

The heterochromatic matchings in edge-colored bipartite graphs *

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Abstract

Let (G, C) be an edge-colored bipartite graph with bipartition (X, Y) . A heterochromatic matching of G is such a matching in which no two edges have the same color. Let $N^c(S)$ denote a maximum color neighborhood of $S \subseteq V(G)$.

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We show that if $|N^c(S)| \geq |S|$ for all $S \subseteq X$, then G has a heterochromatic matching with cardinality at least $\lceil \frac{|X|}{3} \rceil$. We also obtain that if $|X| = |Y| = n$ and $|N^c(S)| \geq |S|$ for all $S \subseteq X$ or $S \subseteq Y$, then G has a heterochromatic matching with cardinality at least $\lceil \frac{3n}{8} \rceil$.

Keywords: heterochromatic matching, color neighborhood

1 Introduction and notation

We use [3] for terminology and notations not defined here and consider simple undirected graphs only.

Let $G = (V, E)$ be a graph. An *edge coloring* of G is a function $C : E \rightarrow N$ (N is the set of nonnegative integers). If G is assigned such a coloring C , then we say that G is an *edge colored graph*, or simply *colored graph*. Denote by (G, C) the graph G together with the coloring C and by $C(e)$ the *color* of the edge $e \in E$. For a subgraph H of G , let $C(H) = \{C(e) : e \in E(H)\}$.

A subgraph H of G is called *heterochromatic*, or *rainbow*, or *colorful* if its any two edges have different colors. There are many publications studying heterochromatic subgraphs. Very often the subgraphs considered are paths, cycles, trees, etc. The heterochromatic hamiltonian cycle or path problems were studied by Hahn and Thomassen (see [9]), Rödl and Winkler (see [7]), Frieze and Reed, Albert, Frieze and Reed (see [1]), and H. Chen and X.L. Li (see [5]). For more references, see [2, 6, 9].

For an uncolored graph the following theorems are well known in matching theory and have been widely used.

Theorem 1 [10]. Let G be a bipartite graph with bipartition (X, Y) . Then G contains a matching that saturates every vertex of X if and only if $|N(S)| \geq |S|$ for all $S \subseteq X$.

Theorem 2 [3]. A bipartite graph G has a perfect matching if and only if $|N(S)| \geq |S|$ for all $S \subseteq V$.

A matching is *heterochromatic* if any two edges of it have different colors. Unlike uncolored matchings for which the maximum matching problem is solvable in polynomial time (see [12]), the maximum heterochromatic matching problem is *NP*-complete, even for bipartite graphs (see [8]). Heterochromatic matchings have been studied for example in [11] in which by defining $N_c(S)$ (see the definition below) Hu and Li gave some sufficient conditions for the existence of perfect heterochromatic matchings in colored graphs.

Let (G, C) be a colored graph. For a vertex v of G , let $CN(v) = \{C(e) : e \text{ is incident with } v\}$ and $CN(S) = \cup_{v \in S} CN(v)$ for $S \subseteq V$. For $S \in V(G)$, denote $N_c(S)$ as one of the minimum set(s) W satisfying $W \subseteq N(S) \setminus S$ and $[CN(S) \setminus C(G[S])] \subseteq CN(W)$.

Theorem 3[11]. Let (B, C) be a colored bipartite graph with bipartition X, Y . Then, B contains a heterochromatic matching that saturates every vertex in X , if $|N_c(S)| \geq |S|$, for all $S \subseteq X$.

Theorem 4[11]. A colored graph (G, C) has a perfect heterochromatic matching, if

- (1) $o(G - S) \leq |S|$, where $o(G - S)$ denotes the number of odd components in the remaining graph $G - S$, and
- (2) $|N_c(S)| \geq |S|$ for all $S \subseteq V$ such that $0 \leq |S| \leq \lfloor \frac{|G|}{2} \rfloor$ and $|N(S) \setminus S| \geq |S|$.

