

Resonance of Plane Bipartite Graphs

Jingrong Chen^{1*}, Heping Zhang²

1. College of Mathematics, Physics and Software Engineering, Lanzhou
Jiaotong University, Lanzhou, Gansu 730070, P. R. China

2. School of Mathematics and Statistics, Lanzhou University, Lanzhou,
Gansu 730000, P. R. China

*E-mail address: jdlbxx@mail.lzjtu.cn

Abstract: A graph G is called resonant if the boundary of each face of G is an F -alternating closed trail with respect to some f -factor F of G . We show that a plane bipartite graph G is resonant if and only if it is connected and each edge of G is contained in an f -factor and not in another f -factor.

Keywords : Plane bipartite graph; Factor-covered graph; Factor-deleted graph; Resonance

In this paper, we consider the resonance of bipartite graphs. The concept for resonance originate from chemical significance of hexagonal systems and was extended to plane bipartite graphs. A *hexagonal system* [11] is a special connected plane graph without cut vertices, each interior face of which is surrounded by a regular hexagon of side length one. Some elegant characterizations for hexagonal system to be resonant were obtained in [12, 13] and there is a survey [15] about this topic. For plane elementary bipartite graphs, it also can be described by resonant faces [14].

We take the basic terminology from [7]. The graphs in this paper will be simple and finite. Let G be a plane bipartite graph with vertex set $V(G)$ and edge set $E(G)$. By a plane graph G we mean an embedding of a planar graph in the plane. This plane graph decomposes the plane into a number of open regions called *faces*. A bipartite graph means a graph for which the vertices are colored by white or black so that two adjacent vertices receive different colors.

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For a vertex x of graph G , the degree of x is denoted by $d_G(x)$. Let g and f be two integer-valued functions defined on $V(G)$ such that $g(x) \leq f(x)$ for each $x \in V(G)$. A (g, f) -factor of G is a spanning subgraph F of G satisfying $g(x) \leq d_F(x) \leq f(x)$ for all $x \in V(G)$. In particular, a (f, f) -factor is called an f -factor. If f is a constant function taking the value k , then an f -factor is said to be a k -factor. A 1-factor of G is also a *perfect matching*. In 1952, Tutte [10] gave a criterion for the existence of an f -factor of a graph. A necessary and sufficient conditions for existence of a (g, f) -factor was obtained by Lovász [8] in 1970. Little [6] introduced the concept of factor-covered graph. Liu [2, 3] extended this concept to general situation. For a graph G with an f -factor is called *f -factor-covered* if each edge of G is contained in some f -factor. G is called *f -factor-deleted* if $G - e$ contains an f -factor for every edge e of G .

We now extend some concepts concerning 1-factor to f -factor, such as “allowed”, “elementary” and “resonance” etc.. If $f \equiv 1$, then these concepts express the general meanings. Let G be a graph with an f -factor. An edge of G is called *f -allowed* if it lies in some f -factor of G and *f -forbidden* otherwise. An *f -fixed single edge* of G is the edge that there is no f -factor of G containing it. And an edge e of G is called *f -fixed double edge* if e belongs to all f -factors of G . G is said to be *f -elementary* if all its f -allowed edges form a connected subgraph of G . If G is 1-elementary, we always say that G is *elementary*. As early as in 1915, König [5] had employed this concept in studying the decomposition of a determinant. After nearly half a century, Hetyei [4] formally used the term “elementary” for this concept and obtained various properties of elementary bipartite graphs. Let F be an f -factor of G . A closed trail P of G is called *F -alternating* if the edges of P appear alternately in $E(F)$ and $E(G) \setminus E(F)$. A face φ of G is said to be *resonant* if G has an f -factor F such that the boundary of φ is an F -alternating closed trail. G is called *resonant* if each face of G is resonant.

From [14] we know the following result: Let G be a plane bipartite graph with more than two vertices, then each face of G is resonant if and only if G is elementary. But for a plane bipartite graph G with f -factors (where $f \neq 1$), the conditions of f -elementary is not enough for G to be resonant. For example (see Fig. 1), there is no f -factor such that the faces φ_1 or φ_2 are resonant.

