

# Orientable Open Domination of Graphs

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## ABSTRACT

An open dominating set for a digraph  $D$  is a set  $S$  of vertices of  $D$  such that every vertex of  $D$  is adjacent from some vertex of  $S$ . The cardinality of a minimum open dominating set for  $D$  is the open domination number  $\rho_1(D)$  of  $D$ . The lower orientable open domination number  $\text{dom}_1(G)$  of a graph  $G$  is the minimum open domination number among all orientations of  $G$ . Similarly, the upper orientable open domination number  $\text{DOM}_1(G)$  of  $G$  is the maximum such open domination number.

For a connected graph  $G$ , it is shown that  $\text{dom}_1(G)$  and  $\text{DOM}_1(G)$  exist if and only if  $G$  is not a tree. A discussion of the upper orientable open domination number of complete graphs is given. It is shown that for each integer  $c$  with  $\text{dom}_1(K_n) \leq c \leq \text{DOM}_1(K_n)$ , there exists an orientation  $D$  of  $K_n$  such that  $\rho_1(D) = c$ .

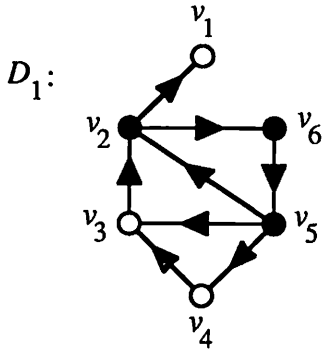
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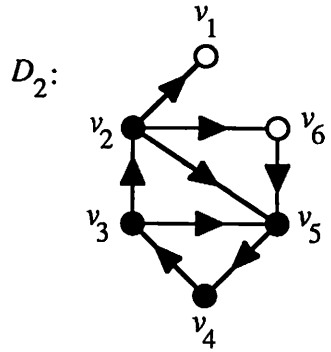
## 1. Orientable Open Domination

A vertex  $v$  in a graph  $G$  is said to *1-step dominate* or *openly dominate* each of its neighbors. A set  $S$  of vertices of  $G$  is an *open dominating set* of  $G$  if every vertex of  $G$  is openly dominated by some vertex of  $S$ . The minimum cardinality among the open dominating sets of  $G$  is called the *open domination number* of  $G$  and is denoted by  $\rho_1(G)$ .

Analogous definitions can be made for digraphs. In particular, for a digraph  $D$ , a vertex  $v$  *openly dominates* all vertices  $w$  with  $(v, w) \in E(D)$ . The *open domination number*  $\rho_1(D)$  of  $D$  is the minimum cardinality among the open dominating sets of  $D$ . Although for a graph  $G$  without isolated vertices, the open domination number  $\rho_1(G)$  always exists, such is not the case for digraphs. However, for a digraph  $D$ , a necessary and sufficient condition for  $\rho_1(D)$  to exist is that  $\text{id } v \geq 1$  for every vertex  $v$  of  $D$ . This condition is satisfied for the digraphs  $D_1$  and  $D_2$  of Figure 1. In  $D_1$ , the vertex  $v_1$  is uniquely openly dominated by  $v_2$ , the vertex  $v_4$  is uniquely openly dominated by  $v_5$ , and  $v_5$  is uniquely openly dominated by  $v_6$ . Hence if  $S$  is an open dominating set of  $D_1$ , then  $\{v_2, v_5, v_6\} \subseteq S$ . Since  $S_1 = \{v_2, v_5, v_6\}$  is itself an open dominating set of  $D_1$ , it follows that  $S_1$  is a (in fact, the unique) minimum open dominating set of  $D_1$ . Thus  $\rho_1(D_1) = 3$ . Similarly,  $S_2 = \{v_2, v_3, v_4, v_5\}$  is the unique minimum open dominating set of  $D_2$  and so  $\rho_1(D_2) = 4$ . Since  $D_1$  and  $D_2$  are orientations of the same graph and  $\rho_1(D_1) \neq \rho_1(D_2)$ , this suggests some definitions.



$$S_1 = \{v_2, v_5, v_6\}$$



$$S_2 = \{v_2, v_3, v_4, v_5\}$$

Figure 1

For a graph  $G$ , we say that  $D$  is a *valid orientation* of  $G$  if every vertex of  $D$  has positive indegree. Let  $D_1, D_2, \dots, D_k$  be the distinct valid orientations of a graph  $G$ . We define the *lower* and *upper orientable open domination numbers* of  $G$ , respectively, as

$$\text{dom}_1(G) = \min\{\rho_1(D_i) \mid 1 \leq i \leq k\} \text{ and } \text{DOM}_1(G) = \max\{\rho_1(D_i) \mid 1 \leq i \leq k\}$$

These concepts were defined and investigated for ordinary domination in digraphs in [1]. In order to present a necessary and sufficient condition for these parameters to be defined for a graph  $G$ , we recall a theorem of Robbins [3].

**Theorem (Robbins)** A graph  $G$  has a strong orientation if and only if  $G$  is 2-edge-connected (connected and has no bridges).

**Theorem 1** Let  $G$  be a connected graph. Then  $\text{dom}_1(G)$  and  $\text{DOM}_1(G)$  exist if and only if  $G$  is not a tree.

**Proof** First assume that  $\text{dom}_1(G)$  and  $\text{DOM}_1(G)$  are defined and suppose, to the contrary, that  $G$  is a tree of order  $n$ . Let  $D$  be any valid orientation of  $G$ . Since  $G$  is a tree, it follows that the sum of the indegrees of the vertices of  $D$  is  $\sum_{v \in V(D)} \text{id } v = |E(D)| = |E(G)| = n - 1$ . However, since every vertex has indegree at least 1, it follows that  $\sum_{v \in V(D)} \text{id } v \geq n$ , producing a contradiction.

Next, suppose that  $G$  is not a tree. We show that there exists a valid orientation  $D$  of  $G$ . Since  $G$  is not a tree, we know that  $G$  contains a cyclic block. Let  $B_1, B_2, \dots, B_k$  denote the cyclic blocks of  $G$ . By Robbins' Theorem, each block  $B_i, 1 \leq i \leq k$ , has a strong orientation. For the desired orientation  $D$ , begin by producing a strong orientation of each block  $B_i, 1 \leq i \leq k$ . Then  $\text{id } v \geq 1$  for every vertex  $v$  of the blocks  $B_1, B_2, \dots, B_k$ . Next, for any shortest path  $P$  between two cyclic blocks  $B_i$  and  $B_j$  ( $i \neq j$ ), orient the edges of  $P$  from one block to the other (say, from  $B_i$  to  $B_j$ ). Then for each vertex  $v$  of  $P$ , we have  $\text{id } v \geq 1$ . Finally, let  $V$  be the set  $\bigcup_{i=1}^k V(B_i)$ , along with the vertices of all shortest paths between two cyclic blocks, and let  $u$  be any vertex of  $G$  not belonging to  $V$ . Then there exists a shortest path from  $u$  to some vertex, say  $v$ , of  $V$ . Direct the edges of the path from  $v$  to  $u$ . The resulting orientation  $D$  has minimum indegree at least 1. Therefore  $\rho_1(D)$  exists, implying that  $\text{dom}_1(G)$  and  $\text{DOM}_1(G)$  exist.  $\square$

Some simple bounds for these parameters are presented next.

