

Extraconnectivity of Cartesian product graphs of paths*

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Abstract. Given a graph G and a non-negative integer g , the g -extraconnectivity of G (written $\kappa_g(G)$) is the minimum cardinality of a set of vertices of G , if any, whose deletion disconnects G , and every remaining component has more than g vertices. The usual connectivity and superconnectivity of G correspond to $\kappa_0(G)$ and $\kappa_1(G)$, respectively. In this paper, we determine $\kappa_g(P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s})$ for $0 \leq g \leq s$, where \times denotes the Cartesian product of graphs. We generalize $\kappa_g(Q_n)$ for $0 \leq g \leq n$, $n \geq 4$, where Q_n denotes the n -cube.

Key words: Network; Extraconnectivity; Cartesian product

1 Introduction

The topology of interconnected network is often modeled by a connected graph of communication links. In the network the connectivity $\kappa(G)$ is an important factor determining the reliability and fault tolerance of the network. Here, we consider the extraconnectivity which corresponds to a kind of conditional connectivity introduced by Harary [3].

Let G be a connected undirected graph, and \mathcal{P} be a graph-theoretic property, Harary [3] defined the conditional connectivity $\kappa(G; \mathcal{P})$ as the minimum cardinality of a set of vertices of G , if any, whose deletion disconnects G and every remaining component has property \mathcal{P} . Let g be a non-negative integer and let \mathcal{P}_g be the property of having more than g vertices. Fàbrega and Fiol [2] defined the g -extraconnectivity $\kappa_g(G)$ of G as $\kappa(G; \mathcal{P}_g)$. Given a graph G and a non-negative integer g , the g -extraconnectivity of G (written $\kappa_g(G)$) is the minimum cardinality of a set of vertices of G , if any, whose deletion disconnects G , and every remaining component has more than g vertices.

The Cartesian product of two graphs G and H , denoted by $G \times H$, is the graph with vertex set $V(G) \times V(H)$ such that two vertices (u_1, u_2) and (v_1, v_2) are adjacent iff either $u_1 = v_1$ with $u_2 v_2 \in E(H)$ or $u_2 = v_2$ with $u_1 v_1 \in E(G)$. Let

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$G = P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}$, where P_{n_i} denotes the path with n_i vertices. If $n_i = 2$ for all i , then G is s -cube, denoted by Q_s , which has been studied by [4,5,6,7,9]. We assign the vertices of each path P_{n_i} a natural ordering by $\{1, 2, \dots, n_i\}$. Thus we can use the s -dimensional array $x_1 x_2 \cdots x_s$ to denote the vertex of G , where $1 \leq x_i \leq n_i, i = 1, 2, \dots, s$. Clearly, any two vertices $x = x_1 x_2 \cdots x_s$ and $y = y_1 y_2 \cdots y_s$ are adjacent iff there is exactly an integer i such that $x_i = y_i + 1$ or $y_i - 1$ and $x_j = y_j$ for $j \neq i$. We use G_i^j to denote the subgraph induced by $\{x_1 \cdots x_{i-1} j x_{i+1} \cdots x_s | 1 \leq x_t \leq n_t, 1 \leq t \leq s, t \neq i\}$. By definition, we see that $G_i^1 \cong G_i^2 \cong \cdots \cong G_i^{n_i} \cong P_{n_1} \times \cdots \times P_{n_{i-1}} \times P_{n_{i+1}} \times \cdots \times P_{n_s}$. Sometimes we express G as $G_i^1 \odot G_i^2 \odot \cdots \odot G_i^{n_i}$. If $v = x_1 x_2 \cdots x_{i-1} 1 x_{i+1} \cdots x_s = v_i^1 \in V(G_i^1)$, we use v_i^j denote vertex $x_1 x_2 \cdots x_{i-1} j x_{i+1} \cdots x_s$. Clearly, $v_i^1 \cdots v_i^{n_i} \cong P_{n_i}$. Let $B \subset G_i^j$, we use B_i^t denote the subgraph of G_i^t which use t to instead of j of the i th coordinate of the vertices of B . Clearly, $B \cong B_i^t$. It is well known that $\kappa(G \times H) \geq \kappa(G) + \kappa(H)$, we have $\kappa(P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}) = s - 1$, and $\kappa_1(P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}) = 2s - 1$ for $n_i \geq 3, i = 1, 2, \dots, s$ (see [8]). In this paper, we derive $\kappa_g(P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s})$ when $g \leq s$ for $n_i \geq 2, i = 1, 2, \dots, s$. We abbreviate $P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}$ as G in the following sections. We use $N_G(v)$ to denote the set of the neighbors of v in G , $N_G(A)$ to denote the set $(\bigcup_{v \in V(A)} N_G(v)) \setminus V(A)$, $C_G(A)$ to denote the set $N_G(A) \cup V(A)$. We follow Bondy [1] for terminologies not given here.

