

# On super connectedness and super restricted edge-connectedness of total graphs \*

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**Abstract** A graph  $G$  is *super-connected*, *super- $\kappa$* , for short if every minimum vertex-cut isolates a vertex of  $G$ . Call  $G$  *super restricted edge-connected*, in short, *super- $\lambda'$* , if every minimum restricted edge-cut isolates an edge. We view the *total graph*  $T(G)$  of  $G$  as the disjoint of  $G$  and the line graph  $L(G)$ , together with the lines of the subdivision graph  $S(G)$ ; for each line  $l = (u, v)$  in  $G$  there are two lines in  $S(G)$ , namely  $(\hat{l}, u)$  and  $(\hat{l}, v)$ . In this paper, we prove that  $T(G)$  is *super- $\kappa$*  if  $G$  is a *super- $\kappa$*  graph with  $\delta(G) \geq 4$ . We also show that  $T(G)$  is *super- $\lambda'$*  if  $G$  is a  $k$ -regular graph with  $\kappa(G) \geq 3$ . Furthermore, we give examples which illustrate that the results are best possible.

**Keywords:** Super connected; Super edge-connected; Total graph

## 1 Introduction

A network can be conveniently modeled as a graph  $G = (V, E)$ , with vertices representing nodes and edges representing links. A classic measure of network reliability is the connectivity  $\kappa(G)$  and the edge-connectivity  $\lambda(G)$ . In general, the larger  $\kappa(G)$  (or  $\lambda(G)$ ) is, the more reliable the network is. For  $\kappa(G) \leq \lambda(G) \leq \delta(G)$ , where  $\delta(G)$  is the minimum degree of  $G$ , a graph  $G$  with  $\kappa(G) = \delta(G)$  ( $\lambda(G) = \delta(G)$ ) is naturally said to be *maximally connected* (*maximally edge-connected*), or  *$\kappa$ -optimal* ( *$\lambda$ -optimal*) for simplicity. As more refined indices of reliability than maximal connectivity and maximal edge-connectivity, super connectivity and super

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edge-connectivity were proposed in [1, 3]. A graph  $G$  is *super-connected*, *super- $\kappa$* , for short (resp. *super edge-connected*, *super- $\lambda$* , for short) if every minimum vertex-cut (resp. edge-cut) isolates a vertex of  $G$ .

For further study, Esfahanian and Hakimi introduce the concept of restricted edge-connectivity [8]. The concept of restricted edge-connectivity is one kind of conditional edge-connectivity proposed by Harary in [9], and has been successfully applied in the further study of tolerance and reliability of networks, see [7, 12, 13]. Let  $F$  be a set of edges in  $G$ . Call  $F$  a *restricted edge-cut* if  $G - F$  is disconnected and contains no isolated vertices. The minimum cardinality over all restricted edge-cuts is called *restricted edge-connectivity* of  $G$ , and denoted by  $\lambda'(G)$ . It was shown by Wang and Li that the larger  $\lambda'(G)$  is, the more reliable the network is [13]. In [8], the authors proved that if a connected graph  $G$  of order  $n \geq 4$  is not a star  $K_{1,n-1}$ , then  $\lambda'(G)$  is well-defined and  $\lambda(G) \leq \lambda'(G) \leq \xi(G)$ , where  $\xi(G) = \min\{d_G(u) + d_G(v) - 2 : uv \in E(G)\}$  is the minimum edge degree of  $G$ . A graph  $G$  with  $\lambda'(G) = \xi(G)$  is called a  *$\lambda'$ -optimal graph*. Call  $G$  *super restricted edge-connected*, in short, *super- $\lambda'$* , if every minimum restricted edge-cut isolates an edge, that is, every minimum restricted edge-cut is a set of edges adjacent to a certain edge with minimum edge degree in  $G$ . By the definitions, a super- $\lambda'$  graph must be a  $\lambda'$ -optimal graph. However, the converse is not true since there are many  $\lambda'$ -optimal graphs not to be super- $\lambda'$ . For example,  $C_n$  ( $n \geq 6$ ), the cycle of length  $n$  is a trivial counterexample.

It should be pointed out that if  $\delta(G) \geq 3$ , then a  $\lambda'$ -optimal graph must be super- $\lambda$ . In fact, a graph  $G$  is super- $\lambda$  if and only if  $\lambda(G) < \lambda'(G)$  [12]. Thus, the concepts of  $\lambda$ -optimal graph, super- $\lambda$  graph,  $\lambda'$ -optimal graph and super- $\lambda'$  graph describe reliable interconnection structure for graphs at different levels.

