

On the order of close to regular graphs without a matching of given size

Sabine Klinkenberg and Lutz Volkmann

Lehrstuhl II für Mathematik, RWTH Aachen University, 52056 Aachen,
Germany
e-mail: volkm@math2.rwth-aachen.de

Abstract

A graph G is a $(d, d+k)$ -graph, if the degree of each vertex of G is between d and $d+k$. Let $p \geq 0$ and $d, k \geq 2$ be integers. If G is a $(d, d+k)$ -graph of order n with at most p odd components and without a matching M of size $2|M| = n - p$, then we show in this paper that

- (i) $n \geq 2d + p + 2$ when $p \leq k - 2$,
- (ii) $n \geq 2\lceil(d(p+2))/k\rceil + p + 2$ when $p \geq k - 1$.

Corresponding results for $0 \leq p \leq 1$ and $0 \leq k \leq 1$ were given by Wallis [6], Zhao [8], and Volkmann [5].

Examples will show that the given bounds (i) and (ii) are best possible.

Keywords: Matching, close to regular graph

We shall assume that the reader is familiar with standard terminology on graphs (see, e.g., Chartrand and Lesniak [2]). In this paper, all graphs are finite and simple. The vertex set of a graph G is denoted by $V(G)$. The neighborhood $N_G(x) = N(x)$ of a vertex x is the set of vertices adjacent with x , and the number $d_G(x) = d(x) = |N(x)|$ is the *degree* of x in the graph G . If $d \leq d_G(x) \leq d+k$ for each vertex x in a graph G , then we speak of a *close to regular graph* or more precisely of a $(d, d+k)$ -graph. If M is a matching in a graph G with the property that every vertex (with exactly one exception) is incident with an edge of M , then M is a *perfect*

matching (an almost perfect matching). We denote by $K_{r,s}$ the complete bipartite graph with partite sets A and B , where $|A| = r$ and $|B| = s$. If G is a graph and $A \subseteq V(G)$, then we denote by $q(G - A)$ the number of odd components in the subgraph $G - A$.

As a generalization of a result by Wallis [6] (see also [7]), Zhao [8] proved in 1991 the following theorem.

Theorem 1 (Zhao [8] 1991) Let $d \geq 2$ be an integer. If a $(d, d+1)$ -graph G has no odd component and no perfect matching, then

$$|V(G)| \geq 3d + 4.$$

Theorem 1 follows easily from the next result by Volkmann [5].

Theorem 2 (Volkmann [5] 2004) Let $d \geq 2$ be an integer, and let G be a $(d, d+1)$ -graph with exactly one odd component and without any almost perfect matching. Then

- 1) $|V(G)| \geq 4(d+1) + 1$,
- 2) $|V(G)| \geq 4(d+1) + 3$ when $d \geq 3$ is odd or $d = 2$ and G is connected,
- 3) $|V(G)| \geq 4(d+1) + 5$ when $d \geq 3$ is odd and G is connected.

In [5] one can also find the following corresponding result for $(d, d+2)$ -graphs.

Theorem 3 (Volkmann [5] 2004) If G is a $(d, d+2)$ -graph with exactly one odd component and without any almost perfect matching, then

$$|V(G)| \geq 3d + 3.$$

Instead of $(d, d+1)$ -graphs or $(d, d+2)$ -graphs, we investigate in this paper the general case of $(d, d+k)$ -graphs for $k \geq 2$. Our main theorem (Theorem 4) is a supplement to Theorems 1 and 2 and an extension of Theorem 3. The proof of our main theorem is based on the following generalization of Tutte's famous 1-factor theorem [3] by Berge [1] in 1958, and we call it the Theorem of Tutte-Berge (for a proof see e.g., [4]).

Theorem of Tutte-Berge (Berge [1] 1958) Let G be a graph of order n . If M is a maximum matching of G , then

$$n - 2|M| = \max_{A \subseteq V(G)} \{q(G - A) - |A|\}.$$

Theorem 4 Let $p \geq 0$ and $d, k \geq 2$ be integers. If G is a $(d, d+k)$ -graph of order n with at most p odd components and without any matching M of size $2|M| = n - p$, then

(i) $n \geq 2d + p + 2$ when $p \leq k - 2$,

(ii) $n \geq 2\lceil(d(p+2))/k\rceil + p + 2$ when $p \geq k - 1$.