**Theorem 2** For every connected graph  $G$  of order  $n \geq 3$  that is not a tree,

$$3 \leq \text{dom}_1(G) \leq \text{DOM}_1(G) \leq n.$$

**Proof** Clearly,  $\text{DOM}_1(G) \leq n$ . Let  $D$  be a valid orientation of  $G$  and let  $S$  be a minimum open dominating set of  $D$ . For  $u \in S$ , there exists some vertex  $v \in S$  that openly dominates  $u$ . So  $(v, u) \in E(D)$ . Also the vertex  $v$  must

be openly dominated by some vertex of  $S$ . Further, this vertex cannot be  $u$  since  $(u, v) \notin E(D)$ . Hence there exists a vertex  $w \in S$  with  $w \neq u, v$  that openly dominates  $v$ . Therefore  $\{u, v, w\} \subseteq S$ , implying  $\rho_1(D) = |S| \geq 3$ . We now know  $\rho_1(D) \geq 3$  for any valid orientation  $D$  of  $G$ . Consequently,  $\text{dom}_1(G) \geq 3$ .  $\square$

The following result shows that the upper orientable domination number of a graph attains the upper bound only in one special case. The proof of this result uses the following theorem of Hall [2].

**Theorem (Hall)** A collection  $S_1, S_2, \dots, S_n, n \geq 1$ , of finite nonempty sets has a system of distinct representatives if and only if the union of any  $k$  of these sets contains at least  $k$  elements, for each  $k$  such that  $1 \leq k \leq n$ .

**Theorem 3** For a connected graph  $G$  of order  $n \geq 3$ ,  $\text{DOM}_1(G) = n$  if and only if  $G = C_n$ .

**Proof** First assume that  $G = C_n$ . Then, we orient the edges so that the corresponding digraph  $D$  is a directed cycle. It follows that every vertex is openly dominated only by the preceding vertex along the cycle. Therefore every vertex must belong to the open dominating set, implying that  $\rho_1(D) = n$ . Consequently,  $\text{DOM}_1(G) = n$ .

For the converse, assume that  $\text{DOM}_1(G) = n$ . Let  $D$  be an orientation of  $G$  such that  $\rho_1(D) = \text{DOM}_1(G) = n$ . Assume that  $V(D) = \{v_1, v_2, \dots, v_n\}$ . Define a digraph  $D'$  by  $V(D') = \{v_1, v_2, \dots, v_n, v'_1, v'_2, \dots, v'_n\}$  and  $(u, v) \in E(D')$  if and only if

- (1)  $u = v_i$  for some  $i \in \{1, 2, \dots, n\}$ ,
- (2)  $v = v'_j$  for some  $j \in \{1, 2, \dots, n\}$ , and
- (3)  $(v_i, v_j) \in E(D)$ .

Let  $S_j = \{v_i \mid (v_i, v'_j) \in E(D')\}$  for  $j = 1, 2, \dots, n$ . Thus  $S_j$  is the set of vertices that openly dominate the vertex  $v_j$  in  $D$ . We claim that  $S_1, S_2, \dots, S_n$  has a system of distinct representatives. Now let  $J \subseteq \{1, 2, \dots, n\}$  such that  $|J| = k$ . Let  $S = \bigcup_{j \in J} S_j$  and suppose, to the contrary, that  $|S| < |J|$ . Let  $W = \{v_j \mid j \in J\}$  and let  $W' = \{v'_1, v'_2, \dots, v'_n\} - W$ . In  $D$  every vertex  $v_j$ , where  $j \in J$ , is openly dominated by some vertex of  $S$ . Now for each  $v'_\ell \in W'$ , we know since  $\text{id}_D v_\ell \geq 1$ , that there exists  $(v_i, v_\ell) \in E(D)$  and hence  $(v_i, v'_\ell) \in E(D')$ . Let  $v_{i_\ell}$  be such a vertex for each  $v'_\ell \in W'$  and let  $S' = \{v_{i_\ell} \mid v'_\ell \in W'\}$ . Then  $|S'| \leq |W'| = n - k$ . Now, in  $D$ , every vertex  $v_j$ , where  $j \notin J$ , is openly dominated by some vertex of  $S'$ . Consequently, the set  $S \cup S'$  is an open dominating set of  $D$ . However  $|S \cup S'| \leq |S| + |S'| < k + n - k = n$ . This implies that  $\rho_1(D) < n$ , producing a contradiction. Hence  $|S| \geq |J|$ , implying that  $S_1, S_2, \dots, S_n$  has a system of distinct representatives. It follows that  $D'$  has an independent set of arcs. This independent set of arcs corresponds to a disjoint union of directed cycles, say  $D_1, D_2, \dots, D_m$ , in  $D$ .

Suppose, to the contrary, that  $\langle V(D_i) \rangle$  is not a directed cycle for some  $i = 1, 2, \dots, m$ . Then there exist vertices  $v_j$  and  $v_k$  in  $V(D_i)$  such that  $(v_j, v_k) \in E(D)$  but  $v_k$  does not follow  $v_j$  on the directed cycle  $D_i$ . Let  $v_\ell$  be the vertex preceding  $v_k$  on the directed cycle  $D_i$ . Then  $V(D_i) - \{v_\ell\}$  is an open dominating set of  $\langle V(D_i) \rangle$ . But this would imply that  $\rho_1(D) < n$ . Hence  $\langle V(D_i) \rangle$  is a directed cycle for each  $i = 1, 2, \dots, m$ .

Finally we show that  $m = 1$ . If  $m \geq 2$ , then since the underlying graph  $G$  is connected, there exist vertices  $v_i$  and  $v_j$  such that  $(v_i, v_j) \in E(D)$  where  $v_i \in V(D_i)$  and  $v_j \in V(D_j)$ ,  $i \neq j$ . Let  $v_k$  be the vertex preceding  $v_j$  on the directed cycle  $D_j$ . Then  $V(D_i) \cup V(D_j) - \{v_k\}$  is an open dominating set of  $\langle V(D_i) \cup V(D_j) \rangle$ . Again, this would imply that  $\rho_1(D) < n$ . Hence  $m = 1$  and  $D = \langle V(D_1) \rangle$ , a directed cycle of length  $n$ . That is,  $G = C_n$ .  $\square$

We next study a class of graphs, the difference of whose lower and upper orientable domination numbers is arbitrarily large. In particular, for  $n \geq 4$ , let  $G_n = K_1 + P_{n-1}$  (see Figure 2).