We define a maximum color neighborhood and study heterochromatic matchings in colored bipartite graphs under a new condition related to maximum color neighborhoods of subsets of vertices.

Let (G, C) be a colored bipartite graph with bipartition (X, Y) . For a vertex set $S \subseteq X$ or Y , a *color neighbourhood* of S is defined as a set $T \subseteq N(S)$ such that there are $|T|$ edges between S and T that are adjacent to distinct vertices of T and have distinct colors. A *maximum color neighborhood* $N^c(S)$ is a color neighborhood of S with maximum size. Given a set S and a color neighborhood T of S , denote by $C(S, T)$ a set of $|T|$ distinct colors on the $|T|$

edges between S and distinct vertices of T . Note that there might be more than one such set $C(S, T)$. If there is no ambiguity, let $C(S, T)$ be a fixed color set in the following.

Let M be a heterochromatic matching of G , we denote $b_M = \{c : c \in C(M) \text{ and there exists an edge } e \in E(G - V_M) \text{ such that } C(e) = c\}$ and denote by $(X_M \cup Y_M)$ with $X_M \subseteq X, Y_M \subseteq Y$, the set of vertices that are incident with the edges in M .

For a given heterochromatic matching M , an *alternating 4-cycle* AC_M is a cycle $\{xy, yx', x'y', y'x\}$ such that $C(xy) = C(x'y')$ and $C(xy') = C(x'y) \notin C(M)$, in which $xy \in E(M), x'y' \in E(G - V_M)$.

The following main results are obtained in this paper.

Theorem 5. Let (G, C) be a colored bipartite graph with bipartition (X, Y) and $|N^c(S)| \geq |S|$ for all $S \subseteq X$, then G has a heterochromatic matching of cardinality at least $\lceil \frac{|X|}{3} \rceil$.

Theorem 6. Let (G, C) be a colored bipartite graph with bipartition (X, Y) and $|X| = |Y| = n$. If $|N^c(S)| \geq |S|$ for all $S \subseteq X$ or $S \subseteq Y$, then G has a heterochromatic matching of cardinality at least $\lceil \frac{3n}{8} \rceil$.

Under the conditions of Theorem 6, the following example shows that the best bound can not be better than $\lceil \frac{n}{2} \rceil$. Let $G = (X, Y)$ with $X = \{x_1, x_1, \dots, x_{2s}\}$ and $Y = \{y_1, y_2, \dots, y_{2s}\}$ be a bipartite graph such that $E(G) = \{x_i y_i | i = 1, 2, \dots, 2s\} \cup \{x_{2i-1} y_{2i} | i = 1, 2, \dots, s\} \cup \{x_{2i} y_{2i-1} | i = 1, 2, \dots, s\}$. The edge coloring C of G is given by $C(x_{2i-1} y_{2i-1}) = C(x_{2i} y_{2i}) = 2i - 1$ and $C(x_{2i-1} y_{2i}) = C(x_{2i} y_{2i-1}) = 2i$ for $i = 1, 2, \dots, s$. Clearly the cardinality of the maximum heterochromatic matching of (G, C) is $s = \lceil \frac{2s}{2} \rceil$. This example shows that the bound in Theorem 6 is not very far away from the best.

2 Proof of Theorem 5

Let M be a maximum heterochromatic matching of G . Put $S = X - X_M$. Let $N^c(S)$ be a maximum color neighborhood of S . And

write $N^c(S) = Y_P \cup Y_Q (Y_P \cap Y_Q = \phi)$, where $C(S, Y_P) \cap C(M) = \phi$ and $C(S, Y_Q) \subseteq C(M)$. Clearly $|Y_Q| \leq |M|$.

If $Y_P \not\subseteq Y_M$, then there is an edge $e \in E(X - X_M, Y - Y_M)$ and $C(e) \notin C(M)$. Hence $M \cup \{e\}$ is a heterochromatic matching with cardinality $|M| + 1$, contrary to the maximality of M .

