

# Implicit degree sum condition for hamiltonian cycles

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**Abstract:** The hamiltonian problem is a classical problem in graph theory. Most of the research on hamiltonian problem is looking for sufficient conditions for a graph to be hamiltonian. For a vertex  $v$  of a graph  $G$ , Zhu, Li and Deng introduced the concept of implicit degree  $id(v)$ , according to the degrees of its neighbors and the vertices at distance 2 with  $v$  in  $G$ . In this paper, we will prove that: Let  $G$  be a 2-connected graph on  $n \geq 3$  vertices. If the maximum value of the implicit degree sums of 2 vertices in  $S$  is more than or equal to  $n$  for each independent set  $S$  with  $\kappa(G) + 1$  vertices, then  $G$  is hamiltonian.

**Keywords:** Implicit degree sum; Independent set; Hamiltonian cycle

## 1 Introduction

In this paper, we consider only finite, undirected and simple graphs. Notation and terminology not defined here can be found in [2]. Let  $G = (V(G), E(G))$  be a graph with vertex set  $V(G)$  and edge set  $E(G)$ , and  $H$  be a subgraph of  $G$ . For a vertex  $u \in V(G)$ , let  $N_H(u) = \{v \in V(H) : uv \in E(G)\}$  and  $d_H(u) = |N_H(u)|$ . If  $G = H$ , we always use  $N(u)$  and  $d(u)$  in place of  $N_G(u)$  and  $d_G(u)$  respectively. Let  $N_2(u) = \{v \in V(G) : d(u, v) = 2\}$ , where  $d(u, v)$  denotes the distance between vertices  $u$  and  $v$  in  $G$ .

Let  $\alpha(G)$  and  $\kappa(G)$  denote the independence number and the connectivity of  $G$ , respectively. For a nonempty set  $S \subset V(G)$ , let  $\Delta_k(S) = \max\{\sum_{x \in X} d(x) : X \text{ is a subset of } S \text{ with } k \text{ vertices}\}$ .

A cycle containing all vertices of  $G$  is called a hamiltonian cycle. A graph  $G$  is called hamiltonian if it contains a hamiltonian cycle. The hamiltonian problem is an important problem in graph theory. Various sufficient conditions for a graph to be hamiltonian have been given in terms of degree conditions. The following two sufficient conditions are essential.

**Theorem 1** ([7]). *Let  $G$  be a graph on  $n \geq 3$  vertices. If  $d(x) + d(y) \geq n$  for every pair of nonadjacent vertices  $x$  and  $y$ , then  $G$  is hamiltonian.*

**Theorem 2** ([5]). *Let  $G$  be a 2-connected graph on  $n \geq 3$  vertices. If  $\max\{d(x), d(y)\} \geq n/2$  for every pair of vertices  $x$  and  $y$  at distance 2, then  $G$  is hamiltonian.*

In 1972, by considering the relationship between the independence number and the connectivity of a graph, Chvátal and Erdős gave a sufficient condition for a 2-connected graph to be hamiltonian.

**Theorem 3** ([3]). *Let  $G$  be a 2-connected graph with  $\alpha(G) \leq \kappa(G)$ . Then  $G$  is hamiltonian.*

Recently, Yamashita improved Ore's Theorem (Theorem 1) as follows.

**Theorem 4** ([8]). *Let  $G$  be a 2-connected graph of on  $n \geq 3$  vertices. If  $\Delta_2(S) \geq n$  for every independent set  $S$  of order  $\kappa(G) + 1$ , then  $G$  is hamiltonian.*

In order to generalize and improve the classic results of hamiltonian problem, Zhu, Li and Deng [9] gave the concept of implicit degree of a vertex.

