

A Cube-covering Problem*

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Abstract

We consider the problem of covering a unit cube with smaller cubes. The size of a cube is given by its side length and the size of a covering is the total size of the cubes used to cover the unit cube. We denote by $g_3(n)$ the smallest size of a minimal covering using n cubes. We present tight results for the upper and lower bounds of $g_3(n)$.

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1 Introduction

In 1932, Erdős defined a function $f(n)$ which denotes the maximum sum of n squares that can be packed into a unit square S . In [1], Erdős and Soifer gave some results of $f(n)$. Inspired by [1], Fan and Zhang [4] discussed the dual version, that is a square-covering problem. In this paper, we generalize this kind of covering problem to the case of cubes. That is, using smaller cubes to cover a unit cube, obtaining corresponding results.

First, we give the definition of a minimal cube-covering. The size of a given cube c is the side length of c and is denoted by $s(c)$. A covering \mathcal{C} is given by a set of cubes S positioned inside a unit cube C in such a way that $0 < s(c) < 1$, for each $c \in S$, and any point of C is covered by at least one of the cubes in S .

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Given a covering \mathcal{C} of the unit cube C using a set of cubes $S = \{c_1, \dots, c_n\}$ of n smaller cubes, where $0 < s(c_i) < 1$, we denote by $s(\mathcal{C})$ the size of the covering \mathcal{C} , which is given by $\sum_{i=1}^n s(c_i)$.

A covering is said to be minimal if there is no other covering of C using a set of cubes S' where $S' = \{c_1, \dots, c_{i-1}, c_{i+1}, \dots, c_n\}$ or $S' = \{c_1, \dots, c_{i-1}, c'_i, c_{i+1}, \dots, c_n\}$, where $s(c'_i) < s(c_i)$. We denote by $g_3(n)$ the smallest size of a minimal covering using a set of n cubes. That is,

$$g_3(n) = \min\{s(\mathcal{C}) : \mathcal{C} \text{ is a minimal covering of the unit cube with } n \text{ cubes}\}.$$

Let \mathcal{C} be a covering of a unit cube C using a set of cubes $S = \{c_1, \dots, c_n\}$. As each corner point of C has to be covered by a cube in S , and the size of a cube $c \in S$ is such that $0 < s(c) < 1$, we have that different corners of C must be covered by different cubes of S . Therefore, the following lemma is valid.

Lemma 1. *If \mathcal{C} is a covering of the unit cube, then \mathcal{C} has at least 8 cubes.*

For example of a covering, consider the case when $n = 8$. It is easy to show that $g_3(8) \leq 4$. To see this, we can use a set S with 8 cubes of size $1/2$, each one positioned in a different corner of the unit cube. This covering is clearly minimal, as we cannot remove a cube from S or replace by a smaller cube to obtain a smaller covering. The next theorem shows that $g_3(n)$ must be at least 4.

Theorem 2. *For any $n \geq 8$, we have that $g_3(n) \geq 4$.*

Proof. Let \mathcal{C} be a covering of the unit cube C with cubes in S . If a top face of C and a cube $c \in S$ have a common point, then c and the bottom face of C have no common points, because $0 < s(c) < 1$. Let a_1, a_2, \dots, a_s be the set of cubes in S which have common points with the top face of C and let b_1, b_2, \dots, b_t be the set of cubes in S which have common points with the bottom face of C , then $s + t \leq n$, and $\{a_1, \dots, a_s\} \cap \{b_1, b_2, \dots, b_t\} = \emptyset$.

For $i = 1, 2, \dots, s$, the projection of a_i in the top face of C , is a square which has the same length with a_i . That is, the projections of a_1, \dots, a_s in the top face leads to a covering of the top face of C . By the known result for the square covering problem, the side length of these projections is no less than 2, so the total size of the cubes a_1, \dots, a_s is no less than 2. In the same way, the sum of the sizes of the cubes b_1, b_2, \dots, b_t is no less than 2. So,

$$g_3(n) \geq 2 + 2 = 4.$$

□

Obviously, the following corollary is valid.

Corollary 3. $g_3(8) = 4$.

