

# On Distance Connected Domination Numbers of Graphs \*

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**Abstract** Let  $k$  be a positive integer and  $G = (V, E)$  be a connected graph of order  $n$ . A set  $D \subseteq V$  is called a  $k$ -dominating set of  $G$  if each  $x \in V(G) - D$  is within distance  $k$  from some vertex of  $D$ . A connected  $k$ -dominating set is a  $k$ -dominating set that induces a connected subgraph of  $G$ . The connected  $k$ -domination number of  $G$ , denoted by  $\gamma_k^c(G)$ , is the minimum cardinality of a connected  $k$ -dominating set. Let  $\delta$  and  $\Delta$  denote the minimum and the maximum degree of  $G$ , respectively. This paper establishes that  $\gamma_k^c(G) \leq \max\{1, n - 2k - \Delta + 2\}$ , and  $\gamma_k^c(G) \leq (1 + o_\delta(1))n \frac{\ln(m(\delta+1)+2-t)}{m(\delta+1)+2-t}$ , where  $m = \lceil \frac{k}{3} \rceil$ ,  $t = 3\lceil \frac{k}{3} \rceil - k$ , and  $o_\delta(1)$  denotes a function that tends to 0 as  $\delta \rightarrow \infty$ . The later generalizes the result of Caro *et al*'s in [Connected domination and spanning trees with many leaves. SIAM J. Discrete Math. 13 (2000), 202-211] for  $k = 1$ .

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## 1 Introduction

For terminology and notation on graph theory not given here, the reader is referred to [2] or [13]. Let  $G = (V, E)$  be a finite simple graph with vertex set  $V = V(G)$  and edge set  $E = E(G)$ . The order, the maximum degree and the minimum degree of vertices of  $G$  are denoted by  $n(G)$ ,  $\Delta(G)$  and  $\delta(G)$ , respectively. The distance  $d_G(x, y)$  between two vertices

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$x$  and  $y$  is the length of a shortest  $xy$ -path in  $G$ . For  $S \subseteq V(G)$ ,  $G[S]$  denotes the subgraph of  $G$  induced by  $S$ , and for  $v \in V(G)$ ,  $d_G(v, S) = \min_{u \in V(S)} \{d_G(v, u)\}$ . The eccentricity  $e_G(v)$  of  $v$  is  $\max_{x \in V(G)} \{d_G(v, x)\}$ . The radius  $\text{rad}(G)$  is the smallest eccentricity of a vertex in  $G$ . Let  $k$  be a positive integer. For every vertex  $x \in V(G)$ , the  $k$ -neighborhood  $N_k(x)$  of  $x$  is defined by  $N_k(x) = \{y \in V(G) : d_G(x, y) \leq k, x \neq y\}$ , and  $N_1(x)$  is usually called the neighborhood of  $x$  in  $G$ .

A set  $D$  of vertices in  $G$  is called a  $k$ -dominating set of  $G$  if every vertex of  $V(G) - D$  is within distance  $k$  from some vertex of  $D$ . A  $k$ -dominating set  $D$  is called to be connected if  $G[D]$  is connected. The minimum cardinality among all  $k$ -dominating sets (resp. connected  $k$ -dominating sets) of  $G$  is called the  $k$ -domination number (resp. connected  $k$ -domination number) of  $G$  and is denoted by  $\gamma_k(G)$  (resp.  $\gamma_k^c(G)$ ). The concept of the  $k$ -dominating set was first introduced by Chang and Nemhauser [4, 5].

Since the distance versions of domination have a strong background of applications, many efforts have been made by several authors to consider the distance parameters (see, for example, [4] ~ [10], [12, 14]).

