The smallest realization of a given vector

Xiaoxiao Duan^{a,b} Kefeng Diao^{a*} Fuliang Lu^a Xiao Zhu^{a,b}

a. School of Science, Linyi University, Linyi, Shandong, 276005, China

b. School of Mathematical Science, Shandong Normal University, Jinan, Shandong, 250014, China

Abstract

For a vector $R = (r_1, r_2, \ldots, r_m)$ of non-negative integers, a mixed hypergraph \mathcal{H} is a realization of R if its chromatic spectrum is R. In this paper, we determine the minimum number of vertices of realizations of a special kind of vectors R_2 . As a result, we partially solve an open problem proposed by Král in 2004.

Key words: mixed hypergraph; strict coloring; feasible set; chromatic spectrum; realization

1 Introduction

A mixed hypergraph is a triple $\mathcal{H} = (X, \mathcal{C}, \mathcal{D})$ where X is a finite set and \mathcal{C} and \mathcal{D} are families of subsets of X called the \mathcal{C} -edges and \mathcal{D} -edges, respectively. If $X' \subset X$, $\mathcal{C}' = \{C \in \mathcal{C} | C \subseteq X'\}$ and $\mathcal{D}' = \{D \in \mathcal{D} | D \subseteq X'\}$, then the hypergraph $\mathcal{H}' = (X', \mathcal{C}', \mathcal{D}')$ is called the *induced sub-hypergraph* of \mathcal{H} on X', denoted by $\mathcal{H}[X']$.

A coloring of the vertices of \mathcal{H} is proper if there are two vertices with a Common color in each \mathcal{C} -edge and there are two vertices with Distinct colors in each \mathcal{D} -edge. A proper coloring using exactly k-colors is called a strict k-coloring and a mixed hypergraph is k-colorable if it has a strict k-coloring. A coloring may also be viewed as a partition (feasible partition) of the vertex set, where the color classes (partition classes) are the sets of vertices assigned to the same color. An edge is said to be monochromatic (resp. polychromatic) if all of its vertices have the same color (resp. different colors). The maximum (resp. minimum) number of colors which can be

^{*}Corresponding author: kfdiao@163.com

used in a strict coloring of \mathcal{H} is called the *upper chromatic number* (resp. lower chromatic number) and denoted by $\chi(\mathcal{H})$ (resp. $\chi(\mathcal{H})$). The study of the colorings of mixed hypergraphs has made a lot of progress since its inception [6]. For more information, we refer readers to [3, 5, 7, 8].

The feasible set $\mathcal{F}(\mathcal{H})$ of a mixed hypergraph \mathcal{H} is the set of all the values k such that \mathcal{H} has a strict k-coloring. For each k, let r_k denote the number of feasible partitions of the vertex set into k nonempty subsets. The vector $R(\mathcal{H}) = (r_1, r_2, \ldots, r_\chi)$ is called the chromatic spectrum of \mathcal{H} , where $\overline{\chi}$ is the upper chromatic number of \mathcal{H} . If S is a set of positive integers, we say that a mixed hypergraph \mathcal{H} is a realization of S if it is a realization of S and all the entries of the chromatic spectrum of S are either 0 or 1. The concept of one-realization was introduced by Jiang et al. in [2] with an extra condition as proper realizations and further studied by Král in [4]. Moreover, for a vector S of non-negative integers, a mixed hypergraph S is called a realization of S if S if S if S if S if S is a set of positive integers.

Bujtás and Tuza [1] gave a necessary and sufficient condition for S to be the feasible set of an r-uniform mixed hypergraph. Jiang et al. [2] proved that the minimum number of vertices of realizations of S is $2 \max(S)$ $\min(S)$ if |S| = 2 and $\max(S) - 1 \notin S$. Moreover, they also mentioned that the question of finding the minimum number of vertices in a mixed hypergraph with feasible set S of size at least 3 remains open. Král [4] proved that there exists a one-realization of S with at most $|S| + 2 \max(S) \min(S)$ vertices and proposed the following problem: what is the number of vertices of the smallest mixed hypergraph whose spectrum is equal to a given spectrum (r_1, r_2, \ldots, r_m) ? In [9], P. Zhao et al. obtained an upper bound on the minimum number of vertices of 3-uniform bi-hypergraphs with a given feasible set. Moreover, P. Zhao et al. [10] proved that the minimum number of vertices of one-realizations of a given set is $2 \max(S)$ – $\min(S)$ if $\max(S) - 1 \notin S$ or $2 \max(S) - \min(S) - 1$ otherwise. In this paper, we generalize this result to the minimum number of vertices of realizations of a special kind of vector.