So $Y_P \subseteq Y_M$. Since $|N^c(S)| = |Y_P| + |Y_Q| \geq |S|$, it follows that $|M| = |Y_M| \geq |Y_P| \geq |S| - |Y_Q| \geq |X| - |M| - |M|$. This gives $|M| \geq \lceil \frac{|X|}{3} \rceil$. \square

3 Proof of Theorem 6

Let M be a maximum heterochromatic matching of G with $t := |M|$ such that $|b_M|$ is maximum. Assume to the contrary that $t < \frac{3n}{8}$.

Suppose that the maximum number of the vertex-disjoint AC_M s in G is l_1 . Clearly $0 \leq l_1 \leq t$. If $l_1 \geq 1$, assume that the i -th alternating cycle $AC_M^i = \{x_i y_i, y_i x'_i, x'_i y'_i, y'_i x_i\}$ and denote that $C(xy) = C(x'_i y'_i) = c_i \in C(M)$, $C(x_i y'_i) = C(x'_i y_i) = c'_i \notin C(M)$, where $xy \in E(M)$, and $x'_i \in X - X_M, y_i \in Y - Y_M$.

Denote

$$\begin{aligned} X_{L_1} &= \{x'_1, x'_2, \dots, x'_{l_1}\}, Y_{L_1} = \{y'_1, y'_2, \dots, y'_{l_1}\}, \\ X_{M_{l_1}} &= \{x_1, x_2, \dots, x_{l_1}\} \subseteq X_M, \\ Y_{M_{l_1}} &= \{y_1, y_2, \dots, y_{l_1}\} \subseteq Y_M, \end{aligned}$$

where $\{x_1 y_1, x_2 y_2, \dots, x_{l_1} y_{l_1}\} = E(M_{l_1}) \subseteq E(M)$.

We define a procedure named *breeding* as follows.

For $i \geq 1$, if there is an edge $x^i y^i \in E(M - M_{l_1+i-1})$ such that x^i is adjacent with $y^{i'}$, y^i is adjacent with $x^{i'}$ and $C(x^{i'} y^{i'}) = C(x^i y^i)$, where $x^{i'} \in X - X_M - X_{L_1+i-1}, y^{i'} \in Y - Y_M - Y_{L_1+i-1}$; Moreover, if $C(x^i y^{i'}) \in C(M)$, then $C(x^i y^{i'}) \in C(M_{l_1+i-1})$ and if $C(x^{i'} y^i) \in C(M)$, then $C(x^{i'} y^i) \in C(M_{l_1+i-1})$; Then we do the i -th breeding (*step i*) as follows.

Denote $c_{l_1+i} = C(x^i y^i), x_{l_1+i} = x^i, y_{l_1+i} = y^i, x'_{l_1+i} = x^{i'}, y'_{l_1+i} = y^{i'}$. And let $X_{M_{l_1+i}} = X_{M_{l_1+i-1}} \cup \{x_{l_1+i}\}, Y_{M_{l_1+i}} = Y_{M_{l_1+i-1}} \cup$

$\{y_{l_1+i}\}$, $X_{L_1+i} = X_{L_1+i-1} \cup \{x^{i'}\}$, $Y_{L_1+i} = Y_{L_1+i-1} \cup \{y^{i'}\}$ and $M_{l_1+i} = M_{l_1+i-1} \cup \{x_{l_1+i}y_{l_1+i}\}$.

Suppose that after k steps, the breeding procedure stops, which means that we can not do the breeding again. Let $l = l_1 + k$ and denote $X_L = X_{L_1+k}$, $Y_L = Y_{L_1+k}$, $X_{M_l} = X_{M_{l_1+k}}$, $Y_{M_l} = Y_{M_{l_1+k}}$. Denote $C_l = \{c_1, c_2, \dots, c_l\}$ and $C_L = \{c : c \in C(AC_M^i) \text{ and } c \notin C_l, 1 \leq i \leq l\}$. Clearly $C(M) - C(M_l) = C(M - M_l)$. Suppose $|C_L| = l'$.

