

Close to regular multipartite tournaments containing a Hamiltonian path

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Abstract

If x is a vertex of a digraph D , then we denote by $d^+(x)$ and $d^-(x)$ the outdegree and the indegree of x , respectively. The global irregularity of a digraph D is defined by $i_g(D) = \max\{d^+(x), d^-(x)\} - \min\{d^+(y), d^-(y)\}$ over all vertices x and y of D (including $x = y$). If $i_g(D) = 0$, then D is regular and if $i_g(D) \leq 1$, then D is called almost regular. The local irregularity is defined as $i_l(D) = \max|d^+(x) - d^-(x)|$ over all vertices x of D . The path covering number of D is the minimum number of directed paths in D that are pairwise vertex disjoint and cover the vertices of D . A semicomplete c -partite digraph is a digraph obtained from a complete c -partite graph by replacing each edge with an arc, or a pair of mutually opposite arcs with the same end vertices. If a semicomplete c -partite digraph D does not contain an oriented cycle of the length two, then D is called a c -partite tournament.

In 2000, Gutin and Yeo [7] proved sufficient conditions for the local irregularity of a semicomplete multipartite digraph to secure a path covering number of at most k . In this paper, we will give a useful supplement to this result by using bounds for the global irregularity that guarantee a path covering number of at most k . As an application, we will present sufficient conditions for close to regular multipartite tournaments containing a Hamiltonian path. Especially,

we will characterize almost regular c -partite tournaments containing a Hamiltonian path.

Keywords: Multipartite tournaments; Semicomplete multipartite digraphs; Hamiltonian path

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1 Terminology and introduction

In this paper all digraphs are finite without loops and multiple arcs. The vertex set and arc set of a digraph D is denoted by $V(D)$ and $E(D)$, respectively. If xy is an arc of a digraph D , then we write $x \rightarrow y$ and say x dominates y , and if X and Y are two disjoint vertex sets or subdigraphs of D such that every vertex of X dominates every vertex of Y , then we say that X dominates Y , denoted by $X \rightarrow Y$. Furthermore, $X \rightsquigarrow Y$ denotes the fact that there is no arc leading from Y to X . By $d(X, Y)$ we denote the number of arcs from the set X to the set Y , i.e., $d(X, Y) = |\{xy \in E(D) : x \in X, y \in Y\}|$. If D is a digraph, then the *out-neighborhood* $N_D^+(x) = N^+(x)$ of a vertex x is the set of vertices dominated by x and the *in-neighborhood* $N_D^-(x) = N^-(x)$ is the set of vertices dominating x . Therefore, if there is the arc $xy \in E(D)$, then y is an *outer neighbor* of x and x is an *inner neighbor* of y . The numbers $d_D^+(x) = d^+(x) = |N^+(x)|$ and $d_D^-(x) = d^-(x) = |N^-(x)|$ are called the *outdegree* and *indegree* of x , respectively. For a vertex set X of D , we define $D[X]$ as the subdigraph induced by X . If we speak of a *cycle* or *path*, then we mean a directed cycle or directed path, and a cycle of length n is called an *n -cycle*. A cycle or path of a digraph D is *Hamiltonian*, if it includes all the vertices of D . The *path covering number* of a digraph D ($pc(D)$) is the minimum number of paths in D that are pairwise vertex disjoint and cover the vertices of D . A *factor* is a spanning subgraph of a digraph. A factor is a *k -path-cycle*, if it consists of a set of vertex disjoint paths and cycles, where k stands for the number of paths in the set.

There are several measures of how much a digraph differs from being regular. In [14], Yeo defines the *global irregularity* of a digraph D by

$$i_g(D) = \max_{x \in V(D)} \{d^+(x), d^-(x)\} - \min_{y \in V(D)} \{d^+(y), d^-(y)\}$$

and the *local irregularity* as $i_l(D) = \max |d^+(x) - d^-(x)|$ over all vertices x of D . Clearly, $i_l(D) \leq i_g(D)$. If $i_g(D) = 0$, then D is *regular* and if $i_g(D) \leq 1$, then D is called *almost regular*.

A *c-partite* or *multipartite tournament* is an orientation of a complete *c-partite* graph. A *tournament* is a *c-partite* tournament with exactly *c* vertices. A *semicomplete multipartite digraph* is obtained by replacing each edge of a complete multipartite graph by an arc or by a pair of two mutually opposite arcs. If V_1, V_2, \dots, V_c are the partite sets of a *c-partite* tournament or semicomplete *c-partite* digraph D and the vertex x of D belongs to the partite set V_i , then we define $V(x) = V_i$. If D is a *c-partite* tournament with the partite sets V_1, V_2, \dots, V_c such that $|V_i| = n_i$ for $i = 1, 2, \dots, c$, then we speak of the *partition-sequence* $(n_i) = n_1, n_2, \dots, n_c$.