The *line graph*  $L(G)$  of  $G$  is that graph whose point set can be put in one-to-one correspondence with the line set of  $G$ , such that two points of  $L(G)$  are adjacent if and only if the corresponding lines of  $G$  are adjacent. We view the *total graph*  $T(G)$  of  $G$  as the disjoint of  $G$  and  $L(G)$ , together with the lines of the subdivision graph  $S(G)$ ; for each line  $l = (u, v)$  in  $G$  there are two lines in  $S(G)$ , namely  $(\hat{l}, u)$  and  $(\hat{l}, v)$ . For convenience, the points of  $T(G)$  belonging to  $L(G)$  will be called *linear points* of  $T(G)$ . We simplify notation by setting  $\kappa = \kappa(G)$ ,  $\kappa_L = \kappa(L(G))$ ,  $\kappa_T = \kappa(T(G))$ , and similarly for  $\lambda$  and  $\delta$ . Except where noted we follow the definitions and notations of Bondy [4].

In 1969, Chartrand and Stewart [5] proved that (i)  $\kappa(L(G)) \geq \lambda(G)$ , if  $\lambda(G) \geq 2$ ; (ii)  $\lambda(L(G)) \geq 2\lambda(G) - 2$ . Furthermore, Hellwig et. [10] showed that

**Theorem 1.1.** *Let  $G$  be a graph with  $|V(G)| \geq 4$  and  $G$  is not a star. Then  $\kappa(L(G)) = \lambda'(G)$ .*

By Theorem 1.1, we easily have

**Corollary 1.2.** *Let  $G$  be a graph with  $|V(G)| \geq 4$  and  $G$  is not a star. If  $G$  is a super- $\lambda$  graph, then  $\kappa(L(G)) = \lambda'(G) > \lambda(G)$ .*

In [2], the authors showed the following two theorems.

**Theorem 1.3.** *If  $\lambda(G) \geq 2$ , then  $T(G)$  is maximally edge-connected.*

**Theorem 1.4.**  $\kappa(G) + \lambda(G) \leq \kappa_T(G) \leq 2\lambda(G)$ .

By Theorem 1.4, we obtain the following two corollaries.

**Corollary 1.5.** *If  $\kappa(G) = \lambda(G)$ , then  $\kappa_T(G) = 2\lambda(G)$ .*

**Corollary 1.6.** *If  $G$  is maximally connected, then  $T(G)$  is maximally connected.*

Chen and Meng [6] proved that for all but a few exceptions, the total graph  $T(G)$  of  $G$  is super- $\lambda$ .

**Theorem 1.7.** *For a given connected graph  $G$ ,  $T(G)$  is super- $\lambda$  if and only if one of the following two conditions applies:*

(i) *If  $\lambda(G) \geq 2$  and  $G$  has a cut vertex  $u$  with  $d_G(u) = \delta(G)$ , then there are at least three edges between  $u$  and any component of  $G \setminus u$ .*

(ii) *If  $\lambda(G) = 1$  and  $e = uv$  is a bridge, then  $\min\{d_G(u), d_G(v)\} \geq 2\delta(G)$ .*

In this paper, we prove that  $T(G)$  is super- $\kappa$  if  $G$  is a super- $\kappa$  graph with  $\delta(G) \geq 4$  in section 2. We also show that  $T(G)$  is super- $\lambda'$  if  $G$  is a  $k$ -regular graph with  $\kappa(G) \geq 3$  in section 3. Furthermore, we give examples which illustrate that the results are best possible.

## 2 Super connected total graphs

In [14], Xu et. proved the following theorem.

**Theorem 2.1.** *Let  $G$  be a super- $\kappa$  graph with  $\delta(G) \geq 4$ . Then  $G$  is super- $\lambda$ .*

Now, we prove the following result.

**Theorem 2.2.** *Let  $G$  be a super- $\kappa$  graph with  $\delta(G) \geq 4$ . Then  $T(G)$  is super- $\kappa$ , and thus  $T(G)$  is super- $\lambda$ .*

**Proof.** Suppose to the contrary that  $T(G)$  is not super- $\kappa$ . Then there exists a minimum vertex-cut  $S$  of  $T(G)$  such that  $|S| \leq \delta(T(G)) = 2\delta(G)$  and every component of  $T(G) - S$  has at least two vertices. Let  $X_1, X_2, \dots, X_t$  ( $t \geq 2$ ) be the components of  $T(G) - S$ . Then we have  $|V(X_i)| \geq 2$  for  $i = 1, 2, \dots, t$ . We consider three cases.

