

On Resolving Acyclic Partitions of Graphs

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Abstract

For a vertex v of a connected graph G and a subset S of $V(G)$, the distance between v and S is $d(v, S) = \min\{d(v, x) : x \in S\}$, where $d(v, x)$ is the distance between v and x . For an ordered k -partition $\Pi = \{S_1, S_2, \dots, S_k\}$ of $V(G)$, the code of v with respect to Π is the k -vector $c_\Pi(v) = (d(v, S_1), d(v, S_2), \dots, d(v, S_k))$. The k -partition Π is a resolving partition if the codes $c_\Pi(v)$, $v \in V(G)$, are distinct. A resolving partition $\Pi = \{S_1, S_2, \dots, S_k\}$ is acyclic if each subgraph $\langle S_i \rangle$ induced by S_i ($1 \leq i \leq k$) is acyclic in G . The minimum k for which there is a resolving acyclic k -partition of $V(G)$ is the resolving acyclic number $a_r(G)$ of G . We study connected graphs with prescribed order, diameter, vertex-arboricity, and resolving acyclic number. It is shown that, for each triple d, k, n of integers with $2 \leq d \leq n-2$ and $3 \leq (n-d+1)/2 \leq k \leq n-d+1$, there exists a connected graph of order n having diameter d and resolving acyclic number k . Also, for each pair a, b of integers with $2 \leq a \leq b-1$, there exists a connected graph with resolving acyclic number a and vertex-arboricity b . We present a sharp lower bound for the resolving acyclic number of a connected graph in terms of its clique number. The resolving acyclic number of the Cartesian product $H \times K_2$ of nontrivial connected graph H and K_2 is studied.

Key Words: distance, resolving partition, acyclic resolving partition.

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1 Introduction

Let G be a nontrivial connected graph. For a set S of vertices of G and a vertex v of G , the *distance* $d(v, S)$ between v and S is defined as

$$d(v, S) = \min\{d(v, x) : x \in S\},$$

where $d(v, x)$ is the distance between v and x . For an ordered k -partition $\Pi = \{S_1, S_2, \dots, S_k\}$ of $V(G)$ and a vertex v of G , the *code of v with respect to Π* is defined as the k -vector

$$c_\Pi(v) = (d(v, S_1), d(v, S_2), \dots, d(v, S_k)).$$

The partition Π is defined in [4] as a *resolving partition* for G if the distinct vertices of G have distinct codes with respect to Π . The minimum k for which there is a resolving k -partition of $V(G)$ is the *partition dimension* $\text{pd}(G)$ of G . A resolving partition of $V(G)$ containing $\text{pd}(G)$ elements is called a *minimum resolving partition*. Resolving partitions in graphs were first introduced and studied in [4]. Resolving partitions that satisfy some additional prescribed properties have been studied [5, 6, 10]

A resolving partition $\Pi = \{S_1, S_2, \dots, S_k\}$ of $V(G)$ is *independent* if each subgraph $\langle S_i \rangle$ induced by S_i ($1 \leq i \leq k$) is independent in G . This topic was first introduced from the point of view of graph coloring in [5, 6]. If Π is an independent partition of $V(G)$, then, by coloring the vertices in S_i by i ($1 \leq i \leq k$), we obtain a proper coloring with k colors that distinguishes all vertices of G in terms of their distances from the color classes. Thus, such a coloring of a graph G is called a *resolving-coloring* (or *locating-coloring*). A *minimum resolving-coloring* uses a minimum number of colors and this number is the *resolving-chromatic number* $\chi_r(G)$ of G . Since every resolving-coloring is a coloring, $\chi(G) \leq \chi_r(G)$ for each connected graph G , where $\chi(G)$ is the chromatic number of G .

A resolving partition $\Pi = \{S_1, S_2, \dots, S_k\}$ of $V(G)$ is defined in [10] to be *acyclic* if each subgraph $\langle S_i \rangle$ induced by S_i ($1 \leq i \leq k$) is acyclic in G . The minimum k for which G contains a resolving acyclic k -partition is the *resolving acyclic number* $a_r(G)$ of G . The *vertex-arboricity* $a(G)$ of G is defined in [1, 2] as the minimum k such that $V(G)$ has an acyclic k -partition. Since every resolving acyclic partition is an acyclic partition, $a(G) \leq a_r(G)$ for each connected graph G .

It was observed in [10] that

$$2 \leq \text{pd}(G) \leq a_r(G) \leq \chi_r(G) \leq n$$

for every nontrivial connected graph G of order n .