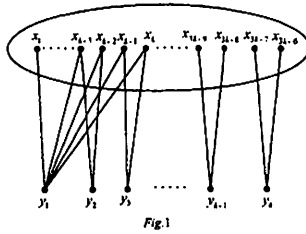
**Definition 1** ([9]). *Let  $v$  be a vertex of a graph  $G$  and  $k = d(v) - 1$ . Set  $M_2 = \max\{d(u) : u \in N_2(v)\}$  and  $m_2 = \min\{d(u) : u \in N_2(v)\}$ . Suppose  $d_1 \leq d_2 \leq d_3 \leq \dots \leq d_k \leq d_{k+1} \leq \dots$  is the degree sequence of vertices in  $N(v) \cup N_2(v)$ . If  $N_2(v) \neq \emptyset$  and  $d(v) \geq 2$ , define*

$$d^*(v) = \begin{cases} m_2, & \text{if } d_k < m_2; \\ d_{k+1}, & \text{if } d_{k+1} > M_2; \\ d_k, & \text{otherwise,} \end{cases}$$

*then the implicit degree of  $v$  is defined as  $id(v) = \max\{d(v), d^*(v)\}$ . If  $N_2(v) = \emptyset$  or  $d(v) \leq 1$ , then  $id(v) = d(v)$ .*

Clearly,  $id(v) \geq d(v)$  for each vertex  $v$  from the definition of implicit degree. For  $S \subseteq V(G)$  with  $S \neq \emptyset$ , let  $i\Delta_k(G) = \max\{\sum_{x \in X} id(x) : X \text{ is a subset of } S \text{ with } k \text{ vertices}\}$ . The authors [9] used implicit degree sum instead of degree sum in Ore's theorem (Theorem 1), and got a sufficient condition for a graph to be hamiltonian.

**Theorem 5** ([9]). *Let  $G$  be a 2-connected graph on  $n \geq 3$  vertices. If  $id(u) + id(v) \geq n$  for each pair of nonadjacent vertices  $u$  and  $v$  in  $G$ , then  $G$  is hamiltonian.*



From Theorem 5, we can easily deduce Fan's Theorem (Theorem 2). Moreover, the authors in [9] gave an example to illustrate Theorem 5 is stronger than Fan's theorem. Motivated by the results of Theorem 4 and Theorem 5, we study implicit degrees and the hamiltonicity of graphs and obtain the following main result.

**Theorem 6.** *Let  $G$  be a 2-connected graph on  $n \geq 3$  vertices. If  $i\Delta_2(S) \geq n$  for every independent set  $S$  of order  $\kappa(G) + 1$ , then  $G$  is hamiltonian.*

We postpone the proof of Theorem 6 in next section. Here we give two remarks. One shows that Theorem 6 is much stronger than Theorem 4 and Theorem 5; the other shows that the condition " $i\Delta_2(S) \geq n$ " could not be weakened to " $i\Delta_2(S) \geq n - 1$ ".

Let  $G_1, \dots, G_k$  be  $k$  vertex disjoint graphs. The union of  $G_1, \dots, G_k$ , denoted by  $G_1 \cup \dots \cup G_k$ , is the graph with vertex set  $\cup_{i=1}^k V(G_i)$  and edge set  $\cup_{i=1}^k E(G_i)$ . We use  $kQ$  instead of  $G_1 \cup \dots \cup G_k$  if each  $G_i$  is isomorphic to  $Q$ . The join of  $G_1, \dots, G_k$ , denoted by  $G_1 \vee \dots \vee G_k$ , is the graph obtained from  $\cup_{i=1}^k G_i$  by joining each vertex of  $G_i$  to each vertex of  $G_j$  for  $i \neq j$ .