2 Preliminaries

Before discussing the $\kappa_g(G)$, we give the following Lemmas.

Lemma 2.1. Assume $n_i \geq 3, i = 1, 2, \dots, s, G = P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}, A \subset G$. If $|V(A)| = g + 1, g \leq s$, then $|N_G(A)| \geq s + (s + 1)g - 2g - \binom{g}{2}$.

Proof. By induction on $|V(A)|$. Clearly, the result holds for $|V(A)| = 1$. Assume that the result holds for all A with $|V(A)| \leq h$. Next we show that the result is true for A with $|V(A)| = h + 1$. We directly use $g + 1$ instead of $h + 1$ in the following.

Let $A_j = G_i^j \cap A$. We first show that G can be decomposed into $G_i^1 \odot G_i^2 \odot \cdots \odot G_i^{n_i}$ for some i such that at least two subgraphs of $\{A_j, j = 1, 2, \dots, n_i\}$ are nonempty. Note that $|V(A)| \geq 2$, let $x = x_1 x_2 \cdots x_s$ and $y = y_1 y_2 \cdots y_s$ be two distinct vertices of A . Without loss of generality, we assume $x_s \neq y_s$, and $x_s = 1, y_s = 2$, thus $x \in V(G_i^1), y \in V(G_i^2)$, that is $V(A_1) \neq \emptyset, V(A_2) \neq \emptyset$.

Without loss of generality, we consider $G = G_s^1 \odot G_s^2 \odot \cdots \odot G_s^{n_s}$ in the following. Assume all the nonempty subgraphs of $\{A_j | A_j = G_s^j \cap A\}$ are A_1, A_2, \dots, A_m , and $|V(A_i)| = N_i$. Clearly, $N_i < |V(A)|$, by induction, $|N_{G_s^i}(A_i)| \geq s - 1 + s(N_i - 1) - 2(N_i - 1) - \binom{N_i - 1}{2}$. Since $N_{G_s^i}(A_i) \cap N_{G_s^j}(A_j) = \emptyset$ if $i \neq j$, it not difficult to see that $|N_G(A)| \geq \sum_{i=1}^m |N_{G_s^i}(A_i)| \geq s + (s + 1)g - 2g - \binom{g}{2}$ for $m \geq 3$, or $m = 2$ and $N_i \geq 2, i = 1, 2$. Next we assume that $m = 2$ and $N_i = 1$ (or $N_j = 1, j \neq i$). Without loss of generality, we consider $N_1 = 1$ (or $N_2 = 1$). If $N_2 = 1$, then $|N_G(A)| \geq 2s - 1 = s + (s + 1)g - 2g - \binom{g}{2}$. Assume that $N_2 \geq 2$. It is easy to see that $|N_{G_s^1}(A_1)| + |N_{G_s^2}(A_2)| \geq s + (s + 1)g - 2g - \binom{g}{2} - 1$ (since

