

# Some New Difference Triangle Sets

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**ABSTRACT.** In [5] Kløve gave tables of the best bounds known on the size of optimal difference triangle sets. In this note we give examples of difference triangle sets found by computer search which improve on the upper bounds in [5]. In four cases these examples are proved to be optimal.

In [5] Kløve defined an  $(I, J)$  difference triangle set,  $\Delta$ , as a set of integers  $\{a_{ij} \mid 1 \leq i \leq I, 0 \leq j \leq J\}$  such that all the differences  $a_{ij} - a_{ik}$ ,  $1 \leq i \leq I$ ,  $0 \leq k \neq j \leq J$  are distinct. Let  $m = m(\Delta)$  be the maximum difference. Kløve defined  $M(I, J) = \min\{m(\Delta) \mid \Delta \text{ is an } (I, J) \text{ difference triangle set}\}$  and gave tables of bounds for  $M$ . In this note we present difference triangle sets found by computer search which improve these bounds.

Table 1 shows the improved upper bounds. The bounds marked with an \* are exact. Examples achieving these bounds are given in Table 2. In [5] Kløve stated  $M(3, 6) = 77$ ,  $M(4, 5) = 65$  and  $M(5, 5) = 83$ . These values apparently should have been upper bounds.

The computer program used was a modified version of the author's Golomb ruler finding program ([8], see also [9]). Golomb rulers are  $(I, J)$  difference triangle sets with  $I = 1$ . The Golomb ruler search program implements a straightforward depth first backtrack search. It attempts to build up a Golomb ruler with  $k$  marks and length  $B$  or less by picking  $0 = a_0 < a_1 < a_2 < \dots < a_k \leq B$  in order ( $a_j$  being picked at level  $j$  of the search tree). At each node of the search tree the program keeps track of which differences have been used (i.e. are formed by pairs of the elements which have already been picked) and of which integers can be added to

the current partial ruler without violating the distinct difference condition. The sons of a node at level  $j$  are formed by setting  $a_{j+1}$  to each of the elements in the level  $j$  eligibility list in turn. The search tree is pruned (i.e. the program backtracks) when too few integers remain eligible to allow completion of the ruler, when not enough small differences remain unused to allow completion of the ruler or when allowed by a symmetry condition. (Since the mirror image of a Golomb ruler is also a Golomb ruler we may require for example that  $a_{k/2} < B/2$ .)

We modified the program to search for  $(I, J)$  difference triangle sets with  $I > 1$  as follows. Such difference triangle sets may be thought of as sets of  $I$  Golomb rulers of  $J + 1$  marks each, such that no difference occurs in more than one ruler. We may assume  $a_{10} = a_{20} = \dots = a_{I0} = 0$ . The modified program attempts to build a  $(I, J)$  difference triangle set,  $\Delta$ , with  $m(\Delta) \leq B$  (for some fixed bound  $B$ ) by picking  $a_{11}, a_{12}, \dots, a_{1J}, a_{21}, a_{22}, \dots, a_{2J}, \dots, a_{I1}, a_{I2}, \dots, a_{IJ}$  in order ( $0 < a_{h1} < a_{h2} < \dots < a_{hJ} \leq B, 1 \leq h \leq I$ ). While building up the  $h$ th ruler the program, at each node of the search tree, keeps track of which differences have been used (in the current and earlier rulers) and of which integers can be added to the current partial ruler without violating the distinct difference condition. The program backtracks when too few integers remain eligible to allow completion of the current ruler, when not enough small differences remain unused to allow completion of the current and any remaining rulers or when allowed by symmetry. Since we may permute the order of the rulers in a difference triangle set as well as reversing rulers we may require  $a_{11} > a_{21} > \dots > a_{I1}$  as well as  $a_{hJ/2} < B/2$  (or if  $J$  is odd  $a_{hJ/2} + a_{hJ/2+1} < B$ )  $1 \leq h \leq I$ .

In most cases it was not feasible to perform a complete search. To handle these cases we modified the program to allow it to search a fraction of the complete search tree by arbitrarily pruning many nodes in the search tree. The complete program amounts to about 150 lines of VS Fortran code. The program was run on an IBM 3090 mainframe. Each partial search was run for at least 5000 seconds of cpu time. For the four cases where a complete search was feasible we give all the examples which achieve the bound. Otherwise we give a single example even if several were found.

This work was done in 1990-1991. Since then some other work has improved on Kløve's bounds ([2], [3] and [1]). However, except for  $M(2, 7) = 70$  the results in Table 1 remain the best known. Chee and Colbourn [1] independently showed by computer search that  $M(2, 7) = 70$ . Their program ran for a week on a network of 30 machines (and did not show uniqueness). The author's program ran for about 3000 seconds (with  $B = 70$ ) on a 3090.