The symmetric difference of two finite sets A and B is defined by $A \oplus B := (A \cup B) \setminus (A \cap B)$. This binary operation is associative and commutative. Let M_1 and M_2 be two different perfect matchings of G . Then the symmetric difference $M_1 \oplus M_2$ consists of mutually disjoint (M_1, M_2) -alternating cycles. But this is different from f -factors.

Lemma 1[9]. *Let G be a connected graph with all its vertices in even*

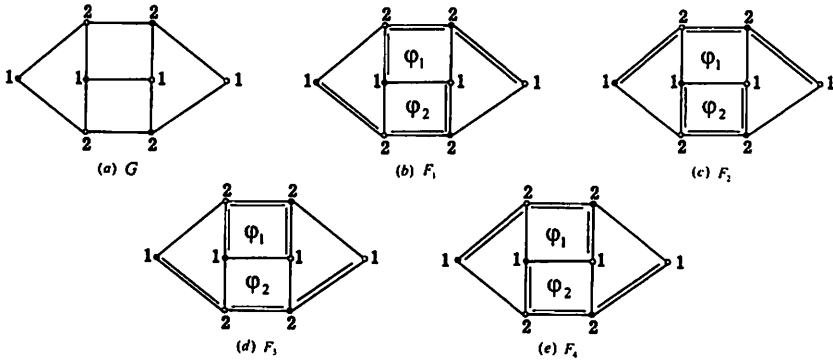


Fig. 1. Bipartite graph G and its f -factors

degree, and G has a 2-edge coloring (not necessary proper) with blue and red colors. If the number of red-color-edges equals to the number of blue-color-edges for every vertex of G , then G has an (red, blue)-alternating Euler tour.

Lemma 2. *If F_1 and F_2 are two different f -factors of G , then the connected components of $G' = (V, E(F_1) \oplus E(F_2))$ are (F_1, F_2) -alternating closed trails.*

Proof: Color the edges of F_1 and F_2 with red and blue respectively. For a vertex x of G , if there are k ($0 \leq k \leq f(x)$) edges incident with x which lie in both F_1 and F_2 , then $d_{G'}(x) = 2(f(x) - k)$. Thus the number of red edges equals to the number of blue edges for every vertex of G' . By Lemma 1, the connected components of G' are (F_1, F_2) -alternating closed trails. \square

Lemma 3[1]. *Let F be an f -factor of a bipartite graph G and P an alternating closed trail relative to F . Then P is the edge-disjoint union of alternating cycles relative to F .*

Corollary 4. *Let F be an f -factor of a bipartite graph G and P an alternating closed trail relative to F . Then the edges which do not belong to any cycle of P form an F -alternating path P' .*

Proof: By Lemma 3, the cycles of P are F -alternating. Assume that P' is not an F -alternating path. Then there are two adjacent edges e_1 and e_2 of P' such that $e_i \in F$ or $e_i \notin F$ ($i = 1, 2$). u is the same end of e_1 and e_2 . Let P_1 be the F -alternating closed trail of P using u as its initial and terminal vertices, and $e_i \notin P_1$. Then $P_2 = e_1 P_1 e_2$ or $P_2 = e_2 P_1 e_1$ is not an F -alternating trail of P , which contradicts $P_2 \subseteq P$. \square

Lemma 5[14]. *Let M be a perfect matching of a graph G and C an M -alternating cycle of G . Then $M \oplus E(C)$ is also a perfect matching of*

G and C is an $(M \oplus E(C))$ -alternating cycle of G .

For f -factors, we have the similar result below.

Lemma 6. *Let F be an f -factor of a bipartite graph G and P an F -alternating closed trail of G . Then $F \oplus E(P)$ is also an f -factor of G and P is an $(F \oplus E(P))$ -alternating closed trail of G .*

Proof: The result is obvious by using Lemmas 3 and 5. \square

If a bipartite graph G is elementary, then G is 2-edge-connected. That is, G has no cut-edge. Conversely, for a bipartite graph G with perfect matchings, if G has a cut-edge, then the cut-edge is certainly forbidden, and G is certainly non-elementary. But for f -factors, this is different.