Claim 1. $l' \leq l$.

Proof. In each step i ($i \geq 1$), we conclude that $C(x^i y^{i'}) \in C(M)$ or $C(x^{i'} y^i) \in C(M)$. Otherwise if $C(x^i y^{i'}) \neq C(x^{i'} y^i)$, then $M^1 = M \cup \{x^i y^{i'}, x^{i'} y^i\} - \{x^i y^i\}$ is a heterochromatic matching with size more than $|M|$, a contradiction. If $C(x^i y^{i'}) = C(x^{i'} y^i)$, then $\{x^i y^i, y^i x^{i'}, x^{i'} y^{i'}, y^{i'} x^i\}$ is an AC_M , it follows that the number of the vertex disjoint AC_M s is at least $l_1 + 1$, a contradiction. Thus, $l' \leq l$. \square

Now denote $S_x = X - X_M - X_L$ and $S_y = Y - Y_M - Y_L$.

Claim 2. (2.1). For any vertex $v_x \in S_x$, if there exists an edge $v_x v_y$ such that $C(v_x v_y) \notin C(M - M_l)$, then $v_y \in Y_M - Y_{M_l}$.

(2.2). For any vertex $y \in S_y$, if there exists an edge xy such that $C(xy) \notin C(M - M_l)$, then $x \in X_M - X_{M_l}$.

Proof. By the symmetry, we only prove (2.1). Suppose, to the contrary, there exists an edge $v_x v_y$ such that $C(v_x v_y) \notin C(M - M_l)$, in which $v_x \in S_x$ and $v_y \notin Y_M - Y_{M_l}$. We distinguish the following three cases, and in each case we get a heterochromatic matching with cardinality more than $|M|$, which is a contradiction.

Case 1. $C(v_x v_y) \notin C(M)$. Then let

$$M^1 = \begin{cases} M \cup \{v_x v_y\} & v_y \in Y - Y_M; \\ M \cup \{v_x v_y, x'_i y'_i\} - \{x_i y_i\} & v_y \in Y_{M_l}, \text{ w.l.o.g, suppose } v_y = y_i. \end{cases}$$

Case 2. $C(v_x v_y) \in C_l$ and $v_y \notin Y_{M_l}$. W.o.l.g, suppose $C(v_x v_y) = c_i, 1 \leq i \leq l$.

If $C(x'_iy_i) \notin C(M)$, then let $M^1 = M \cup \{v_x v_y, x'_iy_i\} - \{x_i y_i\}$. Otherwise, by the breeding rule, we have that $C(x'_iy_i) = c_{j_1}$ with $j_1 < i$. And if $C(x'_{j_1} y_{j_1}) \notin C(M)$, then let $M^1 = M \cup \{v_x v_y, x'_iy_i, x_{j_1} y_{j_1}\} - \{x_i y_i, x_{j_1} y_{j_1}\}$. Otherwise, we have that $C(x'_{j_1} y_{j_1}) = c_{j_2}$ with $j_2 < j_1$. Then consider $C(x'_{j_2} y_{j_2}), \dots$. By successively doing this process, we can find $j_k (1 \leq j_k < j_{k-1} < \dots < j_1 < i)$ such that $C(x'_{j_k} y_{j_k}) \notin C(M)$, then let $M^1 = M \cup \{v_x v_y, x'_iy_i, x'_{j_1} y_{j_1}, \dots, x'_{j_{k-1}} y_{j_{k-1}}, x'_{j_k} y_{j_k}\} - \{x_i y_i, x_{j_1} y_{j_1}, \dots, x_{j_k} y_{j_k}\}$.