Proof In view of the hypotheses, we observe that n and p are of the same parity. Suppose to the contrary that there exists a $(d, d+k)$ -graph G with at most p odd components and without any matching M of size $2|M| = n - p$ such that

a) $n \leq 2d + p + 1$ when $p \leq k - 2$,

b) $n \leq 2\lceil(d(p+2))/k\rceil + p + 1$ when $p \geq k - 1$.

By the hypotheses and the Theorem of Tutte-Berge, there exists a non-empty set $A \subseteq V(G)$ such that $q(G - A) \geq |A| + p + 1$. However, since n and p are of the same parity, it is straightforward to verify that this even leads to the better bound $q(G - A) \geq |A| + p + 2$. We call an odd component of $G - A$ large if it has more than d vertices and small otherwise. If we denote by α and β the number of large and small components, respectively, then we deduce that

$$\alpha + \beta = q(G - A) \geq |A| + p + 2, \quad (1)$$

$$n \geq |A| + \beta + \alpha(d + 1), \quad (2)$$

$$n \geq |A| + \beta + \alpha(d + 2) \text{ when } d \geq 3 \text{ is odd.} \quad (3)$$

Firstly, we show that $\alpha \leq p + 1$. In the case that $p \leq k - 2$, it follows from assumption a) and inequality (2) that

$$2d + p + 1 \geq n \geq |A| + \beta + \alpha(d + 1) \geq 1 + \alpha(d + 1).$$

This leads to $(2 - \alpha)(d + 1) + p - 2 \geq 0$ and thus $\alpha \leq p + 1$. In the other case that $p \geq k - 1$, we conclude from assumption b) and (2) that

$$2\left\lceil\frac{d(p+2)}{k}\right\rceil + p + 1 \geq n \geq |A| + \beta + \alpha(d + 1) \geq 1 + \alpha(d + 1).$$

This inequality chain yields

$$2\left(\frac{d(p+2)}{k} + \frac{k-1}{k}\right) + p - \alpha(d + 1) \geq 0.$$

Because of $k \geq 2$, it is a simple matter to verify that this inequality implies $\alpha \leq p + 1$. Applying (1), we arrive at

$$\beta \geq |A| + 1. \quad (4)$$

Since G is a $(d, d+k)$ -graph, it is easy to show that there are at least d edges of G joining each small component of $G - A$ with A . Therefore it follows from the hypothesis that G has at most p odd components that

$$\alpha - p + d\beta \leq |A|(d+k), \quad (5)$$

$$d\beta \leq |A|(d+k) \text{ when } \alpha \leq p. \quad (6)$$

Case 1. Assume that $p \leq k - 2$.

If $|A| \geq d$, then inequalities (1) and (2) lead to the following contradiction to assumption a):

$$\begin{aligned} n &\geq |A| + \beta + \alpha(d+1) \\ &\geq |A| + |A| + p + 2 - \alpha + \alpha(d+1) \\ &\geq 2d + p + 2 + \alpha d \\ &\geq 2d + p + 2 \end{aligned}$$

Let U be a small component of $G - A$. Since $N(x) \subseteq V(U) \cup A$ for $x \in V(U)$, we observe that $|A| + |V(U)| \geq d + 1$. If $|A| < d$, say $|A| = d - t$ with $1 \leq t \leq d - 1$, then we deduce that each small component U contains at least $t + 1$ vertices. Thus (1) implies that

$$\begin{aligned} n &\geq |A| + \beta(t+1) + \alpha(d+1) \\ &\geq d - t + (|A| + p + 2 - \alpha)(t+1) + \alpha(d+1) \\ &= 2d + p + 2 + (d - t + p - \alpha)t + \alpha d. \end{aligned}$$

Because of $\alpha \leq p + 1$ and $t \leq d - 1$, this leads to $n \geq 2d + p + 2$, a contradiction to assumption a).

Case 2. Assume that $p \geq k - 1$ and $\alpha = 0$.

From inequalities (6) and (1), we deduce that

$$|A|(d+k) \geq d\beta \geq d(|A| + p + 2).$$

This yields $|A|k \geq d(p+2)$ and thus $|A| \geq \lceil (d(p+2))/k \rceil$. Combining this with (1) and (2), we arrive at the following contradiction to assumption b):

$$\begin{aligned} n &\geq |A| + \beta \geq |A| + |A| + p + 2 \\ &\geq 2 \left\lceil \frac{d(p+2)}{k} \right\rceil + p + 2 \end{aligned}$$

Case 3. Assume that $p \geq k - 1$ and $\alpha \geq 1$.