Having a covering C with n cubes does not mean that $g_3(n) \leq s(C)$. This may occur when we can still obtain a covering using smaller cubes or removing cubes in C and therefore, using less than n cubes. For example, consider a covering with $n \geq 8$ cubes having 8 cubes of size $1/2$, each one in a different corner of the unit cube, and the remaining $n - 8$ cubes of size ε , all in one of the corners. This leads to a covering C with size $4 + (n - 8)\varepsilon$, where ε can be made as close to zero as desired. Clearly this covering is not minimal, since the 8 cubes of size $1/2$ already leads to a covering and we can remove the $n - 8$ cubes of size ε . In this example, the smaller covering was easy to detect, but for a general covering, obtaining a smaller covering or proving that it is minimal may be a non-trivial task.

On the other hand, if we have a covering C that is minimal, each cube in C is necessary to obtain the covering of the unit cube. So, the following result is straightforward.

Lemma 4. *If C is a minimal covering of the unit cube and C has n cubes, then $g_3(n) \leq s(C)$.*

2 Main result

In this section, we present the upper and lower bounds for $g_3(n)$.

The next theorem shows that $g_3(n)$ cannot be greater than 4, for any $n \geq 9$.

Theorem 5. *For $n \geq 9$, we have $g_3(n) \leq 4 + \delta$, where δ is a positive value that can be made as close to 0 as desired.*

Proof. We denote the unit cube by C . We present a minimal covering C with n cubes divided in two sets: a set L of large cubes and a set S of small cubes. The set L has three cubes with size $1 - \varepsilon$ and one with size $1 - \varepsilon'$, where $\varepsilon' = (n - 7)\varepsilon$ and ε is a positive value such that $\varepsilon < \frac{1}{9n}$. The set S has two cubes of size ε' and $n - 6$ cubes of size ε .

In the following, we show that these cubes leads to a minimal covering. The following facts are valid.

Fact 1. *The total length of the edges of C that can be covered by the cubes in S is smaller than 1.*

Proof. The result follows, as we have 2 cubes of size $\varepsilon' = (n - 7)\varepsilon$ and $n - 6$ cubes of size ε . Therefore, we have a total size of $2n\varepsilon - 14\varepsilon + n\varepsilon - 6\varepsilon = (3n - 20)\varepsilon < 3n\varepsilon$. If a cube is positioned in a corner point of C , it partially cover 3 edges, and the total edge length covered by the cubes in S is less than $3 \cdot 3n\varepsilon = 9n\varepsilon$. The result follows, as we have $\varepsilon < 1/(9n)$. \square

Fact 2. *If $L \cup S$ can cover C , then each edge of C must be intercepted by a cube in L .*

Proof. From Fact 1, the cubes in S can cover only a total edge length that is smaller than 1. As the length of an edge is 1, the edge must also be intercepted by a cube of L . \square

Fact 3. *If $L \cup S$ can cover C , then each cube in L must cover a different corner point of C and each face of C has exactly two cubes of L covering opposite corners of the face.*

Proof. Let $c_i, i = 1, 2, 3, 4$ be the four cubes in L , and n_i is the number of edges of C which is intercepted by c_i . From Fact 1, $\sum_{i=1}^4 n_i \geq 12$. But each cube cannot intercept more than 3 edges of C , which means that $n_i \leq 3$. So, $n_i = 3, i = 1, 2, 3, 4$. This happens only when c_i covers a corner of C .

Let there are n_{12} edges of C intercepted by c_1 and c_2 simultaneously. Then there are at least $12 - (n_1 + n_2 - n_{12}) = 6 + n_{12}$ edges intercepted by c_3 or c_4 . So, $n_3 + n_4 \geq 6 + n_{12}$, and this cannot happen when $n_{12} > 0$ because $n_3 = n_4 = 3$. This means that each edge of C can be intercepted by only one cubes in L . This leads to a configuration where in each face of C , we have exactly two cubes of L covering opposite corners of the face. \square

So, consider the non-covered space after placing the large cubes of L . In Figure 1, we exemplify the placing of the large cubes: three cubes of size $(1 - \varepsilon)$ in positions $(0, 0, 1)$, $(0, 1, 0)$, $(1, 0, 0)$ and one cube of size $(1 - \varepsilon')$ in position $(1, 1, 1)$. The other possible coverings are symmetries of this one. To help visualize the covering in Figure 1, we only present the covering of the edges of C . The placing of the large cubes of L leads to 4 non-covered cuboid regions: one with dimensions $(\varepsilon, \varepsilon, \varepsilon)$ (at corner point $(0, 0, 0)$ in Figure 1) and three with dimensions $(\varepsilon, \varepsilon, \varepsilon')$ (at corner points $(0, 1, 1)$, $(1, 1, 0)$ and $(1, 0, 1)$, rotating if necessary).