In the rest of this paper, we always assume that n_i and l_i are two sets of integers with $2 \le n_s < \cdots < n_1$, $l_1 = 0$ and $0 \le l_i < n_{i-1} - n_i$ for all $i \in \{2, \ldots, s\}$; moreover, $R_2 = (r_1, r_2, \ldots, r_{n_1})$ is the vector with $r_1 = 0$ and $r_{n_i} = 2^{l_i}$, $i \in \{1, 2, \ldots, s\}$. For any positive integer n, let [n] denote the set $\{1, 2, \ldots, n\}$. The main result of this paper is as follows:

Theorem 1.1 If $\delta(R_2)$ is the minimum number of vertices of realizations of R_2 , then

$$\delta(R_2) = \begin{cases} 2n_1 - n_s, & \text{if } n_1 > n_2 + 1, \\ 2n_1 - n_s - 1, & \text{if } n_1 = n_2 + 1, \end{cases}$$

As a result, we partially solve the open problem proposed by Král.

2 The proof of Theorem 1.1

In this section, we first show that the number $\delta(R_2)$ given in Theorem 1.1 is a lower bound on the number of vertices of the smallest realization of R_2 , then construct two families of mixed hypergraphs which meet the bound in each case. For a mixed hypergraph $\mathcal{H} = (X, \mathcal{C}, \mathcal{D})$ and a strict coloring c of \mathcal{H} , let c(v) denote the color of $v \in X$ under c.

Lemma 2.1

$$\delta(R_2) \ge \begin{cases} 2n_1 - n_s, & \text{if } n_1 > n_2 + 1, \\ 2n_1 - n_s - 1, & \text{if } n_1 = n_2 + 1. \end{cases}$$

Proof. Since the first part was given by Theorem 3 in [2] and Lemma 2.2 in [10] respectively, it suffices to prove the second part. Let $n_1 = n_2 + 1$ and $\mathcal{H} = (X, \mathcal{C}, \mathcal{D})$ be a realization of R_2 . Then $l_2 = 0$, i.e., $r_{n_2} = 1$. Suppose $|X| \leq 2n_1 - (n_s + 2)$. Then for any strict n_1 -coloring $c_1 = \{C_1, C_2, \ldots, C_{n_1}\}$ of \mathcal{H} , there exist at least $n_s + 2$ color classes of size one. Assume that $C_1 = \{\alpha_1\}, C_2 = \{\alpha_2\}, \ldots, C_{n_s+2} = \{\alpha_{n_s+2}\}$. For any strict n_s -coloring c_s of \mathcal{H} , there are the following two possible cases.

Case 1. There exist three vertices in $\{\alpha_1, \alpha_2, \ldots, \alpha_{n_s+2}\}$ with the same color under c_s . Suppose $c_s(\alpha_1) = c_s(\alpha_2) = c_s(\alpha_3)$. Then $\{\alpha_1, \alpha_2\}, \{\alpha_1, \alpha_3\}, \{\alpha_2, \alpha_3\} \notin \mathcal{D}$, which follows that $\{C_1 \cup C_2, C_3, \ldots, C_{n_1}\}$, $\{C_1, C_2 \cup C_3, C_4, \ldots, C_{n_1}\}$ are strict n_2 -colorings of \mathcal{H} , a contradiction to that $r_{n_2} = 1$.

Case 2. There exist two pairs of vertices in $\{\alpha_1, \alpha_2, \ldots, \alpha_{n_s+2}\}$ such that each pair have the same color under c_s . Suppose $c_s(\alpha_1) = c_s(\alpha_2)$ and $c_s(\alpha_3) = c_s(\alpha_4)$. Then $\{\alpha_1, \alpha_2\}, \{\alpha_3, \alpha_4\} \notin \mathcal{D}$, which implies that $\{C_1 \cup C_2, C_3, \ldots, C_{n_1}\}$ and $\{C_1, C_2, C_3 \cup C_4, C_5, \ldots, C_{n_1}\}$ are strict n_2 -colorings of \mathcal{H} , also a contradiction to that $r_{n_2} = 1$.