Hamiltonian cycles in multipartite tournaments are well studied (see e.g. [2, 3, 4, 8, 13, 14]). For example, Yeo [13] presented a result for regular semicomplete multipartite digraphs.

Theorem 1.1 (Yeo [13], 1997) *Every regular semicomplete multipartite digraph D is Hamiltonian.*

On the other hand, it is not paid much attention on the existence of Hamiltonian paths in such digraphs. In 1988, Gutin [5] gave a characterization of semicomplete multipartite digraphs having a Hamiltonian path.

Theorem 1.2 (Gutin [5], 1988) *A semicomplete multipartite digraph D has a Hamiltonian path if and only if it contains a 1-path-cycle factor.*

This result was used to prove another result of Gutin and Yeo [7].

Theorem 1.3 (Gutin, Yeo [7], 2000) *Let D be a semicomplete multipartite digraph with the partite sets V_1, V_2, \dots, V_c such that $|V_1| \leq |V_2| \leq \dots \leq |V_c|$. If there exists a positive integer k such that*

$$i_i(D) \leq \min\{|V(D)| - 3|V_c| + 2k + 1, \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 3k + 2}{2}\},$$

then $pc(D) \leq k$.

In this paper, we will give a useful supplement to Theorem 1.3 by presenting bounds for the global irregularity $i_g(D)$. As an application of this result, we will develop sufficient conditions for close to regular multipartite tournaments containing a Hamiltonian path.

If D is regular, then Theorem 1.1 clearly guarantees the existence of a Hamiltonian path which is already shown by Zhang [15] in 1989. Using

a sufficient condition for multipartite tournaments with an arbitrary large irregularity number containing a Hamiltonian path, we will show that every almost regular c -partite tournament D contains a Hamiltonian path with the exception that $c = 2$ and one partite set consists of two vertices more than the other partite set. Furthermore, we will precisely examine the case that $i_g(D) = 2$. If $c \geq 5$, then D contains a Hamiltonian path also in this case. If $c = 4$, then there is only a finite family of graphs that do not contain any Hamiltonian path. Finally, we will show that there are infinitely many 2-partite and 3-partite tournaments with $i_g(D) = 2$ that have no Hamiltonian path at all.

Further new results about Hamiltonian paths can be found in [11, 12]. For more informations about multipartite tournaments, we recommend to read [1, 6, 10].

2 Preliminary results

The following results play an important role in our investigations.

Lemma 2.1 (Tewes, Volkmann, Yeo [9], 2002) *Let D be a c -partite tournament with the partite sets V_1, V_2, \dots, V_c such that $|V_1| \leq |V_2| \leq \dots \leq |V_c|$. Then $|V_c| \leq |V_1| + 2i_g(D)$.*

Lemma 2.2 (Yeo [14], 1999) *Let D be a semicomplete multipartite digraph with the partite sets V_1, V_2, \dots, V_c . Let $X \subset Y \subseteq V(D)$ and let $y_i = |Y \cap V_i|$ for all $i = 1, 2, \dots, c$. Then*

$$\begin{aligned} & \frac{d(X, Y - X) + d(Y - X, X)}{|X|} + \frac{d(X, Y - X) + d(Y - X, X)}{|Y - X|} \\ & \geq |Y| - \max\{y_i \mid i = 1, 2, \dots, c\}. \end{aligned}$$

A slight reformulation of a result of Yeo [14] is presented in the following lemma.

Lemma 2.3 (Yeo [14], 1999) *If D is a semicomplete c -partite digraph, then the following holds.*

$$i_i(D) \geq \max_{\emptyset \neq X \subseteq V(D)} \left\{ \frac{|d(X, V(D) - X) - d(V(D) - X, X)|}{|X|} \right\}$$

Theorem 2.4 (Gutin, Yeo [7], 2000) *If D is a semicomplete c -partite digraph, then $pc(D) > k$ ($k \geq 1$) if and only if $V(D)$ can be partitioned into subsets Y, Z, R_1, R_2 such that*

$$R_1 \rightsquigarrow Y, (R_1 \cup Y) \rightsquigarrow R_2, Y \text{ is an independent set} \quad (1)$$

and $|Y| > |Z| + k$.

Theorem 2.5 (Volkman [11]) *If D is an almost regular bipartite tournament with the partite sets X, Y such that $1 \leq |X| \leq |Y|$, then every arc of D is contained in a Hamiltonian path if and only if $|Y| \leq |X| + 1$ and D is not isomorphic to $T_{3,3}$, where $T_{3,3}$ is the bipartite tournament presented in Figure 1.*

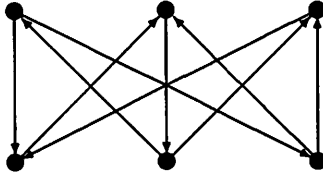


Figure 1: The almost regular bipartite tournament $T_{3,3}$

3 Path covering number and irregularity

The following theorem is a useful supplement to Theorem 1.3, which we will apply in the next section. The proof is similar to the proof of Lemma 4.3 in [14] and Theorem 3.2 in [7].

