

A note on circumferences in k -connected graphs with given independence number

Qing Cui*
Lingping Zhong†

Department of Mathematics
Nanjing University of Aeronautics and Astronautics
Nanjing 210016, P. R. China

Abstract

Fouquet and Jolivet conjectured that if G is a k -connected n -vertex graph with independence number $\alpha \geq k \geq 2$, then G has circumference at least $\frac{k(n+\alpha-k)}{\alpha}$. This conjecture was recently proved by O, West and Wu. In this note, we consider the set of k -connected n -vertex graphs with independence number $\alpha > k \geq 2$ and circumference exactly $\frac{k(n+\alpha-k)}{\alpha}$. We show that all of these graphs have a similar structure.

Keywords: Circumference; Connectivity; Independence number

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1 Introduction

We only consider finite graphs without loops or multiple edges. For a graph G , we use $V(G)$ and $E(G)$ to denote the vertex set and edge set of G , respectively. For any $S \subseteq V(G)$, define $G - S$ to be the subgraph of G with vertex set $V(G) - S$ and edge set $\{e \in E(G) : e \text{ is not incident with any vertex in } S\}$. We use $G[S]$ to denote the subgraph of G induced by S . A graph G is k -connected if $|V(G)| > k$ and every subgraph obtained from G by deleting fewer than k vertices is still connected. The *connectivity* of G , denoted by $\kappa(G)$, is the maximum integer k such that G is k -connected. An *independent*

*E-mail address: cui@nuaa.edu.cn

†E-mail address: zhong@nuaa.edu.cn

set of a graph G is a set of pairwise nonadjacent vertices. The *independence number* of G , denoted by $\alpha(G)$, is the maximum size of such a set. The *circumference* of a graph G , denoted by $c(G)$, is the length of a longest cycle in G . A spanning cycle (respectively, spanning path) is often called a *Hamilton cycle* (respectively, *Hamilton path*).

In 1972, Chvátal and Erdős [2] proved that every k -connected graph with independence number $\alpha \leq k$ contains a Hamilton cycle. There are infinitely many graphs showing that the condition $\alpha \leq k$ is sharp. For $\alpha \geq k$, Fouquet and Jolivet [3] proposed the following conjecture.

Conjecture 1.1 *Let G be a k -connected n -vertex graph with independence number α . If $\alpha \geq k \geq 2$, then $c(G) \geq \frac{k(n+\alpha-k)}{\alpha}$.*

The following two graphs H_1 and H_2 demonstrate the lower bound of $c(G)$ in Conjecture 1.1 is best possible. Let α, k and m be three positive integers such that $\alpha \geq k \geq 2$ and $m \geq 2$. Let G be an arbitrary graph with k vertices, and let K_k and K_m be the complete graphs of order k and m , respectively. Let αK_m be the union of α vertex disjoint copies of K_m . We define H_1 to be the join of the two graphs G and αK_m , and H_2 to be the graph obtained from the join of the two graphs K_k and αK_m by deleting an edge (or even more edges) with one end in K_k and the other end in αK_m . It is easy to check that for each $i = 1, 2$, we have $|V(H_i)| = n = k + \alpha m$, $\kappa(H_i) = k$, $\alpha(H_i) = \alpha$ and $c(H_i) = k(m + 1) = \frac{k(n+\alpha-k)}{\alpha}$.

Kouider [6] proved that if G is a k -connected n -vertex graph with independence number α then the vertices of G can be covered by at most $\lceil \frac{\alpha}{k} \rceil$ cycles. This implies that G contains a cycle of length at least $\frac{n}{\lceil \alpha/k \rceil}$, which is close to the conjectured threshold of $\frac{nk}{\alpha} + \frac{k(\alpha-k)}{\alpha}$. Fournier [4] proved Conjecture 1.1 for the case that $\alpha = k + 1$ or $\alpha = k + 2$ in his doctoral dissertation. In 1984, he [5] further showed that Conjecture 1.1 holds for $k = 2$. The case that $k = 3$ was verified by Manoussakis [7]. Chen, Hu and Wu [1] proved that if $\alpha \geq k \geq 4$ then $c(G) \geq \frac{k(n+\alpha-k)}{\alpha} - \frac{(k-3)(k-4)}{2}$. This shows that Conjecture 1.1 is true for $k = 4$. In the same paper, they also proved this conjecture for $\alpha = k + 3$. Recently, O, West and Wu [8] proved Conjecture 1.1 completely.