In the rest of this section, we shall construct the desired mixed hypergraphs. Our construction here is based on the Construction I in [10]. In order to get the desired r_{n_k} strict n_k -colorings, we define the vertices of γ_{pk}^1 and γ_{pk}^2 in the construction.

Construction I. For $i \in [n_s]$, $p \in [s] \setminus \{1\}$, $j \in \{0, l_p + 1, l_p + 2, \ldots, n_{p-1} - n_p - 1\}$ and $k \in [l_p]$, write

$$\theta_i = (\underbrace{i, i, \dots, i}_{\sum_{i=1}^s 2^{l_i}}), \ \beta_1 = (n_1, \underbrace{n_2, \dots, n_2}_{2^{l_2}}, \dots, \underbrace{n_s, \dots, n_s}_{2^{l_s}})$$

$$\begin{split} &\alpha_{pj} = \underbrace{(n_p + j, \dots, n_p + j, \underbrace{1, \dots, 1}_{\sum_{i=1}^{p-1} 2^{l_i}}, \underbrace{\sum_{i=p}^{s} 2^{l_i}}_{\sum_{i=p} 2^{l_i}} \\ &\beta_{pj} = \underbrace{(n_p + j, \dots, n_p + j, n_p, \dots, n_p, \dots, n_s, \dots, n_s)}_{\sum_{i=1}^{p-1} 2^{l_i}}, \\ &\gamma_{pk}^1 = \underbrace{(n_p + k, \dots, n_p + k, \underbrace{\sum_{i=1}^{p-1} 2^{l_i}}_{2^{k_i}}, \underbrace{\sum_{i=1}^{p-1} 2^{l_i}}_{2^{k_i}}, \underbrace{\sum_{i=1}^{p-1} 2^{l_i}}_{2^{k_i}}, \underbrace{\sum_{i=p+1}^{p-1} 2^{l$$

where $\alpha_{l(j)}$ is the j^{th} entry of the vertex α_l . Then $\mathcal{H} = (X, \mathcal{C}, \mathcal{D})$ is a mixed hypergraph with $2n_1 - n_s$ vertices.

For example, when $S = \{2, 4\}$ and $l_2 = 1$, then $X = \{\theta_1, \alpha_{20}, \beta_{20}, \gamma_{21}^1, \gamma_{21}^2, \beta_1\}$, where $\theta_1 = (1, 1, 1), \theta_2 = (2, 2, 2), \alpha_{20} = (2, 1, 1), \beta_{20} = (2, 2, 2), \gamma_{21}^1 = (3, 2, 1), \gamma_{21}^2 = (3, 1, 2), \beta_1 = (4, 2, 2)$ and $C = \{\{\theta_1, \alpha_{20}, \beta_{20}\}, \{\theta_1, \gamma_{21}^1, \gamma_{21}^2\}\}, \{\alpha_{20}, \gamma_{21}^1, \gamma_{21}^2\}\}, \{\beta_{20}, \gamma_{21}^1, \gamma_{21}^2\}, \{\gamma_{21}^1, \alpha_{20}, \beta_{20}\}, \{\gamma_{21}^2, \alpha_{20}, \beta_{20}\}, \{\beta_1, \alpha_{20}, \beta_{20}\}, \{\beta_1, \gamma_{21}^1, \gamma_{21}^2\}\}.$ $D = \{\{\theta_1, \beta_{20}\}, \{\theta_1, \beta_1\}, \{\alpha_{20}, \beta_1\}, \{\beta_1, \gamma_{21}^1, \gamma_{21}^2\},$

$$\{\alpha_{20},\beta_{20},\gamma_{21}^1\},\{\alpha_{20},\beta_{20},\gamma_{21}^2\},\{\alpha_{20},\gamma_{21}^1,\gamma_{21}^2\},\{\beta_{20},\gamma_{21}^1,\gamma_{21}^2\}\}.$$

Moreover, when $S = \{2, t\}$ and $l_2 = 0$, our construction is similar to the first construction in [2], but our edge-set gives a one-realizer whereas the construction in [2] does not. The later construction in [2] of their proper realizers uses many more than the minimum number of vertices.