**Case 1.** There is one component  $X_i$  ( $1 \leq i \leq t$ ) such that  $V(X_i) \subseteq V(G)$ .

If  $|V(X_i)| = 2$ , then  $|S| \geq 2\delta - 1 + \min\{\kappa(G), n - 2\} > 2\delta = \delta(T(G))$ , a contradiction. If  $|V(X_i)| \geq 3$ , then  $|S| \geq 3\delta - 3 + \min\{\kappa(G), n - |V(X_i)|\} > 2\delta = \delta(T(G))$ , also a contradiction.

**Case 2.** There is one component  $X_i$  ( $1 \leq i \leq t$ ) such that  $V(X_i) \subseteq E(G)$ .

By Case 1, we can assume that  $V(X_j) \cap E(G) \neq \emptyset$  for  $j = 1, 2, \dots, t$ . Let  $Y = V(G[X_i])$ . Then  $|Y| \geq 3$  by  $|V(X_i)| \geq 2$ . Since  $G$  is super- $\kappa$  and  $\delta(G) \geq 4$ , we have that  $G$  is super- $\lambda$  by Theorem 2.1, thus  $\kappa(L(G)) > \lambda(G)$  by Corollary 1.2. If  $|Y| \geq \delta(G)$ , then  $|S| \geq |Y| + \kappa(L(G)) \geq |Y| + \delta(G) + 1 > 2\delta(G)$ , a contradiction. Thus, we assume that  $3 \leq |Y| \leq \delta(G) - 1$ . Then  $|S| \geq |Y|\delta(G) - |Y|(|Y| - 1) + |Y| = |Y|(\delta(G) + 2 - |Y|) \geq 3(\delta(G) - 1) > 2\delta(G)$  by  $\delta(G) \geq 4$ , also a contradiction.

**Case 3.**  $V(X_i) \cap E(G) \neq \emptyset$  and  $V(X_i) \cap V(G) \neq \emptyset$  for  $i = 1, 2, \dots, t$ .

Since  $S \cap V(G)$  is a vertex cut of  $G$  and  $S \cap E(G)$  is a vertex cut of  $L(G)$ , we obtain that  $|S| \geq \kappa(G) + \kappa(L(G)) \geq \delta + (\delta + 1) > 2\delta$  (The inequality  $\kappa(L(G)) \geq \lambda(G) + 1 = \delta(G) + 1$  is obtained by the proof of Case 2), contradicting to  $|S| \leq 2\delta$ .  $\square$

We present a class of graphs, which show that Theorem 2.2 is best possible, in the sense that  $\kappa(G) = \delta(G)$  does not guarantee that  $T(G)$  is super- $\kappa$ .

**Example 2.3.** Let  $n$  and  $\delta$  be arbitrary integers with  $n \geq 2\delta \geq 8$ . Furthermore, let  $H_1 \cong K_\delta$  with vertex set  $V(H_1) = \{x_1, x_2, \dots, x_\delta\}$  and let  $H_2 \cong K_{n-\delta}$  with vertex set  $V(H_2) = \{y_1, y_2, \dots, y_{n-\delta}\}$ . We define the graph  $G$  as the union of  $H_1$  and  $H_2$  together with the  $\delta$  edges  $x_1y_1, x_2y_2, \dots, x_\delta y_\delta$ . Then  $n(G) = n$ ,  $\delta(G) = \delta$ , and  $\kappa(G) = \delta(G)$ . But we have that  $T(G)$  is not super- $\kappa$ , since  $V(H_1) \cup \{x_1y_1, x_2y_2, \dots, x_\delta y_\delta\}$  is a vertex-cut of  $T(G)$  with cardinality  $\delta(T(G))$ .

The following example shows that the condition  $\delta(G) \geq 4$  is necessary in Theorem 2.2.

**Example 2.4.** The graphs defined in Fig.2.1 are super- $\kappa$ , and  $\delta(G_i) = i$  for  $i = 1, 2, 3$ . But we verify that  $T(G_i)$  is not super- $\kappa$  for  $i = 1, 2, 3$ , since  $\{v_{r+1}\} \cup \{v_{r+1}v_{r+2}\}$  is a vertex-cut of  $T(G_1)$  with cardinality 2,  $\{v_3, v_4\} \cup$

$\{v_1v_3, v_2v_4\}$  is a vertex-cut of  $T(G_2)$  with cardinality 4, and  $\{v_4, v_5, v_6\} \cup \{v_1v_4, v_2v_5, v_3v_6\}$  is a vertex-cut of  $T(G_3)$  with cardinality 6.