In this note, we consider the structure of k -connected n -vertex graphs with independence number $\alpha > k \geq 2$ and circumference exactly $\frac{k(n+\alpha-k)}{\alpha}$. We prove the following result.

Theorem 1.2 *Let G be a k -connected n -vertex graph with independence number $\alpha > k \geq 2$. Then $c(G) = \frac{k(n+\alpha-k)}{\alpha}$ if and only if there exists a k -cut S in G such that $G - S$ contains exactly α components, each of which is a complete graph of order m ($m \in \mathbb{N}$).*

The two graphs H_1 and H_2 defined above show that the subgraph of G induced by S may not be a complete graph or some vertex in S may not be adjacent to all the vertices of $G - S$. Hence Theorem 1.2 is best possible in some sense. Note that Theorem 1.2 is not true for $\alpha = k$. Let $k, m_1, m_2, \dots, m_{k-1}$ be k positive integers such that $k \geq 3$ and $m_1 \geq 3$. As before, let K_k and K_{m_i} ($1 \leq i \leq k-1$) be the complete graphs of order k and m_i , respectively. Let H be the union of the $k-1$ vertex disjoint complete graphs K_{m_i} . We define H_3 to be the graph obtained from the join of the two graphs K_k and H by deleting an edge with both ends in K_{m_1} . Clearly, $\kappa(H_3) = \alpha(H_3) = k$ and H_3 has a Hamilton cycle. But $H_3 - V(K_k)$ contains exactly $k-1$ components and some component of $H_3 - V(K_k)$ is not a complete graph.

We conclude this section with some notation and terminology. Let G be a graph. A *block* in G is a maximal subgraph of G containing no cut vertex. Let F and H be two vertex disjoint subgraphs of G . A path from F to H (or an (F, H) -path) in G is a path with one endvertex in $V(F)$, the other endvertex in $V(H)$ and no internal vertex in $V(F) \cup V(H)$.

We write $A := B$ to rename B as A . For any graph G and any $v \in V(G)$, we use $N_G(v)$ to denote the neighborhood of v in G . For any subset S of $V(G)$, we let $N_G(S) := (\bigcup_{v \in S} N_G(v)) - S$. For a specified orientation of a cycle C and for distinct vertices x, y of C , we use xCy to denote the subpath of C from x to y in the given orientation.

2 Proof of Theorem 1.2

In this section, we prove the main result of this paper.

The following result was virtually proved by O, West and Wu [8, Lemma 3.2 and Corollary 3.3]. Since we only consider the set of k -connected n -vertex graphs with independence number $\alpha > k \geq 2$ and circumference $\frac{k(n+\alpha-k)}{\alpha}$, we state this result as follows.

Lemma 2.1 *Let G be a k -connected n -vertex graph with independence number $\alpha > k \geq 2$ and circumference $c(G) = \frac{k(n+\alpha-k)}{\alpha}$, and let $l = \alpha - k$. Then there exist cycles C_0, C_1, \dots, C_l in G such that*

- (i) $\alpha(G - \bigcup_{j=0}^i V(C_j)) = \alpha - k - i$ for $0 \leq i \leq l$,
- (ii) $|V(C_i)| = \frac{k(n+\alpha-k)}{\alpha}$ for $0 \leq i \leq l$, and
- (iii) $|V(C_i) - V(C_0)| = |V(C_i) - \bigcup_{j=0}^{i-1} V(C_j)| = \frac{|V(C_0)|}{k} - 1$ for $1 \leq i \leq l$.

Proof. By [8, Lemma 3.2], there exist cycles C_0, C_1, \dots, C_l in G such that $\alpha(G - \bigcup_{j=0}^i V(C_j)) \leq \alpha - k - i$ for $0 \leq i \leq l$ and $|V(C_i) - \bigcup_{j=0}^{i-1} V(C_j)| \leq$

$\frac{|V(C_0)|}{k} - 1$ for $1 \leq i \leq l$. Since $c(G) = \frac{k(n+\alpha-k)}{\alpha}$, it follows from the proof of [8, Corollary 3.3] that $|V(C_0)| = \frac{k(n+\alpha-k)}{\alpha}$ and $|V(C_i) - \bigcup_{j=0}^{i-1} V(C_j)| = \frac{|V(C_0)|}{k} - 1$ for $1 \leq i \leq l$. Then by the construction of C_i ($0 \leq i \leq l$) in [8, Lemma 3.2], we see that $\alpha(G - \bigcup_{j=0}^i V(C_j)) = \alpha - k - i$ for $0 \leq i \leq l$, $|V(C_i) - V(C_0)| = |V(C_i) - \bigcup_{j=0}^{i-1} V(C_j)|$ and $|V(C_i)| = |V(C_0)|$ for $1 \leq i \leq l$. Hence the assertion of Lemma 2.1 holds. ■

We can now prove Theorem 1.2. The idea is partially inspired by O, West and Wu [8, Theorem 3.1].