Note the inspiration for our construction is so that for each $i \in [s]$ and $h \in [2^{l_i}]$, $c_{ih} = \{X_{ih1}, X_{ih2}, \dots, X_{ihn_i}\}$ is a strict n_i -coloring of \mathcal{H} , where X_{iht} consists of vertices

$$(x_1^1, x_2^1, \dots, x_2^{2^{l_2}}, \dots, x_i^1, \dots, x_i^{h-1}, t, x_i^{h+1}, \dots, x_i^{2^{l_i}}, \dots, x_s^1, \dots, x_s^{2^{l_s}}) \in X.$$

Lemma 2.2 Let $c = \{C_1, C_2, \ldots, C_m\}$ be a strict coloring of \mathcal{H} . Then we may reorder the color classes such that the following conditions hold:

- (i) $\theta_i \in C_i$, $i \in [n_s]$;
- (ii) $\alpha_{pj} \notin C_i$, $i \in [n_s 1] \setminus \{1\}$;
- (iii) $\beta_1, \beta_{pj} \notin C_i, i \in [n_s 1];$
- (iv) $\gamma_{pk}^1, \gamma_{pk}^2 \notin C_i, i \in [n_s 1] \setminus \{1\};$
- (v) $\alpha_{s0} \in C_1 \cup C_{n_s}$

Proof. The \mathcal{D} -edge $\{\theta_i, \theta_j\}$ implies that $c(\theta_i) \neq c(\theta_j)$ if $i \neq j$. Hence, we may reorder the color classes such that $\theta_i \in C_i$, $i \in [n_s]$. It follows that (i) holds.

For $i \in [n_s - 1] \setminus \{1\}$, from the \mathcal{D} -edges $\{\alpha_{pj}, \theta_i\}$, $\{\gamma_{pk}^1, \theta_i\}$, $\{\gamma_{pk}^2, \theta_i\}$, one gets $\alpha_{pj}, \gamma_{pk}^1, \gamma_{pk}^2 \notin C_i$. Hence, (ii) and (iv) hold.

Since $\{\beta_1, \theta_i\}$, $\{\beta_{pj}, \theta_i\}$ are \mathcal{D} -edges, we have $\beta_1, \beta_{pj} \notin C_i$ for $i \in [n_s - 1]$, which implies that (iii) holds. The \mathcal{C} -edge $\{\alpha_{s0}, \theta_{n_s}, \theta_1\}$ implies that $\alpha_{s0} \in C_1 \cup C_{n_s}$, as desired.

Lemma 2.3 Let $c = \{C_1, C_2, \dots, C_m\}$ be a strict coloring of \mathcal{H} satisfying the conditions (i)-(v) in Lemma 2.2.

- (i) Suppose $c(\alpha_{tj_t}) \neq c(\beta_{tj_t})$ for some $t \in [s] \setminus \{1\}$ and $j_t \in \{0, l_t + 1, \ldots, n_{t-1} n_t 1\}$. Then there exists an $a \in [m] \setminus [n_s 1]$ such that $\alpha_{pj} \in C_1$ and $\beta_1, \beta_{pj} \in C_a$ for all $p \in [t] \setminus \{1\}$;
- (ii) Suppose $c(\alpha_{qj_q}) = c(\beta_{qj_q})$ for some $q \in [s] \setminus \{1\}$ and $j_q \in \{0, l_q + 1, \ldots, n_{q-1} n_q 1\}$. Then $c(\alpha_{pj}) = c(\beta_{pj})$ and $c(\gamma_{pk}^1) = c(\gamma_{pk}^2)$ for all $p \in \{q, \ldots, s\}$.

Proof. (i) From the C-edge $\{\alpha_{tj_t}, \beta_{tj_t}, \theta_1\}$, we have $\alpha_{tj_t} \in C_1$. From Lemma 2.2, there is an $a \in [m] \setminus [n_s - 1]$ with $\beta_{tj_t} \in C_a$. For $p \in [t] \setminus \{1\}$, if $n_p + j \neq n_t + j_t$, then the C-edges $\{\beta_1, \alpha_{tj_t}, \beta_{tj_t}\}$ and $\{\beta_{pj}, \alpha_{tj_t}, \beta_{tj_t}\}$ imply that $\beta_1, \beta_{pj} \in C_a$; moreover, from the C-edge $\{\alpha_{pj}, \beta_{pj}, \theta_1\}$ and the D-edge $\{\alpha_{pj}, \beta_{tj_t}\}$, we have $\alpha_{pj} \in C_1$, as desired.