**Table 1**

<i>J\I</i>	2	3	4	5	6	7	8	9	10	11	12	13	14	15
4								93		116	127	138	150	160
5			*64	81	100	118	137	157	170	191	209			
6			*72	98	125	153	178	204	232	264				
7			*70	104	144	182	218	256	302	334				
8			*94	145	199	251	302	353	405					
9			126	196	269	329	405							
10			254											

**Table 2**

I   J      Difference Triangle Set

2	7	0	6	11	18	28	37	62	70					
		0	1	4	24	40	54	67	69					
2	8	0	10	15	29	45	51	77	85	94				
		0	1	3	24	28	61	74	81	92				
2	8	0	14	16	22	34	43	62	79	94				
		0	4	7	37	42	68	81	91	92				
2	9	0	7	10	22	52	68	99	100	118	123			
		0	4	17	38	44	73	87	98	124	126			
3	6	0	5	19	21	50	57	70						
		0	4	10	32	43	58	66						
		0	3	12	30	47	71	72						
3	6	0	7	22	30	62	66	71						
		0	3	19	20	48	54	72						
		0	2	13	27	39	60	70						
3	6	0	10	15	27	61	70	72						
		0	7	13	33	49	63	71						
		0	1	25	29	48	66	69						
3	6	0	13	16	31	56	67	68						
		0	7	21	26	48	65	71						
		0	2	10	30	34	63	72						
3	7	0	17	31	36	57	65	102	103					
		0	4	16	27	55	79	97	104					
		0	2	22	32	35	76	82	91					
3	8	0	6	40	50	54	110	111	126	135				
		0	3	21	38	66	103	115	134	145				
		0	2	22	29	55	91	130	138	143				

3	9	0	3	20	26	86	95	117	162	173	194
3	10	0	16	43	90	94	116	154	231	236	242
3	10	0	16	43	90	94	116	154	231	236	242
3	10	0	16	43	90	94	116	154	231	236	242
4	5	0	12	13	31	58	60	82	92	92	92
4	6	0	5	6	25	59	72	89	107	170	200
4	7	0	21	43	59	69	125	143	144	142	142
4	8	0	21	62	73	77	149	159	193	199	192
4	9	0	20	75	76	118	140	157	207	234	269
5	5	0	16	23	43	68	81	113	119	219	222
5	6	0	18	28	48	92	111	118	179	179	181
5	7	0	24	43	69	103	149	179	179	181	181
0	0	23	65	81	82	120	170	182	182	182	192
0	0	0	35	48	92	156	166	177	177	177	196
0	0	0	5	41	61	114	128	165	180	180	196
0	3	31	99	146	146	153	153	171	171	171	196

5	8	0	26	36	74	96	125	175	220	244
5	9	0	32	82	119	152	163	289	306	310
5	9	0	32	82	119	152	163	289	306	324
6	5	0	14	35	54	92	96	157	214	247
6	6	0	18	45	59	123	146	153	214	251
6	7	0	2	18	67	70	97	100	107	125
6	7	0	60	61	89	111	163	197	216	244
6	8	0	38	46	90	110	222	250	287	300
6	9	0	28	49	83	102	144	269	286	336
7	0	0	26	70	103	190	201	316	353	375
7	0	0	10	46	62	64	199	205	212	335
7	0	0	8	38	86	119	133	233	265	306
7	0	0	4	24	84	113	129	176	247	348
7	0	0	0	28	49	83	102	144	269	382
7	0	0	0	26	70	103	190	201	316	383
7	0	0	0	10	46	62	64	199	205	405
7	0	0	0	8	38	86	119	133	233	387
7	0	0	0	4	24	84	113	129	176	383
7	0	0	0	0	28	49	83	102	144	395

7	5	0	23	28	65	85	115
7	6	0	13	32	66	91	113
7	7	0	11	46	60	104	114
7	8	0	8	39	75	102	118
7	9	0	6	21	76	105	117
7	10	0	4	7	52	90	116
7	11	0	1	18	74	98	107
7	12	0	30	36	84	106	172
7	13	0	27	47	82	108	160
7	14	0	24	58	77	150	163
7	15	0	18	39	56	135	146
7	16	0	10	50	75	124	168
7	17	0	8	12	41	110	153
7	18	0	5	14	37	109	169
7	19	0	37	71	82	118	207
7	20	0	23	77	108	127	156
7	21	0	16	41	96	111	203
7	22	0	9	51	68	73	199
7	23	0	4	12	30	105	132
7	24	0	3	63	87	154	216
7	25	0	2	99	112	119	177
7	26	0	0	31	68	88	147
7	27	0	0	28	49	78	130
7	28	0	0	27	67	92	122
7	29	0	0	24	80	121	174
7	30	0	0	11	71	146	154
7	31	0	0	10	15	58	148
7	32	0	0	2	35	119	163
8	5	0	28	48	67	111	129
8	6	0	26	34	61	99	136
8	7	0	17	24	70	121	124
8	8	0	6	31	80	95	125
8	9	0	5	21	76	108	117
8	10	0	4	47	60	126	137
8	11	0	2	42	52	120	134
8	12	0	1	59	116	128	