To illustrate these concepts, consider the graph G of Figure 1. Let $\Pi = \{S_1, S_2, S_3\}$, where $S_1 = \{x\}$, $S_2 = \{u\}$, and $S_3 = \{v, y, z\}$. Then the

corresponding codes of the vertices of G are

$$\begin{aligned} c_{\Pi}(u) &= (1, 0, 1) & c_{\Pi}(v) &= (1, 2, 0) & c_{\Pi}(x) &= (0, 1, 1) \\ c_{\Pi}(y) &= (1, 1, 0) & c_{\Pi}(z) &= (2, 1, 0). \end{aligned}$$

Since the codes of the vertices of G with respect to Π are distinct, Π is a resolving partition of G . Because no 2-partition is a resolving partition of G , it follows that Π is a minimum resolving partition of G and so $\text{pd}(G) = 3$. However, Π is not acyclic since $\langle S_3 \rangle = K_3$. On the other hand, let $\Pi' = \{S'_1, S'_2, S'_3, S'_4\}$, where $S'_1 = \{x\}$, $S'_2 = \{u\}$, $S'_3 = \{v, y\}$, and $S'_4 = \{z\}$. It can be verified that Π' is a resolving acyclic partition of G and no 3-partition is a resolving acyclic partition of G . Thus $a_r(G) = 4$. Furthermore, $\chi_r(G) = 5$ (see [6]). Therefore, $\text{pd}(G) < a_r(G) < \chi_r(G)$ for the graph G of Figure 1. The example just described also illustrates an important point. When determining whether a given partition Π is a resolving partition of a connected graph G , we need only verify that vertices of G belonging to the same subset of $V(G)$ in Π have distinct codes since the codes of two vertices in different subsets in Π have 0 in different coordinates.

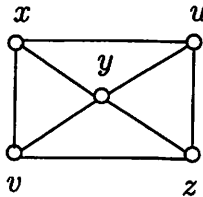


Figure 1: A graph G with $\text{pd}(G) < a_r(G) < \chi_r(G)$

The concept of resolvability in graphs has appeared in the literature. Slater introduced and studied these ideas with different terminology in [11, 12]. Slater described the usefulness of these ideas when working with U.S. sonar and coast guard Loran (Long range aids to navigation) stations. Harary and Melter [7] discovered these concepts independently. Recently, these concepts were rediscovered by Johnson [8, 9] of the Pharmacia Company while attempting to develop a capability of large datasets of chemical graphs. We refer to the book [3] for graph theory notation and terminology not described here.

The resolving acyclic numbers of some familiar classes of graphs have been determined in [10].

Theorem A *Let $n \geq 2$ and let G be a connected graph of order n . Then*

- (a) $a_r(G) = 2$ if and only if $G = P_n$;

(b) $a_r(G) = n$ if and only if $G = K_n$;

(c) $a_r(G) = n - 1$ if and only if

$$G \in \{C_4 + K_1, K_{1,n-1}, K_n - e, K_1 + (K_1 \cup K_{n-2})\}$$

for $n \geq 5$.

Theorem B *Let r, s be positive integers. If G is a connected bipartite graph with partite sets of cardinalities r and s , then*

$$a_r(G) \leq \begin{cases} r + 1 & \text{if } r = s \\ \max\{r, s\} & \text{if } r \neq s. \end{cases}$$

Moreover, equality holds if G is a complete bipartite graph.

2 Graphs with Prescribed Resolving Acyclic Number and Other Parameters

In this section, we study connected graphs with prescribed resolving acyclic number and other parameters, such as diameter, order, vertex-arboricity, clique number, and etc. We begin with the connected graphs with prescribed resolving acyclic number, order, and diameter. The *diameter* of a connected graph G is the largest distance between two vertices in G . Bounds for the resolving acyclic number of a connected graph was established in terms of its order and diameter in [10].

Theorem C *If G is a connected graph of order $n \geq 3$ and diameter $d \geq 2$, then*

$$\log_{d+1} n \leq a_r(G) \leq n - d + 1.$$

The lower bound in Theorem C is not sharp. In order to show this, we present a result that gives an upper bound for the order of a connected graph in terms of its diameter and resolving acyclic number.

Proposition 2.1 *If G is a nontrivial connected graph of order n , diameter d , and resolving acyclic number k , then $n \leq kd^{k-1}$.*

Proof. If $d = 1$, then $G = K_n$. By Theorem A, $a_r(G) = n = k$, and so the result holds for $d = 1$. Thus we may assume that $d \geq 2$ and so $n \geq 3$. Let Π be a resolving acyclic k -partition of $V(G)$. Since (1) exactly one coordinate of the code of a vertex in G with respect to Π is 0 and there are k choices for the zero coordinates in the code of a vertex, (2) each of the $k - 1$ nonzero coordinate of the code of a vertex is a positive integer not

exceeding d , and (3) all codes of the n vertices of G are distinct, it follows that $kd^{k-1} \geq n$. ■

Certainly, for each pair d, k of positive integers, $(d+1)^k > kd^{k-1}$. Thus, if G is a nontrivial connected graph of order n , diameter d , and resolving acyclic number k , then, by Proposition 2.1, $(d+1)^k > kd^{k-1} \geq n$, and so $a_r(G) > \log_{d+1} n$. Thus, the following corollary is an immediate consequence of Proposition 2.1.

Corollary 2.2 *If G is a connected graph of order $n \geq 3$ and diameter $d \geq 2$, then*

$$a_r(G) \geq 1 + \log_{d+1} n.$$

Notice that if $G = P_n$, then $d = n - 1$ and $\log_{d+1} n = 1$. By Theorem A, the equality in Corollary 2.2 holds for $G = P_n$. The upper bound in Theorem C is sharp. In fact, more can be said, as we will see later. First, we present a useful lemma, whose proof is straightforward and is therefore omitted.

