

# A Note on MAD Spanning Trees

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**ABSTRACT.** The question whether every connected graph  $G$  has a spanning tree  $T$  of minimum average distance such that  $T$  is distance preserving from some vertex is answered in the negative. Moreover it is shown that, if such a tree exists, it is not necessarily distance preserving from a median vertex.

Over the past two decades the average distance of graphs has sparked considerable interest (see e.g. [2, 3, 7, 8]). One particular problem that received some attention is that of finding a minimum average distance (MAD) spanning tree of a given connected graph  $G$ , i.e. a spanning tree  $T$  such that  $T$  has minimum average distance among all spanning trees of  $G$ . It was proved by Johnson, Lenstra, and Rinnooy-Kan [6] that this problem is NP-complete. Entringer, Kleitman, and Szekely [4] showed that every connected graph  $G$  contains a spanning tree  $T$  that is distance preserving from some vertex in  $G$  such that the average distance of  $T$  is less than twice the average distance of  $G$ . This prompted the following questions by Entringer [5]:

1. Does every graph have a MAD spanning tree that is also distance preserving from some vertex?
2. If the graph  $G$  has a MAD spanning tree  $T$  that is distance preserving, is it distance preserving from a median vertex?

In this note it is shown that both questions have a negative answer.

We employ the notation and terminology of [1]. In particular, if  $G$  is a connected graph with vertex set  $V$ , then the *distance* of  $G$ ,  $d(G)$ , is defined as

$$d(G) = \sum_{\{v,w\} \subset V} d_G(v,w),$$

where  $d_G(v, w)$  denotes the distance between  $v$  and  $w$ . The *average distance* of  $G$ ,  $\mu(G)$ , is the average of all distances in  $G$ , i.e.  $\mu(G) = \binom{|V|}{2}^{-1} \sum_{\{v,w\} \subset V} d_G(v, w)$ .

**Definition 1.** For positive integers  $k, r$  let  $H$  be the graph obtained from a cycle  $C = (v_1, v_2, \dots, v_{2k-1}, v_1)$  and  $2r$  vertices  $w_1, w_2, \dots, w_r, x_1, x_2, \dots, x_r$  by joining the  $w_i$  with  $v_k$  and  $x_i$  with  $v_{k+1}$  for  $i = 1, \dots, r$ . Let  $H'$  be a disjoint copy of  $H$  with vertex set  $\{v'_1, \dots, v'_{2k-1}, w'_1, \dots, w'_r, x'_1, \dots, x'_r\}$ . Define  $G(k, r)$  as the graph obtained from the union of  $H$  and  $H'$  by adding the edge  $v_1 v'_1$ .

**Theorem 1.** Let  $v$  be a vertex of  $G(k, r)$  and let  $T$  be a spanning tree distance preserving from  $v$ . Then for  $k \geq 3$  and  $r \geq 6$ ,  $T$  is not a MAD spanning tree.

**Proof:** Without loss of generality we can assume that  $v$  is in  $H'$ . Then  $T$  does not contain the edge  $e = v_k v_{k+1}$  but the edge  $f = v_{k+1} v_{k+2}$ . Let  $T_1 = T - f + e$ . It is easy to verify that  $T_1$  is also a spanning tree of  $G(k, r)$ .

We now prove that replacing the edge  $f$  by the edge  $e$  decreases the distance of  $T$  (and thus  $T$  is not a MAD spanning tree). The only distances that are changed are distances between the  $w_i$  and the  $x_j$  (each goes down by  $2k - 3$ ), between  $w_i$  and  $v_{k+1}$  (down by  $2k - 3$ ), between each vertex in  $\{x_i\} \cup \{v_{k+1}\}$  and every vertex in  $H'$  (up by 1). It is easy to check that all the other changes in the distances cancel out. Hence we have, for  $k \geq 3$  and  $r \geq 6$

$$\begin{aligned} d(T_1) &= d(T) - r^2(2k - 3) - r(2k - 3) + (r + 1)(2k - 1 + 2r) \\ &= d(T) - (r + 1)((2k - 5)(r - 1) - 4) \\ &< d(T). \end{aligned}$$

Thus  $T$  is not a MAD spanning tree. □

**Definition 2.** Let  $H = H(k, r)$  be the graph obtained from a cycle  $C = (v_1, v_2, \dots, v_{2k-1}, v_1)$  of order  $2k - 1$  and  $5r$  isolated vertices  $x_1, x_2, \dots, x_r, y_1, y_2, \dots, y_r, z_1, z_2, \dots, z_{3r}$ , by joining  $x_i$  to  $v_k$  and  $y_i$  to  $v_{k+1}$  for  $i = 1, \dots, r$ , and  $z_j$  to  $v_1$  for  $j = 1, \dots, 3r$ .

**Theorem 2.** (i) The graph  $H$  contains a MAD spanning tree which is distance preserving from some vertex.

(ii) Let  $T$  be a spanning tree of  $H$  distance preserving from a vertex of minimum distance in  $H$ . Then  $T$  is not a MAD spanning tree.

**Proof:** (i) Every spanning tree of  $H$  is obtained by deleting a cycle edge  $e$ . It is easy to see that it is distance preserving from the vertex on the cycle opposite  $e$ . Hence  $H$  has a MAD spanning tree which is distance preserving from some vertex.

(ii) It is easy to check that the median of  $H$  consists only of  $v_1$ . Let  $T$  be the spanning tree of  $H$  distance preserving from  $v_1$ . Then  $T = H - e$  where  $e = v_k v_{k+1}$ . We now show that the spanning tree  $T_1 = H - f$  with  $f = v_{k+1} v_{k+2}$ , which is distance preserving from  $v_2$ , has a lower distance than  $T$ , hence  $T$  is not a MAD spanning tree.

As above we show that replacing the edge  $f$  in  $T$  by the edge  $e$  decreases the distance of  $T$  (and thus  $T$  is not a MAD spanning tree). The only distances that are changed are distances between the  $x_i$  and the  $y_j$  (each goes down by  $2k - 3$ ), between  $x_i$  and  $v_{k+1}$  (down by  $2k - 3$ ), between the vertices in  $\{y_i\} \cup \{v_{k+1}\}$  and  $z_j$  (up by 1). It is easy to check that all the other changes in the distances cancel out. Hence we have, for  $k \geq 4$

$$\begin{aligned} d(T_1) &= d(T) - r^2(2k - 3) - r(2k - 3) + (r + 1)3r \\ &= d(T) - r(r + 1)(2k - 6) \\ &< d(T). \end{aligned}$$

Thus  $T$  is not a MAD spanning tree. □

## References

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