Case 3. $C(v_x v_y) \in C_l$ and $v_y \in Y_{M_l}$. Suppose $v_y = y_j, 1 \leq j \leq l$. Let $M' = M \cup \{x'_j y'_j\} - \{x_j y_j\}$, then use the same method as in Case 2, we can get a heterochromatic matching M^1 with size more than $|M|$. \square

If $b_M \setminus C_l \neq \phi$, w.o.l.g, we assume that $b_M \setminus C_l = \{c_{l+1}, c_{l+2}, \dots, c_{l+d}\}$ and moreover assume that $C(x_{l+i} y_{l+i}) = c_{l+i}, 1 \leq i \leq d$. And denote $X_{M_d} = \{x_{l+1}, x_{l+2}, \dots, x_{l+d}\}$, $Y_{M_d} = \{y_{l+1}, y_{l+2}, \dots, y_{l+d}\}$ and $M_d = \{x_{l+1} y_{l+1}, \dots, x_{l+d} y_{l+d}\}$.

Claim 3. If there exists an edge $v_x y_i$ such that $C(v_x y_i) \notin C(M - M_l)$, where $v_x \in S_x$ and $y_i \in Y_{M_d}$, then the following two conclusions hold.

(3.1). The edges in $E(G - V_M)$ with color c_i are incident with v_x .

(3.2). If there is an edge $x_i v_y$, where $v_y \in S_y$, then $C(x_i v_y) \in C(M - M_l)$.

Proof. Since $y_i \in Y_{M_d}$, by the definition of Y_{M_d} , there exists an edge $e \in E(G - V_M)$ such that $C(e) = C(x_i y_i)$. If e is not incident with v_x , we distinguish the following two cases.

Case 1. $C(v_x y_i) \notin C(M)$. Then let $M^1 = M \cup \{e, v_x y_i\} - \{x_i y_i\}$. Then M^1 is a heterochromatic matching M^1 with $|M^1| > |M|$, a contradiction.

Case 2. $C(v_x y_i) \in C_l$. Suppose $C(v_x v_y) = c_j, 1 \leq j \leq l$. Similarly as in Claim 2, we can find j_k such that $C(x_{j_k} y'_{j_k})$ (or $C(x'_{j_k} y_{j_k})$) $\notin C(M)$, then let $M^1 = M \cup \{e, v_x y_i, x'_j y_j$ (or $x_j y'_j$), $x'_{j_1} y_{j_1}$ (or $x_{j_1} y'_{j_1}$), $\dots, x'_{j_k} y_{j_k}$ (or $x_{j_k} y'_{j_k}$) $\} - \{x_i y_i, x_j y_j, x_{j_1} y_{j_1}, \dots, x_{j_k} y_{j_k}\}$. Here if e is incident with x'_{j_r} , we choose the edge $x_{j_r} y'_{j_r}$, where $0 \leq r \leq k$ and $j_0 = j$. Thus, we get a heterochromatic matching

M^1 such that $|M^1| > |M| + 1$, a contradiction. This completes the proof of (3.1).

For (3.2), suppose, to the contrary, $C(x_i v_y) \notin C(M - M_l)$. Similarly as in (3.1), the edges in $E(G - V_M)$ with color $C(x_i y_i)$ are incident with v_y . Thus $C(v_x v_y) = C(x_i y_i)$.

If $C(x_i v_y), C(v_x y_i) \notin C(M)$ and $C(x_i v_y) \neq C(v_x y_i)$, then $M^1 = M \cup \{x_i v_y, v_x y_i\} - \{x_i y_i\}$ is a heterochromatic matching such that $|M^1| > |M|$, a contradiction.

If $C(x_i v_y), C(v_x y_i) \notin C(M)$ and $C(x_i v_y) = C(v_x y_i)$, then $\{x_i y_i, y_i v_x, v_x v_y, v_y x_i\}$ is an AC_M . Thus the number of the vertex disjoint AC_M is at least $l_1 + 1$, a contradiction.

