# A New $\sigma_3$ Type Condition for Heavy Cycles in Weighted Graphs

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Abstract. A weighted graph is one in which every edge e is assigned a nonnegative number w(e), called the weight of e. The weight of a cycle is defined as the sum of the weights of its edges. The weighted degree of a vertex is the sum of the weights of the edges incident with it. In this paper, motivated by a recent result of Fujisawa, we prove that a 2-connected weighted graph G contains either a Hamilton cycle or a cycle of weight at least 2m/3 if it satisfies the following conditions: (1) The weighted degree sum of every three pairwise nonadjacent vertices is at least m; (2) In each induced claw and each induced modified claw of G, all edges have the same weight. This extends a theorem of Zhang, Broersma and Li.

## 1 Terminology and notation

We use Bondy and Murty [3] for terminology and notation not defined here and consider finite simple graphs only.

Let G = (V, E) be a simple graph. G is called a weighted graph if each edge e is assigned a nonnegative number w(e), called the weight of e. For a subgraph H of G, V(H) and E(H) denote the sets of vertices and edges of H, respectively. The weight of H is defined by

$$w(H) = \sum_{e \in E(H)} w(e).$$

For a vertex  $v \in V$ ,  $N_H(v)$  denotes the set, and  $d_H(v)$  the number, of vertices in H that are adjacent to v. We define the weighted degree of v in

H by

$$d_H^w(v) = \sum_{h \in N_H(v)} w(vh).$$

When no confusion occurs, we will denote  $N_G(v)$ ,  $d_G(v)$  and  $d_G^w(v)$  by N(v), d(v) and  $d^w(v)$ , respectively.

An unweighted graph can be regarded as a weighted graph in which each edge e is assigned weight w(e) = 1. Thus, in an unweighted graph,  $d^w(v) = d(v)$  for every vertex v, and the weight of a subgraph is simply the number of edges of it.

An (x, y)-path is a path connecting the two vertices x and y. The distance between two vertices x and y, denoted by d(x, y), is the length of a shortest (x, y)-path. If u and v are two vertices on a path P, P[u, v] denotes the segment of P from u to v.

The number of vertices in a maximum independent set of G is denoted by  $\alpha(G)$ . If G is noncomplete, then for a positive integer  $k \leq \alpha(G)$  we denote by  $\sigma_k(G)$  the minimum value of the degree sum of any k pairwise nonadjacent vertices, and by  $\sigma_k^w(G)$  the minimum value of the weighted degree sum of any k pairwise nonadjacent vertices. If G is complete, then both  $\sigma_k(G)$  and  $\sigma_k^w(G)$  are defined as  $\infty$ .

We call the graph  $K_{1,3}$  a *claw*, and the graph obtained by joining a pendant edge to some vertex of a triangle a *modified claw*.

#### 2 Results

There have been many results on the existence of long cycles in graphs. The following theorem is well-known.

**Theorem A (Pósa [7]).** Let G be a 2-connected graph such that  $\sigma_2(G) \geq s$ . Then G contains either a Hamilton cycle or a cycle of length at least s.

This result was generalized by the following two theorems along different lines.

**Theorem B (Fan [4]).** Let G be a 2-connected graph such that  $\max\{d(x), d(y)|d(x,y)=2\} \ge c/2$ . Then G contains either a Hamilton cycle or a cycle of length at least c.

**Theorem C** (Fournier & Fraisse [5]). Let G be a k-connected graph where  $2 \le k < \alpha(G)$ , such that  $\sigma_{k+1}(G) \ge m$ . Then G contains either a Hamilton cycle or a cycle of length at least 2m/(k+1).

Bondy et al. [2] gave a weighted generalization of Theorem A as follows.

**Theorem 1 (Bondy et al. [2]).** Let G be 2-connected weighted graph such that  $\sigma_2^w(G) \geq s$ . Then G contains either a Hamilton cycle or a cycle of weight at least s.