$|N_{G_s^i}(A_i)| \geq s - 1 + s(N_i - 1) - 2(N_i - 1) - \binom{N_i - 1}{2}$). Note that $(A_2)_s^3 = N_{G_s^3}(A_2)$. Since $V((A_2)_s^3) \subset N_G(A)$ and $N_2 \geq 2$, we have that $|N_G(A)| \geq |N_{G_1^1}(A_1)| + |N_{G_2^2}(A_2)| + |V((A_2)_s^3)| \geq s + (s + 1)g - 2g - \binom{g}{2} - 1 + 2 > s + (s + 1)g - 2g - \binom{g}{2}$ (if $N_2 = 1$, we have that $|N_G(A)| \geq |N_{G_1^1}(A_1)| + |N_{G_2^2}(A_2)| + |V((A_2)_s^3)| \geq s + (s + 1)g - 2g - \binom{g}{2} - 1 + 1 = s + (s + 1)g - 2g - \binom{g}{2}$). \square

Corollary 2.2. Assume $G = P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}$, $n_i \geq 2$, $A \subset G$, $|V(A)| = g + 1$, $g \leq s$. Let k be the number of K_2 of $\{P_{n_i}, i = 1, 2, \dots, s\}$, where K_2 is the complete graph of two vertices, then $|N_G(A)| \geq s + (s + 1)g - 2g - \binom{g}{2} - \min\{g, k\}$.

Proof. Combining the structure of G and a similar argument of Lemma 2.1, the result holds. \square

Remark 2.3. Let $h_s(g) = s + (s + 1)g - 2g - \binom{g}{2}$, $h_s(g)$ is increasing when $0 \leq g \leq s - 1$, the maximum of $h_s(g)$ is $h_s(s - 1) = h_s(s) = \frac{s(s+1)}{2}$. By the proof of Lemma 2.1, we have $h_{s-1}(g_1) + \cdots + h_{s-1}(g_t) > h_s(g)$ when $0 \leq g_1, \dots, g_t \leq s$, $t \geq 3$ and $(g_1 + 1) + \cdots + (g_t + 1) \geq g + 1$. In particular, if $t = 2$, then $h_{s-1}(g_1) + h_{s-1}(g_2) > h_s(g)$ for $g_i \geq 2$, $i = 1, 2$ and $(g_1 + 1) + (g_2 + 1) \geq g + 1$ (see [9] for the detail).

Lemma 2.4. Let $G = G_s^1 \odot \cdots \odot G_s^{n_s}$, $B \subseteq G$, $|V(B)| \geq s$, $n_i \geq 3$, $i = 1, 2, \dots, s$. If there exists a subgraph G_s^m such that $V(B \cap G_s^m) = \emptyset$, then $|N_G(B)| \geq \frac{s(s+1)}{2}$.

Proof. Let $B_i = B \cap G_s^i$, $i = 1, 2, \dots, n_s$, without loss of generality, assume $B \cap G_s^{n_s} = \emptyset$. We verify the result by considering two cases.

Case 1. There exists a B_i such that $|V(B_i)| \geq s$.

Let T_{B_i} be a subgraph of B_i such that $|V(T_{B_i})| = s$. By Lemma 2.1, we have $|N_{G_s^i}(B_i)| \geq (s - 1) + s(s - 1) - 2(s - 1) - \binom{s-1}{2}$, that is $|C_{G_s^i}(B_i)| \geq (s - 1) + s(s - 1) - 2(s - 1) - \binom{s-1}{2} + s = \frac{s(s+1)}{2}$. Clearly, there are at least $\frac{s(s+1)}{2}$ internally disjoint paths between $C_{G_s^i}(B_i)$ and $G_s^{n_s}$. Therefore $|N_G(B)| \geq \frac{s(s+1)}{2}$.

Case 2. All B_i satisfy $|V(B_i)| < s$.

