

# An Introduction to Maximum Sum Permutations of Graphs

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**ABSTRACT:** This paper introduces the problem of finding a permutation  $\phi$  on the vertex set  $V(G)$  of a graph  $G$  such that the sum of the distances from each vertex to its image under  $\phi$  is maximized. We let  $\mathcal{S}(G) = \max_{\phi \in \Pi} \sum_{v \in V(G)} d(v, \phi(v))$ , where the maximum is taken over all permutations  $\phi$  of  $V(G)$ . Explicit formulae for several classes of graphs as well as general bounds are presented.

**1. Introduction.** We assume  $G$  is a connected graph of order  $n$  with vertex set  $V(G) = \{v_1, v_2, \dots, v_n\}$  and edge set  $E(G)$ . As usual, the distance between two vertices  $v_i$  and  $v_j$  is denoted by  $d(v_i, v_j)$  and equals the minimum number of edges in a  $v_i$ - $v_j$  path. Because  $G$  is connected, each  $d(v_i, v_j)$  is finite. The eccentricity of a vertex  $v_i$  is  $e(v_i) = \max \{d(v_i, v_j) : 1 \leq j \leq n\}$ , the maximum distance from  $v_i$  to another vertex. The neighborhood of a vertex  $v_i$  is  $N(v_i) = \{v_j : v_i v_j \in E(G)\}$ , and the degree of  $v_i$  is  $\deg(v_i) = |N(v_i)|$  which is the number of vertices at distance one from  $v_i$ . The distance of vertex  $v_i$  is  $d(v_i) = \sum \{d(v_i, v_j) : 1 \leq j \leq n\}$ , the sum of the distances from  $v$  to all of the vertices in  $V(G)$ .

The radius of  $G$  is the minimum eccentricity of a vertex,  $r(G) = \min \{e(v) : v \in V(G)\}$ , and the center of  $G$  is the set of vertices of minimum eccentricity,  $C(G) = \{v_i \in V(G) : e(v_i) = r(G)\}$ . The median of  $G$  is  $M(G) = \{v_i \in V(G) : d(v_i) \leq d(v_j) \text{ for } 1 \leq j \leq n\}$ , the set of vertices of minimum distance.

Here we are interested in permutations  $\phi: V(G) \rightarrow V(G)$  which maximize the sum of the distances from each vertex to its image under  $\phi$ . Let  $\Pi$  denote the set of all permutations from  $V(G)$  onto  $V(G)$ . We define the maximum sum permutation distance  $\mathcal{S}(G)$  to be

$$\mathcal{S}(G) = \max_{\phi \in \Pi} \sum_{i=1}^n d(v_i, \phi(v_i)).$$

We identify permutation  $\phi: V(G) \rightarrow V(G)$  with the permutation  $\phi: \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$  by equating having  $\phi(v_i) = v_j$  with  $\phi(i) = j$ .

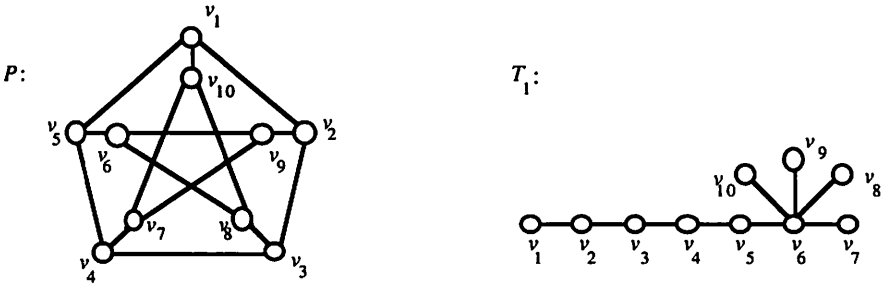


Figure 1. Petersen graph  $P$  and a tree  $T_1$ .