**Remark 1.** *The graph in Fig.1 shows that Theorem 6 is much stronger than Theorem 4 and Theorem 5. Let  $k \geq 4$ ,  $V(K_{3k-6}) = \{x_1, x_2, \dots, x_{3k-6}\}$ ,  $V(kK_1) = \{y_1, y_2, \dots, y_k\}$ . We construct a graph  $G$ , with vertex set  $V(G) = V(K_{3k-6} \cup kK_1)$ , and edge set  $E(G) = E(K_{3k-6}) \cup \{y_1x_i : i = 1, 2, \dots, k\} \cup \{y_i x_{2i+k-7}, y_i x_{2i+k-6} : i = 2, 3, \dots, k\}$ . It is easy to check that  $G$  is a 2-connected hamiltonian graph on  $n = 4k - 6$  vertices. Choose  $S = \{y_1, y_2, y_3\}$ , then  $\Delta_2(S) = k + 2 < n$  and  $G$  does not satisfy the condition in Theorem 4. Since  $id(y_1) = 3k - 6$ ,  $id(y_2) = id(y_3) = k$ ,  $id(y_j) = 3k - 6$  for  $j = 4, 5, \dots, k$ ,  $id(x_2) = 2k + 1$  and  $id(x_l) \geq d(x_l) \geq 3k - 6$  for  $l = 1, 2, \dots, 3k - 6$ . Since  $y_2$  and  $y_3$  are two nonadjacent vertices of  $G$  and  $id(y_2) + id(y_3) = 2k < n$ ,  $G$  does not satisfy the condition in Theorem 5. But it is easy to check that  $G$  satisfies the condition in Theorem 6.*

**Remark 2.** *Let  $G = K_k \vee (k + 1)K_1$ . Clearly,  $G$  is a  $k$ -connected non-hamiltonian graph of order  $n = 2k + 1$ . Let  $S = V((k + 1)K_1)$ . It is*

easy to check that  $i\Delta_2(S) = 2k = n - 1$ , This implies that the condition “ $i\Delta_2(S) \geq n$ ” in Theorem 6 is best possible.

## 2 Proof of Theorem 6

A path  $P$  connecting  $x$  and  $y$  is called an  $H$ -path if  $V(P) \cap V(H) = \{x, y\}$  and  $E(P) \cap E(H) = \emptyset$ . For a cycle  $C$  in  $G$  with a given orientation and a vertex  $x$  in  $C$ ,  $x^+$  and  $x^-$  denote the successor and the predecessor of  $x$  in  $C$ , respectively. For any  $I \subseteq V(C)$ , let  $I^- = \{x : x^+ \in I\}$  and  $I^+ = \{x : x^- \in I\}$ . For two vertices  $x$  and  $y$  in  $C$ , we define  $xCy$  to be the path of  $C$  from  $x$  to  $y$ .  $y\bar{C}x$  denotes the path from  $y$  to  $x$  in the reversed direction of  $C$ . A similar notation is used for paths.

For a path  $P = x_1x_2 \dots x_p$  of a graph  $G$ , let  $l_P(x_1) = \max\{i : x_i \in V(P) \text{ and } x_ix_1 \in E(G)\}$  and  $l_P(x_p) = \min\{i : x_i \in V(P) \text{ and } x_ix_p \in E(G)\}$ . Set  $L_P(x_1) = x_{l_P(x_1)}$  and  $L_P(x_p) = x_{l_P(x_p)}$ . Our proof of Theorem 6 is based on the following lemmas.

**Lemma 1** ([1]). *Let  $G$  be a 2-connected graph on  $n$  vertices and  $C$  be a longest cycle of  $G$  with length at most  $n - 1$ . If  $P$  is a path connecting  $x$  and  $y$  in  $G$  such that  $|V(C)| < |V(P)|$ , then  $d(x) + d(y) < n$ .*

**Lemma 2** ([4]). *Let  $G$  be a 2-connected graph and  $X$  be a subset of  $V(G)$ . If  $|X| \leq \kappa(G)$ , then  $G$  has a cycle that includes every vertex of  $X$ .*

**Lemma 3** ([6]). *Let  $G$  be a 2-connected graph and  $P = x_1x_2 \dots x_p$  be a path of  $G$ . If  $x_1x_p \notin E(G)$ , and  $d(u) < id(x_1)$  for any  $u \in N_{G-V(P)}(x_1) \cup \{x_1\}$ , then either*