In [9], it was showed that if one wants to give a generalization of Theorem B to weighted graphs some extra conditions cannot be avoided. By adding two extra conditions, the authors gave a weighted generalization of Theorem B.

**Theorem 2 (Zhang et al. [9]).** Let G be a 2-connected weighted graph which satisfies the following conditions:

- (1)  $max\{d^w(x), d^w(y)|d(x,y)=2\} \geq s/2;$
- (2) w(xz) = w(yz) for every vertex  $z \in N(x) \cap N(y)$  with d(x,y) = 2;
- (3) In every triangle T of G, either all edges of T have different weights or all edges of T have the same weight.

Then G contains either a Hamilton cycle or a cycle of weight at least s.

Motivated by this result, Zhang et al. [8] gave a weighted generalization of Theorem C in the case k = 2.

**Theorem 3 (Zhang et al.** [8]). Let G be a 2-connected weighted graph which satisfies the following conditions:

- (1)  $\sigma_3^w(G) \geq m$ ;
- (2) w(xz) = w(yz) for every vertex  $z \in N(x) \cap N(y)$  with d(x,y) = 2;
- (3) In every triangle T of G, either all edges of T have different weights or all edges of T have the same weight.

Then G contains either a Hamilton cycle or a cycle of weight at least 2m/3.

Theorem B was further extended by the following result.

**Theorem D** (Bedrossian et al. [1]). Let G be a 2-connected graph. If  $\max\{d(x), d(y)\} \ge c/2$  for each pair of nonadjacent vertices x and y, which are vertices of an induced claw or an induced modified claw of G, then G contains either a Hamilton cycle or a cycle of length at least c.

Fujisawa [6] gave a weighted generalization of Theorem D. The result also generalizes Theorem 2.

**Theorem 4 (Fujisawa** [6]). Let G be a 2-connected weighted graph which satisfies the following conditions:

- (1) For each induced claw and each induced modified claw of G, all its nonadjacent pair of vertices x and y satisfy  $\max\{d^w(x), d^w(y)\} \ge s/2$ ;
- (2) For each induced claw and each induced modified claw of G, all of its edges have the same weight.

Then G contains either a Hamilton cycle or a cycle of weight at least s.

It is clear that Condition (2) of Theorem 4 is weaker than Conditions (2) and (3) of Theorem 3. Our main result in this paper is a further weighted generalization of Theorem C in the case k=2. It turns out that Conditions (2) and (3) of Theorem 3 can be replaced by Condition (2) of Theorem 4.

**Theorem 5.** Let G be a 2-connected weighted graph which satisfies the following conditions:

- (1)  $\sigma_3^w(G) \geq m$ ;
- (2) For each induced claw and each induced modified claw of G, all of its edges have the same weight.

Then G contains either a Hamilton cycle or a cycle of weight at least 2m/3.

We postpone the proof of Theorem 5 to the next section.

#### 3 Proof of Theorem 5

In the proof of Theorem 5, we call a path P a heaviest longest path if P has the following properties

- $\bullet$  P is a longest path of G, and
- w(P) is maximum among all longest paths in G.

To prove Theorem 5, we need the following lemmas. The proof of Lemma 1 is implicit in [2].

**Lemma 1 (Bondy** et al. [2]). Let G be a non-hamiltonian 2-connected weighted graph and  $P = v_1 v_2 \cdots v_p$  be a heaviest longest path in G. Then there is a cycle C in G with  $w(C) \ge d^w(v_1) + d^w(v_p)$ .

**Lemma 2 (Fujisawa** [6]). Let G be a weighted graph satisfying Condition (2) of Theorem 5. If  $x_1yx_2$  is an induced path with  $w(x_1y) \neq w(x_2y)$  in G, then each vertex  $x \in N(y) \setminus \{x_1, x_2\}$  is adjacent to both  $x_1$  and  $x_2$ .