Thus we conclude that $C(x_i v_y) \in C_l$ and $C(v_x y_i) \notin C(M)$ or $C(v_x y_i) \in C_l$ and $C(x_i v_y) \notin C(M)$ or $C(x_i v_y), C(v_x y_i) \in C_l$, it follows that we can continue the breeding procedure, a contradiction. This completes the proof of (3.2). \square

Let $Y_{M_{d_1}} = \{v : v \in Y_{M_d} \text{ and } v \text{ is adjacent with a vertex } v_x \in S_x \text{ such that } C(v_x v) \notin C(M - M_l)\}$ and $|Y_{M_{d_1}}| = d_1$. And define $X_{d_1}' = \{v : v \in S_x \text{ and } v \text{ is adjacent with a vertex } v_y \in Y_{M_d} \text{ such that } C(v v_y) \notin C(M - M_l)\}$ and $|X_{d_1}'| = d_1'$. By Claim 3, $d_1' \leq d_1$.

For a vertex $v \in V(G - V_M)$, let $b_{M_d}(v) = \{c : c \in C(M_d) \text{ and there exists an edge } e \in E(G - V_M) \text{ incident with } v \text{ such that } C(e) = c\}$. And for a subset $V_1 \subseteq V(G - V_M)$, denote $b_{M_d}(V_1) = \{b_{M_d}(v) : v \in V_1\}$. Let Y' denote the set $V \subseteq S_y$ with minimum size satisfying $b_{M_d}(V) = b_{M_d}(S_y)$. Clearly, $|Y'| \leq d$. Now denote $Y'' = Y - Y_M - Y_L - Y'$.

Let $N^c(Y'')$ be a maximum color neighborhood of Y'' . And assume that $N^c(Y'') = X_P \cup X_Q (X_P \cap X_Q = \phi)$, where $C(Y'', X_P) \cap (C(M) \cup C_L) = \phi$ and $C(Y'', X_Q) \subseteq C(M) \cup C_L$. Clearly $|X_Q| \leq t + l'$, then $|X_P| \geq |Y''| - |X_Q| \geq n - t - l - d - (t + l') \geq n - 2t - 2l - d$.

Let $N^c(S_x)$ be a maximum color neighborhood of S_x . And assume that $N^c(S_x) = Y_P \cup Y_Q (Y_P \cap Y_Q = \phi)$, where $C(S_x, Y_P) \cap$

$C(M - M_l) = \phi$ and $C(S_x, Y_Q) \subseteq C(M - M_l)$. Clearly $|Y_Q| \leq t - l$, then $|Y_P| \geq |S_x| - |Y_Q| \geq n - t - l - (t - l) \geq n - 2t$.

Claim 4. $d_1 \geq 2n - 5t - l$.

Proof. By Claim 2, we know that $X_P \subseteq X_M - X_{M_l}$ and $Y_P \subseteq Y_M - Y_{M_l}$. By the definition of Y'' and Claim 3, it holds that $X_P \cap X_{M_d} = \phi$. Thus $X_P \subseteq X_M - X_{M_l} - X_{M_d}$. By Claim 3 and the definition of $Y_{M_{d_1}}$, we conclude that $Y_P \cap (Y_{M_d} \setminus Y_{M_{d_1}}) = \phi$. So we have that $Y_P \subseteq Y_M - Y_{M_l} - (Y_{M_d} \setminus Y_{M_{d_1}})$.

If $|X_P| + |Y_P| > |Y_M - Y_{M_l} - (Y_{M_d} \setminus Y_{M_{d_1}})|$, then there is an edge $x_j y_j \in E(M - M_l - M_d)$ such that $x_j \in X_P$ and $y_j \in Y_P$. Then there exist edges $x_j v_y, v_x y_j$ such that $C(x_j v_y) \notin C(M) \cup C_L$ and $C(v_x y_j) \notin C(M - M_l)$, in which $v_x \in S_x$ and $v_y \in Y''$.