Lemma 2.3 *Let Π be a resolving partition in a connected graph G . If u and v are distinct vertices of G such that $d(u, w) = d(v, w)$ for all $w \in V(G) - \{u, v\}$, then u and v belong to distinct elements of Π . In particular, if G is a connected graph containing a vertex that is adjacent to k end-vertices of G , then $a_r(G) \geq k$.*

We now study those triples of positive integers that are realizable as the diameter, resolving acyclic number, and order of some connected graph.

Theorem 2.4 *For each triple d, k, n of integers with $2 \leq d \leq n - 2$ and $3 \leq (n - d + 1)/2 \leq k \leq n - d + 1$, there exists a connected graph of order n having diameter d and resolving acyclic number k .*

Proof. First, assume that $k = n - d + 1$. Let G be the graph obtained from K_{n-d+1} and the path $P_{d-1} : v_1, v_2, \dots, v_{d-1}$ by joining v_1 to a vertex u in K_{n-d+1} . Then the order of G is n and the diameter of G is d . We show that $a_r(G) = n - d + 1$. Since every element of any resolving acyclic partition of $V(G)$ contains at most one vertex of $V(K_{n-d+1}) - \{u\}$ by Lemma 2.3, it follows that $a_r(G) \geq n - d$. Assume, to the contrary, that $a_r(G) = n - d$. Let $V(K_{n-d+1}) - \{u\} = \{u_1, u_2, \dots, u_{n-d}\}$ and let $\Pi = \{S_1, S_2, \dots, S_{n-d}\}$ be a resolving acyclic partition of $V(G)$. Assume, without loss of generality, that $u_i \in S_i$ and $u \in S_1$. However, then, $c_\Pi(u) = (0, 1, 1, \dots, 1) = c_\Pi(u_1)$, which is a contradiction. Thus, $a_r(G) \geq n - d + 1 = k$. On the other hand, $\Pi^* = \{S_1^*, S_2^*, \dots, S_{n-d+1}^*\}$, where $S_i^* = \{u_i\}$ ($1 \leq i \leq n - d$) and $S_{n-d+1}^* = \{u\} \cup V(P_{d-1})$, is a resolving acyclic partition of $V(G)$ and so $a_r(G) \leq |\Pi^*| = n - d + 1$. Therefore, $a_r(G) = n - d + 1 = k$.

Next, assume that $(n - d + 1)/2 \leq k \leq n - d$. We consider two cases.

Case 1. $d = 2$. Then $(n - 1)/2 \leq k \leq n - 2$. If $(n + 1)/2 \leq k \leq n - 2$, then $k > n - k$ and $G = K_{k,n-k}$ has the desired property by Theorem B. If $(n - 1)/2 \leq k < (n + 1)/2$, then $k = n/2$ if n is even and $k = (n - 1)/2$ if n is odd. We consider these two subcases.

Subcase 1.1. $k = n/2$. Let $G = (K_{k-1} \cup kK_1) + K_1$. Let $V(K_1) = \{u_1\}$, $V(K_{k-1}) = \{u_2, u_3, \dots, u_k\}$, and $V(kK_1) = \{v_1, v_2, \dots, v_k\}$. Since u_1 is adjacent to k end-vertices, $a_r(G) \geq k$ by Lemma 2.3. On the other hand, $\Pi = \{S_1, S_2, \dots, S_k\}$, where $S_i = \{u_i, v_i\}$ ($1 \leq i \leq k$), is a resolving acyclic partition of $V(G)$ and so $a_r(G) = k$.

Subcase 1.2. $k = (n - 1)/2$. Let $G = (P_k \cup kK_1) + K_1$. Let $V(K_1) = \{w\}$, $V(P_k) = \{u_1, u_2, \dots, u_k\}$, and $V(kK_1) = \{v_1, v_2, \dots, v_k\}$. Since w is adjacent to k end-vertices, $a_r(G) \geq k$ by Lemma 2.3. On the other hand, $\Pi = \{S_1, S_2, \dots, S_k\}$, where $S_1 = \{u_1, v_1, w\}$ and $S_i = \{u_i, v_i\}$ ($2 \leq i \leq k$), is a resolving acyclic partition of $V(G)$ and so $a_r(G) = k$.

Case 2. $3 \leq d \leq n - 2$. Again, we consider two subcases.

Subcase 2.1. $k = n - d$. Let G be the graph obtained from the path $P_d : u_1, u_2, \dots, u_d$ by adding the $n - d$ pendant edges u_1v_i ($1 \leq i \leq n - d$). Then the order of G is n and the diameter of G is d . Since u_1 is adjacent to $n - d$ end-vertices, it then follows by Lemma 2.3 that $a_r(G) \geq n - d$. On the other hand, let $\Pi = \{S_1, S_2, \dots, S_{n-d}\}$ be a partition of $V(G)$ where $S_1 = \{u_3, v_1\}$, $S_2 = \{v_2\} \cup (V(P_{n-d} - \{u_3\}))$, and $S_i = \{v_i\}$ for $3 \leq i \leq n - d$. Since Π is a resolving acyclic partition of G , it follows that $a_r(G) \leq |\Pi| = n - d$. Thus $a_r(G) = n - d$.