Assume that all B_1, B_2, \dots, B_m are nonempty. If $m \geq 2$, $|V(B_i)| \geq 2$, $i = 1, 2$, by Remark 2.3, we have $\sum_{i=1}^{n_s} |V(B_i)| \geq s$. It is sufficient to show that $|N_G(B)| \geq \frac{s(s+1)}{2}$ if there exists a B_i such that $|V(B_i)| = 1$. By a similar argument of the last part of Lemma 2.1, we complete this proof. \square

Corollary 2.5. Let $G = G_s^1 \odot \cdots \odot G_s^{n_s}$, $B \subseteq G$, $|V(B)| \geq s$, $n_i \geq 2$, $i = 1, 2, \dots, s$. If there exists a subgraph G_s^m such that $V(B \cap G_s^m) = \emptyset$, and the number of K_2 of $\{P_{n_1}, P_{n_2}, \dots, P_{n_s}\}$ is k , then $|N_G(B)| \geq \frac{s(s+1)}{2} - k$.

Proof. By a similar argument of Lemma 2.4 and Corollary 2.2, the result holds. \square

Lemma 2.6. Let $B \subseteq G$, $|V(B)| \geq s + 1$, $n_i \geq 3$, $i = 1, 2, \dots, s$. If $|V(G) \setminus$

$V(C_G(B)) \geq s + 1$, then $|N_G(B)| \geq \frac{s(s+1)}{2}$.

Proof. By induction on s . In particular, $s = 2$, then G is a grid, it is easy to see that $|N_G(B)| \geq \kappa_1(G) = 2s - 1 = 3 = \frac{s(s+1)}{2}$. Assume that the result is true for $s < h$. Next we show that it holds for $s = h$. We directly use s instead of h in the following.

Let $G = G_s^1 \odot \cdots \odot G_s^{n_s}$, we verify the result by considering two cases.

Case 1. There exists a G_s^i such that $V(B \cap G_s^i) = \emptyset$.

By the Lemma 2.4, we have $|N_G(B)| \geq \frac{s(s+1)}{2}$.

Case 2. All $V(B_i) \neq \emptyset$.

If there exists a G_s^i such that $|V(B_i)| \geq s - 1$ and $|V(G_s^i) \setminus V(C_{G_s^i}(B_i))| \geq s - 1$, by induction, we have $|N_{G_s^i}(B_i)| \geq \frac{s(s-1)}{2}$. Since $n_s \geq 3$, and $\kappa(G_s^i) = s - 1, j \neq i$, we have $|N_G(B)| \geq |N_{G_s^i}(B_i)| + \sum_{i \neq j} \kappa(G_j) \geq \frac{s(s-1)}{2} + 2(s-1) > \frac{s(s+1)}{2}$ for $s \geq 3$. If each pair $\{|V(B_i)|, |V(G_s^i) \setminus V(C_{G_s^i}(B_i))|\}$ has one element less than $s - 1$, without loss of generality, we assume $|V(B_i)| < s - 1$. Since $|V(B)| \geq s$, by Remark 2.3, we have $|N_G(B)| \geq \sum_{i=1}^{n_s} |N_{G_s^i}(B_i)| \geq \frac{s(s+1)}{2}$. \square

Corollary 2.7. Let $B \subseteq G, |V(B)| \geq s + 1, n_i \geq 2, i = 1, 2, \dots, s$, and the number of K_2 of $\{P_{n_1}, P_{n_2}, \dots, P_{n_s}\}$ is k . If $|V(G) \setminus V(C_G(B))| \geq s + 1$, then $|N_G(B)| \geq \frac{s(s+1)}{2} - k$.

Lemma 2.8. Any two vertices x, y of $G = P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s}$ have at most two common neighbors for $s \geq 2$ if they have any. Furthermore, if x and y exactly have two distinct coordinates, then x and y exactly have two common neighbors if they have any.

