

Note on the number of rooted complete N -ary trees

Stephan Dominique Andres
Zentrum für angewandte Informatik Köln
Weyertal 80, 50931 Köln, Germany
andres@zpr.uni-koeln.de

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Abstract

We determine a recursive formula for the number of rooted complete N -ary trees with n leaves, which generalizes the formula for the sequence of Wedderburn-Etherington numbers. The diagonal sequence of our new sequences equals to the sequence of numbers of rooted trees with $N + 1$ vertices.

Key words: parenthesis structure enumeration, complete N -ary rooted tree, N -ary operation, adder topology

1 Introduction

A problem occurring in hardware design is the following: Given an n -operand addition that has to be realized by a set of binary adders, how many possibilities are there to arrange the adders? [1] To be precise we do not care about commutative operations, which can be executed on the adders without changing the arrangement. In a mathematical language, we seek for the number of interpretations of x^n (or the number of ways to insert parentheses) when multiplication is commutative but not associative, or, from another point of view, we are looking for the number of isomorphism classes of n -leaf complete binary rooted trees (where every vertex has either 0 or 2 children). These numbers are known as Wedderburn-Etherington numbers, their sequence $(T_n)_n$ has the key A001190 in the On-Line Encyclopedia of Integer Sequences [2] and its generating function $B(x)$ satisfies the functional

equation

$$B(x) = x + \frac{1}{2} (B(x)^2 + B(x^2)),$$

cf. [3]. So, a recursive method to calculate the sequence is given by

$$\begin{aligned} T_1 &= 1, \\ T_{2n+2} &= \sum_{i=1}^n T_i \cdot T_{2n+2-i} + \frac{T_{n+1}(T_{n+1} + 1)}{2} \quad \text{for } n \geq 0, \\ T_{2n+1} &= \sum_{i=1}^n T_i \cdot T_{2n+1-i} \quad \text{for } n \geq 1. \end{aligned}$$

The aim of this paper is to generalize this formula from complete binary to complete N -ary trees (or from binary adders to N -ary adders as building blocks). Thus we want to determine the number $T_n^{(N)}$ of isomorphism classes of rooted trees with n leaves with the property that each vertex has either N or 0 children.

2 Recursive formula

To abbreviate notation, we call such a rooted complete N -ary tree with n leaves an n -tree whenever N is fixed. Let

$$p_i^{(D)} = \#\{(h_1, \dots, h_D) \mid \forall l : h_l \geq 1, \sum_{l=1}^D h_l = i\} = \binom{i-1}{D-1} \quad (1)$$

be the number of unordered partitions of i into D pieces. Then we may formulate

Theorem 1 For $N \geq 2$ and $n \geq 2$, $T_n^{(N)}$ can be calculated recursively via

$$\begin{aligned} T_1^{(N)} &= 1, \\ T_n^{(N)} &= \sum_{b=1}^N \sum_{\substack{(i_1, \dots, i_b) \\ \sum_{j=1}^b i_j = N \\ i_j \geq 1 \forall j}} \sum_{\substack{(k_1, \dots, k_b) \\ 1 \leq k_1 < k_2 < \dots < k_b \leq n \\ \sum_{j=1}^b i_j k_j = n}} \prod_{j=1}^b \sum_{D=1}^{i_j} \binom{i_j-1}{D-1} \binom{T_{k_j}^{(N)}}{D}. \end{aligned}$$

Proof. Clearly, $T_1^{(N)} = 1$. So consider an n -tree, $n > 1$. Then its root has N children that are roots of r_1 -, r_2 -, ..., r_N -trees, respectively, with $r_s \geq 1$ for

all s , and $\sum_{s=1}^N r_s = n$. r_s is called *input size* of the s -th child's tree. Since we do not care about permutations we may assume without loss of generality that $r_1 \leq r_2 \leq \dots \leq r_N$. We call the set of those r_j -trees with the same input size a *block*, and denote the number of blocks by b . The *block size* is the number of elements of a block. So we have blocks of (positive) block sizes i_1, \dots, i_b with respective input sizes k_1, \dots, k_b where w.l.o.g. $1 \leq k_1 < \dots < k_b \leq n$. We state that