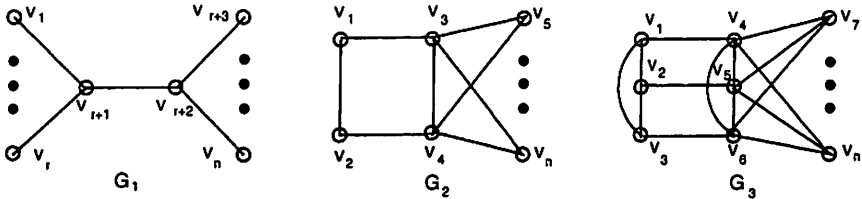


Fig.2.1.

### 3 Super restricted-edge-connected total graphs

For a vertex  $x \in V(G)$  and two vertex sets  $X, Y \subseteq V(G)$  with  $X \cap Y = \emptyset$ , let  $E_G(x) = \{e \in E(G) : e \text{ is incident with } x\}$ ,  $[X, Y] = \{e = xy \in E(G) : x \in X, y \in Y\}$  and  $E_G(X) = [X, \bar{X}]$ .

In the proof of Theorem 3.2, we will use the following theorem, which was given by König [11] in 1931.

**Theorem 3.1.** *If  $G$  is a bipartite graph, then the maximum size of a matching in  $G$  equals the minimum size of a vertex cover of  $G$ .*

Now, we are ready to prove the following main result.

**Theorem 3.2.** *Let  $G$  be a  $k$ -regular graph with  $\kappa(G) \geq 3$ . Then  $T(G)$  is super- $\lambda'$ .*

**Proof.** Suppose to the contrary that  $T(G)$  is not super- $\lambda'$ . Then there exists a minimum restricted edge-cut  $F$  such that  $|F| \leq \xi(T(G)) = 4k - 2$  and each component of the two components of  $T(G) - F$  has at least three vertices. Let  $X$  and  $Y$  be the two vertex sets of the two components of  $T(G) - F$ . We consider three cases.

**Case 1.** One of  $X$  and  $Y$  belongs to  $V(G)$ .

Assume, without loss of generality, that  $X \subseteq V(G)$ . If  $|X| = 3$ , then  $|F| \geq 3k + (3k - 6) = 6(k - 1) > 4k - 2$ , a contradiction. If  $|X| \geq 4$ , then  $|F| \geq 4k > 4k - 2$ , also a contradiction.

**Case 2.** One of  $X$  and  $Y$  belongs to  $E(G)$ .

Assume, without loss of generality, that  $X \subseteq E(G)$ . Let  $Z = V(G[X])$ . Since  $|X| \geq 3$ , we have  $|Z| \geq 3$  and equality holds if and only if  $G[X] \cong K_3$ .

If  $|Z| \geq 4$ , then there are at least four vertices  $u_1, u_2, u_3, u_4 \in V(G) \subseteq Y$ , such that  $E_G(u_i) \cap X \neq \emptyset$  for  $i=1,2,3,4$ . Let  $|E_G(u_i) \cap X| = n_i \geq 1$  and  $|E_G(u_i) \cap Y| = d_G(u_i) - n_i$ . For any  $e_i \in E_G(u_i) \cap X$  and  $e_j \in E_G(u_i) \cap Y$ ,  $e_i$  and  $e_j$  are adjacent in  $T(G)$ , they contribute one edge to  $F$ . Thus,

$$|F| \geq \sum_{i=1}^4 [n_i + n_i(d_G(u_i) - n_i)] = \sum_{i=1}^4 n_i(k+1 - n_i) \geq \sum_{i=1}^4 k > 4k - 2,$$

a contradiction.

If  $|Z| = 3$ , then  $G[X] \cong K_3$ , and let  $Z = \{v_1, v_2, v_3\} \subseteq V(G) \subseteq Y$ . It is easy to see that  $|E_G(v_i) \cap X| = n_i = 2$  and  $|E_G(v_i) \cap Y| = k - n_i = k - 2$ .

Thus  $|F| \geq \sum_{i=1}^3 n_i(k - n_i + 1) = 6(k - 1) > 4k - 2$ , also a contradiction.

**Case 3.**  $X \cap V(G) \neq \emptyset$ ,  $X \cap E(G) \neq \emptyset$ ,  $Y \cap V(G) \neq \emptyset$  and  $Y \cap E(G) \neq \emptyset$ .

Let  $V_1 = V(G) \cap X$ ,  $V_2 = V(G) \cap Y$ ,  $E_1 = E(G) \cap X$  and  $E_2 = E(G) \cap Y$ .