It is quite difficult to determine the value of  $\gamma_k(G)$  or  $\gamma_k^c(G)$  for any given graph  $G$ . In this paper, we prove that for any nontrivial connected graph  $G$  with order  $n$ ,  $\gamma_k^c(G) = \min \gamma_k^c(T)$ , where the minimum is taken over all spanning trees  $T$  of  $G$ . We also get two upper bounds for  $\gamma_k^c(G)$  in terms of the maximum degree  $\Delta = \Delta(G)$ , that is,

$$\gamma_k^c(G) \leq \max\{1, n - 2k - \Delta + 2\},$$

and the minimum degree  $\delta = \delta(G)$ , that is,

$$\gamma_k^c(G) \leq (1 + o_\delta(1))n \frac{\ln[m(\delta + 1) + 2 - t]}{m(\delta + 1) + 2 - t},$$

where  $m = \lceil \frac{k}{3} \rceil$ ,  $t = 3\lceil \frac{k}{3} \rceil - k$ , and  $o_\delta(1)$  denotes a function that tends to 0 as  $\delta \rightarrow \infty$ . The later generalizes the result of Caro *et al*'s [3] for  $k = 1$ , that is,

$$\gamma_1^c(G) \leq (1 + o_\delta(1))n \frac{\ln(\delta + 1)}{\delta + 1}.$$

The method used here is a generalization and refinement of theirs.

## 2 Elementary Results

**Theorem 1** Let  $G$  be a nontrivial connected graph, and  $k$  be a positive integer. Then  $\gamma_k^c(G) = \min \gamma_k^c(T)$ , where the minimum is taken over all spanning trees  $T$  of  $G$ .

*Proof* Let  $G$  be a nontrivial connected graph and  $T$  be a spanning tree of  $G$ . Then any connected  $k$ -dominating set of  $T$  is also a connected  $k$ -dominating set of  $G$ . Therefore,  $\gamma_k^c(G) \leq \gamma_k^c(T)$ . Thus we have that  $\gamma_k^c(G) \leq \min \gamma_k^c(T)$ , where the minimum is taken over all spanning trees  $T$  of  $G$ .

Now we show the reverse inequality. If  $G$  is a tree, then the theorem holds trivially. So we may assume that  $G$  is a connected graph containing cycles. Let  $D$  be a minimum connected  $k$ -dominating set of  $G$  and  $C$  be a cycle in  $G$ . If we can prove that  $D$  is also a connected  $k$ -dominating set of  $G - e$  for some cycle edge  $e \in E(C)$ , then  $\gamma_k^c(G - e) \leq |D| = \gamma_k^c(G)$ . By applying this process a finite number of times, we have  $\gamma_k^c(T) \leq \gamma_k^c(G)$  for some spanning tree  $T$  of  $G$ . Thus, we have that  $\min \gamma_k^c(T) \leq \gamma_k^c(G)$ , where the minimum is taken over all spanning trees  $T$  of  $G$ .

If  $V(C) \subseteq V(D)$ , then obviously  $G[D] - e$  for any  $e \in E(C)$  is also connected and the vertices in  $V(G) - D$  are also all within distance  $k$  to  $D$ .

If  $V(C) \not\subseteq V(D)$ , then we select an edge  $xy$  in  $C$  such that  $d_G(x, D) + d_G(y, D) = \max\{d_G(u, D) + d_G(v, D) : uv \in E(C)\}$ . Now we will show that  $D$  is a connected  $k$ -dominating set of  $G - \{xy\}$ .

First for any two adjacent vertices  $u$  and  $v$  in  $G$ , we have  $|d_G(u, D) - d_G(v, D)| \leq 1$ . Then if  $w$  is a vertex in  $V(C)$  such that  $d_G(w, D) = \max\{d_G(v, D) : v \in V(C)\}$ , we have that  $w = x$  or  $w = y$ . Without loss of generality, suppose that  $d_G(x, D) = \max\{d_G(v, D) : v \in V(C)\}$ .