(ii) If $c(\alpha_{tj_t}) \neq c(\beta_{tj_t})$ holds for some $t \in \{q, \dots, s\}$ and $j_t \in \{0, l_t + 1, \dots, n_{t-1} - n_t - 1\}$, then by (i) we have $c(\alpha_{pj}) \neq c(\beta_{pj})$ for all $p \in [t] \setminus \{1\}$. It follows that $c(\alpha_{qj_q}) \neq c(\beta_{qj_q})$, a contradiction. Hence, $c(\alpha_{pj}) = c(\beta_{pj})$ for $p \in \{q, \dots, s\}$. Moreover, for $p \in \{q, \dots, s\}$, from the \mathcal{D} -edges $\{\gamma_{pk}^1, \alpha_{p0}, \beta_{p0}\}$ and $\{\gamma_{pk}^2, \alpha_{p0}, \beta_{p0}\}$, we have $c(\gamma_{pk}^1), c(\gamma_{pk}^2) \neq c(\alpha_{p0})$; and the \mathcal{C} -edge $\{\gamma_{pk}^1, \gamma_{pk}^2, \alpha_{p0}\}$ implies that $c(\gamma_{pk}^1) = c(\gamma_{pk}^2)$.

Lemma 2.4 Let $c = \{C_1, C_2, \ldots, C_m\}$ be a strict coloring of \mathcal{H} satisfying that $c(\alpha_{t0}) \neq c(\beta_{t0})$ for some $t \in [s] \setminus \{1\}$. Then there exists an $a \in [m] \setminus [n_s - 1]$ such that

- (i) $\gamma_{tk}^1, \gamma_{tk}^2 \in C_1 \cup C_a \text{ and } c(\gamma_{tk}^1) \neq c(\gamma_{tk}^2);$
- (ii) $\gamma_{pk}^1 \in C_1$ and $\gamma_{pk}^2 \in C_a$ for each $p \in [t-1] \setminus \{1\}$.

Proof. For $p \in [t-1] \setminus \{1\}$, from Lemma 2.3, there is an $a \in [m] \setminus [n_s-1]$ with $\alpha_{p0}, \alpha_{t0} \in C_1$ and $\beta_{p0}, \beta_{t0} \in C_a$. For $p \in [t] \setminus \{1\}$, since $\{\gamma_{pk}^1, \alpha_{p0}, \beta_{p0}\}$ and $\{\gamma_{pk}^2, \alpha_{p0}, \beta_{p0}\}$ are \mathcal{C} -edges, $\gamma_{pk}^1, \gamma_{pk}^2 \in C_1 \cup C_a$; and the \mathcal{D} -edges $\{\gamma_{pk}^1, \gamma_{pk}^2, \beta_{p0}\}$ and $\{\gamma_{pk}^1, \gamma_{pk}^2, \alpha_{p0}\}$ imply that $c(\gamma_{pk}^1) \neq c(\gamma_{pk}^2)$. Specially, $\gamma_{tk}^1, \gamma_{tk}^2 \in C_1 \cup C_a$ and $c(\gamma_{tk}^1) \neq c(\gamma_{tk}^2)$. Hence, (i) holds.

For $p \in [t-1] \setminus \{1\}$, from the \mathcal{C} -edge $\{\gamma_{pk}^1, \alpha_{t0}, \beta_{t0}\}$ and the \mathcal{D} -edge $\{\gamma_{pk}^1, \beta_{t0}\}$, one gets $\gamma_{pk}^1 \in C_1$. Then by (i) we have $\gamma_{pk}^2 \in C_a$, which implies that (ii) holds.

Theorem 2.5 \mathcal{H} is a realization of R_2 .

Proof. It suffices to prove that $c_{11}, c_{21}, \ldots, c_{22^{l_2}}, \ldots, c_{s_1}, \ldots, c_{s_2^{l_s}}$ are all of the strict colorings of \mathcal{H} . Suppose $c = \{C_1, C_2, \ldots, C_m\}$ is a strict coloring of \mathcal{H} satisfying the conditions (i)-(v) in Lemma 2.2. In particular, $\alpha_{s0} \in C_1 \cup C_{n_s}$.

Case 1 $\alpha_{s0} \in C_1$.

Note that $\beta_{s0} \in C_{n_s}$. By Lemma 2.3 we have $\alpha_{pj} \in C_1$ and $\beta_1, \beta_{pj} \in C_{n_s}$ for all $p \in [s] \setminus \{1\}$. Then by Lemma 2.4 one gets that

- (i) $\gamma_{sk}^1, \gamma_{sk}^2 \in C_1 \cup C_{n_s}$ and $c(\gamma_{sk}^1) \neq c(\gamma_{sk}^2)$;
- (ii) $\gamma_{pk}^1 \in C_1$ and $\gamma_{pk}^2 \in C_{n_s}$ for $p \in [s-1] \setminus \{1\}$.