**Subcase 3.1.**  $|V_1| = 1$  or  $|V_2| = 1$ .

Assume, without loss of generality, that  $|V_1| = 1$ . Let  $V_1 = \{u\}$ ,  $N(u) = \{u_1, u_2, \dots, u_k\}$ ,  $e_i = uu_i$  ( $i = 1, 2, \dots, k$ ) and  $t = |\{e_1, e_2, \dots, e_k\} \cap E_1|$ . Then  $|E_G(u) \cap E_2| = n_u = k - t$  and  $|E_G(u) \cap E_1| = k - n_u = t$ .

If  $t = 1$ , assume that  $\{e_1, e_2, \dots, e_k\} \cap E_1 = \{e_1\}$ . Since  $|X| \geq 3$  and  $T(G)[X]$  is connected, there exists a vertex  $u_{k+1} \in N(u_1) \setminus \{u\}$  such that  $e_{k+1} = u_1u_{k+1} \in E_1$ . Let  $|E_G(u_1) \cap E_1| = n_1 \geq 1$ ,  $|E_G(u_1) \cap E_2| = k - n_1$ ,  $|E_G(u_{k+1}) \cap E_1| = n_{k+1} \geq 1$  and  $|E_G(u_{k+1}) \cap E_2| = k - n_{k+1}$ . Then

$$\begin{aligned} |F| &\geq n_1 + n_1(k - n_1) + n_{k+1} + n_{k+1}(k - n_{k+1}) + n_u + n_u(k - n_u) + |N_G(u)| \\ &\geq k + k + 2(k - 1) + k = 5k - 2 > 4k - 2, \end{aligned}$$

a contradiction.

If  $2 \leq t \leq k-1$ , assume, without loss of generality, that  $\{e_1, e_2, \dots, e_k\} \cap E_1 = \{e_1, e_2, \dots, e_t\}$ . Let  $|E_G(u_i) \cap E_1| = n_i \geq 1$  and  $|E_G(u_i) \cap E_2| = k - n_i$  for  $i = 1, 2, \dots, t$ . Then

$$|F| \geq \sum_{i=1}^t [n_i + n_i(k - n_i)] + n_u + n_u(k - n_u) + |N_G(u)| \geq tk + k + k > 4k - 2,$$

a contradiction.

If  $t = k$ , then  $E_G(u_i) \cap E_1 \neq \emptyset$  for  $i = 1, 2, \dots, k$ . Let  $|E_G(u_i) \cap E_1| = n_i \geq 1$  and  $|E_G(u_i) \cap E_2| = k - n_i$  for  $i = 1, 2, \dots, k$ . Then

$$|F| \geq \sum_{i=1}^k [n_i + n_i(k - n_i)] + |N_G(u)| \geq k \times k + k > 4k - 2,$$

also a contradiction.

**Subcase 3.2.**  $|V_1| = 2$  and  $|V_2| \geq 2$ , or  $|V_1| \geq 2$  and  $|V_2| = 2$ .

Suppose that  $|V_1| = 2$  and  $|V_2| \geq 2$ . Denote  $V_1 = \{u, v\}$ .

**Subcase 3.2.1.**  $uv \notin E(G)$ .

Let  $N(u) = \{u_1, u_2, \dots, u_k\}$  and  $e_i = uu_i$  for  $i = 1, 2, \dots, k$ . Since  $|X| \geq 3$  and  $T(G)[X]$  is connected, assume, without loss of generality, that  $e_1 \in E_1$ .

If  $e_2$  is also in  $E_1$ , then  $E_G(u_i) \cap E_1 \neq \emptyset$  for  $i = 1, 2$ , and let  $|E_G(u_i) \cap E_1| = n_i \geq 1$  and  $|E_G(u_i) \cap E_2| = k - n_i$  for  $i = 1, 2$ . Thus  $|F| \geq \sum_{i=1}^2 [n_i + n_i(k - n_i)] + |E_G(\{u, v\})| \geq 2k + 2k > 4k - 2$ , a contradiction.

If  $e_2$  is not in  $E_1$ , then  $E_G(u) \cap E_2 \neq \emptyset$ , and let  $|E_G(u) \cap E_2| = n_u \geq 1$  and  $|E_G(u) \cap E_1| = k - n_u$ . Since  $E_G(u_1) \cap E_1 \neq \emptyset$ , let  $|E_G(u_1) \cap E_1| = n_1 \geq 1$  and  $|E_G(u_1) \cap E_2| = k - n_1$ . Thus  $|F| \geq n_u + n_u(k - n_u) + n_1 + n_1(k - n_1) + |E_G(\{u, v\})| \geq 2k + 2k > 4k - 2$ , also a contradiction.