Proof. Let $x = x_1 x_2 \cdots x_s, y = y_1 y_2 \cdots y_s$ be two vertices of G . By definition, if x and y have common neighbor, then their coordinate have one or two is distinct. Clearly, they have one or two common neighbors. Furthermore, if x and y exactly have two distinct coordinates, say i and j , then $y_i = x_i + 1$ or $x_i - 1, y_j = x_j + 1$ or $x_j - 1$. Without loss of generality, we assume that $y_i = x_i + 1, y_j = x_j + 1$, then $x_1 \cdots x_{i-1} y_i x_{i+1} \cdots x_s$ and $x_1 \cdots x_{j-1} y_j x_{j+1} \cdots x_s$ are two common neighbors of x and y . \square

Lemma 2.9. Let $n_i \geq 3, i = 1, 2, \dots, s, G = P_{n_1} \times P_{n_2} \times \cdots \times P_{n_s} = G_s^1 \odot \cdots \odot G_s^{n_s}$ and $u = 11 \cdots 1 \in V(G_s^1), u_i = 11 \cdots 2 \cdots 1 \in N(u), A = G[u, u_1, u_2, \dots, u_g], 1 \leq i \leq g \leq s$, where u_i denotes the neighbor of v which the i th coordinate is different from w 's. Then $|N_G(A)| = s + (s + 1)g - 2g - \binom{g}{2}$ and $G - C_G(A)$ is a connected subgraph of G with property \mathcal{P}_g .

Proof. By lemma 2.8, it is easy to see that $|N_G(A)| = s + (s + 1)g - 2g - \binom{g}{2}$.

Next we show $G - C_G(A)$ is connected. By induction on s . If $s = 2, G$ is a grid, the result is clear. Assume that it is true for $s - 1 \geq 2$. Next we verify the Lemma for s .

Clearly, $A \subset G_s^1 \cup G_s^2$ and $|V(C_G(A) \cap G_s^i)| \leq 1$. By induction, $G_s^i - C_G(A)$ is connected for all i . That is, if $x, y \in V(G_s^i - C_G(A))$, then x connects to y . Next we assume that $x_s^i \in V(G_s^i), y_s^j \in V(G_s^j), i < j$. Since $n_s \geq 3$, then $x_s^i x_s^{i+1} \dots x_s^{n_s}$ is a path of $G - C_G(A)$ from x_s^i to $x_s^{n_s} \in V(G_s^{n_s})$. Similarly, $y_s^j y_s^{j+1} \dots y_s^{n_s}$ is a path of $G - C_G(A)$ from y_s^j to $y_s^{n_s} \in V(G_s^{n_s})$. Since $x_s^{n_s}$ connects $y_s^{n_s}$, we have x_s^i connects to y_s^j . \square

Corollary 2.10. Let $n_i \geq 2, i = 1, 2, \dots, s, G = P_{n_1} \times P_{n_2} \times \dots \times P_{n_s} = G_s^1 \odot \dots \odot G_s^{n_s}$ and $u = 11 \dots 1 \in V(G_s^1), u_i = 11 \dots 2 \dots 1 \in N(u), A = G[u, u_1, u_2, \dots, u_g], 1 \leq i \leq g$, where u_i denotes the neighbor of u which the i th coordinate is different from u 's. The number of K_2 of $\{P_{n_1}, P_{n_2}, \dots, P_{n_s}\}$ is k . Then $|N_G(A)| \geq s + (s + 1)g - 2g - \binom{g}{2} - \min\{g, k\}$, for $0 \leq g \leq s$ and $G - C_G(A)$.

Proof. By a similar argument of Lemma 2.9. \square

3 Main results

Now we come to our main results.

Theorem 3.1. Let $G = P_{n_1} \times P_{n_2} \times \dots \times P_{n_s}, n_i \geq 3, i = 1, 2, \dots, s, 0 \leq g \leq s$, then $\kappa_g(G) = s + (s + 1)g - 2g - \binom{g}{2}$.