- (1) *there exists a vertex  $x_j \in N_P(x_1)^-$  such that  $d(x_j) \geq id(x_1)$ ; or*
- (2)  *$N_P(x_1)^- = N_P(x_1) \cup \{x_1\} - \{L_P(x_1)\}$ ,  $d(x_j) < id(x_1)$  for any vertex  $x_j \in N_P(x_1)^-$  and  $id(x_1) = \min\{d(v) : v \in N_2(x_1)\}$ .*

**Proof of Theorem 6.** Suppose to the contrary that  $G$  is a graph satisfying the condition of Theorem 6 and  $G$  is not hamiltonian. Let  $C$  be a longest cycle of  $G$  and give  $C$  a clockwise orientation. Then  $|V(C)| < n$ .

Let  $H$  be a component of  $G - V(C)$  and  $y_0 \in V(H)$ . Set  $k = \kappa(G)$ . By Lemma 2,  $|V(C)| \geq k$ . Since the connectivity of  $G$  is  $k$ , there are  $k$  paths  $P_1(y_0, y_1), P_2(y_0, y_2), \dots, P_k(y_0, y_k)$  from  $y_0$  to  $C$  having only  $y_0$  in common pairwise and  $V(P_i) \cap V(C) = \{y_i\}$  for each  $1 \leq i \leq k$ . Without loss of generality, we orient  $P_i$  from  $y_0$  to  $y_i$ . Assume without loss of generality that  $y_1, y_2, \dots, y_k$  occur in this order along  $C$ . Let  $x_i = y_i^+$  for each  $i = 1, 2, \dots, k$ . Let  $x_0$  and  $z_0$  be the predecessor of  $y_2$  on the path  $P_2$  and the predecessor of  $y_1$  on the path  $P_1$ , respectively.

**Claim 1.**  $\{x_0, x_1, x_2, \dots, x_k\}$  is an independent set of  $G$ . Similarly,  $\{z_0, x_1, x_2, \dots, x_k\}$  is also an independent set of  $G$ .

**Proof.** If  $x_0x_2 \in E(G)$ , then  $C' = x_0x_2Cy_2x_0$  is a cycle longer than  $C$ , a contradiction.

If  $x_0x_i \in E(G)$  for  $i \neq 2$ , then  $C' = x_ix_0\bar{P}_2y_0P_iy_i\bar{C}x_i$  is a cycle longer than  $C$ , a contradiction.

If  $x_ix_j \in E(G)$  for  $1 \leq i < j \leq k$ , then  $C' = x_iCy_j\bar{P}_jy_0P_iy_i\bar{C}x_jx_i$  is a cycle longer than  $C$ , a contradiction. Therefore,  $\{x_0, x_1, x_2, \dots, x_k\}$  is an independent set of  $G$ . Similarly,  $\{z_0, x_1, x_2, \dots, x_k\}$  is also an independent set of  $G$ .  $\square$

By Claim 1 and the hypothesis of Theorem 6, there exist at least two vertices in  $\{x_0, x_1, x_2, \dots, x_k\}$  with implicit degree sum more than or equal to  $n$  and there exist at least two vertices in  $\{z_0, x_1, x_2, \dots, x_k\}$  with implicit degree sum more than or equal to  $n$ .

**Case 1.** There exist some  $i$  and  $j$  with  $1 \leq i < j \leq k$  such that  $id(x_i) + id(x_j) \geq n$ .

Set  $P = x_iCy_j\bar{P}_jy_0P_iy_i\bar{C}x_j$ . Clearly,  $|V(P)| > |V(C)|$ .

**Claim 2.**  $N_P(x_i)^- \neq N_P(x_i) \cup \{x_i\} - \{L_P(x_i)\}$  and  $N_P(x_j)^+ \neq N_P(x_j) \cup \{x_j\} - \{L_P(x_j)\}$ .