**Subcase 3.2.2.**  $uv \in E(G)$ .

Let  $N(u) = \{v, u_1, \dots, u_{k-1}\}$  and  $N(v) = \{u, v_1, \dots, v_{k-1}\}$ . If  $e = uv \in E_2$ , assume, without loss of generality, that  $e_1 = uu_1 \in E_1$  by  $|X| \geq 3$  and  $T(G)[X]$  is connected. Let  $|E_G(u) \cap E_2| = n_u \geq 1$ ,  $|E_G(v) \cap E_2| = n_v \geq 1$  and  $|E_G(u_1) \cap E_1| = n_1 \geq 1$ . Then  $|F| \geq n_u + n_u(k - n_u) + n_v + n_v(k - n_u) + n_1 + n_1(k - n_1) + |E_G(\{u, v\})| \geq 3k + 2k - 2 > 4k - 2$ , a contradiction. Therefore, we assume that  $e = uv \in E_1$  in the following.

If  $|X| = 3$ , then  $|E_G(u) \cap E_2| = n_u = k - 1$  and  $|E_G(v) \cap E_2| = n_v = k - 1$ . Thus  $|F| \geq n_u + n_u(k - n_u) + n_v + n_v(k - n_v) + |E_G(\{u, v\})| = 2(k - 1) + 2(k - 1) + 2(k - 1) > 4k - 2$ , also a contradiction.

If  $|X| \geq 4$ , then we can assume, without loss of generality, that  $e_1 = uu_1 \in E_1$  by  $T(G)[X]$  is connected. Let  $|E_G(u_1) \cap E_1| = n_1 \geq 1$  and  $|E_G(u_1) \cap E_2| = k - n_1$ .

We claim that  $E_G(v) \cap E_2 = \emptyset$ . Suppose that  $E_G(v) \cap E_2 \neq \emptyset$ , let  $|E_G(v) \cap E_2| = n_v \geq 1$  and  $|E_G(v) \cap E_1| = k - n_v$ . If  $E_G(u) \cap E_2 \neq \emptyset$ , let  $|E_G(u) \cap E_2| = n_u \geq 1$  and  $|E_G(u) \cap E_1| = k - n_u$ . Thus  $|F| \geq n_u + n_u(k - n_u) + n_v + n_v(k - n_v) + n_1 + n_1(k - n_1) + |E_G(\{u, v\})| \geq 3k + 2(k - 1) > 4k - 2$ ,

a contradiction. If  $E_G(u) \cap E_2 = \emptyset$ , then  $e_2 = uu_2 \in E_1$  and  $E_G(u_2) \cap E_1 \neq \emptyset$ , let  $|E_G(u_2) \cap E_1| = n_2 \geq 1$  and  $|E_G(u_2) \cap E_2| = k - n_2$ . Thus  $|F| \geq n_v + n_v(k - n_v) + n_1 + n_1(k - n_1) + n_2 + n_2(k - n_2) + |E_G(\{u, v\})| \geq 3k + 2(k - 1) > 4k - 2$ , also a contradiction. Thus, we assume that  $E_G(v) \cap E_2 = \emptyset$ .

If  $N(v) \setminus \{u, u_1, u_2\} \neq \emptyset$ , let  $v_1 \in \{u, u_1, u_2\}$ . Thus  $E_G(v_1) \cap E_1 \neq \emptyset$ , and let  $|E_G(v_1) \cap E_1| = n_{v_1} \geq 1$  and  $|E_G(v_1) \cap E_2| = k - n_{v_1}$ . If  $E_G(u) \cap E_2 \neq \emptyset$ , let  $|E_G(u) \cap E_2| = n_u \geq 1$  and  $|E_G(u) \cap E_1| = k - n_u$ . Thus  $|F| \geq n_u + n_u(k - n_u) + n_{v_1} + n_{v_1}(k - n_{v_1}) + n_1 + n_1(k - n_1) + |E_G(\{u, v\})| \geq 3k + 2(k - 1) > 4k - 2$ , a contradiction. If  $E_G(u) \cap E_2 = \emptyset$ , then  $e_2 = uu_2 \in E_1$  and  $E_G(u_2) \cap E_1 \neq \emptyset$ , let  $|E_G(u_2) \cap E_1| = n_2 \geq 1$  and  $|E_G(u_2) \cap E_2| = k - n_2$ . Thus  $|F| \geq n_{v_1} + n_{v_1}(k - n_{v_1}) + n_1 + n_1(k - n_1) + n_2 + n_2(k - n_2) + |E_G(\{u, v\})| \geq 3k + 2(k - 1) > 4k - 2$ , also a contradiction. Thus, we assume that  $N(v) = \{u, u_1, u_2\}$ . Since  $\kappa(G) \geq 3$ , we obtain that  $G \cong K_4$ . It is easy to verify that  $|F| > 10$ , a contradiction.