It follows that $c \in \{c_{sh} \mid h \in [2^{l_s}]\}.$

Case 2 $\alpha_{s0} \in C_{n_s}$.

In this case, we shall prove that $c \in \{c_{ih} \mid i \in [s-1], h \in [2^{l_i}]\}$. Let $t \in [s] \setminus \{1\}$ be the minimum number such that $c(\alpha_{t0}) = c(\beta_{t0})$. By Lemma 2.3 we have $c(\alpha_{pj}) = c(\beta_{pj})$ and $c(\gamma_{pk}^1) = c(\gamma_{pk}^2)$ for each $p \in \{t, \ldots, s\}$. For $p_1, p_2 \in \{t, \ldots, s\}$ and $p_1 \neq p_2$, from the \mathcal{D} -edges $\{\alpha_{p_1j_1}, \beta_{p_2j_2}\}$, $\{\alpha_{p_1j_1}, \beta_{p_1j_1}, \gamma_{p_2k_2}^2\}$ and $\{\alpha_{p_2j_2}, \beta_{p_2j_2}, \gamma_{p_1k_1}^1\}$, $\{\gamma_{p_1k_1}^1, \gamma_{p_2k_2}^1, \gamma_{p_2k_2}^2\}$, one gets $c(\alpha_{p_1j_1}) \neq c(\beta_{p_2j_2})$, $c(\alpha_{p_1j_1}) \neq c(\gamma_{p_2k_2}^2)$ and $c(\gamma_{p_1k_1}^1) \neq c(\beta_{p_2j_2})$, $c(\gamma_{p_1k_1}^1) \neq c(\gamma_{p_2k_2}^2)$ for $j_i \in \{0, l_{p_i} + 1, \ldots, n_{p_i-1} - n_{p_i} - 1\}$, $k_i \in [l_{p_i}]$, i = 1, 2. Moreover, for $p \in \{t, \ldots, s\}$, the \mathcal{D} -edge $\{\beta_{pj_1}, \alpha_{pj_2}\}$ implies that $c(\alpha_{pj}) \neq c(\gamma_{pk}^1)$; and the \mathcal{D} -edge $\{\gamma_{pk_1}^1, \gamma_{pk_1}^1, \gamma_{pk_2}^2\}$ implies that $c(\gamma_{pk_1}^1) \neq c(\gamma_{pk_2}^1)$ if $k_1 \neq k_2$. Hence, we may assume that $\alpha_{pj}, \beta_{pj} \in C_{n_p+j}$ and $\gamma_{pk}^1, \gamma_{pk}^2 \in C_{n_p+k}$ for $p \in \{t, \ldots, s\}$, $j \in \{0, l_p + 1, \ldots, n_{p-1} - n_p - 1\}$ and $k \in [l_p]$.

Case 2.1 t > 2. That is to say, $c(\alpha_{t-1,0}) \neq c(\beta_{t-1,0})$. By Lemma 2.3 we have $\alpha_{t-1,0} \in C_1$ and $\beta_{t-1,0} \notin C_1$. For $p \in \{t, \ldots, s\}$, from the \mathcal{D} -edges $\{\beta_{t-1,0}, \alpha_{pj}\}$ and $\{\beta_{t-1,0}, \gamma_{pk}^1\}$, we have $\beta_{t-1,0} \notin C_{n_p+j} \cup C_{n_p+k}$. Hence, we may assume that $\beta_{t-1,0} \in C_{n_{t-1}}$. By Lemma 2.3 we have $\alpha_{pj} \in C_1$ and $\beta_1, \beta_{pj} \in C_{n_{t-1}}$ for $p \in [t-1] \setminus \{1\}$. Moreover, by Lemma 2.4 we have

(i)
$$\gamma_{t-1,k}^1, \gamma_{t-1,k}^2 \in C_1 \cup C_{n_{t-1}}$$
 and $c(\gamma_{t-1,k}^1) \neq c(\gamma_{t-1,k}^2)$;

(ii) $\gamma_{pk}^1 \in C_1$ and $\gamma_{pk}^2 \in C_{n_{t-1}}$ for any $p \in [t-2] \setminus \{1\}$.