If $C(v_x y_j) \in C_l$, suppose that $C(v_x y_j) = c_i (1 \leq i \leq l)$. Use the same method as in Claim 2, we can find j_k such that $C(x_{j_k} y'_{j_k}) \notin C(M)$, then let $M^1 = M \cup \{x_j v_y, v_x y_j, x'_i y_i, x'_{j_1} y_{j_1}, \dots, x'_{j_k} y_{j_k}\} - \{x_j y_j, x_i y_i, x_{j_1} y_{j_1}, \dots, x_{j_k} y_{j_k}\}$. Thus M^1 is a heterochromatic matching such that $|M^1| > |M|$, which is a contradiction.

If $C(v_x y_j) \notin C_l$ and $C(v_x y_j) \neq C(x_j v_y)$. Then $M^1 = M \cup \{v_x y_j, x_j v_y\} - \{x_j y_j\}$ is a heterochromatic matching with cardinality $|M| + 1$, a contradiction.

So we conclude that $C(v_x y_j) = C(x_j v_y)$. Then let $M^1 = M \cup \{x_j v_y\} - \{x_j y_j\}$. It follows that $|M^1| = |M|$ and $|b_{M^1}| > |b_M|$, which contradicts with the choice of M .

Thus $|X_P| + |Y_P| \leq |Y_M - Y_{M_l} - (Y_{M_d} \setminus Y_{M_{d_1}})|$. It follows that $n - 2t - 2l - d + n - 2t \leq t - l - (d - d_1)$, then $d_1 \geq 2n - 5t - l$. \square

Similarly, we define $X_{M_{d_2}} = \{v : v \in X_{M_d} \text{ and } v \text{ is adjacent with a vertex } v_y \in S_y \text{ such that } C(v_y v) \notin C(M - M_l)\}$ and $|X_{M_{d_2}}| = d_2$. Similarly as Claim 4, we have that $d_2 \geq 2n - 5t - l$. By Claim 3, we know that $d_1 + d_2 \leq d$.

Let $S'_x = X - X_M - X_L - X_{d'_1}$ and $N^c(S'_x)$ be a maximum color neighborhood of S'_x . And assume that $N^c(S'_x) = Y'_P \cup Y'_Q (Y'_P \cap$

$Y'_Q = \phi$), in which $C(S'_x, Y'_P) \cap C(M - M_l) = \phi$ and $C(S'_x, Y'_Q) \subseteq C(M - M_l)$. Clearly, $|Y'_Q| \leq t - l$, then $|Y'_P| \geq |S'_x| - |Y'_Q| \geq n - t - l - (t - l) - d'_1 = n - 2t - d'_1 \geq n - 2t - d_1$.

Claim 5. $|Y_M - Y_{M_l} - Y_{M_d}| \geq n - 2t - d_1$.

Proof. By Claim 2, it holds that $Y'_P \subseteq Y_M - Y_{M_l}$. And by Claim 3 and the definition of S'_x , it holds that $Y'_P \cap Y_{M_d} = \phi$. Then it follows that $Y'_P \subseteq Y_M - Y_{M_l} - Y_{M_d}$. Thus we have that $|Y_M - Y_{M_l} - Y_{M_d}| \geq |Y'_P| \geq n - 2t - d_1$. \square

Now we have that

$$\begin{aligned} t &\geq |Y_M - Y_{M_l} - Y_{M_d}| + |Y_{M_d}| + |Y_{M_l}| \\ &\geq n - 2t - d_1 + d + l \\ &\geq n - 2t + 2n - 5t - l + l \\ &\geq 3n - 7t. \end{aligned}$$

Thus $t \geq \frac{3n}{8}$, which is a contradiction. The proof of Theorem 6 is complete.

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