8	6	0	33	58	87	155	191	202
8	7	0	47	60	68	142	257	259
8	8	0	44	62	90	162	195	258
9	4	0	27	34	86	91	371	405
9	5	0	45	62	83	125	157	
0	8	0	23	25	64	111	154	
0	9	0	20	34	85	137	142	
0	10	0	18	49	58	118	145	
0	11	0	15	19	92	121	147	
0	12	0	12	71	78	150	153	
0	13	0	11	44	81	105	135	
0	14	0	10	16	114	136	149	
0	15	0	8	36	84	151	152	
0	16	0	0	0	3	11	69	92

9	6	0	38	49	112	148	225	232
9	7	0	33	58	137	141	215	320
9	7	0	27	71	84	118	238	274
10	5	0	22	73	134	196	206	283
10	5	0	16	54	86	175	218	294
10	5	0	15	23	65	193	242	245
10	5	0	7	35	116	162	255	256
10	5	0	6	17	37	231	276	306
10	5	0	5	69	122	151	282	303
10	5	0	2	14	99	138	268	286
10	6	0	57	66	73	152	160	
10	6	0	37	48	107	156	158	
10	6	0	34	39	97	123	164	
10	6	0	29	72	93	155	168	
10	6	0	19	50	90	118	170	
10	6	0	18	54	74	165	166	
10	6	0	10	12	35	81	136	169
10	6	0	10	14	116	141	163	
10	6	0	6	38	82	143	167	
10	6	0	3	30	45	145	162	
10	6	0	29	49	113	149	221	244
10	6	0	28	59	83	193	239	240
10	6	0	18	44	106	169	243	260
10	6	0	14	66	85	141	184	245
10	6	0	12	15	60	167	206	264
10	6	0	8	73	89	166	176	217
10	6	0	7	57	124	220	253	262
10	6	0	5	27	119	188	228	241
10	6	0	4	34	102	182	207	261
10	6	0	2	37	123	224	235	256

11	4	0	38	53	88	112
		0	27	40	89	98
		0	25	42	73	107
		0	23	52	55	116
		0	18	37	84	104
		0	11	54	110	114
		0	10	36	87	115
11	5	0	0	0	0	0
		0	1	46	76	109
		0	31	41	91	180
		0	30	42	97	156
		0	21	57	58	185
		0	18	25	53	136
		0	14	65	84	159
		0	13	33	137	160
		0	0	0	11	106
		0	0	0	74	123
		0	0	4	73	100
		0	3	46	128	168
		0	2	26	87	134
12	4	0	55	59	121	127
		0	31	34	94	126
		0	27	38	111	116
		0	20	21	96	122
		0	19	42	90	119
		0	18	53	103	117
		0	16	52	97	109
		0	15	58	98	123
		0	13	46	74	125
		0	10	17	54	124
		2	49	88	118	115

12	5	0	43	97	105	180	209
0	0	39	50	63	142	201	
0	0	36	66	123	179	194	
0	0	0	32	41	102	157	197
0	0	0	23	33	129	173	207
0	0	0	16	38	98	206	208
0	0	0	14	49	107	135	183
13	4	0	0	0	0	0	
0	0	0	1	53	100	177	
0	0	0	29	59	117	130	
0	0	0	22	30	120	138	
0	0	0	15	62	103	137	
0	0	0	14	56	80	135	
0	0	0	12	50	96	123	
0	0	0	10	45	102	128	
0	0	0	7	51	100	119	
0	0	0	6	17	95	131	
0	0	0	4	32	85	133	
0	0	0	3	63	94	127	
0	0	0	2	23	99	136	
0	0	0	1	40	105	110	
14	4	0	0	0	0	0	
0	0	31	35	126	134		
0	0	25	65	97	150		
0	0	24	27	101	139		
0	0	21	57	79	130		
0	0	20	69	128	144		
0	0	18	28	84	145		
0	0	15	67	137	148		
0	0	12	41	88	143		
0	0	13	30	113	136		
0	0	13	30	113	136		
0	0	6	45	116	135		
0	0	5	14	68	146		
0	0	2	48	98	140		
0	1	34	94	120			

15	4	0	57	74	156	157
		0	39	53	98	149
		0	29	60	115	150
		0	28	70	141	159
		0	27	32	105	143
		0	24	61	67	130
		0	23	33	152	155
		0	21	47	112	148
		0	20	66	124	154
		0	16	50	118	158
		0	13	25	120	139
		0	11	52	144	146
		0	8	64	136	145
		0	7	22	84	160
		0	4	48	97	151

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