We note that inequality (5) is equivalent to

$$\beta \leq |A| + \frac{p + k\beta - \alpha}{d+k}. \quad (7)$$

Subcase 3.1. Assume that $p + k\beta - \alpha \leq d + k - 1$. It follows from (7) that $\beta \leq |A|$, a contradiction to (4).

Subcase 3.2. Assume that $p + k\beta - \alpha \geq d + k$. This implies that $k\beta \geq d + k - p + \alpha$. Combining this with (5), we obtain

$$|A| \geq \frac{\alpha - p + d\beta}{d + k} \geq \frac{\alpha - p + \frac{d}{k}(d + k - p + \alpha)}{d + k}.$$

For $\alpha = p + 1$, this yields $|A| \geq (d + 1)/k$ and hence (2) leads to

$$\begin{aligned} n &\geq |A| + \beta + \alpha(d + 1) \\ &\geq \frac{d + 1}{k} + \frac{d + k + 1}{k} + (p + 1)(d + 1) \\ &= \frac{2d + 2}{k} + d(p + 1) + p + 2. \end{aligned} \quad (8)$$

Since $p \geq 0$ and $d, k \geq 2$, we observe that $k(pd + d - 2) \geq 2(pd + d - 2)$, and this is equivalent to

$$\frac{2d + 2}{k} + d(p + 1) \geq 2 \frac{d(p + 2) + (k - 1)}{k}.$$

Combining this inequality with (8), we arrive at a contradiction to assumption b) as follows:

$$\begin{aligned} n &\geq \frac{2d + 2}{k} + d(p + 1) + p + 2 \\ &\geq 2 \frac{d(p + 2) + (k - 1)}{k} + p + 2 \\ &\geq 2 \left\lceil \frac{d(p + 2)}{k} \right\rceil + p + 2 \end{aligned}$$

Assume next that $1 \leq \alpha \leq p$, and let $\alpha = p - s$ with $0 \leq s \leq p - 1$. We deduce from (1) that $\beta \geq |A| + s + 2$, and this yields together with (6) the inequality

$$|A| \geq \left\lceil \frac{d(s + 2)}{k} \right\rceil \quad (9)$$

Subcase 3.2.1. Assume that $k = 2$ and that $d \geq 2$ is even. In this case (2) and (9) lead to the following contradiction to b):

$$\begin{aligned} n &\geq |A| + \beta + \alpha(d + 1) \\ &\geq 2|A| + s + 2 + (p - s)(d + 1) \\ &\geq 2 \frac{d(s + 2)}{2} + s + 2 + (p - s)(d + 1) \\ &= d(p + 2) + p + 2 \\ &= 2 \left\lceil \frac{d(p + 2)}{2} \right\rceil + p + 2 \end{aligned}$$

Subcase 3.2.2. Assume that $k = 2$ and that $d \geq 3$ and s are odd. Since d is odd, (3) and (9) yield the following contradiction to b):

$$\begin{aligned}
 n &\geq |A| + \beta + \alpha(d+2) \\
 &\geq 2|A| + s + 2 + (p-s)(d+2) \\
 &\geq 2\frac{d(s+2)+1}{2} + s + 2 + (p-s)(d+2) \\
 &= d(p+2) + p + 3 + p - s \\
 &\geq d(p+2) + p + 4 \\
 &\geq 2\left\lceil \frac{d(p+2)}{2} \right\rceil + p + 2
 \end{aligned}$$

Subcase 3.2.3. Assume that $k = 2$ and that $d \geq 3$ is odd and that s is even. Combining (3) and (9), we arrive at the following contradiction to b):

$$\begin{aligned}
 n &\geq |A| + \beta + \alpha(d+2) \\
 &\geq 2|A| + s + 2 + (p-s)(d+2) \\
 &\geq d(s+2) + s + 2 + (p-s)(d+2) \\
 &= d(p+2) + 2p + 2 - s \\
 &\geq d(p+2) + p + 3 \\
 &= 2\frac{d(p+2)+1}{2} + p + 2 \\
 &\geq 2\left\lceil \frac{d(p+2)}{2} \right\rceil + p + 2
 \end{aligned}$$