Subcase 2.2. $k \leq n - d - 1$. Since $k \leq n - d - 1$, it follows that $n - d - k + 1 \geq 2$. Let G be the graph obtained from the path $P_{d-1} : u_1, u_2, \dots, u_{d-1}$ by adding the $n - d + 1$ pendant edges u_1v_i ($1 \leq i \leq k$), u_2w_1 , and $u_{d-1}w_j$ ($2 \leq j \leq n - d - k + 1$). Then the order of G is n and the diameter of G is d . We show that $a_r(G) = k$. Since $k \geq (n - d + 1)/2$, it follows that $k \geq n - d - k + 1 > n - d - k$. Thus $a_r(G) \geq k$ by Lemma 2.3. On the other hand, let $\Pi = \{S_1, S_2, \dots, S_k\}$, where $S_1 = \{v_1, w_1\}$, $S_2 = \{u_2, v_2, w_2\}$, $S_3 = (V(P_{d-1}) - \{u_2\}) \cup \{v_3, w_3\}$, $S_i = \{v_i, w_i\}$ for $4 \leq i \leq n - d - k + 1$, and $S_i = \{v_i\}$ for $n - d - k + 2 \leq i \leq k$. Since G is a tree, it suffices to show that Π is a resolving partition of G . Observe that (1) $c_\Pi(v_1) = (0, 2, \dots)$ and $c_\Pi(w_1) = (0, 1, \dots)$, (2) $c_\Pi(u_2) = (1, 0, 1, \dots)$, $c_\Pi(v_2) = (2, 0, 1, \dots)$, and $c_\Pi(w_2) = (2, 0, 2, \dots)$ if $d = 3$ and $c_\Pi(w_2) = (d - 1, 0, 1, \dots)$ if $d \geq 4$, (3) $c_\Pi(v_3) = (2, 2, 0, \dots)$, and $c_\Pi(w_3) = (2, 1, 0, \dots)$ if $d = 3$ and $c_\Pi(w_3) = (d - 1, 2, 0, \dots)$ if $d \geq 4$, $c_\Pi(u_i) = (1, *, 0, \dots)$ and $c_\Pi(w_i) = (i - 1, *, 0, \dots)$ for $3 \leq i \leq d - 1$, (4) $c_\Pi(v_i) = (2, *, 1, \dots)$, $c_\Pi(w_i) = (2, *, 2, \dots)$ if $d = 3$ and $c_\Pi(w_i) = (d - 1, *, 1, \dots)$ if $d \geq 4$, where $*$ represents an irrelevant

coordinate. Thus all codes of the vertices of G are distinct, and so Π is resolving, implying that $a_r(G) \leq |\Pi| = k$. Therefore, $a_r(G) = k$. ■

We have seen that if G is a connected graph with $a(G) = a$ and $a_r(G) = b$, then $1 \leq a \leq b$ and $b \geq 2$. Next we study those pairs a, b of integers with $1 \leq a \leq b$ and $b \geq 2$ that are realizable as vertex-arboricity and resolving acyclic number of some connected graph. Since trees are the only connected graphs having vertex-arboricity 1, it follows that $a(T) \neq a_r(T)$ for all trees T . Thus we may assume that $a \geq 2$. It is known that if G is a connected graph of order n , then $a(G) \leq \lceil n/2 \rceil$. Thus if $a_r(G) > \lceil n/2 \rceil$, then $a(G) \neq a_r(G)$. On the other hand, we show next that each pair a, b of integers with $2 \leq a \leq b - 1$ is realizable as vertex-arboricity and resolving acyclic number of some connected graph.

Theorem 2.5 *For each pair a, b of integers with $2 \leq a \leq b - 1$, there exists a connected graph G with $a(G) = a$ and $a_r(G) = b$.*

Proof. For $a = 2$, let $G = K_{a,b}$. It is known that $a(G) = 2$. Since $a < b$, it follows by Theorem B that $a_r(G) = b$ and so G has the desired properties. For $a \geq 3$, let G be the graph obtained from K_{2a} with $V(K_{2a}) = \{u_1, u_2, \dots, u_{2a}\}$ by adding the $b + a - 1$ new vertices v_1, v_2, \dots, v_a and w_1, w_2, \dots, w_{b-1} , joining each vertex v_i to u_i ($1 \leq i \leq a$) and joining each vertex w_j ($1 \leq j \leq b - 1$) to u_a . Thus $a(G) = a$.

Next, we show that $a_r(G) = b$. Since u_a is adjacent to b end-vertices, namely, v_a and w_j ($1 \leq j \leq b - 1$), it then follows by Lemma 2.3 that $a_r(G) \geq b$. On the other hand, let $\Pi = \{S_1, S_2, \dots, S_b\}$ be a partition of $V(G)$, where $S_i = \{u_i, u_{a+i}, w_i\}$, $1 \leq i \leq a$, $S_i = \{w_i\}$, $a + 1 \leq i \leq b - 1$, and $S_b = \{v_1, v_2, \dots, v_a\}$. Thus Π is acyclic. Notice that (1) $d(u_i, S_j) = 1$ for $1 \leq i \neq j \leq a$ and $d(u_i, S_b) = 1$, (2) $d(u_{a+i}, S_j) = 1$ for $1 \leq i \neq j \leq a$ and $d(u_{a+i}, S_b) = 2$, and (3) $d(w_i, S_j) = 2$ for $1 \leq i \neq j \leq a$ and $j \neq a$ and $d(w_i, S_a) = 1$ for $1 \leq i \leq a$. Since $a \geq 3$, it follows that $c_{\Pi}(u_i)$, $c_{\Pi}(u_{a+i})$, and $c_{\Pi}(w_i)$ are distinct for $1 \leq i \leq a$. Furthermore, the i th coordinate of $c_{\Pi}(v_i)$ is 1, the b th coordinate of $c_{\Pi}(v_i)$ is 0, and the remaining coordinates of $c_{\Pi}(v_i)$ are 2 or 3, implying that the codes $c_{\Pi}(v_i)$, $1 \leq i \leq a$, are distinct. Hence Π is a resolving acyclic partition of $V(G)$ and so $a_r(G) \leq |\Pi| = b$. Therefore, $a_r(G) = b$. ■