Lemma 7. *Assume that a bipartite graph G is f -elementary. If G has cut-edges, then all the cut-edges of G are f -fixed double edges.*

Proof: At first we show that all the cut-edges of G are allowed. Otherwise, suppose a cut-edge e is a fixed single edge. Then the union of all f -factors in G forms a disconnected subgraph of G , which contradicts that G is f -elementary.

Further we show that a cut edge ab is a fixed double edge. Otherwise ab is not a fixed double edge. Then there is an f -factor F of G such that $ab \notin F$. Let $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ be the two connected components of $G - ab$. Then $F_1 := F \cap G_1$ is an f -factor of G_1 . We have

$$\sum_{x \in V(G_1)} f(x) = \sum_{x \in V(G_1)} d_{F_1}(x) \equiv 0 \pmod{2}.$$

Since ab is an allowed edge, there is an f -factor F' of G such that $ab \in F'$. Then

$$\sum_{x \in V(G_1)} f(x) = \sum_{x \in V(G_1)} d_{F'}(x) = \sum_{x \in V(G_1)} d_{F' \cap G_1}(x) + 1.$$

So

$$\sum_{x \in V(G_1)} d_{F' \cap G_1}(x) \equiv 1 \pmod{2},$$

contradicting the Handshaking Lemma. The theorem is proved. \square

Corollary 8. *If a connected graph G is both f -factor-covered and f -factor-deleted, then G is 2-edge-connected.*

Proof: If G has a cut-edge e , by Lemma 7 e is a f -fixed double edge since G is f -elementary. Hence $G - e$ has no f -factor, which contradicts that G is f -factor-deleted. \square

Lemma 9. *Suppose a graph G is both f -factor-covered and f -factor-deleted, and e is an edge of an f -factor F of G . Then there are an edge e' adjacent to e and an f -factor F' of G such that $e' \notin F$, $e \notin F'$ and $e' \in F'$.*

Proof: Let u be an endpoint of e . Assume that e_1, e_2, \dots, e_k are the edges incident with u but not belong to F . Because G is f -factor-deleted, we have $k = d_G(u) - f(u) \geq 1$ and $G - e$ has an f -factor F' such that F' contains some e_i . Otherwise, only $f(u) - 1$ edges incident with u belonging to F' . The result is proved. \square

Theorem 10. *A connected plane bipartite graph G is both f -factor-covered and f -factor-deleted if and only if each face of G is resonant.*

Proof: Assume that each face of G is resonant. For any edge e of G , let P be a boundary of the faces containing e . Then there is an f -factor F such that P is an F -alternating closed trail, and by Lemma 6 $F \oplus E(P)$ is another f -factor. If $e \in F$, then $e \notin F \oplus E(P)$; if $e \notin F$, then $e \in F \oplus E(P)$. Hence G is both f -factor-covered and f -factor-deleted.

Conversely, assume that G is both f -factor-covered and f -factor-deleted. Let φ be a face of G . We will show that φ is resonant. For an edge e on the boundary of φ , there are two f -factors F_1 and F_2 such that $e \in F_1$ and $e \notin F_2$. By Lemmas 2 and 3, there is an (F_1, F_2) -alternating cycle C of $F_1 \oplus F_2$ such that $e \in E(C)$. Without loss of generality, assume that φ lies in the interior of C . Denote by $I[C]$ the subgraph of G consisting of C together with the interior. Obviously, $I[C]$ is connected. We now show that the every face of $I[C]$ is resonant by induction on the number m of edges contained in the interior of C .

If $m = 0$, the result is trivially true.

Now we suppose that the result is true for $0 \leq m \leq k$. That is: if the number of edges in the interior of an f -alternating cycle C is no more than k , then the every face of $I[C]$ is resonant. On the basis of this, we give the proof of $m = k + 1$ in the following.

In the interior of C , let e_1 be an edge with an endpoint on C . If the two ends of e_1 are both on C , then C and e_1 form two cycles C_1 and C_2 . Whenever e_1 belongs to F_1 (or F_2) or not, C_1 and C_2 are f -alternating cycles relative to some f -factors of G . And the number of edges in the interior of C_1 and C_2 are no more than k respectively. So by the induction hypothesis, each face of $I[C_1]$ and $I[C_2]$ is resonant.