Proof. Let $u = (1, 1, \dots, 1), u_i = (1, 1, \dots, 2, \dots, 1) (1 \leq i \leq g)$ be the neighbors of u , where u_i denotes the neighbor of u which the i th coordinate is different from u 's. Then $G[u, u_1, u_2, \dots, u_g] \cong K_{1,g}$. By Lemma 2.9, we have $|N_G(A)| = s + (s + 1)g - 2g - \binom{g}{2}$ and $G - C_G(A)$ has property \mathcal{P}_g . Thus we have $\kappa_g(G) \leq s + (s + 1)g - 2g - \binom{g}{2}$.

Next we show that $\kappa_g(G) \geq s + (s + 1)g - 2g - \binom{g}{2}$. Suppose F is a vertex cut set of G such that every component of $G - F$ has property \mathcal{P}_g and $|F| \leq s + (s + 1)g - 2g - \binom{g}{2} - 1$. We will show that it is impossible. If there exists a component A of $G - F$ such that $|V(A)| \leq s$, we have $g + 1 \leq |V(A)| \leq s$. By Lemma 2.1, we have $|F| \geq s + (s + 1)g - 2g - \binom{g}{2}$, a contradiction. If every component of $G - F$ has size more than s , we derive a contradiction by Lemma 2.6. \square

Corollary 3.2. Let $G = P_{n_1} \times P_{n_2} \times \dots \times P_{n_s}, n_i \geq 2 (i = 1, 2, \dots, s)$. The number of K_2 of $\{P_{n_1}, P_{n_2}, \dots, P_{n_s}\}$ is k .

(i) If $0 < k < \frac{s}{2}$, then $\kappa_g(G) \geq s + (s + 1)g - 2g - \binom{g}{2} - \min\{g, k\}$, for $0 \leq g \leq s - 2; \kappa_g(G) = \frac{s(s+1)}{2} - k$, for $s - 1 \leq g \leq s$.

(ii) If $\frac{s}{2} \leq k < s$, then $\kappa_g(G) \geq s + (s + 1)g - 2g - \binom{g}{2} - \min\{g, k\}$, for $0 \leq g \leq s - 3; \kappa_g(G) = \frac{s(s+1)}{2} - k$, for $s - 2 \leq g \leq s$.

(iii) If $k = s$, then $\kappa_g(G) \geq s + (s + 1)g - 2g - \binom{g}{2} - \min\{g, k\}$, for $0 \leq g \leq s - 4; \kappa_g(G) = \frac{s(s+1)}{2} - s$, for $s - 3 \leq g \leq s$.

Proof. We only give the proof of (iii). Assume that A is the same subgraph as Theorem 3.1. By Corollary 2.10, we have that $|N_G(A)| = \frac{s(s+1)}{2} - s$ if $g = s$, and by Corollary 2.7, we have that $|N_G(B)| \geq \frac{s(s+1)}{2} - s$ for any subgraph B with order more than s . Thus $\kappa_g(G) \leq \frac{s(s+1)}{2} - s$ for $g \leq s$.

Let $f(g) = s + (s+1)g - 2g - \binom{g}{2} - \min\{g, k\}$. It is easy to see that $f(1) < f(2) < \dots < f(s-4) < f(s-3) = f(s) < f(s-2) = f(s-1)$. By Corollary 2.2, we have $\kappa_g(G) = f(g)$ for $g \leq s-4$ and $\kappa_g(G) = f(s-3) = f(s) = \frac{s(s+1)}{2} - s$. \square

Corollary 3.3. $\kappa_g(Q_n) = (g+1)n - 2g - \binom{g}{2}$ for $0 \leq g \leq n-4$ and $\kappa_g(Q_n) = \frac{n(n+1)}{2} - n$ for $n-3 \leq g \leq n$, where $Q_n = P_2 \times P_2 \times \dots \times P_2$ is a well known networks (see [9] for a detail proof).

Proof. By (iii) of Corollary 3.2, Corollary 3.3 is clearly true. \square

Remark 3.4. There are also some other known results can be obtained directly by Theorem 3.1 and Corollary 3.3, such as the results about hypercube of [6] and the results of [8].

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