$$\sum_{j=1}^b i_j k_j = n, \quad \sum_{j=1}^b i_j = N.$$

We will count at first the number of possibilities for a block with block size i and input size k , a number which we denote by $B(k, i)$. Then we derive the number $T(b, (i_1, \dots, i_b), (k_1, \dots, k_b))$ of trees for fixed (r_1, \dots, r_N) , i.e., for fixed b , block sizes i_1, \dots, i_b , and input sizes k_1, \dots, k_b . Note that for distinct values of the 3-tupel $(b, (i_1, \dots, i_b), (k_1, \dots, k_b))$ trees cannot be isomorphic, since isomorphic trees must have the same number of blocks, corresponding block sizes, and corresponding input sizes, as we assume $k_1 < \dots < k_b$. Therefore we obtain the total number of n -trees simply by adding $T(b, (i_1, \dots, i_b), (k_1, \dots, k_b))$ over all possibilities for $(b, (i_1, \dots, i_b), (k_1, \dots, k_b))$, i.e.,

$$T_n^{(N)} = \sum_{b=1}^N \sum_{\substack{(i_1, \dots, i_b) \\ \sum_{j=1}^b i_j = N \\ i_j \geq 1 \forall j}} \sum_{\substack{(k_1, \dots, k_b) \\ 1 \leq k_1 < k_2 < \dots < k_b \leq n \\ \sum_{j=1}^b i_j k_j = n}} T(b, (i_1, \dots, i_b), (k_1, \dots, k_b)). \quad (2)$$

Now we consider a block with i elements, each one having input size k . Let D be the number of distinct k -trees occurring in the block. We have

$$\binom{T_k^{(N)}}{D}$$

possibilities to choose D such structures, and $p_i^{(D)}$ possibilities to partition the k -trees of the block in D subblocks each one having equal k -trees as elements. These choices are independent in the sense of non-isomorphism, so in total we have

$$p_i^{(D)} \binom{T_k^{(N)}}{D}$$

possibilities, and

$$B(k, i) = \sum_{D=1}^i p_i^{(D)} \binom{T_k^{(N)}}{D}. \quad (3)$$

Now, what happens in different blocks is independent of each other block again, thus we conclude

$$T(b, (i_1, \dots, i_b), (k_1, \dots, k_b)) = \prod_{j=1}^b B(k_j, i_j). \quad (4)$$

Combining (1), (2), (3), and (4) yields the theorem. Note that in the right-hand-side of the recursive formula the expression $T_n^{(N)}$ does not occur, since $k_b < n$ whenever $N \geq 2$. \square

3 Final remarks

A very interesting sequence is the diagonal sequence

$$(T_{N^2}^{(N)})_{N=2,3,4,\dots} = 2, 4, 9, 20, 48, 115, \dots$$

Theorem 2 $T_{N^2}^{(N)}$ equals to the number of rooted trees with $N + 1$ vertices.

Proof. We construct an isomorphism between rooted complete N -ary trees with N^2 leaves and rooted trees with $N + 1$ vertices, $N \geq 2$, in the following way. Let T be a rooted complete N -ary tree with N^2 leaves. Delete all leaves to obtain a rooted tree with $N + 1$ vertices. (Note that there are always $N + 1$ inner vertices in T .) On the other hand, let R be a rooted tree with $N + 1$ vertices. So every vertex has at most N children. Construct a complete N -ary tree from R by adding children in such a way that every vertex from R has exactly N children, and every new vertex has no children. Obviously, these mappings are one-to-one. \square

Theorem 3 $\lim_{N \rightarrow \infty} T_{(n+1)(N-1)+1}^{(N)} = T_{n^2}^{(n)}$

Proof. Let T be a rooted complete N -ary tree with $(n + 1)(N - 1) + 1$ leaves, and $N > n$. This means that T has exactly $n + 1$ inner vertices, hence every inner vertex has at most n children which are inner vertices. So for every inner vertex we may delete $N - n$ of its children which are leaves to obtain a rooted complete n -ary tree with $(n + 1)((N - 1) - (N - n)) + 1 = n^2$ leaves. This construction can be reversed by adding children in an appropriate way. \square

References

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