**Proof.** By the choices of  $C$ , we know  $x_iy_i \in E(G)$ ,  $x_iy_0 \notin E(G)$  and  $x_jy_j \in E(G)$ ,  $x_jy_0 \notin E(G)$ . Since  $y_0$  is before  $y_i$  on the path  $P$ ,  $N_P(x_i)^- \neq N_P(x_i) \cup \{x_i\} - \{L_P(x_i)\}$ . Since  $y_j$  is before  $y_0$  on the path  $P$ ,  $N_P(x_j)^+ \neq N_P(x_j) \cup \{x_j\} - \{L_P(x_j)\}$ .  $\square$

**Claim 3.** There exists a path  $W(w_1, w_2)$  such that (i)  $V(P) \subseteq V(W)$ , and (ii)  $d(w_1) \geq id(x_i)$  and  $d(w_2) \geq id(x_j)$ .

**Proof.** For convenience, set  $P = u_1u_2 \dots u_p$  with  $u_1 = x_i$  and  $u_p = x_j$ . By the choice of  $C$ , we have  $N_{G-V(P)}(x_i) \cap N_{G-V(P)}(x_j) = \emptyset$ .

**Case 1.1.** There is a vertex  $u \in N_{G-V(P)}(x_i) \cup \{x_i\}$  such that  $d(u) \geq id(x_i)$ .

If there is a vertex  $v \in N_{G-V(P)}(x_j) \cup \{x_j\}$  such that  $d(v) \geq id(x_j)$ , then set

$$W(w_1, w_2) = ux_iPx_jv, \text{ where } w_1 = u \text{ and } w_2 = v.$$

If  $d(v) < id(x_j)$  for each vertex  $v \in N_{G-V(P)}(x_j) \cup \{x_j\}$ , then by Claim 2 and Lemma 3, there exists some vertex  $u_l \in N_P(x_j)^+$  such that  $d(u_l) \geq id(x_j)$ . Set

$$W(w_1, w_2) = ux_i P u_{l-1} x_j \bar{P} u_l, \text{ where } w_1 = u \text{ and } w_2 = u_l.$$

**Case 1.2.**  $d(u) < id(x_i)$  for each vertex  $u \in N_{G-V(P)}(x_i) \cup \{x_i\}$ .

By Claim 2 and Lemma 3, there is some vertex  $u_h \in (N_P(x_i))^-$  such that  $d(u_h) \geq id(x_i)$ . If there is a vertex  $v \in N_{G-V(P)}(x_j) \cup \{x_j\}$  such that  $d(v) \geq id(x_j)$ , then set

$$W(w_1, w_2) = u_h \bar{P} x_i u_{h+1} P x_j v, \text{ where } w_1 = u_h \text{ and } w_2 = v.$$

Next, we suppose  $d(v) < id(x_j)$  for any vertex  $v \in N_{G-V(P)}(x_j) \cup \{x_j\}$ . If  $h+1 \leq l_P(x_j)$  (where  $h$  is the index of  $u_h$  on the path  $P$ ), then there is some  $u_l \in (N_P(x_j))^+$  such that  $d(u_l) \geq id(x_j)$  by Claim 2 and Lemma 3. Set

$$W(w_1, w_2) = u_h \bar{P} x_i u_{h+1} P u_{l-1} x_j \bar{P} u_l, \text{ where } w_1 = u_h \text{ and } w_2 = u_l.$$

If  $h+1 > l_P(x_j)$ , set

$$A = \{u_s : u_{s+1} \in N_P(x_j) \text{ and } s < h\},$$

$$B = \{u_s : u_{s-1} \in N_P(x_j) \text{ and } s > h+1\}, \text{ and}$$

$$C = \{u_s : u_{s+1} \in N_{P_i, j}(x_j) \text{ } s \geq h+1 \text{ and is as small as possible}\}.$$