Proof of Theorem 1.2. If there exists a k -cut S in G such that $G - S$ contains exactly α components, each of which is a complete graph of order m ($m \in \mathbb{N}$), then $n = k + \alpha m$. Now it is easy to check that $c(G) = k(m + 1) = \frac{k(n+\alpha-k)}{\alpha}$. So we need only to consider the opposite direction.

Suppose $c(G) = \frac{k(n+\alpha-k)}{\alpha}$. Then by Lemma 2.1, there exist cycles C_0, C_1, \dots, C_l in G satisfying (i), (ii) and (iii). By (ii) and (iii), $|V(C_i) - V(C_0)| = \frac{|V(C_0)|}{k} - 1 = \frac{n+\alpha-k}{\alpha} - 1 = \frac{n-k}{\alpha}$ for $1 \leq i \leq l$. Since $\frac{n-k}{\alpha}$ is a positive integer, we have $n = k + \alpha m$ ($m \in \mathbb{N}$). For each $1 \leq i \leq l$, let $V_i = V(C_i) - V(C_0)$ and let $G_i = G[V_i]$. Then $|V(C_0)| = \frac{k(n+\alpha-k)}{\alpha} = k(m + 1)$ and $|V_i| = \frac{n-k}{\alpha} = m$. By (iii), we see that $V_i \cap V_j = \emptyset$ for $1 \leq i < j \leq l$ if $l \geq 2$.

We first consider G_l . By (i), we have $\alpha(G_l) = \alpha(G - \bigcup_{j=0}^{l-1} V(C_j)) = 1$. This implies that G_l is a complete graph of order m . Let B be the block of $G - V(C_0)$ containing G_l . Since G is a k -connected graph, by Menger's Theorem, for each $b \in V(B)$, there exist k paths from b to C_0 in G that pairwise share only b (since $|V(C_0)| \geq k$). We call these k paths a (b, C_0) -fan. For a fixed orientation of C_0 , let $S = \{s_1, s_2, \dots, s_p\}$ be the set of endvertices on C_0 of all the (B, C_0) -paths, and assume that s_1, s_2, \dots, s_p occur on C_0 in the given orientation. Then $p \geq k$. By the maximality of B , any two (B, C_0) -paths with distinct endvertices in B are internally disjoint.

We claim that

$$(1) \quad p = k \text{ and } |V(s_i C_0 s_{i+1})| = m + 2 \text{ for } 1 \leq i \leq k.$$

First, suppose that $|V(B)| = 1$. Then $B = G_l$ and $m = 1$. Let b be the only vertex in B . By the maximality of B , we have $N_G(b) = S$. Choose $s_i \in S$ such that $|V(s_i C_0 s_{i+1})|$ is minimum, where $s_{p+1} := s_1$. Then by the pigeonhole principle, we have $|V(s_i C_0 s_{i+1})| \leq m + 2 = 3$. Let $C := s_i C_0 s_i \cup \{b s_i, b s_{i+1}\}$. If $|V(s_i C_0 s_{i+1})| = 2$, then $|V(C)| = |V(C_0)| + 1$, which contradicts the assumption that $c(G) = |V(C_0)|$. Hence $|V(s_i C_0 s_{i+1})| = 3$. By the choice of s_i , we deduce that $p = k$ and $|V(s_i C_0 s_{i+1})| = 3$ for $1 \leq i \leq k$.