Let  $z$  be another neighbor of  $x$  different from  $y$  in  $V(C)$ . So we immediately have that  $d_G(z, D) \leq d_G(y, D)$ . Thus, we get the distance between a vertex in  $V(G) - D$  and  $D$  is not influenced by deleting the edge  $\{xy\}$ . That is to say,  $d_{G-\{xy\}}(v, D) = d_G(v, D)$  for all vertices  $v$  in  $V(G)$ . Hence,  $D$  is also a connected  $k$ -dominating set of  $G - e$  for some cycle edge  $e$ . ■

**Proposition 2** Let  $G = (V, E)$  be a nontrivial connected graph, and  $k$  be a positive integer. If  $\text{rad}(G) \leq k$ , then  $\gamma_k^c(G) = 1$ .

### 3 Main Results

**Theorem 3** Let  $G$  be a connected graph of order  $n \geq 2$  with maximum degree  $\Delta = \Delta(G)$ , and  $k$  be a positive integer, then

$$\gamma_k^c(G) \leq \max\{1, n - 2k - \Delta + 2\}.$$

*Proof* By Theorem 1, it is sufficient to show that  $\gamma_k^c(T) \leq \max\{1, n - 2k - \Delta + 2\}$ , for any spanning tree  $T$  with maximum degree  $\Delta = \Delta(T)$ .

If  $\text{rad}(T) \leq k$ , then by Theorem 2, we get  $\gamma_k^c(T) = 1$ . So we may assume that  $\text{rad}(T) > k$ . Let  $P$  be a longest path in  $T$  with end-vertices  $u$  and

$v$ . Then there exists two vertices  $x$  and  $y$  of  $P$  such that  $d_T(x, u) = k$  and  $d_T(y, v) = k$ . Let  $P_{xy}$  be the  $xy$ -subpath of  $P$ , and let  $D' = V(P) - V(P_{xy})$ . Let  $D = V(T) - (D' \cup \mathcal{L}(T))$ , where  $\mathcal{L}(T)$  is the set of leaves of  $V(T)$ . Thus  $D$  must contain a connected  $k$ -dominating set of  $T$ . Since  $u, v \in D' \cap \mathcal{L}(T)$ , and  $\mathcal{L}(T) \geq \Delta$ , we have

$$\begin{aligned} \gamma_k^c(T) &\leq |V(T)| - |D' \cup \mathcal{L}(T)| \\ &\leq |V(T)| - |D'| - |\mathcal{L}(T)| + |D' \cap \mathcal{L}(T)| \\ &\leq n - 2k - \Delta + 2 \end{aligned}$$

as required. ■

We use probabilistic method to give an upper bound of  $\gamma_k^c(G)$  in terms of the minimum degree  $\delta = \delta(G)$  below. This bound improves the results of Caro *et al* [3] for  $k = 1$  and the method is a generalization and refinement of theirs.

For an event  $A$  and for a random variable  $Z$  of an arbitrary probability space,  $P[A]$  and  $E[Z]$  denote the probability of  $A$ , the expectation of  $Z$ , respectively.

**Lemma 4** (Xu, Tian and Huang [14]) Let  $S$  be a  $k$ -dominating set of a connected graph  $G$ . If  $G[S]$  has  $h$  components, then

$$\gamma_k^c(G) \leq |S| + 2(h - 1)k.$$

**Theorem 5** Let  $G$  be a nontrivial connected graph of order  $n$  with minimum degree  $\delta$ , then

$$\gamma_k^c(G) < n \frac{72k + 20km^2 + 17 + 0.5\sqrt{\ln q} + \ln q}{q}, \quad (1)$$

where  $q = m(\delta + 1) + 2 - t$ ,  $m = \lceil \frac{k}{3} \rceil$  and  $t = 3\lceil \frac{k}{3} \rceil - k$ .

