# The Properties of Middle Digraphs\*

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Abstract: Let D be a strongly connected digraph with order at least two, M(D) denote the middle digraph of D,  $\kappa(D)$  and  $\lambda(D)$  denote the connectivity and arcconnectivity of D, respectively. In this paper we study super-arc-connected and
super-connected middle digraphs and the spectral of middle digraphs.

Key words: Middle digraph; Super-arc-connected; Super-connected; Spectral

#### 1 Introduction

For graph-theoretical terminology and notation not defined here we follow Bondy and Murty [2]. Let D=(V(D),A(D)) be a digraph with vertex set V(D) and arc set A(D). For a vertex  $v\in V(D)$ , we denote the indegree, the outdegree of v, the minimum indegrees and outdegrees in D by  $d_D^-(v)$ ,  $d_D^+(v)$ ,  $\delta^-(D)$  and  $\delta^+(D)$ , respectively. We denote the minimum degree of D by  $\delta(D)=\min\{\delta^-(D),\delta^+(D)\}$ .  $\overrightarrow{K_n}$  denotes the complete digraph of order n.

Let D = (V(D), A(D)) be a digraph, |V(D)| = n, |A(D)| = m,  $V(D) = \{v_1, v_2, \dots, v_n\}$ . In 1960, Harary and Norman [4] introduced the concept of the line digraph. For a digraph D, the line digraph of D, denoted by L(D), is the digraph with vertex set  $V(L(D)) = \{a_{ij} | a_{ij} = (v_i, v_j) \text{ is an } \{v_i, v_j \in A(D)\}$ 

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arc in D}, and a vertex  $a_{ij}$  is adjacent to a vertex  $a_{st}$  in L(D) if and only if  $v_j = v_s$  in D. In 1977, Zamfirescu [8] introduced the concept of the middle digraph. For a digraph D, the middle digraph of D, denoted by M(D), is the digraph with vertex set  $V(M(D)) = V(D) \cup A(D)$ , there is an arc  $(a,b) \in A(M(D))$  from vertex a to vertex b in M(D) if and only if one of the following cases holds: (1). If  $a \in V(D)$  and  $b \in A(D)$ , then a is the tail of arc b in D. (2). If  $a \in A(D)$  and  $b \in V(D)$ , then b is the head of arc a in a. (3). If  $a \in A(D)$  and a0, then the head of arc a2 is the tail of arc a3 in a5. The middle digraph has been discussed in [7, 8, 9].

In fact, the middle digraph M(D) can be viewed as  $V(M(D)) = V(D) \cup V(L(D))$  and  $A(M(D)) = A(L(D)) \cup A(D, L(D))$ , where A(D, L(D)) denotes the arcs with one end in V(D) and the other end in V(L(D)). For each vertex  $v \in V(D)$ ,  $d_{M(D)}^+(v) = d_D^+(v)$ , since there are  $d_D^+(v)$  out-arcs from v to vertices in V(L(D)). For each vertex  $a_{ij} \in V(L(D))$ ,  $d_{M(D)}^+(a_{ij}) = d_D^+(v_j) + 1$ , since  $d_{L(D)}^+(a_{ij}) = d_D^+(v_j)$  and there are exactly one out-arc  $(a_{ij}, v_j)$  from  $a_{ij}$  to vertices in V(D). Similarly,  $d_{M(D)}^-(v) = d_D^-(v)$ ,  $d_{M(D)}^-(a_{ij}) = d_D^-(v_i) + 1$ .

An arc-cut of a strongly connected digraph D with order at least two is a set of arcs whose remove makes D no longer strongly connected. The  $arc\text{-}connectivity\ \lambda(D)$  is the minimum cardinality of an arc-cut over all arc-cuts of D. The inequality  $\lambda(D) \leq \delta(D)$  is wellknown. We call a digraph D maximally arc-connected, for short  $max\text{-}\lambda$ , if  $\lambda(D) = \delta(D)$ . For a vertex  $v \in V(D)$ , denote by  $N_D^+(v)$  the set of out-neighbors of v,  $N_D^-(v)$  the set of in-neighbors of v,  $E_D^-(v)$  the set of in-arcs of v. A strongly connected digraph D is super-arc-connected, for short  $super-\lambda$ , if every minimum arc-cut is either  $E_D^+(v)$  or  $E_D^-(v)$  for some vertex v. The  $connectivity\ \kappa(D)$ ,  $max\text{-}\kappa$ ,  $super\text{-}\kappa$  are defined similarly.