So we may assume that $|V(B)| \geq 2$. Let $S_1 = \{s_{i_1}, s_{i_2}, \dots, s_{i_q}\} \subseteq S$ be the set of endvertices on C_0 such that for each $s_{i_j} \in S_1$, there exist two disjoint (B, C_0) -paths P_{i_j} and $P_{i_{j+1}}$ in G with $V(P_{i_j}) \cap V(C_0) = \{s_{i_j}\}$ and $V(P_{i_{j+1}}) \cap V(C_0) = \{s_{i_{j+1}}\}$. Clearly, $S_1 \neq \emptyset$. Let $b_{i_j} \in V(B)$ and $b_{i_{j+1}} \in V(B)$ be the other endvertices of P_{i_j} and $P_{i_{j+1}}$, respectively. Since B is a block of $G - V(C_0)$ and G_l is a complete graph, by Menger's Theorem, we can always find a path P in B from b_{i_j} to $b_{i_{j+1}}$ containing V_l (by considering $|\{b_{i_j}, b_{i_{j+1}}\} \cap V_l| = 0, 1$ or 2). Now let $C := s_{i_{j+1}}C_0s_{i_j} \cup P_{i_j} \cup P \cup P_{i_{j+1}}$. If $|V(s_{i_j}C_0s_{i_{j+1}})| \leq m + 1$, then we have $|V(C)| \geq |V(C_0)| + 1$, which contradicts the assumption that $c(G) = |V(C_0)|$. So we know that $|V(s_{i_j}C_0s_{i_{j+1}})| \geq m + 2$ for $1 \leq j \leq q$, and hence $q \leq k$.

Let $S_2 = S - S_1$. By the definition of S_1 , for each $s_j \in S_2$, there exist two internally disjoint paths P_j and P_{j+1} in G from the same vertex of $G - V(C_0)$, say v_j , to C_0 such that $V(P_j) \cap V(C_0) = \{s_j\}$ and $V(P_{j+1}) \cap V(C_0) = \{s_{j+1}\}$. It is possible that $v_j \notin V(B)$. Let $C := s_{j+1}C_0s_j \cup P_j \cup P_{j+1}$. If $|V(s_jC_0s_{j+1})| = 2$, then $|V(C)| \geq |V(C_0)| + 1$, contradicting the assumption that $c(G) = |V(C_0)|$. Hence we have $|V(s_jC_0s_{j+1})| \geq 3$ for each $s_j \in S_2$. For each $b \in V(B)$, a (b, C_0) -fan has k endvertices on C_0 . Since $q \leq k$, at least $k - q$ of them are contained in S_2 . Moreover, again by the definition of S_1 , for distinct vertices of B , these endvertices contained in S_2 are pairwise distinct. Therefore, we see that $|S_2| = p - q \geq |V(B)|(k - q)$.

Then we have $k(m + 1) = |V(C_0)| \geq q(m + 1) + 2(p - q) \geq q(m + 1) + 2|V(B)|(k - q) \geq q(m + 1) + 2m(k - q) = (2k - q)m + q$. Since $|V(B)| \geq 2$, it is easy to check the above inequality holds if and only if $p = q = k$. So we have $|V(s_iC_0s_i + 1)| = m + 2$ for $1 \leq i \leq k$. This proves (1).

We also claim that

$$(2) \quad B = G_l.$$

Suppose for a contradiction that $B \neq G_l$. Then $|V(B)| \geq 2$. By (iii), $G - V_l$ is 2-connected. Then by Menger's Theorem, there exist two internally disjoint $(B - V_l, C_0)$ -paths P_1 and P_2 in $G - V_l$ with distinct endvertices on C_0 . Without loss of generality, let $b_1 \in V(B) - V_l$ and s_1 be the two endvertices of P_1 , and let $b_2 \in V(B) - V_l$ and $s_i \neq s_1$ be the two endvertices of P_2 .

First, assume that $k = 2$. Then $s_i = s_2$. Suppose $b_1 \neq b_2$. Since B is a block of $G - V(C_0)$ and G_l is a complete graph, by Menger's theorem, we can always find a path P in B from b_1 to b_2 containing V_l . Let $C := s_2C_0s_1 \cup P_1 \cup P \cup P_2$. Since $b_1, b_2 \in V(B) - V_l$, we have $|V(C)| \geq |V(C_0)| + 2$, which contradicts the assumption that $c(G) = |V(C_0)|$. So we may assume that $b_1 = b_2$. Then there must exist a path P_3 (disjoint with P_1) in G from $B - \{b_1\}$ to C_0 with $V(P_3) \cap V(C_0) = \{s_2\}$ or a path P'_3 (disjoint with P_2) in G from $B - \{b_1\}$ to C_0 with $V(P'_3) \cap V(C_0) = \{s_1\}$; for otherwise, b_1 would

be a cut vertex in G , a contradiction. By symmetry between s_1 and s_2 , let P_3 be a path (disjoint with P_1) in G from $B - \{b_1\}$ to C_0 with endvertices $b_3 \in V(B) - \{b_1\}$ and s_2 . As before, let P' be a path in B from b_1 to b_3 containing V_i , and let $C' := s_2 C_0 s_1 \cup P_1 \cup P' \cup P_3$. Since $b_1 \in V(B) - V_i$, $|V(C')| \geq |V(C_0)| + 1$, a contradiction.