The *clique number* of a graph is the maximum order among the complete subgraphs of the graph. We state a lower bound for the resolving acyclic number of a connected graph in terms of its clique number.

Proposition 2.6 *If the clique number of a connected graph G is ω , then*

$$a_r(G) \geq \left\lceil \frac{\omega}{2} \right\rceil + 1.$$

Moreover, for each integer $\omega \geq 2$, there exists a connected graph G_ω having clique number ω such that $a_r(G_\omega) = \lceil \omega/2 \rceil + 1$.

3 Cartesian Products

We show next that the resolving acyclic number of the Cartesian product of K_2 and a nontrivial connected graph H is bounded above by $a_r(H) + a(H)$.

Theorem 3.1 *For every nontrivial connected graph H ,*

$$a_r(H \times K_2) \leq a_r(H) + a(H)$$

Proof. Let $G = H \times K_2$, where H_1 and H_2 are the two copies of H in the construction of G . Suppose that $a_r(H) = k$ and $a(H) = a$. Let $\Pi = \{S_1, S_2, \dots, S_k\}$ be a resolving acyclic partition of $V(H_1)$, and let $\{W_1, W_2, \dots, W_a\}$ be an acyclic partition of $V(H_2)$. Then

$$\Pi^* = \{S_1, S_2, \dots, S_k, W_1, W_2, \dots, W_a\}$$

is an acyclic partition of $V(G)$. We show that Π^* is a resolving partition of $V(G)$. Let x and y be vertices of G such that $c_{\Pi^*}(x) = c_{\Pi^*}(y)$. We show that $x = y$. Assume, to the contrary, that $x \neq y$. We consider three cases.

Case 1. Both x and y belong to H_1 . Then $d_G(x, S_i) = d_{H_1}(x, S_i)$ and $d_G(y, S_i) = d_{H_1}(y, S_i)$ ($1 \leq i \leq k$). Since $d_G(x, S_i) = d_G(y, S_i)$ for all $1 \leq i \leq k$, it follows that $d_{H_1}(x, S_i) = d_{H_1}(y, S_i)$ for $1 \leq i \leq k$ and so $c_\Pi(x) = c_\Pi(y)$, contradicting the fact that Π is a resolving acyclic partition of H_1 .

Case 2. Both x and y belong to H_2 . Let x' and y' be the vertices of H_1 that correspond to x and y in H_2 , respectively. Since $x \neq y$, we have $x' \neq y'$. Notice that $d_G(x, S_i) = d_G(x', S_i) + 1 = d_{H_1}(x', S_i) + 1$ and $d_G(y, S_i) = d_G(y', S_i) + 1 = d_{H_1}(y', S_i) + 1$ for $1 \leq i \leq k$. Thus $d_{H_1}(x', S_i) = d_G(x, S_i) - 1$ and $d_{H_1}(y', S_i) = d_G(y, S_i) - 1$ for $1 \leq i \leq k$. Since $d_G(x, S_i) = d_G(y, S_i)$ for $1 \leq i \leq k$, it follows that $d_{H_1}(x', S_i) = d_{H_1}(y', S_i)$ for $1 \leq i \leq k$ and so $c_\Pi(x') = c_\Pi(y')$, a contradiction.

Case 3. Either $x \in V(H_1)$ and $y \in V(H_2)$, or $x \in V(H_2)$ and $y \in V(H_1)$, say the former. Suppose that $y \in W_i$, where $1 \leq i \leq a$ and so $d_G(y, W_i) = 0$. However, $x \notin V(H_2)$ and so $x \notin W_i$, which implies that $d_G(x, W_i) > 0$. Thus $c_{\Pi^*}(x) \neq c_{\Pi^*}(y)$, contradicting the assumption.

Therefore, Π^* is a resolving acyclic partition of $V(G)$ and so $a_r(G) \leq |\Pi^*| = a_r(H) + a(H)$, as desired. \blacksquare

Equality in Theorem 3.1 can hold. For example, if $H = P_n$, where $n \geq 2$, then $a_r(P_n) = 2$ and $a(P_n) = 1$. By Theorem A, $a_r(P_n \times K_2) \geq 3$. On the

other hand, let $H_1 : u_1, u_2, \dots, u_n$ and $H_2 : v_1, v_2, \dots, v_n$ be two copies of P_n in $P_n \times K_2$. Since $\Pi = \{S_1, S_2, S_3\}$, where $S_1 = V(H_1)$, $S_2 = \{v_1\}$, and $S_3 = V(H_2) - \{v_1\}$, is a resolving acyclic partition of $V(P_n \times K_2)$, it follows that $a_r(P_n \times K_2) = 3$. Therefore, $a_r(P_n \times K_2) = 3 = a_r(P_n) + a(P_n)$ for $n \geq 2$.