*Proof* Let  $k = 3m - t$ , where  $m \geq 1$ ,  $0 \leq t \leq 2$ . For  $\delta(G) < 72\lfloor \frac{k}{m} \rfloor + 20km$ , we immediately have  $\gamma_k^c(G) \leq n$ , and the theorem holds. We assume that  $\delta(G) \geq 72\lfloor \frac{k}{m} \rfloor + 20km \geq 92$  below. Let  $p = \frac{\ln q}{q}$ , where  $q = m(\delta + 1) + 2 - t$ , and let us pick, randomly and independently, each vertex of  $V$  with probability  $p$ . Let  $X$  be the set of vertices picked. Let  $Y$  be the random set of all vertices that are not picked and have no  $k$ -neighbors in  $X$ . By the choice of  $Y$ ,  $X \cup Y$  is a  $k$ -dominating set of  $G$ .

**Claim 1**  $d_G(X, Y) = k + 1$ .

*Proof of Claim 1.* It is clear from the choice of  $Y$  that  $d_G(X, Y) \geq k + 1$ . Now let  $a \in X$ ,  $b \in Y$  be two vertices whose distance in  $G$  is the smallest, that is,  $d_G(a, b) = d_G(X, Y)$ . Let  $P$  be any shortest path from  $a$  to  $b$  and let  $v$  be the second-last vertex on  $P$ . Then  $v \notin Y$ . If  $d_G(a, b) \geq k + 2$ , then  $v$  has no  $k$ -neighbors in  $X$ . By definition of  $Y$ , we should get  $v \in Y$ , a contradiction. ■

Let  $\alpha = |X|$ ,  $\beta = |Y|$  and  $P_{XY}$  denote one shortest path from  $X$  to  $Y$ . By Claim 1, we have  $|V(P_{XY})| = k + 2$ . Let  $\mu$  denote the number of components in  $G[X]$ . Then  $X \cup Y \cup V(P_{XY})$  is a subgraph of  $G$  having at most  $\mu + \beta - 1$  components. By Lemma 4, we have

$$\gamma_k^c(G) \leq \alpha + \beta + k + 2(\mu + \beta - 1 - 1)k = \alpha + (2k + 1)\beta + 2k\mu - 3k.$$

In order to prove (1), it therefore suffices to show that with positive probability,

$$\alpha + (2k + 1)\beta + 2k\mu - 3k < n \frac{72k + 20km^2 + 17 + 0.5\sqrt{\ln q} + \ln q}{q}. \quad (2)$$

**Claim 2**  $|N_k(v)| \geq m(\delta + 1) + 1 - t$  for any  $v \in V(G)$ .

*Proof of Claim 2* Let  $X_i(v) = \{u \in V(G) : d_G(u, v) = i\}$ .

If  $v \in X \cup Y$ , then by  $d_G(X, Y) = k + 1$  and  $G$  is connected,  $X_i(v) \neq \emptyset$  for  $i = 1, \dots, k$ . Clearly,  $|X_1(v)| \geq \delta$ . For  $2 \leq i \leq k - 2$ , we have that  $|X_i(v)| + |X_{i+1}(v)| + |X_{i+2}(v)| \geq \delta + 1$ . In fact, for any  $u \in X_{i+1}(v)$ ,  $N_1(u) \subseteq X_i(v) \cup X_{i+1}(v) \cup X_{i+2}(v)$ , thus,  $|X_i(v)| + |X_{i+1}(v)| - 1 + |X_{i+2}(v)| \geq \delta$ . So, we have

$$\begin{aligned} |N_k(v)| &= |X_1(v)| + |X_2(v)| + \dots + |X_k(v)| \\ &\geq \delta + \left\lfloor \frac{k-1}{3} \right\rfloor (\delta + 1) + \left( k - 1 - 3 \left\lfloor \frac{k-1}{3} \right\rfloor \right) \\ &= \delta + (m-1)(\delta + 1) + (2-t) \\ &= m(\delta + 1) + 1 - t. \end{aligned}$$

Let  $v \in V(G) - (X \cup Y)$ . If  $d_G(v, Y) \geq k$  or  $d_G(v, X) \geq k$ , using the same discussion as above we get  $|N_k(v)| \geq m(\delta + 1) + 1 - t$ . Now suppose that  $d_G(v, Y) < k$  and  $d_G(v, X) < k$ . Since  $d_G(X, Y) = k + 1$ , there must exist a shortest path between a vertex  $a \in X$  and a vertex  $b \in Y$  through  $v$  such that  $d_G(a, b) \geq k + 1$ ,  $d_G(v, b) < k$  and  $d_G(a, v) < k$ . We only consider the worst case  $d_G(a, b) = k + 1$ , and let  $P_{ab}$  denote the shortest path from  $a$  to  $b$  passing through  $v$ .