Hence we may assume that $k \geq 3$. Then there must exist a path P_4 (disjoint with P_1) in G from $B - \{b_1\}$ to C_0 with $V(P_4) \cap V(C_0) = \{s_2\}$ or $V(P_4) \cap V(C_0) = \{s_k\}$; for otherwise, $S' := (S - \{s_2, s_k\}) \cup \{b_1\}$ would be a $(k-1)$ -cut in G , a contradiction. By symmetry between s_2 and s_k , let $b_4 \in V(B) - \{b_1\}$ and s_2 be the two endvertices of P_4 . By the same argument as before, we can find a path in B from b_1 to b_4 containing V_i and construct a cycle of length at least $|V(C_0)| + 1$ in G (since $b_1 \in V(B) - V_i$), again a contradiction. So (2) holds.

We further claim that

$$(3) \quad N_G(V_i) = S.$$

Suppose to the contrary that $N_G(V_i) \neq S$. Then by the argument in (1) and by (2), we have $|V_i| = |V(B)| \geq 2$. Since G is a k -connected graph, $|N_G(V_i)| \geq k = |S|$, and hence $N_G(V_i) - S \neq \emptyset$. Let uv be an edge in G with $v \in V_i$ and $u \in N_G(V_i) - S$. By (iii), $G - V_i$ is 2-connected. Then by Menger's Theorem, there exist two internally disjoint (u, C_0) -paths Q'_1 and Q'_2 in $G - V_i$ with distinct endvertices on C_0 . Then $Q_1 := Q'_1 \cup \{uv\}$ and $Q_2 := Q'_2 \cup \{uv\}$ are two paths in G from G_l to C_0 with $|V(Q_1)| \geq 3$ and $|V(Q_2)| \geq 3$. Without loss of generality, let s_1 and $s_i \neq s_1$ be the other endvertices of Q_1 and Q_2 , respectively.

Suppose that $k = 2$. Therefore, $s_i = s_2$. Then there must exist a path Q_3 (disjoint with Q_1) in G from $G_l - \{v\}$ to C_0 with $V(Q_3) \cap V(C_0) = \{s_2\}$ or a path Q'_3 (disjoint with Q_2) in G from $G_l - \{v\}$ to C_0 with $V(Q'_3) \cap V(C_0) = \{s_1\}$; for otherwise, v would be a cut vertex in G , a contradiction. By symmetry between s_1 and s_2 , let Q_3 be a path (disjoint with Q_1) in G from $G_l - \{v\}$ to C_0 with endvertices $w \in V_i - \{v\}$ and s_2 . Let Q be a Hamilton path in G_l from v to w , and let $C := s_2 C_0 s_1 \cup Q_1 \cup Q \cup Q_3$. Since $|V(Q_1)| \geq 3$, we have $|V(C)| \geq |V(C_0)| + 1$, a contradiction.

So we may assume that $k \geq 3$. Then there must exist a path Q_4 (disjoint with Q_1) in G from $G_l - \{v\}$ to C_0 with $V(Q_4) \cap V(C_0) = \{s_2\}$ or $V(Q_4) \cap V(C_0) = \{s_k\}$; for otherwise, $S' := (S - \{s_2, s_k\}) \cup \{v\}$ would be a $(k-1)$ -cut in G , a contradiction. By symmetry between s_2 and s_k , let $w' \in V(B) - \{b_1\}$ and s_2 be the two endvertices of Q_4 . By the same argument as before, since $|V(Q_1)| \geq 3$, we can find a cycle of length at least $|V(C_0)| + 1$ in G , again a contradiction. So we have (3).

