

# A square-covering problem

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## Abstract

Erdős and Soifer [3] and later Campbell and Staton [1] considered a problem which was a favorite of Erdős [2]: Let  $S$  be a unit square. Inscribe  $n$  squares with no common interior point. Denote by  $e_1, e_2, \dots, e_n$  the side lengths of these squares. Put  $f(n) = \max \sum_{i=1}^n e_i$ . And they discussed the bounds for  $f(n)$ . In this paper, we consider its dual problem— covering a unit square with squares.

**keywords:** packing, minimal covering

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## 1. Introduction

Erdős and Soifer [3] and later Campbell and Staton [1] considered a problem which was a favorite of Erdős [2]: Let  $S$  be a unit square. Inscribe  $n$  squares with no common interior point. Denote by  $e_1, e_2, \dots, e_n$  the side length of these squares. Put  $f(n) = \max \sum_{i=1}^n e_i$ . They discussed the bounds for  $f(n)$ . Inspired by [3], we discuss a problem on square-covering and give corresponding functions  $g_i(n)(i = 1, 2)$ .

First, we give the definition of a minimal square-covering.

**Definition 1.1.** Let  $S$  be a unit square. If  $n$  squares  $S_1, S_2, \dots, S_n$  cover  $S$ , in such a way which satisfies:

(1) Each  $S_i$  has side of length  $s_i(0 < s_i < 1)$  and is placed so that its sides are parallel to those of  $S$ ;

(2) Each  $S_i$  can not be smaller; that is, there does not exist any  $S_{i1} \subset S_i$  such that  $\{S_j, j = 1, 2, \dots, i-1, i+1, \dots, n\} \cup \{S_{i1}\}$  can cover  $S$  admitting translation.

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We call this kind of covering a minimal square-covering.

With this definition of minimal square covering, define:

$$g_1(n) = \min \sum_{i=1}^n s_i, \quad g_2(n) = \max \sum_{i=1}^n s_i.$$

When  $n \leq 3$ , since  $0 < s_i < 1$ , each  $S_i (i = 1, 2, \dots, n)$  can only cover one corner of a unit square, but it has four corners, so  $S_1, S_2, \dots, S_n$  can not cover  $S$ . That is, when  $n \leq 3$ ,  $g_i(n) (i = 1, 2)$  has no meaning. So in the following, let  $n \geq 4$ .

## 2. The bounds of $g_1(n)$

**Proposition 2.2.**  $g_1(n) \geq 2$ .

*Proof.* Obviously, if  $n$  squares cover the unit square, they must cover its two opposite sides. Since  $0 < s_i < 1$ , no square can cover the points of two opposite sides in the same time, so  $g_1(n) \geq 2$  must hold.  $\square$

**Theorem 2.3.** When  $n$  is even,  $g_1(n) \leq 3 - \frac{4}{n}$ .

*Proof.* Consider a minimal square-covering of a unit square  $S$  with  $n - 2$  squares  $S_2, S_3, \dots, S_{n-1}$  each of which has side of length  $x$ , a square  $S_1$  with side of length  $1 - x$ , and a square  $S_n$  with side of length  $1 - (\frac{n}{2} - 1)x$ . Since  $s_1 + s_n \geq 1$ , we have  $x \leq \frac{2}{n}$ . When  $n = 6$ , see Figure 1 for the placement. It's easy to see this is a minimal-square covering. So by the definition of  $g_1(n)$ ,  $g_1(n) \leq (1 - x) + (n - 2)x + [1 - (\frac{n}{2} - 1)x] = 2 + (\frac{n}{2} - 2)x \leq 3 - \frac{4}{n}$ .

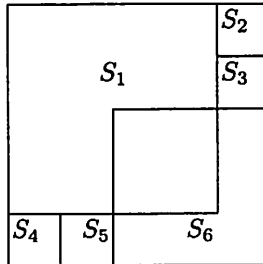


Figure 1: A unit square covered by six squares

$\square$

**Corollary 2.4.**  $g_1(4) = 2$ .

**Theorem 2.5.** When  $n$  is odd,  $g_1(n) < \frac{5}{2} - \frac{1}{2(2n-7)}$ .

*Proof.* Consider a minimal square-covering of a unit square  $S$  with  $n - 3$  squares  $S_2, S_3, \dots, S_{n-2}$  each of which has side of length  $x$ , a square  $S_1$  with side of length  $1 - x$ , a square  $S_n$  with side of length  $1 - (n - 3)x$  and a square with side of length  $(n - 3)x$ . Since  $s_1 + s_n \geq 1$  and  $s_{n-1} < s_{n-2} + s_n$ , we have  $x < \frac{1}{2n-7}$ . When  $n = 7$ , see Figure 2 for the placement. It's easy to see this is a minimal-square covering. So by the definition of  $g_1(n)$ ,  $g_1(n) \leq (1 - x) + (n - 3)x + 1 - (n - 3)x + (n - 3)x = 2 + (n - 4)x < 2 + \frac{n-4}{2n-7} = \frac{5}{2} - \frac{1}{2(2n-7)}$ .