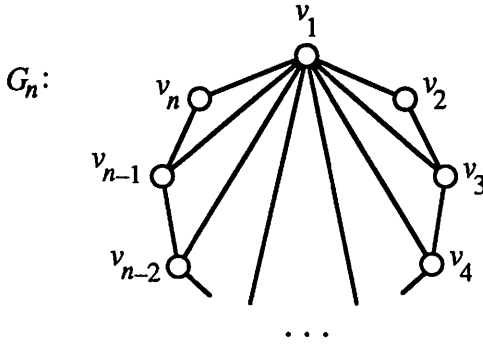
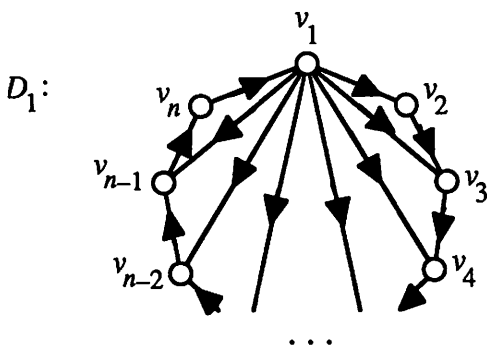


Figure 2

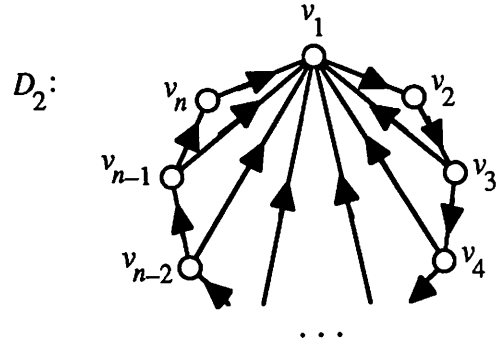
It can be shown that  $\text{dom}_1(G_n) = 3$  and  $\text{DOM}_1(G_n) = n - 1$  with the aid of the orientations  $D_1$  and  $D_2$  of  $G_n$  shown in Figure 3.

For  $D_1$ , observe that  $\text{id } v_2 = 1$  and thus  $v_1$  must belong to every open dominating set of  $D_1$ . Similarly, since  $\text{id } v_n = 1$  and  $\text{id } v_1 = 1$ , it follows that  $v_{n-1}$  and  $v_n$  must also belong to each open dominating set of  $D_1$ . Further,  $S_1 = \{v_1, v_{n-1}, v_n\}$  is an open dominating set of  $D_1$ . Thus  $\rho_1(D_1) = |S_1| = 3$ . Also  $3 \leq \text{dom}_1(G_n) \leq \rho_1(D_1) = 3$ , which implies that  $\text{dom}_1(G_n) = 3$ .

In a similar manner, observe that  $\text{id}_{D_2} v_i = 1$  for  $i = 2, 3, \dots, n$ . Thus  $v_1, v_2, \dots, v_{n-1}$  must belong to every open dominating set. Since  $S_2 = \{v_1, v_2, \dots, v_{n-1}\}$  is an open dominating set, it follows that  $\text{DOM}_1(G_n) \geq \rho_1(D_2) = |S_2| = n - 1$ . Further, by Theorem 3,  $\text{DOM}_1(G_n) \neq n$ . Thus  $\text{DOM}_1(G_n) = n - 1$ .



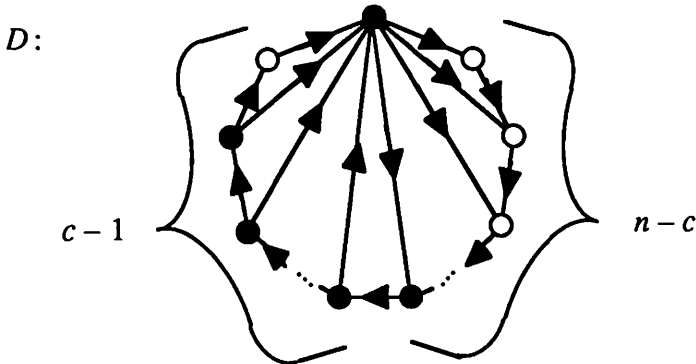
$$S_1 = \{v_1, v_{n-1}, v_n\}$$



$$S_2 = \{v_1, v_2, \dots, v_{n-1}\}$$

Figure 3

This class of graphs also has the property that for each integer  $c$ ,  $3 \leq c \leq n - 1$ , there exists an orientation of  $G_n$ , resulting in a digraph  $D$ , such that  $\rho_1(D) = c$ . (See Figure 4.)



$$\rho_1(D) = c$$

Figure 4



## 2. Upper Orientable Open Domination Numbers of Complete Graphs

In this section we investigate the growth of the function  $\text{DOM}_1(K_n)$ . First, we show that  $\text{DOM}_1(K_n)$  is a nondecreasing function.

**Lemma 4** For all  $n \geq 3$ ,  $\text{DOM}_1(K_n) \leq \text{DOM}_1(K_{n+1})$ .

**Proof** Choose a valid orientation  $D_n$  of  $K_n$  such that  $\rho_1(D_n) = \text{DOM}_1(K_n)$ . Next define an orientation  $D_{n+1}$  of  $K_{n+1}$  by adding a vertex  $w$  to  $D_n$  and adding arcs  $(u, w)$  for every  $u \in V(D_n)$ . Then  $D_{n+1}$  is a valid orientation, and

$$\text{DOM}_1(K_{n+1}) \geq \rho_1(D_{n+1}) = \rho_1(D_n) = \text{DOM}_1(K_n). \quad \square$$

**Corollary 5** For  $3 \leq m \leq n$ ,  $\text{DOM}_1(K_m) \leq \text{DOM}_1(K_n)$ .

In the next two lemmas, we investigate how slowly the function  $\text{DOM}_1(K_n)$  grows. Of course, every orientation of  $K_n$  is a tournament of order  $n$ .

**Lemma 6** If  $\text{DOM}_1(K_{n+1}) > \text{DOM}_1(K_n)$ , then

- (a) each tournament  $T_{n+1}$  with  $\rho_1(T_{n+1}) = \text{DOM}_1(K_{n+1})$  is strong, and
- (b)  $\text{DOM}_1(K_{n+1}) = \text{DOM}_1(K_n) + 1$ .