Strict inequality in Theorem 3.1 can also hold. As an example, we study $a_r(K_n \times K_2)$ for $n \geq 3$, beginning with $n = 3, 4$.

Proposition 3.2 $a_r(K_3 \times K_2) = 3$ and $a_r(K_4 \times K_2) = 4$.

Proof. For $n = 3, 4$, let H_1 and H_2 be two copies of K_n in G , where $V(H_1) = \{u_1, u_2, \dots, u_n\}$, $V(H_2) = \{v_1, v_2, \dots, v_n\}$, and $u_i v_i \in E(G)$ for $1 \leq i \leq n$. For $n = 3$, let $\Pi = \{S_1, S_2, S_3\}$, where $S_1 = \{u_1, u_2, v_1\}$, $S_2 = \{v_2, v_3\}$, and $S_3 = \{u_3\}$. Since Π is a resolving acyclic partition, it follows by Theorem A that $a_r(K_3 \times K_2) = 3$.

For $n = 4$, let $\Pi^* = \{S_1^*, S_2^*, S_3^*, S_4^*\}$, be a partition of $V(G)$, where $S_1^* = \{u_1, u_3\}$, $S_2^* = \{u_2, v_2, v_3\}$, $S_3^* = \{v_1, v_4\}$, and $S_4^* = \{u_4\}$. Since Π^* is a resolving acyclic partition, $a_r(G) \leq |\Pi^*| = 4$. Assume, to the contrary, that $a_r(G) = 3$. Let $\Pi = \{S_1, S_2, S_3\}$ be a minimum resolving partition of G . Since $\langle S_i \rangle$, $1 \leq i \leq 3$, is acyclic, S_i contains at most two vertices of H_1 . Thus at least two of S_1, S_2 , and S_3 contain vertices of H_1 . If each of S_1, S_2 , and S_3 contains some vertex of H_1 , then assume, without loss of generality, that $u_i \in S_i$ for $1 \leq i \leq 3$. Since Π is a partition of $V(G)$, it follows that $u_4 \in S_i$ for some i with $1 \leq i \leq 3$, say $u_4 \in S_1$. Then $c_\Pi(u_1) = c_\Pi(u_4) = (0, 1, 1)$, a contradiction. Thus exactly two of S_1, S_2 , and S_3 contains vertices of H_1 , say S_1 and S_2 contain vertices of H_1 . Moreover, each of S_1 and S_2 contains exactly two vertices of H_1 . We may assume that $u_1, u_2 \in S_1$ and $u_3, u_4 \in S_2$. Similarly, exactly two of S_1, S_2 , and S_3 contains vertices of H_2 , each of which contains exactly two vertices of H_2 . Since $S_3 \neq \emptyset$ and $S_3 \cap V(H_1) = \emptyset$, it follows that S_3 consists of exactly two vertices of H_2 . This implies that exactly one of S_1 and S_2 contains the two vertices in $V(H_2) - S_3$, say $V(H_2) - S_3 \subseteq S_1$ and so $S_2 = \{u_3, u_4\}$. Since $\langle S_1 \rangle$ is acyclic and $u_1, u_2 \in S_1$, it follows that S_1 contains at most one of v_1 and v_2 . We consider two cases.

Case 1. S_1 contains exactly one of v_1 and v_2 , say $v_1 \in S_1$. Thus $S_1 = \{u_1, u_2, v_1, v_j\}$, where $j = 3$ or $j = 4$. If $S_1 = \{u_1, u_2, v_1, v_3\}$, then $S_2 = \{u_3, u_4\}$, and $S_3 = \{v_2, v_4\}$. Thus $c_\Pi(u_2) = c_\Pi(v_3) = (0, 1, 1)$, a contradiction. If $S_1 = \{u_1, u_2, v_1, v_4\}$, then $S_2 = \{u_3, u_4\}$, and $S_3 = \{v_2, v_3\}$. Thus $c_\Pi(u_2) = c_\Pi(v_4) = (0, 1, 1)$, a contradiction.

Case 2. S_1 contains neither v_1 nor v_2 . Thus $S_1 = \{u_1, u_2, v_3, v_4\}$, $S_2 = \{u_3, u_4\}$, and $S_3 = \{v_1, v_2\}$. Then $c_\Pi(u_1) = c_\Pi(u_2) = (0, 1, 1)$, a contradiction. ■

Thus $a_r(K_n \times K_2) = a_r(K_n) = n$ for $n = 3, 4$. However, $a_r(K_n \times K_2) < n$ for $n \geq 5$. We present an upper bound for $a_r(K_n \times K_2)$ in terms of n without a proof.

