## A Note on Edge Coloring of Graphs

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## Abstract

Let G be a graph of minimum degree  $\delta(G)$ . R.P. Gupta proved the two following interesting results: 1) A bipartite graph G has a  $\delta$ -edge-coloring in which all  $\delta$  colors appear at each vertex. 2) If G is a simple graph with  $\delta(G)>1$ , then G has a  $(\delta-1)$ -edge-coloring in which all  $(\delta-1)$  colors appear at each vertex. Let t be a positive integer. In this paper, we extend the first result by showing that for every bipartite graph, there exists a t-edge coloring such that, at each vertex v,  $\min\{t,d(v)\}$  colors appear. Also, we show that if G is a graph, then the edges of G can be colored using t colors in which for each vertex v, the number of colors appear at v is at least  $\min\{t,d(v)-1\}$ , which generalizes the second result.

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## Introduction.

Throughout this paper, all graphs are simple. Let G be a graph and let E(G) and V(G) denote the edge set and vertex set of G, respectively. We

denote the minimum and the maximum degree of G by  $\delta(G)$  and  $\Delta(G)$ , respectively. For every  $v \in V(G)$ ,  $d_G(v)(=d(v))$  denotes the degree of v. A k-edge-coloring of a graph G is a mapping  $c: E \longrightarrow S$ , where S is a set of k colors. An edge coloring of G is called proper if the adjacent edges receive distinct colors. The edge chromatic number,  $\chi'(G)$ , of a graph G is the minimum positive integer k for which G has a proper k-edge-coloring. Consider an edge coloring of G. For a vertex  $v \in V(G)$ , let  $s_G(v)$  (= s(v)) denote the number of different colors appearing on the edges incident with v. A celebrated theorem due to Vizing says that for every graph G,  $\Delta(G) \leq \chi'(G) \leq \Delta(G) + 1$  (see Theorem 17.4 of [1]).

R.P. Gupta proved that a bipartite graph G has a (not necessarily proper)  $\delta(G)$ -edge-coloring in which all  $\delta(G)$  colors appear at each vertex (see Exercise 17.1.16 of [1]). The following theorem generalizes this result.

**Theorem 1.** Let G be a bipartite graph and let t be a positive integer. Then all edges of G can be colored using t colors such that for each vertex  $v, s(v) \ge \min\{t, d(v)\}$ .

**Proof.** We apply induction on m = |E(G)|. For m = 1, the assertion is trivial. If  $\Delta(G) > t$ , then consider a vertex of degree  $\Delta(G)$ , say u. Let v be one of its neighbors and let e be the edge between u and v. By the induction hypothesis,  $H = G \setminus \{e\}$  has an edge coloring with t colors such that  $s_H(w) \ge \min\{t, d_H(w)\}$  for each vertex w. Since  $\min\{\Delta(G) - 1, t\} = t$ , by the induction hypothesis, t colors appear at u. If  $\min\{t, d(v) - 1\} = t$ , then by an arbitrary coloring of e we obtain the result. Now, if  $\min\{t, d(v) - 1\} = d(v) - 1$ , then we can color e by a color not appearing at v. Obviously,  $s(w) \ge \min\{t, d(w)\}$  for each vertex w of G. So assume that  $\Delta(G) \le t$ . By Theorem 17.2 of [1], since G is a bipartite graph,  $\chi'(G) = \Delta(G)$ . Now, by considering a proper  $\Delta(G)$ -edge-coloring of G we are done.  $\Box$ 

We note that this theorem is not true for a non-bipartite graph. Consider the complete graph  $G = K_{t+1}$  for even t. Then G has no t-edge-coloring with the desired property (see Exercise 17.1.15 [1]).

**Theorem 2.** Let G be a graph with no odd cycle component. Then all edges of G can be colored by 2 colors such that for each vertex v with degree at least 2,  $s(v) \geq 2$ .

**Proof.** Without loss of generality, we may assume that G is a connected graph. First, suppose that G is an Eulerian graph. If G is a cycle, then consecutively color the edges of G by the colors 1 and 2. If G is not a cycle, then it has a vertex v of degree at least 3. It suffices to consider an Eulerian circuit in G with starting vertex v and color its edges alternatively by colors 1 and 2. If G is not Eulerian, then we add a new vertex u and join u to all vertices of odd degree. Then the new graph is an Eulerian graph and we proceed as before. Now, by removing vertex v the proof is complete.

R.P. Gupta proved that if G is a simple graph with  $\delta(G) > 1$ , then G has a  $(\delta(G)-1)$ -edge-coloring (necessarily improper) in which all  $\delta(G)-1$  colors appear at each vertex (see [1, p.461]). The next theorem generalizes Gupta's result.

**Theorem 3.** If G is a graph and t is a positive integer, then all edges of G can be colored using t colors such that for each vertex v,  $s(v) \ge \min\{t, d(v) - 1\}$ .

**Proof.** First, assume that  $t \leq \delta(G) - 1$ . Obviously, by removing  $\delta(G) - t + 1$  edges of G, we find a spanning subgraph G' of G with  $\delta(G') = t + 1$ . By Gupta's result, we conclude that there exists a  $(\delta(G') - 1)$ -edge-coloring of G' such that s(v) = t for every vertex v. Now, we arbitrary color all edges of  $E(G) \setminus E(G')$  and in this case we are done. So assume that  $t \geq \delta(G)$ . There exists a graph H such that G is an induced subgraph of H and  $\delta(H) = t + 1$ . To see this, consider two copies of G and join the corresponding vertices of degree  $\delta(G)$  and repeat this procedure  $t + 1 - \delta(G)$  times. By Gupta's result, there exists a  $(\delta(H) - 1)$ -edge-coloring of H such that s(v) = t for every  $v \in V(H)$ . Now, consider the restriction of this

coloring to the graph G. In this coloring of G, for every  $v \in V(G)$  we have  $s(v) \ge t - \max\{t - d(v) + 1, 0\}$ , which implies that  $s(v) \ge \min\{t, d(v) - 1\}$  and the proof is complete.

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