Therefore, $c \in \{c_{t-1,h} \mid h \in [2^{l_{t-1}}]\}.$

Case 2.2 t=2. From the \mathcal{D} -edges $\{\beta_1, \alpha_{pj}\}$ and $\{\beta_1, \gamma_{pk}^1, \gamma_{pk}^2\}$, we have $\beta_1 \notin C_{n_p+j} \cup C_{n_p+k}$ for all $p \in [s] \setminus \{1\}$. It follows that $c=c_{11}$. The proof is completed.

For the case of $n_1 = n_2 + 1$, we have the following construction.

Construction II. Let $X' = X \setminus \{\alpha_{20}\}$ and $\mathcal{H}' = \mathcal{H}[X']$. Then for each $i \in [s], h \in [2^{l_i}], c'_{ih} = \{X'_{ih1}, X'_{ih2}, \dots, X'_{ihn_i}\}$ is a strict n_i -coloring of \mathcal{H}' , where $X'_{iht} = X_{iht} \cap X'$.

Theorem 2.6 If $n_1 = n_2 + 1$, then \mathcal{H}' is a realization of $R_2 = (0, r_2, \dots, r_{n_2-1}, 1, 1)$.

Proof. Let $\mathcal{H}'' = (X'', \mathcal{C}'', \mathcal{D}'')$ be the realization of $R_2'' = (0, r_2, \dots, r_{n_2-1}, 1)$ with $n_1'' = n_2$ given by Construction I. Let $Y = \{\alpha \mid \alpha \in X, \alpha_{(1)} = \alpha_{(2)}\}$, where $\alpha_{(j)}$ is the j^{th} entry of the vertex α . Then

$$\phi: Y \to X''
(x_2^1, x_2^1, x_3^1, \dots, x_3^{2^{l_3}}, \dots, x_s^1, \dots, x_s^{2^{l_s}}) \to (x_2^1, x_3^1, \dots, x_3^{2^{l_3}}, \dots, x_s^1, \dots, x_s^{2^{l_s}})$$

is an isomorphism from $\mathcal{G} = \mathcal{H}[Y]$ to \mathcal{H}'' . By Theorem 2.5, all of the strict colorings of \mathcal{G} are as follows:

$$e_{ih} = \{Y_{ih1}, Y_{ih2}, \dots, Y_{ihn_i}\}, i \in [s] \setminus \{1\}, h \in [2^{l_i}],$$

where Y_{iht} consists of vertices

$$(x_2^1, x_2^1, x_3^1, \dots, x_3^{2^{l_3}}, \dots, x_i^1, \dots, x_i^{h-1}, t, x_i^{h+1}, \dots, x_i^{2^{l_i}}, \dots, x_s^1, \dots, x_s^{2^{l_s}}) \in X.$$

Note that the restriction on Y of any strict coloring of \mathcal{H}' corresponds to a strict coloring of \mathcal{G} . For any strict coloring $c = \{C_1, C_2, \ldots, C_m\}$ of \mathcal{H}' , there are the following two possible cases.

Case 1 $c|_{Y} = e_{21}$.

By the proof of Theorem 2.5, we have $\beta_1 \notin C_j$ for any $j \in [n_2 - 1]$. Then, it is immediate that $c = c'_{21}$ if $\beta_1 \in C_{n_2}$, or $c = c'_{11}$ if $\beta_1 \notin C_{n_2}$.

Case 2 $c|_Y = e_{ih}$ for some $i \in [s] \setminus \{1, 2\}$ and $h \in [2^{l_i}]$.

Note that $\beta_{i0} \in C_{n_i}$ and $\alpha_{i0} \in C_1$. From the \mathcal{C} -edge $\{\alpha_{i0}, \beta_{i0}, \beta_1\}$ and the \mathcal{D} -edge $\{\beta_1, \theta_1\}$, we observe that $\beta_1 \in C_{n_i}$. Therefore, $c = c'_{ih}$. Hence, the desired result follows.

Combining Lemma 2.1, Theorem 2.5 and Theorem 2.6, the proof of Theorem 1.1 is completed.

Acknowledgment

The research is supported by NSF of Shandong Province (ZR2013AL009), promotive research fund for excellent young and middle-aged scientists of Shandong province(Grant No. BS2013DX026) and AMEP of Linvi University.

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