Proposition 3.3 *If $n \geq 5$, then*

$$a_r(K_n \times K_2) \leq \begin{cases} \frac{3n}{4} & \text{if } n \equiv 0 \pmod{4} \\ \frac{3(n-1)}{4} + 2 & \text{if } n \equiv 1 \pmod{4} \\ \frac{3(n-2)}{4} + 2 & \text{if } n \equiv 2 \pmod{4} \\ \frac{3(n-3)}{4} + 3 & \text{if } n \equiv 3 \pmod{4}. \end{cases}$$

Since $a(K_n) = \lceil n/2 \rceil$ for $n \geq 3$, it follows by Proposition 3.3 that strict inequality in Theorem 3.1 can hold.

4 Acyclic Partition Ratios

If H is an induced subgraph of a graph G , then it is known that $a(H) \leq a(G)$. Thus

$$0 < \frac{a(H)}{a(G)} \leq 1$$

for each induced subgraph H of a graph G . However, it is not true, in general, for the resolving acyclic numbers. Let G be a nontrivial connected graph and H a connected induced subgraph of G . We defined *the acyclic partition ratio* of G and H , respectively, by

$$r_a(H, G) = \frac{a_r(H)}{a_r(G)}$$

By Theorem B $a_r(K_{1,m}) = m$ for all $m \geq 2$. Hence for $G = K_{1,m}$ and $H = K_2$, we can make the ratio $r_a(H, G)$ as small as we wish by choosing m arbitrarily large. Although this may not be surprising, it may be unexpected that, in fact, we can make $r_a(H, G)$ as large as we wish. We now establish the truth of this statement.

For $n \geq 3$, we label the vertices of the star $K_{1,2^{n+1}}$ with $v_0, v_1, v_2, \dots, v_{2^n}, v'_1, v'_2, \dots, v'_{2^n}$, where v_0 is the central vertex. Then we add two new vertices x and x' and 2^{n+1} edges xv_i and $x'v'_i$ for $1 \leq i \leq 2^n$. Next, we add two sets $W = \{w_1, w_2, \dots, w_n\}$ and $W' = \{w'_1, w'_2, \dots, w'_n\}$ of vertices, together with the edges w_ix and $w'_i x'$ for $1 \leq i \leq n$. Finally, we add edges between W and $\{v_0, v_1, v_2, \dots, v_{2^n}\}$ so that each of the 2^n possible n -tuples of 1s and 2s appears exactly once such that the representations $(d(v_i, w_1), d(v_i, w_2), \dots, d(v_i, w_n))$ are distinct for $1 \leq i \leq 2^n$.

2^n . Similarly, edges are added between W' and $\{v'_1, v'_2, \dots, v'_{2^n}\}$ so that $(d(v'_i, w'_1), d(v'_i, w'_2), \dots, d(v'_i, w'_n))$ are distinct for $1 \leq i \leq 2^n$. Denote the resulting graph by G . The graph G for $n = 3$ is shown in Figure 2.

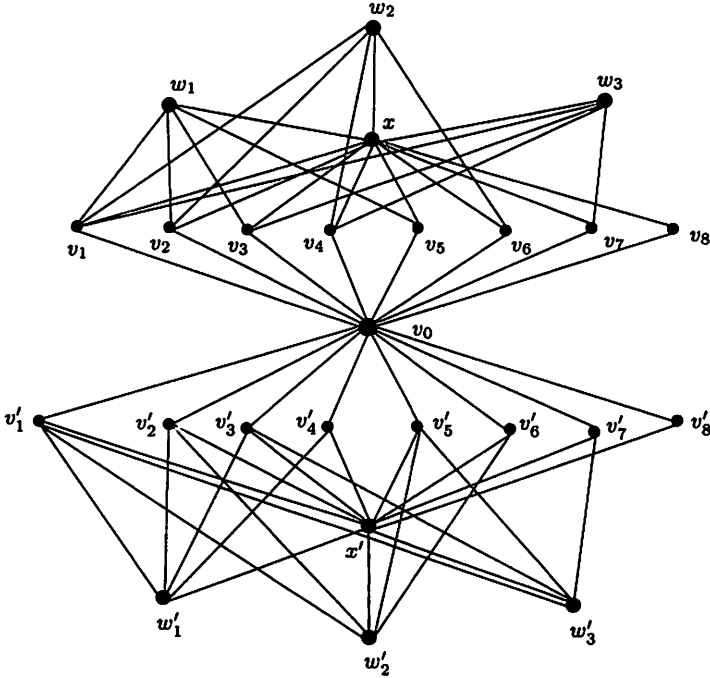


Figure 2: The graph G for $n = 3$

Let $\Pi = \{S_1, S_2, \dots, S_{2n+2}\}$ be the partition of $V(G)$ where $S_i = \{w_i\}$ ($1 \leq i \leq n$), $S_{n+j} = \{w'_j\}$ ($1 \leq j \leq n$), and $S_{2n+1} = \{x, x'\}$ and $S_{2n+2} = V(K_{1,2^{n+1}})$. Since $\langle S_i \rangle$ is acyclic for $1 \leq i \leq 2n + 2$, it follows that Π is acyclic. Next we show that Π is a resolving partition of $V(G)$. By construction, $c_\Pi(v_i) = c_\Pi(v_j)$ implies that $i = j$ and $c_\Pi(v'_i) = c_\Pi(v'_j)$ implies that $i = j$. Moreover,