**Lemma 3** (Fujisawa [6]). Let G be a weighted graph satisfying Condition (2) of Theorem 5. Suppose  $x_1yx_2$  is an induced path such that  $w_1 = w(x_1y)$  and  $w_2 = w(x_2y)$  with  $w_1 \neq w_2$ , and  $yz_1z_2$  is a path such that  $\{z_1, z_2\} \cap \{x_1, x_2\} = \emptyset$  and  $x_2z_2 \notin E(G)$ . Then

- (i)  $\{z_1x_1, z_1x_2, z_2x_1\} \subseteq E(G)$ , and  $yz_2 \notin E(G)$ . Moreover, all edges in the subgraph induced by  $\{x_1, y, x_2, z_1, z_2\}$ , other than  $x_1y$ , have the same weight  $w_2$ .
- (ii) Let Y be the component of  $G \{x_2, z_1, z_2\}$  with  $y \in V(Y)$ . For each vertex  $v \in V(Y) \setminus \{x_1, y\}$ , v is adjacent to all of  $x_1$ ,  $x_2$ , y and  $z_2$ . Furthermore,  $w(vx_1) = w(vx_2) = w(vy) = w(vz_2) = w_2$ .

**Proof of Theorem 5.** Let G be a 2-connected weighted graph satisfying the conditions of Theorem 5. Suppose that G does not contain a Hamilton

cycle. Then it suffices to prove that G contains a cycle of weight at least 2m/3.

Choose a path  $P = v_1 v_2 \cdots v_n$  in G such that

- (a) P is a heaviest longest path in G;
- (b)  $d^w(v_1) + d^w(v_p)$  is as large as possible, subject to (a).

From the choice of P, we can immediately see that  $N(v_1) \cup N(v_p) \subseteq V(P)$ . And it is not difficult to prove that there exists no cycle of length p.

It follows from Lemma 1 that there exists a cycle C of weight  $w(C) \ge d^w(v_1) + d^w(v_p)$ . Without loss of generality, we assume  $d^w(v_1) \le w(C)/2$ .

Claim 0. Let  $P_1$  and  $P_2$  be two heaviest longest paths such that  $P_1$  has v' and  $v_p$  as its end-vertices, and  $P_2$  has v'' and  $v_p$  as its end-vertices. If  $v'v'' \notin E(G)$ , then w(C) > 2m/3.

Proof. Since  $P_1$  and  $P_2$  are heaviest longest paths,  $v'v_p \notin E(G)$  and  $v''v_p \notin E(G)$ . Then v', v'' and  $v_p$  are pairwise nonadjacent. By the choice of the path P in (b),  $d^w(v') \leq d^w(v_1) \leq w(C)/2$  and  $d^w(v'') + d^w(v_p) \leq w(C)$ . So we have  $d^w(v') + d^w(v'') + d^w(v_p) \leq 3w(C)/2$ . It follows from Condition (1) of the theorem that  $w(C) \geq 2m/3$ .

Since G is 2-connected,  $v_1$  is adjacent to at least one vertex on P other than  $v_2$ . Choose  $v_k \in N(v_1)$  such that k is as large as possible. It is clear that  $3 \le k \le p-1$ .

Since  $G - v_k$  is connected, there must be a path Q such that

- Q has end-vertices  $v_r$  and  $v_s$ , such that r < k < s, and
- $\bullet \ V(Q) \cap V(P) = \{v_r, v_s\}.$

We assume that such a path Q was chosen so that

- (i) s is as large as possible;
- (ii) r is as large as possible, subject to (i).

Case 1.  $v_1v_i \in E(G)$  for every i with  $r \leq i \leq k$ .

Claim 1.1.  $v_r v_s \in E(G)$ .

*Proof.* Since r < k, we have  $v_1v_{r+1} \in E(G)$ . If there exists a vertex  $u \notin \{v_r, v_s\}$  on Q, then the path  $Q[u, v_r]v_rv_{r-1}\cdots v_1v_{r+1}v_{r+2}\cdots v_p$  is longer than P, a contradiction.

Case 1.1. s > k + 1.