Let  $v_1$  and  $v_2$  be two neighbors of  $v$  on  $P_{ab}$  from  $b$  to  $v$  and from  $a$  to  $v$ , respectively. Let  $d_G(b, v_1) = \ell_1$ ,  $d_G(a, v_2) = \ell_2$ . Thus,  $\ell_1 + \ell_2 = k - 1$ . We only consider three cases. The other one are analogue or immediate by symmetry.

If  $\ell_1 \equiv 1 \pmod{3}$ ,  $\ell_2 \equiv 1 \pmod{3}$ , then  $k \equiv 0 \pmod{3}$ , that is,  $k = 3m$ ,  $t = 0$ .

$$\begin{aligned} |N_k(v)| &\geq \delta + \left( \left\lfloor \frac{\ell_1}{3} \right\rfloor + \left\lfloor \frac{\ell_2}{3} \right\rfloor \right) (\delta + 1) + 2 \\ &= \delta + \frac{\ell_1 + \ell_2 - 2}{3} (\delta + 1) + 2 \end{aligned}$$

$$\begin{aligned}
&= \delta + \frac{k-3}{3}(\delta+1) + 2 \\
&= \delta + (m-1)(\delta+1) + 2 \\
&= m(\delta+1) + 1
\end{aligned}$$

If  $\ell_1 \equiv 1 \pmod{3}$ ,  $\ell_2 \equiv 2 \pmod{3}$ , then  $k \equiv 1 \pmod{3}$ , that is,  $k = 3m-2$ ,  $t = 2$ . Notice  $\ell_2 \equiv 2 \pmod{3}$  and  $d_G(v, a) < k$ , then  $N_1(a) \subseteq N_k(v)$ , thus  $|X_{\ell_2-1}(v_2)| + |X_{\ell_2}(v_2)| + |X_{\ell_2+1}(v_2)| \geq \delta + 1$ . So we have

$$\begin{aligned}
|N_k(v)| &\geq \delta + \left( \left\lfloor \frac{\ell_1}{3} \right\rfloor + \left\lfloor \frac{\ell_2}{3} \right\rfloor \right) (\delta+1) + 1 + (\delta+1) \\
&= \delta + \frac{\ell_1 - 1 + \ell_2 - 2}{3} (\delta+1) + \delta + 2 \\
&= \delta + \frac{k-4}{3} (\delta+1) + \delta + 2 \\
&= m(\delta+1) \\
&> m(\delta+1) + 1 - t
\end{aligned}$$

If  $\ell_1 \equiv 2 \pmod{3}$ ,  $\ell_2 \equiv 2 \pmod{3}$ , then  $k \equiv 2 \pmod{3}$ , that is,  $k = 3m-1$ ,  $t = 1$ . By the discussion as above, we also get  $|X_{\ell_1-1}(v_1)| + |X_{\ell_1}(v_1)| + |X_{\ell_1+1}(v_1)| \geq \delta + 1$ . Thus, we have,

$$\begin{aligned}
|N_k(v)| &\geq \delta + \left( \left\lfloor \frac{\ell_1}{3} \right\rfloor + \left\lfloor \frac{\ell_2}{3} \right\rfloor \right) (\delta+1) + 2(\delta+1) \\
&= \delta + \frac{\ell_1 - 2 + \ell_2 - 2}{3} (\delta+1) + 2\delta + 2 \\
&= \delta + \frac{k-5}{3} (\delta+1) + 2\delta + 2 \\
&= m(\delta+1) + \delta \\
&> m(\delta+1)
\end{aligned}$$