In this paper, we study super-arc-connected and super-connected middle digraphs and the spectra of middle digraphs.

The following two lemmas will be used in our discussions.

**Lemma 1.1.** [10] Let D be a digraph with order at least two. Then D is strongly connected if and only if the line digraph L(D) is strongly connected.

**Lemma 1.2.** Let D be a stongly connected digraph, then  $\kappa(D) \leq \lambda(D) = \kappa(L(D)) \leq \lambda(L(D))$ .

## 2 The spectra of middle digraphs

Let D=(V(D),A(D)) be a pseudodigraph, i.e. loops and multiple arcs may occur in A(D). Let  $V(D)=\{v_1,v_2,\ldots,v_n\}$  and  $A(D)=\{e_1,e_2,\ldots,e_m\}$ . Its out-incidence and in-incidence matrices (both are (0,1)-matrices), denoted by  $X_o$  and  $X_I$  respectively, are defined as follows.

$$X_o = (x_{ij}^o), \quad X_I = (x_{ij}^I),$$

where

$$x_{ij}^o = \left\{ egin{array}{ll} 1 & ext{if } v_i ext{ is the tail of } e_j; \\ 0 & ext{otherwise;} \end{array} 
ight. \ x_{ij}^I = \left\{ egin{array}{ll} 1 & ext{if } v_i ext{ is the head of } e_j; \\ 0 & ext{otherwise.} \end{array} 
ight.$$

The following lemmas are needed for the proof of Theorem 2.4.

**Lemma 2.1.** [5, 6, 11] Let M,  $M_L$  be the adjacency matrices of digraph D and its line digraph L(D) respectively. Let  $X_o$  and  $X_I$  be the out- and in-incidence matrices of D, then

$$M = X_o X_I^T, \quad M_L = X_I^T X_o.$$

Lemma 2.2. [3]

$$\det \left( \begin{array}{cc} A_{11} & A_{12} \\ A_{21} & A_{22} \end{array} \right) = \left\{ \begin{array}{cc} \det A_{11} \det (A_{22} - A_{21} A_{11}^{-1} A_{12}) & \text{if } A_{11} \text{ is invertible;} \\ \det A_{22} \det (A_{11} - A_{12} A_{22}^{-1} A_{21}) & \text{if } A_{22} \text{ is invertible .} \end{array} \right.$$

**Lemma 2.3.** Let  $\lambda_i (i = 1, 2, ..., n)$  be the eigenvalues of matrix A, let f(x) be a polynomial. Then the eigenvalues of f(A) are  $f(\lambda_i)(i = 1, 2, ..., n)$ .

**Theorem 2.4.** Let D be a pseudodigraph with n vertices and m arcs, let M(D) be the middle digraph of D. Then

$$P_{M(D)}(\lambda) = \lambda^{m-n} \prod_{i=1}^{n} (\lambda^2 - (\lambda+1)\lambda_i),$$

where  $\lambda_i (1 \leq i \leq n)$  are the eigenvalues of D.

*Proof.* By a suitable labeling the vertices of M(D) we can see that the adjacency matrix of M(D) is

$$M_{M(D)} = \left( \begin{array}{cc} 0 & X_o \\ X_I^T & M_L \end{array} \right) = \left( \begin{array}{cc} 0 & X_o \\ X_I^T & X_I^T X_o \end{array} \right).$$

Therefore

$$P_{M(D)}(\lambda) = |\lambda I_{m+n} - M_{M(D)}|$$

$$= \begin{vmatrix} \lambda I_n & -X_o \\ -X_I^T & \lambda I_m - X_I^T X_o \end{vmatrix}.$$

By adding the product of  $-X_I^T$  and the first row of the block matrix to the second row, we have

$$P_{M(D)}(\lambda) = \begin{vmatrix} \lambda I_n & -X_o \\ -(\lambda+1)X_I^T & \lambda I_m \end{vmatrix}.$$