Claim 1.2.  $w(v_1v_{r+1}) = w(v_rv_{r+1}).$ 

*Proof.* First, we consider the case r < k - 1. By the choices of  $v_k$  and  $v_r$ ,  $v_1v_s \notin E(G)$  and  $v_{r+1}v_s \notin E(G)$ . So  $\{v_r, v_{r+1}, v_1, v_s\}$  induces a modified claw. Then we get  $w(v_1v_{r+1}) = w(v_rv_{r+1})$ .

Now consider the case r = k - 1. We need prove  $w(v_1v_k) = w(v_{k-1}v_k)$ .

By the choices of  $v_k$  and  $v_s$ ,  $v_1v_{s+1} \notin E(G)$  and  $v_{k-1}v_{s+1} \notin E(G)$ . So, if  $v_kv_{s+1} \in E(G)$ , then  $\{v_k, v_{k-1}, v_1, v_{s+1}\}$  induces a modified claw. Then we get  $w(v_1v_k) = w(v_{k-1}v_k)$ .

By the choice of  $v_k$ ,  $v_1v_s \notin E(G)$ . So, if  $v_kv_s \notin E(G)$ , then  $\{v_{k-1}, v_k, v_1, v_s\}$  induces a modified claw. Then we get  $w(v_1v_k) = w(v_{k-1}v_k)$ .

Clearly we need only consider the case  $v_k v_{s+1} \notin E(G)$  and  $v_k v_s \in E(G)$ . By the choice of  $v_k$ ,  $v_1 v_{k+1} \notin E(G)$  and  $v_1 v_s \notin E(G)$ . So  $\{v_k, v_{k+1}, v_1, v_s\}$  induces a claw or a modified claw, which implies that  $w(v_1 v_k) = w(v_k v_s)$ ; On the other hand, by the choice of  $v_s$ ,  $v_{k-1} v_{s+1} \notin E(G)$ . So  $\{v_s, v_{k-1}, v_k, v_{s+1}\}$  induces a modified claw, which implies that  $w(v_k v_s) = w(v_{k-1} v_k)$ . Thus we have  $w(v_1 v_k) = w(v_{k-1} v_k)$ .

### Claim 1.3. $w(v_{s-1}v_s) = w(v_rv_s)$ .

Proof. By the choice of  $v_k, v_1v_{s-1} \notin E(G)$  and  $v_1v_s \notin E(G)$ . So, if  $v_{s-1}v_r \in E(G)$ , then  $\{v_r, v_{s-1}, v_s, v_1\}$  induces a modified claw, which implies that  $w(v_{s-1}v_s) = w(v_rv_s)$ . By the choice of  $v_s, v_rv_{s+1} \notin E(G)$ . So, if  $v_{s-1}v_r \notin E(G)$ , then  $\{v_s, v_{s-1}, v_{s+1}, v_r\}$  induces a claw or a modified claw. Thus we have  $w(v_{s-1}v_s) = w(v_rv_s)$ .

It follows from Claims 1.2 and 1.3 that  $v_{s-1}v_{s-2}\cdots v_{r+1}v_1v_2\cdots v_rv_s$   $v_{s+1}\cdots v_p$  is a heaviest longest path different from P and with  $v_p$  as one of its end-vertices. At the same time, by the choice of  $v_k$ ,  $v_1v_{s-1}\notin E(G)$ . From Claim 0 we have  $w(C)\geq 2m/3$ .

#### Case 1.2. s = k + 1.

From the choice of the path Q and the connectedness of  $G-v_{k+1}$ , there exists a path R such that

- R has end-vertices  $v_k$  and  $v_t$  with  $k+2 \le t < p$ , and
- $\bullet \ V(R) \cap V(P) = \{v_k, v_t\}.$

Choose R such that t is as large as possible. Similar to the proof of Claim 1.1, we have the following claim:

Claim 1.4. 
$$v_k v_t \in E(G)$$
.

Claim 1.5.  $w(v_1v_{r+1}) = w(v_rv_{r+1})$ .