The Claim 2 follows. ■

**Claim 3**  $P \left[ \beta > 17 \frac{n}{q} \right] < 0.059$ .

*Proof of Claim 3* For each vertex  $v$ , the probability that  $v \in Y$  is that  $P[v \in Y] = (1-p)^{|N_k(v)|+1}$ . By Claim 2, we already have that  $|N_k(v)| \geq m(\delta+1) + 1 - t$  for any  $v \in V(G)$  and since  $\beta$  can be written as a sum of  $n$  indicator random variables  $\chi_v$ , where  $\chi_v = 1$  if  $v \in Y$  and  $\chi_v = 0$  otherwise, it follows that the expectation of  $\beta$  satisfies  $E[\beta] \leq n(1-p)^q$ . By using Taylor's formula,

$$\left( 1 - \frac{\ln q}{q} \right)^q < \left( e^{-\frac{\ln q}{q}} \right)^q = \frac{1}{q},$$

we have  $E[\beta] < \frac{n}{q}$ . By Markov's inequality, for any  $s > 0$ ,  $P[\beta > s] < \frac{E[\beta]}{s}$ , we have,

$$P\left[\beta > 17\frac{n}{q}\right] < \frac{1}{17} < 0.059$$

as required. ■

**Claim 4**  $P\left[\alpha > n\frac{\ln q}{q} + n\frac{0.5\sqrt{\ln q}}{q}\right] < 0.892.$

*Proof of Claim 4* Since  $\alpha$  can also be written as a sum of  $n$  indicator random variables that each having probability  $p$  of success, we also have  $E[\alpha] = np = n\frac{\ln q}{q}$ . We use an inequality attributed to Chernoff in [1], that is, for any  $s \geq 0$ :

$$P[\alpha > E[\alpha] + s] \leq \exp\left\{\frac{-s^2}{2(E[\alpha] + \frac{s}{3})}\right\}.$$

Take  $s = n\frac{0.5\sqrt{\ln q}}{q}$  to this inequality, we have

$$\begin{aligned} & P\left[\alpha > n\frac{\ln q}{q} + n\frac{0.5\sqrt{\ln q}}{q}\right] \\ & \leq \exp\left(-\frac{n}{8q + \frac{4}{3}q\frac{1}{\sqrt{\ln q}}}\right) \\ & < \exp\left(-\frac{1}{8 + 1.34\frac{1}{\sqrt{\ln[93m+2-t]}}}\right) \\ & \leq \exp\left(-\frac{1}{8 + 1.34\frac{1}{\sqrt{\ln 93}}}\right) < 0.892. \end{aligned}$$

Here  $n \geq |N_k(v)| + 1 \geq q$ . The Claim 4 follows. ■

Like [3], we say that a vertex  $v \in V(G)$  is weakly dominated if  $v$  has fewer than  $\frac{1}{8m^2} \ln q$  neighbors in  $X$ . Let  $N_1^X(v)$  denote the set of neighbors of  $v$  in  $X$ . Let  $\mathcal{D}$  denote the set of weakly dominated vertices in  $X$ .

**Claim 5**  $P\left[|\mathcal{D}| > 19n\frac{\ln q}{q^{1.34}}\right] < 0.047.$

*Proof of Claim 5* First we have, for any  $v \in V(G)$ ,

$$\begin{aligned} E[|N_1^X(v)|] &= |N_1(v)|p \geq \delta p \\ &= \frac{\delta}{q} \ln q \\ &\geq \frac{92}{93m+2-t} \ln q \\ &\geq \frac{92}{93m+2} \ln q, \end{aligned}$$

where  $\frac{\delta}{q}$  is an increasing function for  $\delta$ . By using linearity of expectation and another inequality of Chernoff [1], that is, for any  $s \geq 0$ ,