**Proof of (a)** Let  $T_{n+1}$  be any valid tournament of order  $n + 1$  such that  $\rho_1(T_{n+1}) = \text{DOM}_1(K_{n+1})$ . Suppose, to the contrary, that  $T_{n+1}$  is not strong. Then  $V(T_{n+1})$  can be partitioned into sets  $S_1, S_2, \dots, S_k$  ( $k \geq 2$ ) such that (1) each  $\langle S_i \rangle$  is a strong subdigraph and is maximal with respect to the property of being strong, that is, the subdigraphs  $\langle S_1 \rangle, \langle S_2 \rangle, \dots, \langle S_k \rangle$  are the strong components of  $T_{n+1}$ , and (2) for every vertex  $u$  of  $S_1$  and for every vertex  $v$

of  $S_i, i \geq 2$ , the arc  $(u, v)$  belongs to  $T_{n+1}$ . It is clear that  $S_1$  is an open dominating set of  $T_{n+1}$ . Further, a minimum open dominating set of  $\langle S_1 \rangle$  is also a minimum open dominating set of  $T_{n+1}$ . So  $\text{DOM}_1(K_{n+1}) = \rho_1(T_{n+1}) = \rho_1(\langle S_1 \rangle) \leq \text{DOM}_1(K |_{S_1}) \leq \text{DOM}_1(K_n) < \text{DOM}_1(K_{n+1})$ , producing a contradiction. Thus  $T_{n+1}$  is strong.

**Proof of (b)** It suffices to show that  $\text{DOM}_1(K_{n+1}) \leq 1 + \text{DOM}_1(K_n)$ . Let  $T_{n+1}$  be a valid tournament of order  $n + 1$  such that  $\rho_1(T_{n+1}) = \text{DOM}_1(K_{n+1})$ . Choose a vertex  $v$  of  $T_{n+1}$  with  $\text{od } v \geq 1$ . Let  $S = V(T_{n+1}) - N^+(v)$ , where  $N^+(v)$  represents the out-neighborhood of  $v$ . Then  $v$  openly dominates all vertices belonging to  $N^+(v)$ . Also, since  $\text{id } v \geq 1, S \neq \emptyset$ . Since  $\text{od } v \geq 1, |S| \leq n$ . Now it follows that

$$\text{DOM}_1(K_{n+1}) = \rho_1(T_{n+1}) \leq 1 + \rho_1(\langle S \rangle) \leq 1 + \text{DOM}_1(K_n). \quad \square$$

**Theorem 7** If  $\text{DOM}_1(K_n) = m$  and  $\text{DOM}_1(K_{n+1}) = m + 1$ , then  $\text{DOM}_1(K_i) = m + 1$  for  $n + 1 \leq i \leq 2n + 2$ .

**Proof** It suffices to show that  $\text{DOM}_1(K_{2n+2}) \leq m + 1$ . Let  $T$  be a valid tournament of order  $2n + 2$ . Then  $T$  has a vertex  $v$  such that  $\text{od } v \geq n + 1$ . Let  $S = V(T) - (N^+(v) \cup \{v\})$ . Then  $|S| \leq n$ . As in the proof of Lemma 6(b), we have

$$\rho_1(T) \leq 1 + \rho_1(\langle S \rangle) \leq 1 + \text{DOM}_1(K_n) = 1 + m.$$

Since this is true for any such tournament  $T$ , it follows that  $\text{DOM}_1(K_{2n+2}) \leq m + 1$ .  $\square$

From our previous results, we know that  $\text{DOM}_1(K_n)$  is a nondecreasing function and any increase in functional values is a step increment

of 1. However, our results so far have not shown that  $\text{DOM}_1(K_n)$  increases at all. The next result shows that, in fact, the function increases without bound.

**Theorem 8** The function  $\text{DOM}_1(K_n)$  is unbounded as  $n \rightarrow \infty$ .

**Proof** We proceed by a counting argument. First, let  $T_n$  denote the number of labeled tournaments of order  $n$ . Since  $K_n$  has  $\binom{n}{2}$  edges, each of which can be oriented in one of two directions,  $T_n = 2^{\binom{n}{2}}$ .

Next, let  $T_n^v$  denote the number of valid labeled tournaments of order  $n$ . We claim that  $T_n^v = T_n - nT_{n-1}$ . Observe that for any labeled tournament  $T$  that is not valid, there cannot exist two vertices  $v$  and  $w$  with  $\text{id } v = 0$  and  $\text{id } w = 0$  because either  $(v, w) \in E(T)$  or  $(w, v) \in E(T)$ . Hence for any labeled tournament  $T$  that is not valid, there exists a unique vertex  $v$  such that  $\text{id } v = 0$ . Since there are  $n$  choices for the label of the vertex of indegree 0 and, by removing this vertex, we obtain a labeled tournament of order  $n - 1$ , it follows that the number of labeled tournaments that are not valid is  $nT_{n-1}$ . Thus the number of valid labeled tournaments is  $T_n^v = T_n - nT_{n-1}$ .

Next, let  $T_n^k$  denote the number of labeled tournaments of order  $n$  with open domination number equal to  $k$ ,  $3 \leq k \leq n$ . Any such labeled tournament  $T$  has a minimum open dominating set  $S$  such that  $\langle S \rangle$  is a valid subtournament of order  $k$ . There are  $\binom{n}{k}$  ways to label the vertices of  $S$ . So the number of possible orientations of  $\langle S \rangle$  is  $T_k^v$ . Let  $T'$  be the subtournament  $\langle V(T) - S \rangle$ . The number of possible orientations of  $T'$  is  $T_{n-k}$ . Since  $S$  is an open dominating set of  $T$ , it follows that for every vertex  $v$  of  $T'$ , there is at least one vertex of  $S$  adjacent to  $v$ . That is, each of the  $n - k$  vertices of  $T'$  is adjacent from at least one of the  $k$  vertices of  $S$ . Hence the number of possible orientations of edges between  $S$  and  $V(T')$  is