*Proof.* Suppose r < k - 1. By the same proof as for Claim 1.2, we get  $w(v_1v_{r+1}) = w(v_rv_{r+1})$ .

Suppose r = k - 1. By the choices of  $v_k$  and  $v_s$ ,  $v_1v_t \notin E(G)$  and  $v_{k-1}v_t \notin E(G)$ . So  $\{v_k, v_{k-1}, v_1, v_t\}$  induces a modified claw, so we get  $w(v_1v_k) = w(v_{k-1}v_k)$ .

Claim 1.6.  $w(v_k v_t) = w(v_{t-1} v_t)$ .

*Proof.* Suppose  $v_k v_{t-1} \notin E(G)$ . By the choice of  $v_t$ ,  $v_k v_{t+1} \notin E(G)$ , then  $\{v_t, v_{t-1}, v_{t+1}, v_k\}$  induces a claw or a modified claw. So we get  $w(v_k v_t) = w(v_{t-1}v_t)$ .

Suppose  $v_k v_{t-1} \in E(G)$ . By the choice of  $v_k$ ,  $v_1 v_{t-1} \notin E(G)$  and  $v_1 v_t \notin E(G)$ , then  $\{v_k, v_{t-1}, v_t, v_1\}$  induces a modified claw. So we get  $w(v_k v_t) = w(v_{t-1} v_t)$ .

Claim 1.7.  $w(v_r v_{k+1}) = w(v_k v_{k+1})$ .

*Proof.* Suppose  $v_r v_k \notin E(G)$ . By the choice of  $v_s$ ,  $v_r v_{k+2} \notin E(G)$ . So  $\{v_{k+1}, v_k, v_{k+2}, v_r\}$  induces a claw or a modified claw. Thus  $w(v_r v_{k+1}) = w(v_k v_{k+1})$ .

Suppose  $v_rv_k \in E(G)$ . By the choices of  $v_k$  and  $v_t$ ,  $v_1v_{k+1} \notin E(G)$ ,  $v_1v_t \notin E(G)$  and  $v_rv_t \notin E(G)$ . So  $\{v_k, v_{k+1}, v_t, v_1\}$  induces a claw or a modified claw and  $\{v_k, v_1, v_r, v_t\}$  induces a modified claw. Thus  $w(v_1v_k) = w(v_kv_{k+1})$  and  $w(v_1v_k) = w(v_1v_r)$ . So  $w(v_1v_r) = w(v_1v_k) = w(v_kv_{k+1})$ . We conclude that  $w(v_rv_{k+1}) = w(v_1v_r)$ . Otherwise, apply Lemma 3 (ii) to the subgraph induced by  $\{v_1, v_r, v_{k+1}, v_k, v_t\}$ . Since there must be a vertex  $v_{t'} \in V(P[v_{k+2}, v_p]) \setminus \{v_t\}$  such that  $v_{k+1}v_{t'} \in E(G)$ , we have  $v_{t'}$  is in the component of  $G - \{v_1, v_k, v_t\}$ . Thus we have  $v_{t'}v_1 \in E(G)$ , contradicting the choice of  $v_k$ . So we get  $w(v_rv_{k+1}) = w(v_kv_{k+1})$ .

It follows from Claims 1.5, 1.6 and 1.7 that  $v_{t-1}v_{t-2}\cdots v_{k+1}v_rv_{r-1}\cdots v_1$   $v_{r+1}\cdots v_kv_t$   $v_{t+1}\cdots v_p$  is a heaviest longest path different from P and with  $v_p$  as one of its end-vertices. By the choice of  $v_k$ ,  $v_1v_{t-1} \notin E(G)$ . Then from Claim 0, we have  $w(C) \geq 2m/3$ . This completes the proof of Case 1.

Case 2.  $v_1v_i \notin E(G)$  for some i with  $r \leq i < k$ .

