigma$ which cannot exceed the following obvious upper bound on this sum: $t(t - 1)p$ if $X = M$ and $t(t + 1)p$ if $X = P$. This gives (P) and (M). So all four conditions (i)–(iv) are necessary.

To establish sufficiency, let t be the least nonnegative integer satisfying conditions (i)–(iv). If $\sigma = 0$ then G is r -regular whence $t = 0$, which agrees with conditions (i)–(iv). Assume that $\sigma > 0$. Then $r \geq \Delta \geq 0$ and $r > 0$. Moreover, $t > 0$ by (i). Recall that v_i denotes the i th vertex of the X -graph G and $V = V(G)$. Let U be a set of t extra vertices u_1, u_2, \dots, u_t . Let B be the bipartite $V-U$ multigraph comprising a_i edges joining v_i to U for $i = 1, 2, \dots, n$. Thus σ is the number of all $V-U$ edges. Define an X -graph F to be the edge-disjoint union of X -graphs G , B and H where H is induced by the set U . Assume that $V-U$ edges make up a sequence \mathcal{A} such that, for each i , edges incident with v_{i+1} follow all those incident with v_i . In order to establish incidence of $V-U$ edges with vertices in U , consider the sequence of vertices $(\tilde{u}_1, \tilde{u}_2, \dots, \tilde{u}_t, \dots, \tilde{u}_\sigma)$, where $\tilde{u}_j = u_i$ if $j \equiv i \pmod{t}$ for $j = 1, \dots, \sigma$ and $i = 1, \dots, t$. Assume that the j th vertex \tilde{u}_j is made incident to the j th edge of \mathcal{A} . Then multiplicity of the pair v_i, u_j equals at most $\lceil a_i/t \rceil$ which is at most p by (ii). Hence $F = G \cup B$ if equality holds in condition (i) (which is the case if both $X = M$ and $t = 1$).

Otherwise strong inequality holds in (i) and additionally $t > 1$ if $X = M$. We are going to show that adding some $U-U$ edges can complete the X -graph F . Notice that $nr + \sigma$ is even as the sum of degrees of vertices in the X -graph $G \cup B$. Hence and by (iii) $tr - \sigma$ is an even number. Then for $t = 1$ and $X = P$, it is enough to add $(tr - \sigma)/2$ ($\leq p$ by (P)) loops to the vertex u_1 in order to get a required F . For $t \geq 2$ and $X \in \{P, M\}$, define nonnegative integers h and s so that $\sigma/t = h + s/t$ where $s < t$. Hence $h < r$ because strong inequality in (i) is assumed. Then $\deg_B(u_j) = h + 1$ for $j = 1, 2, \dots, s$, otherwise $\deg_B(u_j) = h$. Thus the degrees in B of vertices of U differ mutually by one at most. Moreover, the greatest remaining deficiency $r - h$ can be covered up in H if either $r - h \leq (t + 1)p$ for $X = P$ or $r - h \leq (t - 1)p$ for $X = M$. But the required conditions (on integers) hold because, by definition of h and s ,

$$r - h - s/t = r - \sigma/t \leq \begin{cases} (t + 1)p & \text{due to (P) if } X = P, \\ (t - 1)p & \text{due to (M) if } X = M, \end{cases}$$

where $0 \leq s/t < 1$. A required X -graph H can be extracted from the complete X -graph on t vertices by using edge-disjoint complete subgraphs

K_t and possibly an edge-decomposition of one K_t into Hamiltonian cycles and, additionally for even t , a perfect matching (see Berge [1] for the well-known solution to Kirkman's problem on packing Hamilton cycles into K_t . The solution presented in Berge appears already in Lucas [7] of 1883 wherein no reference to Kirkman is made).

Now we are going to prove the statement on t_0 , t_0 defined in (1).

Claim 1 *If $r/p \geq n$ and $\delta < r + p - p[r/p]$ then conditions (ii) and (iii) imply that $t \geq t_0$.*

Proof. Note that (ii) implies $t \geq t_1 := [r/p]$ if and only if $r - \delta > ([r/p] - 1)p$ or (equivalently) $\delta < r + p - p[r/p]$. However, if r and $n + t_1$ are odd integers and $t_1 > n$ then $t \neq t_1$ by condition (iii). \square

Claim 2 *Conditions (i) and (ii) hold for $t \geq \max\{[r/p], n\}$.*

Proof. Since $nr \geq \sigma$, condition (i) holds for $t \geq n$. Moreover, (ii) holds for $t \geq [r/p]$. \square

Assume that $\sigma > 0$ in the following part of the proof. Hence it is easily seen that (P) holds for $t \geq r/p - 1$. This and Claim 2 imply that conditions (i),(ii) and (P) hold for $t \geq \max\{[r/p], n\}$. Hence by Claim 1 conditions (i)-(iv) hold for $t \geq t_0$ if $X = P$.

Claim 3 *Condition (M) holds for $t \geq n$ if $r \leq 2pn$, otherwise for $t \geq [r/p] + 1 - n$.*

Proof. Let $L(t)$ stand for the left-hand side of the inequality (M). Then $L(n) = (\sigma - n(r - \Delta)) + n((n - 1)p - \Delta) \geq 0$ as the sum of two nonnegative summands. Hence (M) holds for $t = n$. Assume that $r \leq 2pn$ and note that $L(t)$ is the quadratic trinomial in t which attains its minimum at $\tau := (r/p + 1)/2 \leq n + 1/2$. Therefore for any integer $t \geq n + 1$, $L(t) \geq L(n) \geq 0$.

Otherwise $r > 2pn$ whence $\tau > n + 1/2$. Therefore $L(t) \geq 0$ for $t \leq n$. Hence, since $L(t)$ is symmetrical with respect to $t = \tau > n$, $L(t) \geq 0$ for each $t \geq n + 2(\tau - n) = r/p + 1 - n$. \square

Claim 3 implies that, for $n \geq 2$ and $r/p > 2n$, condition (M) holds for $t \geq [r/p] - 1$. Hence and by Claims 1, 2 and 3, conditions (i)-(iv) hold for $t \geq t_0$ if $X = M$. \blacksquare

Remark 1. A specification of Theorem in case G and F are simple graphs is obtainable by straightforwardly substituting therein $X = M$ and $p = 1$. Additionally, for $r = \Delta$ this way we get main body of the classical Erdős and Kelly's Theorem [3, 4].

Remark 2. For $X \in \{M, P\}$ the following n -vertex X -graphs G show that the upper bound t_0 on the smallest t in Theorem is attainable. Assume

that E_m denotes a set of m edges of the complete graph K_n . For $X = M$, $p = 1$ and $r = \Delta(G) > 0$ (which is the Erdős and Kelly's case): $t = n$ if $n \geq 4$ and $G = K_n - E_m$ where $1 \leq m \leq \lfloor (n-2)/2 \rfloor$ (the Erdős and Kelly example $K_n - e$ being included). For any X , any p , and $r/p \leq n$: $t = n$ if G is any X -graph with $e(G) \leq (r-1)/2$ (whence $\Delta < r$); whereas for $r/p > n$: $t = t_0 > n$ if $\delta(G) < r + p - p\lceil r/p \rceil$ (i.e., if $\delta(G) = 0$ for $p = 1$).

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