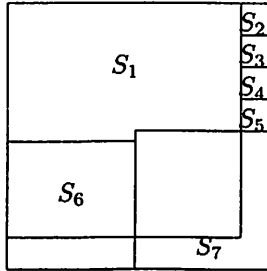


Figure 2: A unit square covered by seven squares

□

### 3. The bounds of $g_2(n)$

**Proposition 4.6.**  $g_2(k^2) \geq k$ .

*Proof.* It is obvious that all the  $S_i (i = 1, 2, \dots, n)$  must be equal, whence  $ne_i = 1$ . That is, for each  $i$ ,  $e_i = \frac{1}{\sqrt{n}}$ . So  $n$  is a perfect square, say,  $n = k^2$ , and the optimal covering is the standard  $n$ -covering, that is, a  $k \times k$  grid of squares each of which has side of length  $\frac{1}{k}$ . Obviously,  $k^2 - 1$  such squares can not cover the unit square. By the definition of  $g_2(n)$ ,  $g_2(k^2) \geq k$ . □

**Proposition 4.7.**  $g_2(k^2 + 1) > k$ .

*Proof.* From a standard  $k^2$ -covering, remove a  $2 \times 2$  grid and replace it with five squares  $S_{i1}, S_{i2}, \dots, S_{i5}$  covering the same area, and which are placed as in Figure 3 such that  $S_{i1}$  is the largest square of  $\{S_{ij} \mid j = 1, 2, \dots, 5\}$  and  $s_{i2} = s_{i3} = \frac{2}{k} - s_{i1}$ ,  $s_{i4} = \frac{2}{k} - 2(\frac{2}{k} - s_{i1}) = 2s_{i1} - \frac{2}{k}$ ,  $s_{i5} = 2(\frac{2}{k} - s_{i1}) = \frac{4}{k} - 2s_{i1}$ ,  $0 < s_{ij} < \frac{2}{k} (j = 1, 2, \dots, 5)$ ,  $s_{i4} + s_{i1} \geq \frac{2}{k}$ ,  $s_{i5} < s_{i4} + s_{i3}$ . The result is a covering with  $k^2 + 1$  squares, the sum of whose lengths is  $s = k - \frac{4}{k} + s_{i1} + 2(\frac{2}{k} - s_{i1}) + (2s_{i1} - \frac{2}{k}) + (\frac{4}{k} - 2s_{i1}) = k + \frac{2}{k} - s_{i1}$ . The inequalities above imply that  $\frac{3}{2k} < s_{i1} < \frac{2}{k}$ , so  $s > k + \frac{2}{k} - \frac{2}{k} = k$ .

Obviously, this covering is a minimal covering, so we have  $g_2(k^2 + 1) > k$ . □

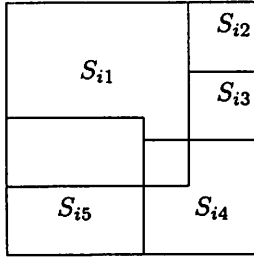


Figure 3: A  $2 \times 2$  grid covered by five squares

**Proposition 4.8.**  $g_2(k^2 - 1) \geq k - \frac{3}{k}$ .

*Proof.* From a standard  $k^2$ -covering, remove a  $3 \times 3$  grid and replace it with eight squares  $S_{i1}, S_{i2}, \dots, S_{i8}$  covering the same area, and which are placed as in Figure 4 such that  $S_{i1}$  is the largest square of  $\{S_{ij} \mid j = 1, 2, \dots, 8\}$  and  $s_{i2} = s_{i3} = s_{i4} = s_{i5} = s_{i6} = s_{i7} = \frac{3}{k} - s_{i1}$ ,  $s_{i8} = \frac{3}{k} - 3(\frac{3}{k} - s_{i1}) = 3s_{i1} - \frac{6}{k}$ ,  $0 < s_{ij} < \frac{3}{k}$  ( $j = 1, 2, \dots, 8$ ),  $s_{i8} + s_{i1} \geq \frac{3}{k}$ . The result is a covering with  $k^2 - 9 + 8 = k^2 - 1$  squares, the sum of whose lengths is  $s = k - \frac{9}{k} + s_{i1} + 6(\frac{3}{k} - s_{i1}) + (3s_{i1} - \frac{6}{k}) = k + \frac{3}{k} - 2s_{i1}$ . The inequalities above imply that  $\frac{9}{4k} \leq s_{i1} < \frac{2}{k}$ , so  $s \geq k + \frac{3}{k} - 2\frac{2}{k} = k - \frac{3}{k}$ .

Obviously, this covering is a minimal covering, so we have  $g_2(k^2 - 1) \geq k - \frac{3}{k}$ .