Choose  $v_l \notin N(v_1)$  with  $r \leq l < k$  such that l is as large as possible. It is clear that  $3 \leq l < k$  and  $v_1v_i \in E(G)$  for every i with  $l < i \leq k$ . Let j be the smallest index such that j > l and  $v_j \notin N(v_1) \cap N(v_l)$ . Since  $v_{l+1} \in N(v_1) \cap N(v_l)$ , we have  $j \geq l+2$ . Also, it is obvious that  $j \leq k+1$ .

If  $w(v_1v_{l+1}) = w(v_lv_{l+1})$ , then  $v_lv_{l-1} \cdots v_1v_{l+1}v_{l+2} \cdots v_p$  is a heaviest longest path different from P and with  $v_p$  as one of its end-vertices. By the choice of  $v_l$ ,  $v_1v_l \notin E(G)$ . Then from Claim 0, we have  $w(C) \geq 2m/3$ . From now on, we have the following assumption.

Assumption 1.  $w(v_1v_{l+1}) \neq w(v_lv_{l+1})$ .

Claim 2.1.  $v_{l+1}v_j \notin E(G)$ .

*Proof.* If  $v_{l+1}v_j \in E(G)$ , then  $v_j \in N(v_1) \cap N(v_l)$  by Lemma 2, contradicting the choice of  $v_j$ .

Claim 2.2.  $v_{l+1}v_{j-1} \notin E(G)$  and  $v_{l+2}v_{j} \notin E(G)$ .

*Proof.* Suppose  $v_{l+1}v_{j-1} \in E(G)$ . If  $j \leq k$ , then by the choice of  $v_j$ , we have  $v_lv_j \notin E(G)$ . Apply Lemma 3 (ii) to the subgraph induced by  $\{v_1,v_{l+1},v_l,v_{j-1},v_j\}$ . Since  $v_s$  is in the component of  $G-\{v_l,v_{j-1},v_j\}$ , we have  $v_sv_1 \in E(G)$ , contradicting the choice of  $v_k$ . So  $v_{l+1}v_{j-1} \notin E(G)$ . The case j=k+1 can be proved similarly.

Similar to the proof of above, we can prove that  $v_{l+2}v_j \notin E(G)$ .

From Claim 2.2, we know that  $l+4 \le j \le k+1$ .

Case 2.1.  $j \leq k$ .

Claim 2.3.  $w(v_{l+1}v_{l+2}) = w(v_1v_{l+2}) = w(v_{j-1}v_j) = w(v_1v_j)$ .

*Proof.* By Claims 2.1 and 2.2, we know that both  $\{v_1, v_{l+1}, v_{l+2}, v_j\}$  and  $\{v_1, v_{j-1}, v_j, v_{l+1}\}$  induce modified claws. Thus  $w(v_{l+1}v_{l+2}) = w(v_1v_{l+2}) = w(v_{j-1}v_j) = w(v_1v_j)$ .

It follows from Claim 2.3 that each of  $v_{l+1}v_l \cdots v_1 v_{l+2}v_{l+3} \cdots v_p$  and  $v_{j-1}v_{j-2} \cdots v_1 v_j \ v_{j+1} \cdots v_p$  is a heaviest longest path with  $v_p$  as one of its end-vertices. By Claim 2.2 and Claim 0, we have  $w(C) \geq 2m/3$ .

Case 2.2. j = k + 1.

Claim 2.4.  $w(v_1v_k) = w(v_lv_k) = w(v_kv_{k+1})$ .

*Proof.* By the choices of  $v_k$  and  $v_l$ ,  $v_1v_{k+1} \notin E(G)$  and  $v_1v_l \notin E(G)$ . So  $\{v_k, v_{k+1}, v_l, v_1\}$  induces a claw or a modified claw, thus  $w(v_1v_k) = w(v_lv_k) = w(v_kv_{k+1})$ .

Claim 2.5. For any vertex  $v \in N(v_1) \cap N(v_l) \setminus \{v_{l+1}, v_k\}$ , we have  $vv_k \in E(G), vv_{l+1} \in E(G)$  and  $vv_{k+1} \notin E(G)$ .