**Subcase 3.3.**  $|V_1| \geq 3$  and  $|V_2| \geq 3$ .

We consider the graph  $H = G[[V_1, V_2]_G]$ , which is induced by the edge set  $[V_1, V_2]_G$  in  $G$ . If  $H$  has a matching  $M$  which contains at least four edges, assume that  $v_{11}v_{21}, v_{12}v_{22}, \dots, v_{1t}v_{2t}$  ( $t \geq 4$ ) are  $t$  edges in  $M$ . Then either  $E_G(v_{1i}) \cap E_2 \neq \emptyset$ , or  $E_G(v_{2i}) \cap E_1 \neq \emptyset$  for  $i = 1, 2, \dots, t$ . Assume, without loss of generality, that  $E_G(v_{1i}) \cap E_2 \neq \emptyset$  for  $i = 1, \dots, r$ , and  $E_G(v_{2j}) \cap E_1 \neq \emptyset$  for  $j = r + 1, \dots, t$ . Let  $|E_G(v_{1i}) \cap E_2| = n_{1i} \geq 1$  and  $|E_G(v_{1i}) \cap E_1| = k - n_{1i}$  for  $i = 1, \dots, r$ ,  $|E_G(v_{2j}) \cap E_1| = n_{2j} \geq 1$  and  $|E_G(v_{2j}) \cap E_2| = k - n_{2j}$  for  $j = r + 1, \dots, t$ . Then  $|F| \geq \sum_{i=1}^r [n_{1i} + n_{1i}(k - n_{1i})] + \sum_{j=r+1}^t [n_{2j} + n_{2j}(k - n_{2j})] \geq rk + (t - r)k > 4k - 2$ , a contradiction. If the maximum size of matching  $\alpha'(H) \leq 2$ , then the minimum size of vertex cover  $\beta(H) = \alpha'(H) \leq 2$  by Theorem 3.1. Since a vertex cover of  $H$  with cardinality at least two is a vertex-cut of  $G$ , thus  $\kappa(G) \leq 2$ , which contradicts to  $\kappa(G) \geq 3$ . Therefore, we assume that  $\beta(H) = \alpha'(H) = 3$  in the following.

Since  $G$  is  $k$ -regular,  $[[V_1, V_2]_G]$  is an even number. Thus  $|E(H)|$  is an even number and  $|E(H)| \geq 4$ . If the number of vertices in  $V_1$  which are adjacent to  $E_2$  in  $T(G)$  and vertices in  $V_2$  which are adjacent to  $E_1$  in  $T(G)$  is at least 4, then we have that  $|F| \geq 4k > 4k - 2$  by a similar argument as above, a contradiction. Since the vertices in  $V_1$  which are adjacent to  $E_2$  in  $T(G)$  and the vertices in  $V_2$  which are adjacent to  $E_1$  in  $T(G)$  constitute a vertex cover of  $H$ ,  $\beta(H) = \alpha'(H) = 3$  and  $\kappa(G) \geq 3$ , we can assume that there are exactly three vertices  $v_1, v_2, v_3$  such that if  $v_i \in V_1$ , then  $E_G(v_i) \cap E_2 \neq \emptyset$ , and if  $v_j \in V_2$ , then  $E_G(v_j) \cap E_1 \neq \emptyset$ . It is easy to verify that  $\{v_1, v_2, v_3\}$  is a minimum vertex cover of  $H$ . Assume, without loss of



generality, that  $v_1, v_2, v_3 \in V_1$ , or  $v_1, v_2 \in V_1$  and  $v_3 \in V_2$ .

**Subcase 3.3.1.**  $v_1, v_2, v_3 \in V_1$ .

Let  $|E_G(v_i) \cap E_2| = n_i \geq 1$  and  $|E_G(v_i) \cap E_1| = k - n_i$  for  $i = 1, 2, 3$ . Since  $v_1, v_2, v_3 \in V_1$ , we have that the linear point set  $\{\{v_1, v_2, v_3\}, V_2\}_G$  of  $T(G)$  is contained in  $E_2$ . If there exists a  $n_i$  such that  $2 \leq n_i \leq k - 1$ , assume, without loss of generality that  $2 \leq n_1 \leq k - 1$ , then  $n_1 + n_1(k - n_1) \geq 2(k - 1)$ . Thus  $|F| \geq \sum_{i=1}^3 [n_i + n_i(k - n_i)] + |E(H)| \geq 2(k - 1) + 2k + 4 > 4k - 2$ , a contradiction. Otherwise, there is a  $n_j$  such that  $n_j = k$ . Assume, without loss of generality that  $n_2 = k$ , then  $|N(v_2) \cap V_2| \geq k - 2$  since  $N(v_2) \cap V_1 \subseteq \{v_1, v_3\}$ . We also have  $|N(v_1) \cap V_2| \geq 1$  and  $|N(v_3) \cap V_2| \geq 1$ . Therefore,  $|F| \geq \sum_{i=1}^3 [n_i + n_i(k - n_i)] + |E(H)| \geq 3k + k - 2 + 2 > 4k - 2$ , also a contradiction.