$$\left[ \binom{k}{1} + \binom{k}{2} + \dots + \binom{k}{k} \right]^{n-k}.$$

Therefore,

$$\begin{aligned}
 T_n^k &\leq \binom{n}{k} \cdot T_k^v \cdot T_{n-k} \cdot [ \binom{k}{1} + \binom{k}{2} + \dots + \binom{k}{k} ]^{n-k} \\
 &= \binom{n}{k} \cdot T_k^v \cdot T_{n-k} \cdot (2^k - 1)^{n-k} \\
 &= \binom{n}{k} \cdot T_k^v \cdot 2^{\frac{(n-k)(n+k-1)}{2}} \cdot \left( \frac{2^k - 1}{2^k} \right)^{n-k} \\
 &= \binom{n}{k} \cdot \left[ 2^{\frac{k(k-1)}{2}} - k 2^{\frac{(k-1)(k-2)}{2}} \right] \cdot 2^{\frac{(n-k)(n+k-1)}{2}} \cdot \left( \frac{2^k - 1}{2^k} \right)^{n-k} \\
 &= \binom{n}{k} \cdot \left[ 2^{\frac{n(n-1)}{2}} - k 2^{\frac{n(n-1)-2(k-1)}{2}} \right] \cdot \left( \frac{2^k - 1}{2^k} \right)^{n-k} \\
 &= \binom{n}{k} \cdot 2^{\binom{n}{2}} \cdot \left[ 1 - \frac{k}{2^{k-1}} \right] \cdot \left( \frac{2^k - 1}{2^k} \right)^{n-k}
 \end{aligned}$$

Also observe that

$$T_n^v = 2^{\binom{n}{2}} - n 2^{\binom{n-1}{2}} = 2^{\binom{n}{2}} \cdot \left[ 1 - \frac{n}{2^{n-1}} \right].$$

Now let  $g(n) = T_n^v - \sum_{i=3}^k T_n^i$ , which counts the number of valid tournaments of order  $n$  with minimum open dominating set of cardinality exceeding  $k$ . Then we have

$$\begin{aligned}
 g(n) &= T_n^v - \sum_{i=3}^k T_n^i \\
 &\geq 2^{\binom{n}{2}} \left[ 1 - \frac{n}{2^{n-1}} \right] - \sum_{i=3}^k \binom{n}{i} \cdot 2^{\binom{n}{2}} \cdot \left[ 1 - \frac{i}{2^{i-1}} \right] \left[ \frac{2^i - 1}{2^i} \right]^{n-i} \\
 &= 2^{\binom{n}{2}} \left( \left[ 1 - \frac{n}{2^{n-1}} \right] - \sum_{i=3}^k \binom{n}{i} \cdot \left[ 1 - \frac{i}{2^{i-1}} \right] \left[ \frac{2^i - 1}{2^i} \right]^{n-i} \right).
 \end{aligned}$$

Since  $\lim_{n \rightarrow \infty} \sum_{i=3}^k \binom{n}{i} \cdot [1 - \frac{i}{2^{i-1}}] [\frac{2^i - 1}{2^i}]^{n-i} = 0$ , it follows that  $\lim_{n \rightarrow \infty} g(n) = \infty$ . Thus, for any fixed  $k$ ,  $3 \leq k < n$ , there exists a sufficiently large integer  $n$  such that for some valid labeled tournament  $T$  of order  $n$ ,  $\rho_1(T) > k$ , implying that  $\text{DOM}_1(K_n) > k$ .  $\square$

Using the method described above, the first value of  $n$  for which  $g(n) > 0$  is 77. Thus  $\text{DOM}_1(K_{77}) > 3$ . However this is not the least integer  $n$  for which  $\text{DOM}_1(K_n) > 3$ . It can be verified that  $\text{DOM}_1(K_n) = 3$  for  $n = 3, 4, \dots, 10$ , the most difficult of which is  $n = 10$ . We verify this.

**Theorem 9**  $\text{DOM}_1(K_{10}) = 3$ .

**Proof** Let  $D$  be an orientation of  $K_{10}$  such that  $\rho_1(D) = \text{DOM}_1(K_{10})$ . There exists a vertex, say  $v_1$ , with maximum outdegree. Necessarily,  $\text{od } v_1 \geq 5$ .

*Case 1*  $\text{od } v_1 = 5$ . Let  $N^+(v_1) = \{u_1, u_2, u_3, u_4, u_5\}$ , and let  $N^-(v_1) = \{v_2, v_3, v_4, v_5\}$ . Assume, without loss of generality, that  $v_2$  is adjacent to at least two of  $\{v_3, v_4, v_5\}$ . In particular, assume the arcs  $(v_2, v_4)$  and  $(v_2, v_5)$  belong to  $D$ .

*Subcase 1.1*  $(v_2, v_3) \in E(D)$ . Then, since  $\text{id } v_2 \geq 1$ , it follows that  $(u_i, v_2) \in E(D)$  for some  $i \in \{1, 2, 3, 4, 5\}$ . It follows that  $S = \{u_i, v_2, v_1\}$  is an open dominating set. Thus  $\rho_1(D) = 3$ .

*Subcase 1.2*  $(v_3, v_2) \in E(D)$ . Since the maximum outdegree of  $D$  is 5, it follows that at least 3 of the vertices of  $N^+(v_1)$  are adjacent to  $v_2$ . Suppose  $u_1, u_2$ , and  $u_3$  are among the vertices adjacent to  $v_2$ . Assume that none of the sets  $\{u_i, v_2, v_1\}$ ,  $1 \leq i \leq 3$ , is an open dominating set. Then there must exist some vertex not openly dominated by any of these sets. In particular,  $v_3$  is the

only such vertex. So  $(v_3, u_i) \in E(D)$  for  $i = 1, 2, 3$ . Now since the maximum outdegree of  $D$  is 5, it follows that  $(v_4, v_3) \in E(D)$  and  $(v_5, v_3) \in E(D)$ . Without loss of generality, assume  $(v_4, v_5) \in E(D)$ . Since the maximum outdegree of  $D$  is 5, it follows that at least 3 of the vertices of  $N^+(v_1)$  are adjacent to  $v_4$ . Hence there is at least one vertex, say  $u_i$ , for some  $i$ , of  $N^+(v_1)$  which is adjacent to both  $v_2$  and  $v_4$ . Then  $S = \{u_i, v_1, v_4\}$  is an open dominating set, implying  $\rho_1(D) = 3$ .

*Case 2*  $\text{od } v_1 = 6$ . Let  $N^+(v_1) = \{u_1, u_2, u_3, u_4, u_5, u_6\}$ , and let  $N^-(v_1) = \{v_2, v_3, v_4\}$ . Assume, without loss of generality, that  $v_2$  is adjacent to at least one of  $v_3$  and  $v_4$ .

*Subcase 2.1*  $\langle v_2, v_3, v_4 \rangle$  is a directed 3-cycle. If  $\{v_2, v_3, v_4\}$  is not an open dominating set, then there exists  $u_i$  ( $1 \leq i \leq 6$ ) such that  $(u_i, v_2), (u_i, v_3), (u_i, v_4) \in E(D)$ . Then  $S = \{u_i, v_1, v_2\}$  is an open dominating set and  $\rho_1(D) = 3$ .