Proof. From Assumption 1 and Claim 2.4, we know that  $w(v_lv_{l+1}) = w(v_lv_k)$  and  $w(v_lv_{l+1}) = w(v_1v_k)$  cannot hold at the same time. Suppose  $w(v_lv_{l+1}) \neq w(v_lv_k)$ . Then applying Lemma 2 to the induced path  $v_kv_lv_{l+1}$  and  $v \in N(v_l) \setminus \{v_k, v_{l+1}\}$ , we get  $vv_k \in E(G)$  and  $vv_{l+1} \in E(G)$ . Now if  $vv_{k+1} \in E(G)$ , then apply Lemma 3 (ii) to the subgraph induced by  $\{v_k, v_l, v_{l+1}, v, v_{k+1}\}$ . Since  $v_l$  is in the component of  $G - \{v_{l+1}, v, v_{k+1}\}$ , we have  $v_lv_l \in E(G)$ , contradicting the choice of  $v_l$ . So  $vv_{k+1} \notin E(G)$ . The case  $w(v_lv_{l+1}) \neq w(v_lv_k)$  can be proved similarly.

Claim 2.6.  $w(v_1v_{k-1}) = w(v_{k-1}v_k)$ .

*Proof.* By Claim 2.5,  $v_{k-1}v_{k+1} \notin E(G)$ . So  $\{v_k, v_{k-1}, v_1, v_{k+1}\}$  induces a modified claw, thus  $w(v_1v_{k-1}) = w(v_{k-1}v_k)$ .

Suppose  $v_1v_{l-1} \in E(G)$ . Then by Claim 2.5, we have  $v_{l-1}v_k \in E(G)$ ,  $v_{l-1}v_{l+1} \in E(G)$  and  $v_{l-1}v_{k+1} \notin E(G)$ . So  $\{v_k, v_1, v_{l-1}, v_{k+1}\}$  induces a modified claw, thus  $w(v_1v_{l-1}) = w(v_{l-1}v_k)$ . We claim that  $w(v_1v_{l-1}) = w(v_{l-1}v_l)$ . Otherwise, apply Lemma 3 (ii) to the subgraph induced by  $\{v_l, v_{l-1}, v_1, v_k, v_{k+1}\}$ . Since  $v_{l+1}$  is in the component of  $G - \{v_1, v_k, v_{k+1}\}$ , we have  $v_{l+1}v_{k+1} \in E(G)$ , contradicting Claim 2.1. So we get  $w(v_{l-1}v_l) = w(v_{l-1}v_k)$ . Therefore, from Claim 2.6 we know that  $v_lv_{l+1} \cdots v_{k-1}v_1v_2 \cdots v_{l-1}v_kv_{k+1} \cdots v_p$  is a heaviest longest path different from P and with  $v_p$  as one of its end-vertices. By the choice of  $v_l, v_1v_l \notin E(G)$ . Then from Claim 0, we have  $w(C) \geq 2m/3$ .

Suppose  $v_1v_{l-1} \notin E(G)$ , we have  $v_{l-1}v_{l+1} \notin E(G)$ . Otherwise,  $\{v_{l+1}, v_l, v_{l-1}, v_1\}$  induces a modified claw, contradicting the assumption  $w(v_1v_{l+1}) \neq w(v_lv_{l+1})$ . From Claim 2.2, we know that  $\{v_l, v_{l-1}, v_k, v_{l+1}\}$  induces a claw or a modified claw. So  $w(v_{l-1}v_l) = w(v_lv_k)$ . Therefore, from Claim 2.6, we know that  $v_{l-1}v_{l-2}\cdots v_1v_{k-1}$   $v_{k-2}\cdots v_lv_kv_{k+1}\cdots v_p$  is a heaviest longest path different from P and with  $v_p$  as one of its end-vertices. Then from Claim 0, we have  $w(C) \geq 2m/3$ .

The proof is complete.

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