**Subcase 3.3.2.**  $v_1, v_2 \in V_1$  and  $v_3 \in V_2$ .

Let  $|E_G(v_i) \cap E_2| = n_i \geq 1$  and  $|E_G(v_i) \cap E_1| = k - n_i$  for  $i = 1, 2$ ,  $|E_G(v_3) \cap E_1| = n_3 \geq 1$  and  $|E_G(v_3) \cap E_2| = k - n_3$ . If there exists a  $n_i$  such that  $2 \leq n_i \leq k - 1$ , then we can obtain a contradiction by a similar argument as the proof of the subcase 3.3.1. Otherwise, there exists a  $n_j$  such that  $n_j = k$ . Since  $T(G)[Y]$  is connected and there exists exactly one vertex  $v_3$  in  $V_2$  such that  $E_G(v_3) \cap E_1 \neq \emptyset$ , we have that  $n_3 \neq k$ . Thus  $n_1 = k$  or  $n_2 = k$ . Assume, without loss of generality that  $n_1 = k$ , then  $|N(v_1) \cap V_2| \geq k - 1$  since  $N(v_1) \cap V_1 \subseteq \{v_2\}$ . We also have  $|N(v_2) \cap V_2| \geq 1$ . Therefore,  $|F| \geq \sum_{i=1}^3 [n_i + n_i(k - n_i)] + |E(H)| \geq 3k + k - 1 + 1 > 4k - 2$ , also a contradiction.  $\square$

If  $G$  is a connected 1-regular graph, then  $G \cong K_2$  and  $T(G) \cong K_3$ . If  $G$  is a connected 2-regular graph, then  $T(G)$  is not super- $\lambda'$  since  $E_{T(G)}(\{u, v, e\})$  is a restricted edge-cut with cardinality 6 for every edge  $e = uv \in E(G)$ . If  $G$  is a  $k(\geq 3)$ -regular with  $\kappa(G) = 1$ , then  $k$  is even. The following examples show that the condition  $\kappa(G) \geq 3$  is necessary in the Theorem 3.2.

**Example 3.3.** Let  $k$  be an even integer with  $k = 2t \geq 4$ . Furthermore, let  $G_1, G_2, \dots, G_t$  be  $t$  connected  $k$ -regular graphs, and  $e_1 = u_1v_1 \in E(G_1), e_2 = u_2v_2 \in E(G_2), \dots, e_t = u_tv_t \in E(G_t)$ . We define the graph  $G$  as the union of  $G_1 - e_1, G_2 - e_2, \dots, G_t - e_t$  together with the  $k$  edges  $wu_1, wv_1, wu_2, wv_2, \dots, wu_t, wv_t$ , where  $w$  is an other vertex not in  $V(G_1) \cup V(G_2) \cup \dots \cup V(G_t)$ . Then  $G$  is a  $k$ -regular graph with  $\kappa(G) = 1$ . But  $T(G)$  is not super- $\lambda'$  since  $E_{T(G)}(V_1 \cup E(G_1) \setminus \{e_1\})$  is a restricted edge-cut with cardinality  $2k + 2$ .

**Example 3.4.** Let  $G_1$  and  $G_2$  be two  $k(\geq 3)$ -regular graphs with  $\kappa(G_1) \geq 2$  and  $\kappa(G_2) \geq 2$ ,  $e_1 = u_1v_1 \in E(G_1)$  and  $e_2 = u_2v_2 \in E(G_2)$ . We define the graph  $G$  as the union of  $G_1 - e_1$  and  $G_2 - e_2$  together with the two edges  $u_1u_2$  and  $v_1v_2$ . Then  $G$  is a  $k$ -regular graph with  $\kappa(G) = 2$ . But  $T(G)$  is not super- $\lambda'$  since  $E_{T(G)}(V(G_1) \cup E(G_1) \setminus \{e_1\})$  is a restricted edge-cut with cardinality  $2k + 2$ .

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