$$\begin{aligned} c_\Pi(x) &= (1, 1, \dots, 1, 4, 4, \dots, 4, 0, 1), \\ c_\Pi(v_i) &= (*, *, \dots, *, 3, 3, \dots, 3, 1, 0), \quad 1 \leq i \leq 2^n, \\ c_\Pi(v_0) &= (2, 2, \dots, 2, 2, 2, \dots, 2, 2, 0), \\ c_\Pi(v'_i) &= (3, 3, \dots, 3, *, *, \dots, *, 1, 0), \quad 1 \leq i \leq 2^n, \\ c_\Pi(x') &= (4, 4, \dots, 4, 1, 1, \dots, 1, 0, 1), \end{aligned}$$

where $*$ represents an irrelevant coordinate. Thus Π is a resolving acyclic $(2n + 2)$ -partition of $V(G)$. Observe that G contains $H = K_{1,2^{n+1}}$ as an

induced subgraph and

$$\frac{a_r(H)}{a_r(G)} \geq \frac{2^{n+1}}{2n+2}.$$

Since

$$\lim_{n \rightarrow \infty} \frac{2^{n+1}}{2n+2} = \infty,$$

there exists a connected graph G and an induced subgraph H of G such that $r_a(H, G) = a_r(H)/a_r(G)$ is arbitrary large.

5 Open Questions and Topics for Study

The lower and upper bounds for the resolving acyclic number of a connected graph in terms of its resolving-chromatic number were established in [10], which we state next.

Theorem D *For every nontrivial connected graph G ,*

$$\frac{\chi_r(G)}{2} \leq a_r(G) \leq \chi_r(G).$$

Thus, if G is a connected graph G with $a_r(G) = a$ and $\chi_r(G) = b$, then $a \geq 2$ and $b/2 \leq a \leq b$. On the other hand, it is easy to show that, for each pair a, b of integers with $a \geq 2$ and $b/2 < a \leq b$, there exists a connected graph G such that $a_r(G) = a$ and $\chi_r(G) = b$. It was shown in [5] that if G is a nontrivial complete graph or complete bipartite graph, then $\chi_r(G) = n$, where n is the order of G . Thus, if $a = b \geq 2$, then the graph K_a has the property that $a_r(G) = \chi_r(G) = a$ by Theorem A. If $a \geq 2$ and $b/2 < a < b$, then $K_{a, b-a}$ has the property that $a_r(G) = a$ by Theorem B and $\chi_r(G) = b$. These observations yield the following.

Proposition 5.1 *For every pair a, b of integers with $a \geq 2$ and $b/2 < a \leq b$, there exist a connected graph G with $a_r(G) = a$ and $\chi_r(G) = b$.*

However, the following question remains open.

Problem 5.2 *Does there exist a connected graph G such that $\chi_r(G) = 2a_r(G)$?*

We have seen in Theorem 2.4 that, for each triple d, k, n of integers with $2 \leq d \leq n-2$ and $3 \leq (n-d+1)/2 \leq k \leq n-d+1$, there exists a connected graph of order n having diameter d and resolving acyclic number k . On the other hand, for those triples d, k, n of integers with $2 \leq d \leq n-2$ and $3 \leq k \leq (n-d-1)/2$, the following question is open.

Problem 5.3 For which triples d, k, n of integers with $2 \leq d \leq n - 2$ and $3 \leq k \leq (n - d - 1)/2$, does there exist a connected graph of order n having diameter d and resolving acyclic number k ?

Theorem 2.5 shows, for each pair a, b of integers with $2 \leq a \leq b - 1$, that there exists a connected graph G with $a(G) = a$ and $a_r(G) = b$. However, we don't know any connected graph G with $a(G) = a_r(G)$. This suggests the following question.

Problem 5.4 Does there exist a connected graph G such that $a(G) = a_r(G)$?

We have seen in Theorem 3.1 that if H is a nontrivial connected graph, then $a_r(H \times K_2) \leq a_r(H) + a(H)$, which suggests the following question.

Problem 5.5 For which triples a, b, c of positive integers with $a \geq 2$ and $3 \leq c \leq a + b$, does there exist a connected graph H such that $a_r(H) = a$ and $a(H) = b$, and $a_r(H \times K_2) = c$?

We conclude this paper by describing some topics for further study. If G is a connected graph with $V(G) = \{v_1, v_2, \dots, v_n\}$, then the ordered partition $\Pi = \{S_1, S_2, \dots, S_n\}$, where $S_i = \{v_i\}$ for $1 \leq i \leq n$, into singleton subsets of $V(G)$ is always a resolving partition of $V(G)$. Since $\langle S_i \rangle$ is trivially acyclic for each i ($1 \leq i \leq n$), it follows that Π is a resolving acyclic partition of $V(G)$ as well, and, consequently, $a_r(G)$ is defined. A similar argument shows that $\chi_r(G)$ is defined for every connected graph G .

This suggests a variety of concepts to study. If P is any graphical property possessed by a trivial subgraph of a connected graph G , then the ordered partition Π of $V(G)$ described above is said to satisfy property P and the P -partition number is defined. Among the various properties P , in addition to the property of being independent or being acyclic, are (1) the property of being a specified graph, say path, star, cycle, a linear forest (every component is a path), or a galaxy (every component is a star), (2) the property of being planar or hamiltonian, and (3) the property of having maximum degree or girth at most k for a fixed nonnegative integer k .

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