Clearly,  $x_j \in B$  and  $|C| = 1$ . Then  $|A| + |B \setminus \{x_j\}| + |C| + |N_{G-V(P)}(x_j)| = d(x_j)$ ,  $u_{l_P(x_j)-1} \in A \cap N_2(x_j)$ . (Since  $d(x_j) = d_P(x_j) + d_{G-V(P)}(x_j) = |N_P(x_j)| + |N_{G-V(P)}(x_j)|$  and  $|N_P(x_j)| = |N_{V(u_1 P u_h)}(x_j)| + |N_{V(u_{h+1} P u_p)}(x_j)| = |A| + |B \setminus \{x_j\}| + |C|$ .) Since  $u_h \notin N(x_j)$ ,  $C \subseteq N_2(x_j)$ . By the definition of  $id(x_j)$ , there is some vertex  $u_l \in (A \cup B) - \{x_j\}$  such that  $d(u_l) \geq id(x_j)$ . When  $u_l \in B \setminus \{x_j\}$ , set

$$W(w_1, w_2) = u_h \bar{P} x_i u_{h+1} P u_{l-1} x_j \bar{P} u_l, \text{ where } w_1 = u_h \text{ and } w_2 = u_l.$$

When  $u_l \in A$ , set

$$W(w_1, w_2) = u_h \bar{P} u_{l+1} x_j \bar{P} u_{h+1} x_i P u_l, \text{ where } w_1 = u_h \text{ and } w_2 = u_l.$$

Now we complete the proof of Claim 3.  $\square$

By Claim 3, there exists a path  $W(w_1, w_2)$  such that  $|V(W)| \geq |V(P)| > |V(C)|$  and  $d(w_1) + d(w_2) \geq id(x_i) + id(x_j) \geq n$ . This contradicts Lemma 1.

**Case 2.** There exists some  $i$  with  $1 \leq i \leq k$  and  $i \neq 2$  such that  $id(x_0) + id(x_i) \geq n$ .

Assume, without loss of generality, that  $id(x_0) + id(x_1) \geq n$  and set  $Q = x_0 \bar{P}_2 y_0 P_1 y_1 \bar{C} x_1$ . Clearly,  $|V(Q)| > |V(C)|$ .

**Claim 4.**  $N_Q(x_0)^- \neq N_Q(x_0) \cup \{x_0\} - \{L_Q(x_0)\}$  and  $N_Q(x_1)^+ \neq N_Q(x_1) \cup \{x_1\} - \{L_Q(x_1)\}$ .

**Proof.** By Claim 1,  $x_0x_2 \notin E(G)$ . Since  $x_0y_2 \in E(G)$  and  $x_2$  is the predecessor of  $y_2$  on the path  $P$ ,  $N_Q(x_0)^- \neq N_Q(x_0) \cup \{x_0\} - \{L_Q(x_0)\}$ . Since  $x_1y_1 \in E(G)$ ,  $x_1x_2 \notin E(G)$  and  $y_1$  is before  $x_2$  on the path  $Q$ ,  $N_Q(x_1)^+ \neq N_Q(x_1) \cup \{x_1\} - \{L_Q(x_1)\}$ .  $\square$

By similar argument as in Claim 3 to the path  $Q$ , we have

**Claim 5.** There exists a path  $W'(w'_1, w'_2)$  such that (i)  $V(Q) \subseteq V(W')$ , and (ii)  $d(w'_1) \geq id(x_0)$  and  $d(w'_2) \geq id(x_1)$ .  $\square$

By Claim 5, there exists a path  $W'(w'_1, w'_2)$  such that  $|V(W')| \geq |V(Q)| > |V(C)|$  and  $d(w'_1) + d(w'_2) \geq id(x_0) + id(x_1) \geq n$ . This contradicts Lemma 1.

**Case 3.**  $id(x_0) + id(x_2) \geq n$ .

Set  $R = x_0y_2\bar{C}x_2$ . Clearly,  $|V(R)| > |V(C)|$ . For convenience, set  $R = r_1r_2 \dots r_s$ .