For example,  $\phi = (1, 3, 5, 2, 4) (6, 7, 8, 9, 10)$  satisfies  $\phi(v_1) = v_3$ ,  $\phi(v_2) = v_4$ ,  $\phi(v_3) = v_5$ ,  $\phi(v_4) = v_1$ ,  $\phi(v_5) = v_2$ ,  $\phi(v_6) = v_7$ ,  $\phi(v_7) = v_8$ ,  $\phi(v_8) = v_9$ ,  $\phi(v_9) = v_{10}$ , and  $\phi(v_{10}) = v_6$ . For the Petersen graph  $P$  in Figure 1,  $\sum_{i=1}^{10} d(v_i, \phi(v_i)) = 10 \cdot 2 = 20 = \mathcal{S}(P)$ , while for tree  $T_1$  of Figure 1,  $\sum_{i=1}^{10} d(v_i, \phi(v_i)) = 2 + 2 + 2 + 3 + 3 + 1 + 2 + 2 + 2 + 1 = 20$ . Note that for  $\psi = (1, 6, 3, 7, 5, 9) (2, 10, 4, 8)$  and tree  $T_1$ , we have  $\sum_{i=1}^{10} d(v_i, \psi(v_i)) = 5 + 5 + 4 + 3 + 2 + 3 + 2 + 5 + 6 + 3 = 38$ . As will be shown,  $\mathcal{S}(T_1) = 38$ .

**2. Examples and bounds.** We can precisely bound  $\mathcal{S}(G)$  for graphs  $G$  of order  $n$ .

**Theorem 1.** For every connected graph  $G$  of order  $n \geq 2$ , we have  $n \leq \mathcal{S}(G) \leq \lfloor n^2/2 \rfloor$  and these bounds are sharp.

**Proof.** Let  $\phi$  be any derangement of  $V(G)$ , that is,  $\phi(v_i) \neq v_i$  for  $1 \leq i \leq n$ . Then  $\mathcal{S}(G) \geq \sum_{i=1}^n d(v_i, \phi(v_i)) \geq n \cdot 1 = n$ . For a median vertex  $v \in M(G)$  we have  $d(v) \leq n^2/4$ . Let  $\phi$  be a maximum sum permutation of  $G$ . Then  $\mathcal{S}(G) = \sum_{i=1}^n d(v_i, \phi(v_i)) \leq \sum_{i=1}^n (d(v_i, v) + d(v, \phi(v_i))) \leq 2 \cdot d(v) \leq 2(n^2/4) = n^2/2$ .

Clearly for complete graph  $K_n$  we have  $\mathcal{S}(K_n) = n$  for  $n \geq 2$ . For path  $P_n$  let  $\phi$  be the permutation with  $\phi(v_i) = v_{n-i+1}$  for  $1 \leq i \leq n$ , and it is easily verified that  $\sum_{i=1}^n d(v_i, \phi(v_i)) = \lfloor n^2/2 \rfloor$ .  $\square$

We note that for cycles  $C_n$  we have  $\mathcal{S}(C_n) = n^2/2$  when  $n$  is even and  $\mathcal{S}(C_n) = n^2/2 - n/2$  if  $n$  is odd. Complete graph  $K_n$  is the only order  $n$  graph  $G$  with  $\mathcal{S}(G) = n$ .

For the eccentricity  $e(v_i)$  we clearly have  $d(v_i, \phi(v_i)) \leq e(v_i)$ .

Theorem 2.  $\mathcal{S}(G) \leq \sum_{i=1}^n e(v_i)$ .  $\square$

For example, the  $k$ -cube  $Q_k$  has  $V(Q_k)$  equal to the set of all binary  $k$ -tuples with two vertices adjacent if and only if they differ in exactly one position. Each vertex has a unique vertex farthest from it,  $d((e_1, e_2, \dots, e_k), (1-e_1, 1-e_2, \dots, 1-e_k)) = k$ . Hence, for  $Q_k$  we have  $n = 2^k$ , and  $\mathcal{S}(Q_k) = k \cdot 2^k$ .

For each edge  $e$  in a tree  $T$ , let  $w(e)$  be the order of the smaller of the two components of  $T - e$ .

Theorem 3. For any tree  $T$ , we have  $\mathcal{S}(T) = \sum_{e \in E(T)} 2 \cdot w(e)$ .

Proof. Note that an edge  $e$  is on the unique path from  $v_i$  to  $\phi(v_i)$  for a permutation  $\phi: V(T) \rightarrow V(T)$  if and only if  $v_i$  and  $\phi(v_i)$  are in different components of  $T - e$ . It follows that the maximum number of  $v_i - \phi(v_i)$  paths containing  $e$  is  $2 \cdot w(e)$ . Thus  $\mathcal{S}(T) = \max_{\phi \in \Pi} \sum_{i=1}^n d(v_i, \phi(v_i)) \leq \sum_{e \in E(T)} 2 \cdot w(e)$ .