*Subcase 2.2*  $\langle v_2, v_3, v_4 \rangle$  is not a directed 3-cycle. Then one of the vertices, say  $v_2$ , is adjacent to the other two,  $v_3$  and  $v_4$ . And since the maximum outdegree of  $D$  is 6, it follows that at least 3 of  $N^+(v_1)$  are adjacent to  $v_2$ . Suppose  $u_1, u_2$ , and  $u_3$  are among the vertices adjacent to  $v_2$ . Then  $S = \{u_1, v_1, v_2\}$  is an open dominating set, implying  $\rho_1(D) = 3$ .

*Case 3*  $\text{od } v_1 = 7$ . Let  $N^+(v_1) = \{u_i \mid 1 \leq i \leq 7\}$  and  $N^-(v_1) = \{v_2, v_3\}$ . Assume, without loss of generality, that  $(v_2, v_3) \in E(D)$ . Since the maximum outdegree of  $D$  is 7, there exist at least two vertices of  $N^+(v_1)$  adjacent to  $v_2$ . Suppose  $u_i$  is one of these vertices for some  $i$ . Then  $S = \{u_i, v_1, v_2\}$  is an open dominating set. So  $\rho_1(D) = 3$ .

*Case 4*  $\text{od } v_1 = 8$ . Let  $N^+(v_1) = \{u_i \mid 1 \leq i \leq 8\}$  and  $N^-(v_1) = \{v_2\}$ . Since the maximum outdegree of  $D$  is 8, there exists at least one vertex, say  $u_i$ , of

$N^+(v_1)$  adjacent to  $v_2$ . Then  $S = \{u_i, v_1, v_2\}$  is an open dominating set. Therefore  $\rho_1(D) = 3$ .  $\square$

Finally, to show that the least integer  $n$  for which  $\text{DOM}_1(K_n) = 4$  is  $n = 11$ , it is required to verify that  $\text{DOM}_1(K_{11}) \neq 3$ .

**Theorem 10**  $\text{DOM}_1(K_{11}) = 4$ .

**Proof** We need only show the existence of an orientation of  $K_{11}$  with open domination number 4. We describe such an orientation  $D$  as follows. Let  $V(D) = \{u_1, u_2, u_3, u_4, u_5, v_1, v_2, v_3, v_4, v_5, w\}$ . The tournament  $D$  contains two strong subtournaments of order 5, say  $T_1 = \langle \{u_1, u_2, \dots, u_5\} \rangle$  and  $T_2 = \langle \{v_1, v_2, \dots, v_5\} \rangle$ . Also the vertices of  $T_2$  are adjacent to  $w$ , and the vertices of  $T_1$  are adjacent from  $w$ . See Figure 5.

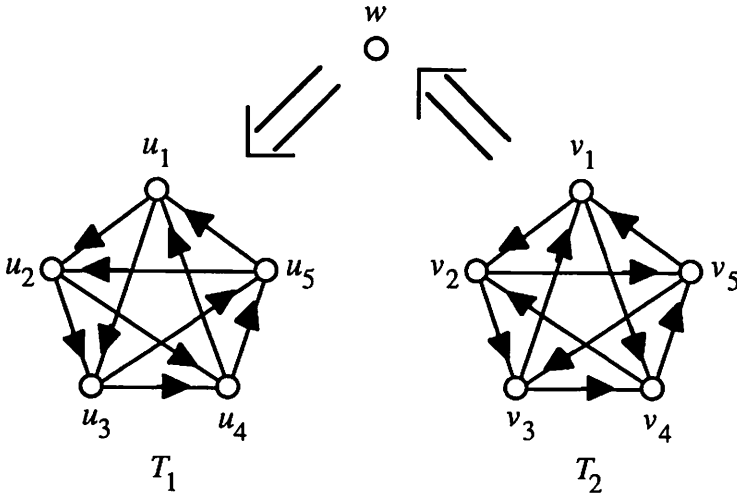


Figure 5

The orientations of edges between  $V(T_1)$  and  $V(T_2)$  remain to be described. We determine these orientations in the following way. There are

precisely 5 directed 3-cycles in  $T_2$ . Let  $S_1 = \{v_1, v_2, v_3\}$ ,  $S_2 = \{v_3, v_4, v_5\}$ ,  $S_3 = \{v_1, v_2, v_5\}$ ,  $S_4 = \{v_2, v_3, v_4\}$ , and  $S_5 = \{v_1, v_4, v_5\}$  be the 5 sets of vertices which determine the directed 3-cycles of  $T_2$ . Now for each  $i = 1, 2, 3, 4, 5$ , let  $u_i$  be adjacent to  $V(S_i)$ . All other arcs not mentioned thus far are directed from  $V(T_2)$  to  $V(T_1)$ . The resulting orientation is the desired orientation  $D$ .

It remains to show that there is no minimum open dominating set of 3 vertices. Suppose, to the contrary, that  $\rho_1(D) = 3$ . Then there exists a directed 3-cycle, say  $\langle S \rangle$ , which openly dominates all vertices of  $D$ . Observe that  $\langle S \rangle$  cannot be one of the triangles of  $T_1$  since  $w$  is adjacent to all vertices of  $V(T_1)$ . Also, by construction, each of the triangles of  $T_2$  is openly dominated by some vertex of  $V(T_1)$ . So  $\langle S \rangle$  cannot be one of the triangles of  $T_2$ . Hence we consider 3-cycles which contain vertices from both  $T_1$  and  $T_2$ .

*Case 1*  $S = \{w, u_i, v_j\}$  for some  $i$ , some  $j$ . Then  $\langle S \rangle$  is the cycle  $w, u_i, v_j, w$ . Now  $u_i$  is directed toward a unique 3-cycle, say  $C$ , of  $T_2$  where  $v_j$  belongs to  $C$ . Of the two vertices of  $V(T_2) - V(C)$ , one is adjacent to  $v_j$  and one is adjacent from  $v_j$ . Let  $v_k$  be the vertex adjacent to  $v_j$ . Then  $v_k$  openly dominates  $S$ . So  $S$  cannot be an open dominating set of  $D$ .

*Case 2*  $S = \{u_i, u_j, v_k\}$  for some  $i, j$  and  $k$ . There are 15 triangles of this type. However, each triangle is openly dominated by some vertex. (See Figure 6.) So  $S$  cannot be an open dominating set of  $D$ .