Subcase 3.2.4. Assume that $k \geq 3$. It follows from (2) and (9) that

$$\begin{aligned}
 n &\geq |A| + \beta + \alpha(d+1) \\
 &\geq 2|A| + s + 2 + (p-s)(d+1) \\
 &\geq 2\left\lceil \frac{d(s+2)}{k} \right\rceil + (p-s)d + p + 2. \tag{10}
 \end{aligned}$$

To receive a contradiction to assumption b), it thus remains to show that

$$2\frac{d(s+2)}{k} + (p-s)d \geq 2\frac{d(p+2) + k - 1}{k},$$

and this is equivalent to

$$d \geq \frac{2(k-1)}{(p-s)(k-2)}. \tag{11}$$

If $s \leq p - 2$, then $k \geq 3$ implies

$$\frac{2(k-1)}{(p-s)(k-2)} \leq \frac{k-1}{k-2} \leq 2 \leq d$$

and (11) is valid. If $s = p - 1$ and $d \geq 4$, then it is easy to see that (11) is also true. In the remaining case that $k \geq 3$, $s = p - 1$, and $2 \leq d \leq 3$, we deduce that

$$\begin{aligned} 2 \left\lceil \frac{d(p+2)}{k} \right\rceil &= 2 \left\lceil \frac{d(p+1)}{k} + \frac{d}{k} \right\rceil \\ &\leq 2 \left\lceil \frac{d(p+1)}{k} \right\rceil + 2 \left\lceil \frac{d}{k} \right\rceil \\ &= 2 \left\lceil \frac{d(p+1)}{k} \right\rceil + 2 \\ &\leq 2 \left\lceil \frac{d(p+1)}{k} \right\rceil + d \\ &= 2 \left\lceil \frac{d(s+2)}{k} \right\rceil + (p-s)d. \end{aligned}$$

Combining this inequality chain with (10), we finally obtain a contradiction to b).

Since we have discussed all possible cases, the proof of Theorem 4 is complete. \square

The following examples show that the bounds in Theorem 4 are best possible.

Example 5 Case 1. Assume that $p + 2 \leq k$. In this case, the complete bipartite graph $K_{d, d+p+2}$ is a $(d, d+k)$ -graph of order $n = 2d + p + 2$ without a matching M of size $2|M| = n - p$. Consequently, Condition (i) is best possible.

Case 2. Assume that $p + 2 \geq k + 1$. Let H be a d -regular bipartite graph with the partite sets $X = \{x_1, x_2, \dots, x_{\lceil \frac{d(p+2)}{k} \rceil}\}$ and $Y = \{y_1, y_2, \dots, y_{\lceil \frac{d(p+2)}{k} \rceil}\}$. Now let G consists of H and $p + 2$ additional vertices u_1, u_2, \dots, u_{p+2} , which are connected with X by $d(p + 2)$ edges such that $d_G(u_i) = d$ for $i = 1, 2, \dots, p + 2$ and $|d_G(x_i) - d_G(x_j)| \leq 1$ for $1 \leq i, j \leq \lceil \frac{d(p+2)}{k} \rceil$. Now G is a $(d, d+k)$ -graph of order $2\lceil \frac{d(p+2)}{k} \rceil + p + 2$ without a matching M of size $2|M| = n - p$. This example shows that Condition (ii) is also best possible.

References

- [1] C. Berge, Sur le couplage maximum d'un graphe, *C. R. Acad. Sci. Paris Math.* **247** (1958), 258-259.
- [2] G. Chartrand, L. Lesniak, Graphs and Digraphs, 3rd Edition, Chapman and Hall, London, 1996.
- [3] W.T. Tutte, The factorizations of linear graphs, *J. London Math. Soc.* **22** (1947), 459-474.
- [4] L. Volkmann, Foundations of Graph Theory, Springer-Verlag, Wien New York (1996) (in German).
- [5] L. Volkmann, On the size of odd order graphs with no almost perfect matching, *Australas. J. Combin.* **29** (2004), 119-126.
- [6] W.D. Wallis, The smallest regular graphs without one-factors, *Ars Combinat.* **11** (1981), 295-300.
- [7] W.D. Wallis, One-Factorizations, Kluwer Academic Publishers, Dordrecht Boston London, Mathematics and Its Applications 390 (1997).
- [8] C. Zhao, The disjoint 1-factors of $(d, d+1)$ -graphs, *J. Combin. Math. Combin. Comput.* **9** (1991), 195-198.