**Claim 6.**  $N_C(x_0) \cap N_C(x_0)^+ = \emptyset$  and  $N_R(x_2)^+ \neq N_R(x_2) \cup \{x_2\} - \{L_R(x_2)\}$ .

**Proof.** If  $x \in N_C(x_0) \cap N_C(x_0)^+$ , then  $x_0xCx^-x_0$  is a cycle longer than  $C$ , contrary to the choice of  $C$ . So  $N_C(x_0) \cap N_C(x_0)^+ = \emptyset$ . Since  $x_2y_2 \in E(G)$ ,  $x_2x_1 \notin E(G)$  and  $y_2$  is before  $x_1$  on the path  $R$ ,  $N_R(x_2)^+ \neq N_R(x_2) \cup \{x_2\} - \{L_R(x_2)\}$ .  $\square$

**Case 3.1.** There is a vertex  $x \in N_{G-V(R)}(x_0) \cup \{x_0\}$  such that  $d(x) \geq id(x_0)$ .

If there is a vertex  $y \in N_{G-V(R)}(x_2) \cup \{x_2\}$  such that  $d(y) \geq id(x_2)$ , then  $R' = xx_0Rx_2y$  is a path such that  $|V(R')| > |V(C)|$  and  $d(x) + d(y) \geq id(x_0) + id(x_2) \geq n$ . This contradicts Lemma 1.

If  $d(y) < id(x_2)$  for each vertex  $y \in N_{G-V(R)}(x_2) \cup \{x_2\}$ , then by Claim 6 and Lemma 3, there exists some vertex  $r_t \in N_R(x_2)^+$  such that  $d(r_t) \geq id(x_2)$ . Then  $R' = xx_0Rr_{t-1}x_2\bar{R}r_t$  is a path such that  $|V(R')| > |V(C)|$  and  $d(x) + d(r_t) \geq id(x_0) + id(x_2) \geq n$ . This contradicts Lemma 1.

**Case 3.2.**  $d(x) < id(x_0)$  for each vertex  $x \in N_{G-V(R)}(x_0) \cup \{x_0\}$ .

By Claim 6,  $N_C(x_0)^+ \subseteq N_2(x_0)$ . Then  $|N_{G-V(R)}(x_0)| + |N_C(x_0)^+| \geq d_R(x_0) + d_{G-V(R)}(x_0) = d(x_0)$ . If  $d_C(x_0) \geq 2$ , then by the definition of  $id(x_0)$ , there exists at least one vertex  $z \in N_C(x_0)^+ \setminus \{x_2\}$  such that  $d(z) \geq id(x_0)$ . Then  $id(z) + id(x_2) \geq id(x_0) + id(x_2) \geq n$ . By similar argument as in Case 1, we can get a contradiction.

Next, suppose  $d_C(x_0) = 1$ . Then by the definition of  $id(x_0)$ ,  $id(x_0) = \min\{d(u) : u \in N_2(x_0)\}$ . If there exists a vertex  $y \in N_2(x_0) \cap (V(G) - V(R))$ , then by similar argument to the path  $yy'x_0Rx_2$  with  $y' \in N(x_0) \cap N(x)$  as in Case 3.1, we can get a contradiction. So suppose  $N_2(x_0) \cap (V(G) - V(R)) = \emptyset$ .

By Claim 1 and the hypothesis of Theorem 6, there exist at least two vertices in  $\{z_0, x_1, x_2, \dots, x_k\}$  with implicit degree sum more than or equal to  $n$ . If there exists some  $i$  with  $1 \leq i \leq k$  and  $i \neq 1$  such that  $id(z_0) + id(x_i) \geq n$ , then by similar argument as in Case 2, we can get a contradiction. So suppose  $id(z_0) + id(x_1) \geq n$ . Then  $id(x_0) + id(x_2) + id(z_0) + id(x_1) \geq 2n$ . Therefore,  $id(z_0) + id(x_2) \geq n$  or  $id(x_0) + id(x_1) \geq n$ . Thus, by similar argument as in Case 2, we can get a contradiction.

Now the proof of Theorem 6 is completed.

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