$$P[|N_1^X(v)| < E[|N_1^X(v)|] - s] \leq \exp\left(-\frac{s^2}{2E[|N_1^X(v)|]}\right),$$

we have,

$$\begin{aligned} & P\left[|N_1^X(v)| < \frac{1}{8m^2} \ln q\right] = P\left[m|N_1^X(v)| < \frac{1}{8m} \ln q\right] \\ & \leq P\left[m|N_1^X(v)| < \frac{(93m+2)}{8m \times 92} E[|N_1^X(v)|]\right] \\ & = P\left[m|N_1^X(v)| - E[m|N_1^X(v)|] < -\left(m - \frac{(93m+2)}{8m \times 92}\right) E[|N_1^X(v)|]\right] \\ & < \exp\left(-\frac{\left(m - \frac{(93m+2)}{8m \times 92}\right)^2 E^2[|N_1^X(v)|]}{2mE[|N_1^X(v)|]}\right) \\ & = \exp\left(-\frac{\left(m - \frac{(93m+2)}{8m \times 92}\right)^2 E[|N_1^X(v)|]}{2m}\right) \\ & \leq \exp\left(-\frac{46}{93m^2+2m} \left(m - \frac{(93m+2)}{8m \times 92}\right)^2 \ln q\right) \\ & \leq \exp\left(-\frac{46m^2}{93m^2+2m} \left(1 - \frac{(93m+2)}{8m^2 \times 92}\right)^2 \ln q\right) \\ & \leq \exp\left(-\frac{46}{95} \left(1 - \frac{95}{736}\right)^2 \ln q\right) \\ & \leq \left(\frac{1}{q}\right)^{0.367}. \end{aligned}$$

Since the event that a vertex  $v$  is picked into  $X$  is independent of the event that  $v$  is a weakly dominated vertex. Hence, the probability that a vertex is in  $X$  and is weakly dominated is,

$$\begin{aligned} & P\left[v \in X; |N_1^X(v)| < \frac{1}{8m^2} \ln q\right] \\ & = P[v \in X] \cdot P\left[|N_1^X(v)| < \frac{1}{8m^2} \ln q\right] \\ & \leq p \left(\frac{1}{q}\right)^{0.367}. \end{aligned}$$

Thus, we have

$$E[|\mathcal{D}|] \leq np \left(\frac{1}{q}\right)^{0.367} = n \frac{\ln q}{q^{1.367}}.$$

By Markov's inequality,

$$P\left[|\mathcal{D}| > 19n \frac{\ln q}{q^{1.34}}\right] < \frac{1}{19q^{0.027}} < \frac{1}{19 \times 93^{0.027}} < 0.047$$



as required. ■

From Claim 3, Claim 4 and Claim 5, we find that all of these events that

$$\begin{aligned}\alpha &\leq n \frac{\ln q}{q} + n \frac{0.5\sqrt{\ln q}}{q} \\ \beta &\leq 17 \frac{n}{q} \\ |\mathcal{D}| &\leq 19n \frac{\ln q}{q^{1.34}}\end{aligned}$$

could happen simultaneously with positive probability, that is,

$$1 - 0.892 - 0.059 - 0.047 = 0.002 > 0.$$

Now we choose a set  $X$  satisfying all of these events simultaneously. Every component of  $X$  that contains no weakly dominated vertex has size at least  $\frac{1}{8m^2} \ln q$ , and  $\mathcal{D}$  has at most  $|\mathcal{D}|$  components. Thus, we have the number of components in  $G[X]$  satisfies,

$$\mu \leq \frac{\alpha}{\frac{1}{8m^2} \ln q} + 19n \frac{\ln q}{q^{1.34}}.$$

Since  $f(\delta) = \frac{\ln q}{q^{0.34}}$  is a decreasing function for  $\delta \geq 72 \lfloor \frac{k}{m} \rfloor + 20km \geq 92$ , we obtain