By Lemma 2.2 and Lemma 2.3, we have

$$P_{M(D)}(\lambda) = |\lambda I_m| |\lambda I_n - \frac{\lambda+1}{\lambda} X_o X_I^T|$$

$$= \lambda^m |\lambda I_n - \frac{\lambda+1}{\lambda} M|$$

$$= \lambda^{m-n} |\lambda^2 I_n - (\lambda+1) M|$$

$$= \lambda^{m-n} \prod_{i=1}^n (\lambda^2 - (\lambda+1)\lambda_i).$$

**Corollary 2.5.** Let D be a pseudodigraph with n vertices and m arcs  $(m \ge n)$ . Then M(D) has m - n zero eigenvalues, and the following 2n eigenvalues:

$$\frac{\lambda_i \pm \sqrt{\lambda_i^2 + 4\lambda_i}}{2} \qquad (i = 1, 2, \dots, n)$$

where  $\lambda_i (1 \leq i \leq n)$  are the eigenvalues of D.

## 3 Super-arc-connected middle digraphs

In this section, we consider strict digraph D (digraph having no loops and no parallel arcs are allowed). If |V(D)| = 1, then M(D) is an isolated vertex. So we consider the case that  $|V(D)| \ge 2$ .

**Theorem 3.1.** Let D be a digraph with order at least two. Then M(D) is strongly connected if and only if D is strongly connected.

Proof. If D is strongly connected, then L(D) is strongly connected by Lemma 1.1. According to the definition of middle digraph, we know that the vertices of V(D) have both in-arcs and out-arcs in M(D). Hence, M(D) is strongly connected. If D is not strongly connected, then there exist two subsets  $X_1$  and  $X_2$  in V(D) such that  $X_1 \cup X_2 = V(D)$  and there are no arcs from  $X_1$  to  $X_2$  in D. It is easy to see that, in M(D), there are no directed paths from vertices in  $X_1$  to vertices in  $X_2$ .

**Theorem 3.2.** Let D be a strongly connected digraph with order at least two. Then  $\lambda(M(D)) \ge \min\{\delta(D), 2\lambda(D)\}.$ 

*Proof.* Clearly,  $\delta(M(D)) = \delta(L(D)) = \delta(D) \geq 1$ . Let S be a minimum arc-cut of M(D), then there exists a nonempty proper vertex subset  $X \subseteq V(M(D))$  such that there is no arc from X to  $\overline{X}$  in M(D) - S, where  $\overline{X} = V(M(D)) \setminus X$ .

We consider three cases.

Case 1.  $X \subseteq V(D)$ .

If |X| = 1, then  $|S| \ge \delta(M(D))$ . If  $n \ge |X| \ge 2$ , since D - S is no longer strongly connected, and every vertex  $v \in X$  has  $d_D^+(v) \ge \delta(D)$  out-neighbors in V(L(D)), we have

$$|S| \ge |X| \cdot \delta(D) > \delta(D).$$

Case 2.  $X \subseteq V(L(D))$ .

If |X| = 1, then  $|S| \ge \delta(M(D))$ . If  $\delta(D) \ge |X| \ge 2$ , since every vertex  $a \in X$  has at least  $\delta(L(D)) - (|X| - 1) = \delta(D) - |X| + 1$  out-neighbor in  $V(L(D)) \setminus X$  and exactly one out-neighbors in V(D), we have

$$|S| \ge |X|(\delta(L(D)) - |X| + 1) + |X| > \delta(D).$$

If  $m \geq |X| > \delta(D)$ , then

$$|S| \ge |X| > \delta(D)$$
.

Case 3.  $X \cap V(L(D)) \neq \emptyset$  and  $X \cap V(D) \neq \emptyset$ .