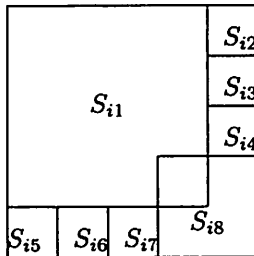


Figure 4: A  $3 \times 3$  grid covered by eight squares

□

When neither  $n - 1$  nor  $n + 1$  is a perfect square, we have the following result:

**Theorem 4.9.** *If neither  $n - 1$  nor  $n + 1$  is a perfect square, then  $g_2(n) > \sqrt{n - 1}$ .*

In the proof of Theorem 4.9, we borrow the main idea of [1].

*Proof.* When  $n = k^2$ , by Theorem 4.6,  $g_2(n) \geq \sqrt{n} > \sqrt{n-1}$ .

When  $n \neq k^2$ ,  $k$  must lie between two perfect squares of different parity. That is, there is an integer  $k$  such that  $k^2 < n < (k+1)^2$ ,  $n - k^2$  and  $(k+1)^2 - n$  have different parity. Consider the values of  $n$  where  $k^2 + 1 < n < (k+1)^2 - 1$ . There are two cases which provide upper bounds for all  $n$  on the interval  $[k^2 + 2, (k+1)^2 - 2]$ :

Case 1.  $(k+1)^2 - n$  is odd. Say,  $(k+1)^2 - n = 2a + 1$  ( $a \geq 1$ ),  $k^2 < n \leq (k+1)^2 - 3$ . From a standard  $(k+1)^2$ -covering of  $S$ , remove an  $(a+1) \times (a+1)$  grid and replace it with an  $a \times a$  grid covering the same area. The result is a covering with  $(k+1)^2 - (a+1)^2 + a^2 = n$  squares, the sum of whose lengths is

$$[(k+1)^2 - (a+1)^2] \frac{1}{k+1} + a^2 \left(\frac{a+1}{a}\right) \left(\frac{1}{k+1}\right) = k+1 - \frac{a+1}{k+1}.$$

Obviously, no one of these  $n$  squares can be smaller. So

$$g_2(n) \geq k+1 - \frac{a+1}{k+1}, g_2^2(n) \geq (k+1 - \frac{a+1}{k+1})^2 > n-1.$$

That is,  $g_2(n) > \sqrt{n-1}$ .

Case 2.  $n - k^2$  is odd. Say,  $n - k^2 = 2a - 1$  ( $a \geq 2$ ),  $k^2 + 3 \leq n < (k+1)^2$ . From a standard  $k^2$ -covering of  $S$ , remove an  $(a-1) \times (a-1)$  grid and replace it with an  $a \times a$  grid covering the same area. The result is a covering with  $k^2 - (a-1)^2 + a^2 = k^2 + 2a - 1 = n$  squares of the unit square  $S$ . The sum of the lengths of sides is

$$[k^2 - (a-1)^2] \frac{1}{k} + a^2 \left(\frac{a-1}{a}\right) \left(\frac{1}{k}\right) = k + \frac{a-1}{k}.$$

Obviously, no  $n-1$  squares of these  $n$  squares can cover  $S$ . So

$$g_2(n) \geq k + \frac{a-1}{k}, g_2^2(n) \geq (k + \frac{a-1}{k})^2 = k^2 + 2a - 1 + (\frac{a-1}{k})^2 - 1 > n-1.$$

That is,  $g_2(n) > \sqrt{n-1}$ . □

The following lemma is a well-known result [4]:

**Lemma 4.10.** *Finitely many squares whose total area is equal to 3 can cover a unit square.*

**Theorem 4.11.**  $g_2(n) \leq 3\sqrt{n}$ .

*Proof.* Let  $\{S_i\}_{i=1}^n$  be a minimal covering of the unit square  $S$ , and  $s_i$  denote the length of the side of  $S_i$  ( $i = 1, 2, \dots, n$ ). We first prove that  $\sum_{i=1}^n s_i^2 \leq 3$ . Otherwise, if  $\sum_{i=1}^n s_i^2 > 3$ , there exists a  $S_{i_1} \subset S_i$ , such that

$s_{i_1} < s_i$  and  $s_{i_1}^2 + \sum_{j=1}^{i-1} s_j^2 + \sum_{j=i+1}^n s_j^2 \geq 3$ . By Lemma 4.10, it is obvious that  $S_1, S_2, \dots, S_{i-1}, S_{i_1}, S_{i+1}, \dots, S_n$  can cover the unit square  $S$ , which contradicts the definition of a minimal covering. So  $\sum_{i=1}^n s_i^2 \leq 3$ .

Let  $\mathbf{s}$  be the vector  $(s_1, s_2, \dots, s_n)$ , and let  $\mathbf{v}$  be the vector  $(1, 1, \dots, 1)$ . Now  $\sum_{i=1}^n s_i \leq \|\mathbf{s}\| \|\mathbf{v}\| \leq \sum_{i=1}^n s_i^2 n^{\frac{1}{2}} = n^{\frac{1}{2}} \sum_{i=1}^n s_i^2 \leq 3\sqrt{n}$ , so  $g_2(n) \leq 3\sqrt{n}$ .

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