$$\frac{\ln q}{q^{0.34}} \leq \frac{\ln(93m + 2 - t)}{(93m + 2 - t)^{0.34}} \leq \frac{\ln(95 - t)}{(95 - t)^{0.34}} \leq \frac{\ln(93)}{(93)^{0.34}} < 1,$$

that is  $19n \frac{\ln q}{q^{1.34}} < 19 \frac{n}{q}$ . Now we take

$$\alpha \leq n \frac{\ln q}{q} + n \frac{0.5\sqrt{\ln q}}{q}$$

to the inequality above, we have

$$\begin{aligned}\mu &< n \frac{8m^2}{q} + n \frac{4m^2}{q} \frac{1}{\sqrt{\ln q}} + \frac{19n}{q} \\ &< n \frac{8m^2}{q} + n \frac{4m^2}{q} \times \frac{1}{2} + n \frac{19}{q} \\ &= n \frac{10m^2 + 19}{q},\end{aligned}$$

where

$$\frac{1}{\sqrt{\ln q}} \leq \frac{1}{\sqrt{\ln(93m + 2 - t)}} \leq \frac{1}{\sqrt{\ln(95 - t)}} < \frac{1}{\sqrt{\ln(93)}} < \frac{1}{2}.$$

Finally, we have

$$\alpha + (2k + 1)\beta + 2k\mu - 3k < n \frac{72k + 20km^2 + 17 + 0.5\sqrt{\ln q} + \ln q}{q}.$$

So, the inequality (2) is proved and the theorem follows. ■

**Remark 1** For  $k=1$ ,

$$\gamma_1^c(G) < n \frac{109 + 0.5\sqrt{\ln(\delta + 1)} + \ln(\delta + 1)}{(\delta + 1)}.$$

It improves the bound in [3], that is,

$$\gamma_1^c(G) < n \frac{145 + 0.5\sqrt{\ln(\delta + 1)} + \ln(\delta + 1)}{(\delta + 1)}.$$

**Remark 2** Since  $X \cup Y$  is also a  $k$ -dominating set of  $G$ , and  $E[\alpha] + E[\beta] \leq n \frac{1 + \ln q}{q}$ , there is at least one choice of  $X \subseteq V(G)$  such that  $\gamma_k(G) \leq |X \cup Y| \leq n \frac{1 + \ln q}{q}$ , where  $q = m(\delta + 1) + 2 - t$ ,  $m = \lceil \frac{k}{3} \rceil$ , and  $t = 3 \lceil \frac{k}{3} \rceil - k$ . It improves the well-known result of Lovász [11], that is,

$$\gamma_1(G) \leq n \frac{1 + \ln(\delta + 1)}{\delta + 1}.$$

**Theorem 6** For any nontrivial connected graph  $G$  with order  $n$  and minimum degree  $\delta$ ,

$$\gamma_k^c(G) \leq (1 + o_\delta(1))n \frac{\ln q}{q},$$

where  $q = m(\delta + 1) + 2 - t$ ,  $m = \lceil \frac{k}{3} \rceil$ , and  $t = 3 \lceil \frac{k}{3} \rceil - k$ .

*Proof* By Theorem 5, we have

$$\gamma_k^c(G) < n \frac{\ln q}{q} \left( 1 + \frac{72k + 20km^2 + 17}{\ln q} + \frac{0.5}{\sqrt{\ln q}} \right).$$

We get the theorem as

$$\lim_{\delta \rightarrow \infty} \left( \frac{72k + 20km^2 + 17}{\ln q} + \frac{0.5}{\sqrt{\ln q}} \right) = 0.$$

**Remark 3** Theorem 6 generalizes the result of Caro *et al* [3] for  $k = 1$ , that is,

$$\gamma_1^c(G) \leq (1 + o_\delta(1))n \frac{\ln(\delta + 1)}{\delta + 1}.$$

For  $\delta$  is sufficiently large, we also find that the upper bound for  $\gamma_k^c(G)$  behaves like the upper bound for  $\gamma_k(G)$ .

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