$S$	openly dominated by vertex
$u_1, u_2, v_5$	$u_5$
$u_1, u_2, v_4$	$u_5$
$u_1, u_3, v_5$	$v_4$
$u_4, u_1, v_1$	$v_5$
$u_5, u_1, v_3$	$u_4$
$u_5, u_1, v_2$	$u_4$
$u_2, u_3, v_1$	$u_1$
$u_2, u_3, v_2$	$u_1$
$u_2, u_4, v_2$	$v_1$
$u_5, u_2, v_3$	$v_2$
$u_3, u_4, v_4$	$u_2$
$u_3, u_4, v_3$	$u_2$
$u_3, u_5, v_4$	$v_3$
$u_4, u_5, v_1$	$u_3$
$u_4, u_5, v_5$	$u_3$

Figure 6

Case 3  $S = \{u_i, v_j, v_k\}$  for some  $i, j$  and  $k$ . There are 15 triangles of this type. However, each triangle is openly dominated by some vertex. (See Figure 7.) Thus  $S$  cannot be an open dominating set of  $D$ .

S	openly dominated by vertex
$v_1, v_2, u_5$	$u_3$
$v_3, v_1, u_2$	$u_1$
$v_3, v_1, u_4$	$v_5$
$v_1, v_4, u_3$	$v_3$
$v_1, v_4, u_1$	$u_5$
$v_5, v_1, u_2$	$u_5$
$v_2, v_3, u_3$	$u_1$
$v_4, v_2, u_2$	$v_1$
$v_4, v_2, u_5$	$u_4$
$v_2, v_5, u_1$	$v_4$
$v_2, v_5, u_4$	$u_3$
$v_3, v_4, u_1$	$u_4$
$v_5, v_3, u_3$	$u_2$
$v_5, v_3, u_5$	$v_2$
$v_4, v_5, u_4$	$u_2$

Figure 7

We have now exhausted all possibilities. Thus there is no such open dominating set of 3 vertices, implying  $\rho_1(D) \neq 3$ . By Lemma 6, it follows that  $\text{DOM}_1(K_{11}) = 4$ .  $\square$

### 3. An Intermediate Value Theorem for Orientable Open Domination in Complete Graphs

By the results of the preceding section, it follows that if  $m, n$ , and  $c$  are positive integers such that  $m \leq n$  and  $\text{DOM}_1(K_m) \leq c \leq \text{DOM}_1(K_n)$ , then there exists an integer  $k$ , with  $m \leq k \leq n$ , such that  $\text{DOM}_1(K_k) = c$ . Hence we have a certain type of Intermediate Value Theorem. In this section, we consider the existence of another type of Intermediate Value Theorem. In particular, for a graph  $G$ , given an integer  $c$  such that  $\text{dom}_1(G) \leq c \leq \text{DOM}_1(G)$ , does there exist an orientation  $D$  of  $G$  such that  $\rho_1(D) = c$ ?

Although such a theorem has not yet been proven for an arbitrary graph  $G$ , the result does hold for complete graphs. We begin by establishing a few lemmas for graphs in general.

**Lemma 11** Let  $G$  be a graph and let  $v$  be any vertex of  $G$ . If  $\text{DOM}_1(G - v)$  is defined, then

$$\text{DOM}_1(G) \geq \text{DOM}_1(G - v) \geq \text{DOM}_1(G) - 1.$$

**Proof** Assume  $\text{DOM}_1(G - v)$  is defined. Let  $D - v$  be an orientation of  $G - v$  such that  $\rho_1(D - v) = \text{DOM}_1(G - v)$ . Define  $D$  by directing all edges of  $G$  incident with  $v$  towards the vertex  $v$ . Notice that a minimum open dominating set of  $D$  also openly dominates  $D - v$ . Hence  $\rho_1(D) \geq \rho_1(D - v)$ . Thus  $\text{DOM}_1(G) \geq \rho_1(D) \geq \rho_1(D - v) = \text{DOM}_1(G - v)$ .

Next let  $D$  be any orientation of  $G$  such that every vertex of  $D$  has positive indegree. Then, for the given vertex  $v$ , there exists some arc, say  $(u, v)$  in  $D$ . Now an open dominating set of  $D$  can be formed from an open dominating set of  $D - v$ , possibly along with the vertex  $u$ . Thus  $\rho_1(D) \leq \rho_1(D - v) + 1 \leq \text{DOM}_1(G - v) + 1$ . Since this is true for any valid orientation  $D$  of  $G$ , it follows that  $\text{DOM}_1(G) \leq \text{DOM}_1(G - v) + 1$ , that is,  $\text{DOM}_1(G - v) \geq \text{DOM}_1(G) - 1$ .  $\square$

**Lemma 12** If  $\text{DOM}_1(G) > \text{dom}_1(G)$ , then there exists a vertex  $v$  of  $G$  such that

- (a)  $\text{dom}_1(G - v)$  and  $\text{DOM}_1(G - v)$  are defined
- (b)  $\text{dom}_1(G) \geq \text{dom}_1(G - v) \geq \text{dom}_1(G) - 1$
- (c) either  $\text{DOM}_1(G - v) = \text{DOM}_1(G)$  or  $\text{DOM}_1(G - v) = \text{DOM}_1(G) - 1$ .

**Proof** Clearly, if  $G$  is a cycle, then  $\text{dom}_1(G) = \text{DOM}_1(G)$ . Thus, assuming  $\text{DOM}_1(G) > \text{dom}_1(G)$ , it follows that  $G$  is not a cycle and  $G$  is not a tree.

Now let  $D$  be an orientation of  $G$  such that every vertex has positive indegree. Let  $S$  be a minimum open dominating set of  $D$ . Since  $G$  is not a cycle, we know  $|S| < n$ . Hence there exists some vertex, say  $v$ , in  $V(G) - S$ . Now since  $S$  is an open dominating set of  $D$ , it follows that  $\langle S \rangle$  contains a cycle. Further every vertex of  $V(G) - S$  is adjacent from some vertex of  $S$ . Consequently  $G - v$  is connected and contains a cycle. In fact, every vertex of  $D - v$  has positive indegree and  $\rho_1(D - v)$  exists. So  $\text{dom}_1(G - v)$  and  $\text{DOM}_1(G - v)$  are defined, proving part (a).

Let  $D$  be an orientation of  $G$  such that  $\rho_1(D) = \text{dom}_1(G)$ . Let  $S$  be a minimum open dominating set of  $D$ . And let  $v$  be a vertex of  $V(G) - S$ , as in the proof of part (a). Then  $\text{dom}_1(G) = \rho_1(D) \geq \rho_1(D - v) \geq \text{dom}_1(G - v)$ , proving the first inequality of part (b).