If one of ends of e_1 is in the interior of C , then there are three cases to be considered.

Case 1: $e_1 \in F_1, e_1 \notin F_2$ (see Fig. 2 (a) and (b)).

By Lemmas 2 and 3, $F_1 \oplus F_2$ has a cycle C' containing e_1 . Let P be an (F_1, F_2) -alternating path of C' such that only end vertices lie on C . If the two end vertices of P are different, then C and P form two cycles C_1 and C_2 such that $C_1 \cap C_2 = P$, see Fig.2(a). Without loss of generality, suppose C_1 is an F_1 -alternating cycle and C_2 is thus an $F_1 \oplus E(C_1)$ -alternating cycle. If the two end vertices of P are the same, then we can get that C_1 is an (F_1, F_2) -alternating cycle and C_2 is an (F_1, F_2) -alternating closed

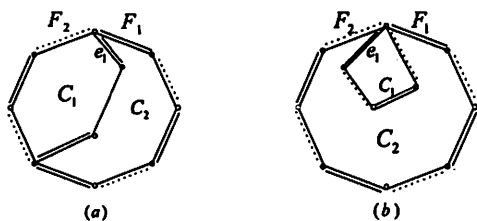


Fig. 2.

trail, see Fig.2(b). By the two narration above, the number of edges in the interior of the C_1 and C_2 are no more than k respectively. So by the induction hypothesis, each face of $I[C_1]$ and $I[C_2]$ is resonant. For $e_1 \notin F_1$ and $e_1 \in F_2$, use the same argument.

Case 2: $e_1 \in F_1 \cap F_2$ (see Fig. 3(a)).

By Lemma 9, there is an edge e_2 adjacent to e_1 and an f -factor F_3 such that $e_1 \notin F_3, e_2 \notin F_1, e_2 \in F_3$. Then by Lemma 2, an (F_1, F_3) -alternating closed trail C' of $F_1 \oplus F_3$ contains e_1 and e_2 . Let P be a part(or the whole) of C' in $I[C]$ such that only end vertices lie on C . If the two end vertices of P are different, then by Corollary 4 we can get an (F_1, F_3) -alternating path by deleting the (F_1, F_3) -alternating cycles of P .

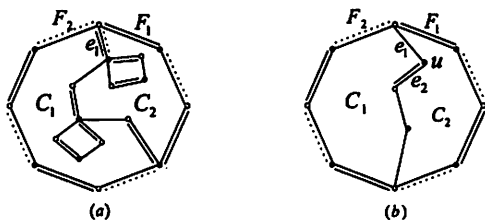


Fig. 3.

If the two end vertices of P are the same, keep an (F_1, F_3) -alternating cycle of P which has a vertex same to the ends of P and delete the other cycles. Then discussed as in Case 1, two cycles C_1 and C_2 (or C_1 is a cycle and C_2 is a closed trail) is obtained, which alternate with respect to some f -factors of G . And the number of edges in the interior of the C_1 and C_2 are no more than k respectively. So by the induction hypothesis, each face of $I[C_1]$ and $I[C_2]$ is resonant.

Case 3: $e_1 \notin F_1, e_1 \notin F_2$ (see Fig. 3(b)).

Let u be an endpoint of e_1 in the interior of C and F_3 be an f -factor such that $e_1 \in F_3$. Among all the edges incident with u except e_1 , there

are $f(u)$ edges in F_1 and $f(u) - 1$ edges in F_3 . Then there are at least one edge e_2 incident with u , such that $e_2 \in F_1$, and $e_2 \notin F_3$. By Lemma 2, there is an alternating closed trail of $F_1 \oplus F_3$ containing e_1 and e_2 . Similar to those methods used in Case 2, we can get an (F_3, F_1) -alternating path or cycle dividing C into two alternating cycles or closed trails C_1 and C_2 . The number of edges in the interior of the C_1 and C_2 are no more than k respectively. So by the induction hypothesis, each face of $I[C_1]$ and $I[C_2]$ is resonant.

Therefore each face of $I[C]$ is resonant over the discussions above, and the theorem is proved. \square

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