We may suppose that  $V(D) \nsubseteq X$  and  $V(L(D)) \nsubseteq X$ . In fact, in the case that  $V(D) \subseteq X$  or  $V(L(D)) \subseteq X$ , by considering  $\overline{X}$ , the proof is analogous to the proof of Case 1 or Case 2. For each arc  $(v_i, v_j)$  from  $X \cap V(D)$  to  $\overline{X} \cap V(D)$  in D, if the corresponding vertex  $a_{ij} \in \overline{X} \cap V(L(D))$ , then  $(v_i, a_{ij})$  is an arc from X to  $\overline{X}$  in M(D); if the corresponding vertex  $a_{ij} \in X \cap V(L(D))$ , then  $(a_{ij}, v_j)$  is an arc from X to  $\overline{X}$  in M(D). Hence,

$$|S| \ge \lambda(D) + \lambda(L(D)) \ge 2\lambda(D).$$

We thus conclude that  $\lambda(M(D)) \ge \min\{\delta(D), 2\lambda(D)\}.$ 

**Remark 3.3.** If the digraph D has the following properties, then  $\lambda(M(D)) = 2\lambda(D)$ .

- (i).  $\lambda(D) = \lambda(L(D))$ .
- (ii). Let S be a minimum arc-cut of D and there exist two vertex set X and  $\bar{X}$  such that there are no arcs from X to  $\bar{X}$  in D-S, where  $\bar{X}=V(D)-X$ .
- (iii). The arc set from  $A(D[X]) \cup [\bar{X}, X] \cup S$  to  $A(D[\bar{X}])$  is a minimum arc-cut of L(D), where  $[\bar{X}, X]$ , D[X] and  $D[\bar{X}]$  denote the arc set from  $\bar{X}$  to X, the vertex-induced subdigraph of X and  $\bar{X}$ , respectively.

Corollary 3.4. Let D be a strongly connected digraph with order at least two. If  $2\lambda(D) \geq \delta(D)$ , then M(D) is max- $\lambda$ .

**Theorem 3.5.** Let D be a strongly connected digraph with order at least two. If  $2\lambda(D) > \delta(D)$ , then M(D) is super- $\lambda$ .

*Proof.* From the proof of Theorem 3.2, only when |X|=1, or in Case 3, the equality  $\lambda(M(D))=\delta(M(D))$  may hold. If |X|=1, then we are done. In Case 3, the equality cannot hold when  $2\lambda(D)>\delta(D)$ . Hence if  $2\lambda(D)>\delta(D)$ , then M(D) is super- $\lambda$ .

## 4 Super-connected middle digraphs

**Theorem 4.1.** Let D be a strongly connected digraph with order at least two. Then  $\kappa(M(D)) = \lambda(D)$ .

Proof. By the definition of middle digraph, in order to destroy the connectivity of M(D), we must destroy all the out- (in-) arc set of some vertex in V(D), or destroy the connectivity of L(D) firstly. Thus,  $\kappa(M(D)) \geq \kappa(L(D)) = \lambda(D)$ . On the other hand, if S is a vertex-cut of L(D), then S is a vertex-cut of M(D). In fact, the corresponding arc set S' of S is an arc-cut in D. Therefore there exists a nonempty proper vertex subset  $X \subseteq V(L(D))$  such that there is no arc from X to  $\overline{X}$  in L(D) - S, and there exists a nonempty proper vertex subset  $X' \subseteq V(D)$  such that there is no arc from X' to  $\overline{X}'$  in D - S', and where  $\overline{X} = V(L(D)) \setminus (X \cup S)$ ,  $\overline{X}' = V(D) \setminus X'$ . It is easy to see that there is no arcs from  $X \cup X'$  to  $\overline{X} \cup \overline{X}'$  in M(D) - S. Hence,  $\kappa(M(D)) = \lambda(D)$ .

By Theorem 4.1, we have the following consequences.

Corollary 4.2. Let D be a strongly connected digraph with order at least two. Then M(D) is max- $\kappa$  if and only if D is max- $\lambda$ .

Corollary 4.3. Let D be a strongly connected digraph with order at least two. Then M(D) is super- $\kappa$  if and only if D is super- $\lambda$ .

*Proof.* Clearly, if M(D) is super- $\kappa$ , then D is super- $\lambda$ . On the other hand, if D is super- $\lambda$ , then L(D) is super- $\kappa$ , For any vertex-cut  $S = \{a_{ij_1}, a_{ij_2}, \ldots, a_{ij_d}\} = N^+(a_{ti})$  of L(D), we find that  $S = \{a_{ij_1}, a_{ij_2}, \ldots, a_{ij_d}\} = N^+(v_i)$  in M(D). Thus M(D) is super- $\kappa$ .

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