Next let  $D - v$  be an orientation of  $G - v$  such that  $\text{dom}_1(G - v) = \rho_1(D - v)$ . Let  $D$  be the orientation of  $G$  formed by directing all edges incident with  $v$  toward the vertex  $v$ . Then  $\text{dom}_1(G) \leq \rho_1(D) \leq \rho_1(D - v) + 1$ . Hence  $\text{dom}_1(G) - 1 \leq \rho_1(D - v) = \text{dom}_1(G - v)$ , proving the second inequality in part (b).

The proof of part (c) follows directly from Lemma 11.  $\square$

**Lemma 13** Let  $G$  be a connected graph such that  $G$  is not a tree, and let  $c$  be an integer such that

$$\text{dom}_1(G) \leq c < \text{DOM}_1(G).$$

Then there exists a sequence  $v_1, v_2, \dots, v_k$  of vertices of  $G$  such that for  $G_i = G - \{v_1, v_2, \dots, v_i\}$ ,  $1 \leq i \leq k$ , we have

- (a)  $\text{dom}_1(G_i)$  and  $\text{DOM}_1(G_i)$  are defined
- (b)  $\text{DOM}_1(G_k) = c$ .

**Proof** The result is obvious if  $c = \text{dom}_1(G)$ . Thus we assume  $\text{dom}_1(G) < c$ . We proceed iteratively.

By Lemma 12, there exists a vertex, say  $v_1$ , such that  $\text{dom}_1(G_1)$  and  $\text{DOM}_1(G_1)$  are defined and either  $\text{DOM}_1(G_1) = \text{DOM}_1(G)$  or  $\text{DOM}_1(G_1) = \text{DOM}_1(G) - 1$ . We consider the following two cases.

*Case 1*  $\text{DOM}_1(G_1) = \text{DOM}_1(G)$ . Then, by Lemma 12,  $\text{dom}_1(G_1) \leq \text{dom}_1(G) < \text{DOM}_1(G) = \text{DOM}_1(G_1)$ . That is,  $\text{dom}_1(G_1) < \text{DOM}_1(G_1)$ .

*Case 2*  $\text{DOM}_1(G_1) = \text{DOM}_1(G) - 1$  and  $c < \text{DOM}_1(G_1)$ . Then, by Lemma 12,  $\text{dom}_1(G_1) \leq \text{dom}_1(G) < c < \text{DOM}_1(G_1)$ . That is,  $\text{dom}_1(G_1) < \text{DOM}_1(G_1)$ .

Observe that for each of these cases, the graph  $G_1$  satisfies the hypothesis of Lemma 12. Thus there exists a vertex, say  $v_2$ , such that  $\text{DOM}_1(G_2)$  is defined, and the process continues.

In general, as long as  $\text{dom}_1(G_i) < \text{DOM}_1(G_i)$ , there exists a vertex  $v_{i+1}$  such that  $\text{DOM}_1(G_{i+1})$  is defined. Further, the process continues as long as

- (1)  $\text{DOM}_1(G_{i+1}) = \text{DOM}_1(G_i)$  or
- (2)  $\text{DOM}_1(G_{i+1}) = \text{DOM}_1(G_i) - 1$  and  $c < \text{DOM}_1(G_{i+1})$

We claim that the process terminates, that is, there exists  $k$  such that  $\text{DOM}_1(G_k) = \text{DOM}_1(G_{k-1}) - 1$  and  $c = \text{DOM}_1(G_k)$ . Let  $k - 1$  be the largest integer such that  $\text{dom}_1(G_{k-1}) \leq c < \text{DOM}_1(G_{k-1})$ . Then, it follows that there exists a vertex  $v_k$  such that  $\text{DOM}_1(G_k)$  is defined and  $\text{DOM}_1(G_k) \leq c$ . But we know either  $\text{DOM}_1(G_k) = \text{DOM}_1(G_{k-1})$  or  $\text{DOM}_1(G_k) = \text{DOM}_1(G_{k-1}) - 1$ . Since  $\text{DOM}_1(G_{k-1}) > c$ , we must have

$\text{DOM}_1(G_k) = \text{DOM}_1(G_{k-1}) - 1$ . Hence  $c \geq \text{DOM}_1(G_k) = \text{DOM}_1(G_{k-1}) - 1 > c - 1$ , implying  $\text{DOM}_1(G_k) = c$ .  $\square$

Finally, we have an Intermediate Value Theorem for the upper orientable open domination number of a complete graph.

**Theorem 14** Let  $c$  be an integer such that  $\text{dom}_1(K_n) \leq c \leq \text{DOM}_1(K_n)$ . Then there exists an orientation  $D$  of  $K_n$  such that  $\rho_1(D) = c$ .

**Proof** Certainly if  $c = \text{dom}_1(K_n)$  or  $c = \text{DOM}_1(K_n)$ , the result is clear. Thus we assume  $\text{dom}_1(K_n) < c < \text{DOM}_1(K_n)$ . By Lemma 13, there exists a set of vertices  $W = \{v_1, v_2, \dots, v_k\}$ ,  $k \geq 1$ , such that  $\text{DOM}_1(K_n - W)$  is defined and  $\text{DOM}_1(K_n - W) = c$ . Let  $D'$  be an orientation of  $K_n - W$  such that  $\rho_1(D') = \text{DOM}_1(K_n - W)$ . Form an orientation  $D$  of  $K_n$  from  $D'$  by letting  $\langle D - W \rangle = D'$  and directing all edges incident with the vertices of  $W$  toward  $W$ . Now let  $S$  be a minimum open dominating set of  $D'$ , that is,  $\rho_1(D') = |S|$ . By construction, the arc  $(u, v_i)$  belongs to  $D$  for every  $u \in S$  and for every  $v_i \in W$ . Thus every vertex of  $W$  is openly dominated by a vertex of  $S$ . Hence,  $S$  is an open dominating set of  $D$ . Further  $S$  is, in fact, a minimum open dominating set, for otherwise we would have  $\rho_1(D') < |S|$ . Therefore

$$\rho_1(D) = |S| = \rho_1(D') = \text{DOM}_1(K_n - W) = c. \quad \square$$

Another consequence of the previous lemmas is the following.

**Corollary 15** Let  $G$  be a connected graph which is not a tree. Let  $c$  be any integer such that  $\text{dom}_1(G) \leq c \leq \text{DOM}_1(G)$ . Then  $G$  contains an induced subgraph  $H$  such that  $\text{DOM}_1(H) = c$ .

We conclude with a conjecture.

**Conjecture** Let  $G$  be a graph. If  $c$  is an integer for which  $\text{dom}_1(G) \leq c \leq \text{DOM}_1(G)$ , then there exists an orientation  $D$  of  $G$  such that  $\rho_1(D) = c$ .

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