We then consider G_{l-1} if $l \geq 2$. By (i), $\alpha(G[V_{l-1} \cup V_l]) = \alpha(G - \bigcup_{j=0}^{l-2} V(C_j)) = 2$. Since $N_G(V_l) \cap V_{l-1} = \emptyset$ (by (3)), we know that G_{l-1}

is also a complete graph of order m . Let B' be the block of $G - V(C_0)$ containing G_{l-1} . Let $T = \{t_1, t_2, \dots, t_{p'}\}$ be the set of endvertices on C_0 of all the (B', C_0) -paths, and assume that $t_1, t_2, \dots, t_{p'}$ occur on C_0 in the given orientation. Then by the same argument as for G_l , we can also conclude that $p' = k$, $|V(t_i C_0 t_{i+1})| = m + 2$ for $1 \leq i \leq k$ and $N_G(V_{l-1}) = T$.

We claim that

$$(4) \quad S = T.$$

Suppose for a contradiction that $S \neq T$. Then $S \cap T = \emptyset$. By symmetry between S and T , we may assume that $s_1, t_1, s_2, t_2, \dots, s_k, t_k$ occur on C_0 in the given orientation. Let $s_1 u_1, s_2 u_2 \in E(G)$ with $u_1, u_2 \in V_l$, and let $t_1 v_1, t_2 v_2 \in E(G)$ with $v_1, v_2 \in V_{l-1}$. If $m \geq 2$, we can choose u_1, u_2, v_1, v_2 so that $u_1 \neq u_2$ and $v_1 \neq v_2$. Let P be a Hamilton path in G_l from u_1 to u_2 , and let Q be a Hamilton path in G_{l-1} from v_1 to v_2 . (If $m = 1$, then let $P := \emptyset$ and $Q := \emptyset$.) Now let $C := t_2 C_0 s_1 \cup P \cup t_1 C_0 s_2 \cup Q \cup \{s_1 u_1, s_2 u_2, t_1 v_1, t_2 v_2\}$. Since $|V(s_1 C_0 t_1)| = |V(s_2 C_0 t_2)| \leq m + 1$, we have $|V(C)| \geq |V(C_0)| + 2$, contradicting the assumption that $c(G) = |V(C_0)|$. This proves (4).

By (i) and by the same arguments as for G_l and G_{l-1} , we know that

$$(5) \quad G_i \text{ is a complete graph of order } m \text{ and } N_G(V_i) = S \text{ for each } 1 \leq i \leq l.$$

We now consider C_0 . For each $1 \leq i \leq k$, let $U_i = V(s_i C_0 s_{i+1}) - \{s_i, s_{i+1}\}$ and let $G_{l+i} = G[U_i]$. Then $|U_i| = m$ (by (1)) and $U_i \cap U_j = \emptyset$ for $1 \leq i < j \leq k$.

We claim that

$$(6) \quad \text{there is no edge in } G \text{ with one end in } U_i \text{ and the other end in } U_j \text{ for } 1 \leq i < j \leq k.$$

For otherwise, let $r_i r_j$ be an edge in G with $r_i \in U_i$ and $r_j \in U_j$. Let $s_i u_1, s_j u_2, s_{i+1} v_1, s_{j+1} v_2 \in E(G)$ such that $u_1, u_2, v_1, v_2 \in V_1$. If $m \geq 2$, we can further choose u_1, u_2, v_1, v_2 so that $u_1 \neq u_2$ and $v_1 \neq v_2$. Let P be a Hamilton path in G_1 from u_1 to u_2 , and let Q be a Hamilton path in G_1 from v_1 to v_2 . (If $m = 1$, then let $P := \emptyset$ and $Q := \emptyset$.) Now let $C := r_i C_0 s_j \cup r_j C_0 s_i \cup P \cup \{s_i u_1, s_j u_2, r_i r_j\}$ and let $C' := s_{i+1} C_0 r_j \cup s_{j+1} C_0 r_i \cup Q \cup \{s_{i+1} v_1, s_{j+1} v_2, r_i r_j\}$. Since $|U_i| = |U_j| = m$, it is easy to check that $|V(C)| + |V(C')| = 2|V(C_0)| + 2$. But this implies either $|V(C)| \geq |V(C_0)| + 1$ or $|V(C')| \geq |V(C_0)| + 1$, which contradicts the assumption that $c(G) = |V(C_0)|$. So (6) holds.

By (5) and (6), we have $G - S$ contains exactly $l + k = \alpha$ components. Since $\alpha(G) = \alpha$, we see that G_{l+i} is also a complete graph of order m for each $1 \leq i \leq k$. This completes the proof of Theorem 1.2. \blacksquare

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