If  $M(T) = \{u, v\}$  let  $T_u$  and  $T_v$  be the components of  $T - uv$  containing  $u$  and  $v$ , respectively. We have  $|V(T_u)| = |V(T_v)|$ , and so there is a permutation  $\phi: V(T) \rightarrow V(T)$  with  $\phi(v_i) \in V(T_v)$  if and only if  $v_i \in V(T_u)$ . For such a  $\phi$  every edge  $e$  is on  $2 \cdot w(e)$  of the  $v_i - \phi(v_i)$  paths, and so  $\mathcal{S}(T) = \sum_{e \in E(T)} 2 \cdot w(e)$ .

If the median consists of one vertex,  $M(T) = \{u\}$ , then each component of  $T - u$  has less than  $n/2$  vertices. We can define permutation  $\phi: V(T) \rightarrow V(T)$  such that every edge is on  $2 \cdot w(e)$  of the  $v_i - \phi(v_i)$  paths iteratively, as follows. Let  $C_1$  and  $C_2$  be the two largest components of  $T - u$ . Select  $x \in V(C_1)$  and  $y \in V(C_2)$  such that  $x$  and  $y$  are endpoints of  $T$ . Let  $\phi(x) = y$  and  $\phi(y) = x$ , replace  $T$  by  $T - x - y$  and iterate this process until only one or

two vertices remain. If it is just  $u$ , let  $\phi(u) = u$ , and if it is  $u$  and  $w$  then let  $\phi(u) = w$  and  $\phi(w) = u$ .  $\square$

**3. Distance-separating edge partitions.** In this concluding section we generalize Theorem 3. An edge set  $F \subseteq E(G)$  is called distance-separating if  $G - F$  has exactly two components, say with vertex sets  $V_1$  and  $V_2$ , and every path from any vertex  $x \in V_1$  to any vertex  $y \in V_2$  of length  $d(x, y)$  contains exactly one edge in  $F$ . Note that in trees edge set  $F$  is distance-separating if and only if  $F$  consist of exactly one edge. We let  $w(F)$  be the order of the smaller of the two components of  $G - F$ . A distance-separating edge partition for  $G$  is a collection  $\{E_1, E_2, \dots, E_t\}$  where each  $E_i$  is a distance-separating edge set,  $\bigcup_{i=1}^t E_i = E(G)$ , and  $1 \leq i < j \leq t$  implies  $E_i \cap E_j = \emptyset$ .

**Theorem 4.** Assume  $\{E_1, E_2, \dots, E_t\}$  is a distance-separating edge partition of graph  $G$ . Then  $\mathcal{S}(G) \leq \sum_{i=1}^t 2 \cdot w(E_i)$ .

**Proof.** Let  $V_i^1$  and  $V_i^2$  be the vertex sets of the components of  $G - E_i$  with  $w(E_i) = |V_i^1| \leq |V_i^2|$ . For any permutation  $\phi: V(G) \rightarrow V(G)$  consider a collection  $\{P_1, P_2, \dots, P_n\}$  where  $P_i$  is a  $v_i$ - $\phi(v_i)$  path of length  $d(v_i, \phi(v_i))$ . For each  $e \in E(G)$ , let  $f(e)$  be the number of paths  $P_i$  that contain  $e$ , and let  $f(E_i) = \sum_{e \in E_i} f(e)$ . We have  $\sum_{j=1}^n d(v_j, \phi(v_j)) = \sum_{e \in E(G)} f(e) = \sum_{i=1}^t f(E_i) \leq \sum_{i=1}^t 2 \cdot w(E_i)$ .

We note that  $\mathcal{S}(G) = \sum_{i=1}^t 2 \cdot w(E_i)$  if and only if there is a permutation  $\phi: V(G) \rightarrow V(G)$  such that for  $1 \leq i \leq t$  we have  $x \in V_i^1$  implying that  $\phi(x) \in V_i^2$ .  $\square$

For example, consider the grid graph  $G_{r,s}$  which is the cartesian product  $P_r \times P_s$  of paths  $P_r$  and  $P_s$ . Theorem 4 can be used to see that  $\mathcal{S}(G_{r,s}) = s \lfloor r^2/2 \rfloor + r \lfloor s^2/2 \rfloor$ .

We have further results for maximum sum permutations of graphs under study.