We can consider these non-covered cuboid regions as one-dimensional bins, considering the largest edge of the cuboid as the size of a one-dimensional bin, that must be covered by one-dimensional items of size ε or ε' (all remaining cubes are the cubes in S , which has $n - 6$ cubes of size ε and 2 cubes of size ε'). In Figure 2, we present these bins with the size of a largest edge.

So, the total size of these bins is $\varepsilon + (n - 7)\varepsilon + 2\varepsilon' = \varepsilon + \varepsilon' + 2\varepsilon' = \varepsilon + 3\varepsilon'$. On the other hand, the total size of cubes in S is equal to $(n - 6)\varepsilon + 2\varepsilon'$ which is also $\varepsilon + 3\varepsilon'$. So, to have a covering of these bins (non-covered cuboids) with the cubes of S , we have to obtain a perfectly covering of the

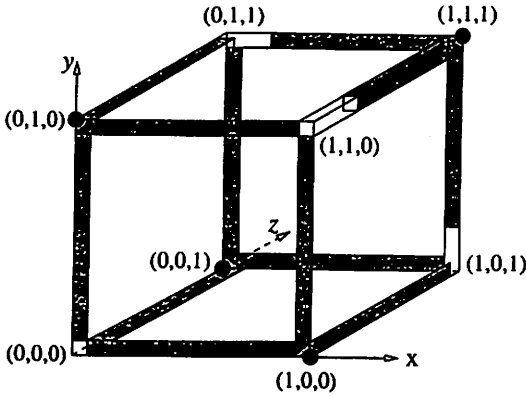


Figure 1: Edges of C covered after placing cubes in L .

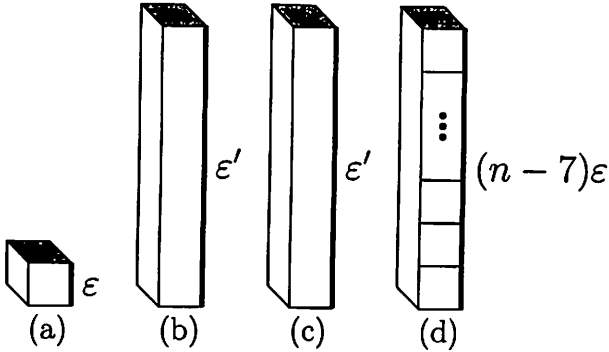


Figure 2: Non-covered cuboids after placing the large cubes of L .

bin size. In fact, the covering is easy to obtain, placing one cube of size ϵ in the cuboid of size ϵ (Figure 2 (a)), two cuboids with dimensions $(\epsilon, \epsilon, \epsilon')$ (Figures 2 (b) and (c)) and the remaining $n - 7$ cuboids of size ϵ covering perfectly the remaining cuboid of dimension $(\epsilon, \epsilon, \epsilon')$ (Figure 2 (d)).

To see that the above covering is minimal, note that we cannot replace one cube of S by a smaller cube, as the small cubes fit perfectly in the total length of the bins (cuboid largest edge). And we also cannot replace one large cube of L by a smaller cube, as there is no more small cubes to be used to cover the new larger cuboid regions.

Now, consider the size of the obtained covering. The cubes in L have total size $3(1 - \epsilon) + (1 - \epsilon') = 3 - 3\epsilon + 1 - (n - 7)\epsilon = 4 - (n - 4)\epsilon$. The cubes in S have total size $2\epsilon' + (n - 6)\epsilon = 2(n - 7)\epsilon + (n - 6)\epsilon = (3n - 20)\epsilon$.

So, the size of the covering is $4 - (n - 4)\varepsilon + (3n - 20)\varepsilon = 4 + (2n - 16)\varepsilon$. As ε can be made as close to 0 as desired, the size of the covering can also be made as close to 4 as desired. \square

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