Theorem 3.1 *Let V_1, V_2, \dots, V_c be the partite sets of the semicomplete multipartite digraph D such that $|V_1| \leq |V_2| \leq \dots \leq |V_c|$. Assume that $pc(D) > k$ for an integer $k \geq 1$. According to Theorem 2.4, $V(D)$ can be partitioned into subsets Y, Z, R_1, R_2 satisfying (1) such that $|Z| + k + 1 \leq |Y| \leq |V_c| - t$ with an integer $t \geq 0$. Let V_i be the partite set with the property that $Y \subseteq V_i$. If $Q = V(D) - Z - V_i$, $Q_1 = Q \cap R_1$ and $Q_2 = Q \cap R_2$, then*

$$i_1(D) \geq |V(D)| - 3|V_c| + 2t + 2k + 2 \quad \text{and}$$

$$i_2(D) \geq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 3k + 3}{2},$$

if $Q_1 = \emptyset$ or $Q_2 = \emptyset$, and

$$i_g(D) \geq i_l(D) \geq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 3k + 3 + t}{2},$$

if $Q_1 \neq \emptyset$ and $Q_2 \neq \emptyset$.

Proof. Let $V(D)$ be partitioned into the subsets Y, Z, R_1, R_2 satisfying (1) such that $|Z| + k + 1 \leq |Y| \leq |V_c| - t$ for integers $k \geq 1$ and $t \geq 0$. Define $Y_1 = R_1 \cap V_i$ and $Y_2 = R_2 \cap V_i$. If Q_1 and Q_2 are defined as above, then we observe that $|Z| \leq |Y| - 1 - k \leq |V_c| - 1 - k - t$, $Q_1 \rightsquigarrow Y \rightsquigarrow Q_2$, $(Q_1 \cup Y_1) \rightsquigarrow (Q_2 \cup Y_2)$ and $Y_1 \cup Y_2 \cup Y \subseteq V_i$. If $i = c$, then let $j = c - 1$ and if $i < c$, then let $j = c$. We now consider the following three cases.

Case 1. Let $Q_1 = \emptyset$. Then $Q_2 = Q$ and we obtain

$$\begin{aligned} & d(Y, V(D) - Y) - d(V(D) - Y, Y) \\ & \geq |Y||Q_2| - |Y||Z| \geq |Y|(|V(D)| - |V_i| - 2|Z|) \\ & \geq |Y|(|V(D)| - |V_c| - 2(|V_c| - 1 - k - t)) \\ & = |Y|(|V(D)| - 3|V_c| + 2 + 2k + 2t). \end{aligned}$$

According to Lemma 2.3, this implies that $i_l(D) \geq |V(D)| - 3|V_c| + 2 + 2k + 2t$, and hence, we have one part of the desired result. We will now show the second part.

Let $\delta^* = \min\{d^-(w) | w \in V_i\}$. Since $Y \subseteq V_i$ and thus $d^-(y) \leq |Z|$ for all $y \in Y$ we observe that $\delta^* \leq |Z| \leq |Y| - k - 1 \leq |V_i| - |Y_2| - 1 - k$. Let $\Delta^* = \max\{d^+(w), d^-(w) | w \in V(D) - V_i\}$ and note that $d^+(w) + d^-(w) \geq |V(D)| - |V_j|$ for all $w \in V(D) - V_i$. The fact that $\sum_{x \in Q_2} (d^-(x) - d^+(x)) \geq |Q_2|(|Y| - |Z| - |Y_2|) \geq |Q_2|(1 + k - |Y_2|)$ implies that there is a vertex $q \in Q_2$ such that $d^-(q) > d^+(q) + k - |Y_2|$. This leads to $2d^-(q) - k + |Y_2| > d^+(q) + d^-(q) \geq |V(D)| - |V_j|$, and thus we conclude that $\Delta^* > \frac{|V(D)| - |V_j| + k - |Y_2|}{2}$. This implies

$$\begin{aligned} i_g(D) & \geq \Delta^* - \delta^* > \frac{|V(D)| - |V_j| + k - |Y_2|}{2} - |V_i| + |Y_2| + k + 1 \\ & = \frac{|V(D)| - |V_j| - 2|V_i| + 3k + 2 + |Y_2|}{2} \\ & \geq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 3k + 2}{2}, \end{aligned}$$

and the second part is proved.

Case 2. Let $Q_2 = \emptyset$. This is analogously to Case 1 (change the orientation of all the arcs in D).

Case 3. Let $Q_1 \neq \emptyset$ and $Q_2 \neq \emptyset$. Since $|V_i| + |V_j| \leq |V_{c-1}| + |V_c|$, we deduce that $|Q| - |V_j| \geq |V(D)| - |V_i| - |Z| - |V_j| \geq |V(D)| - |V_{c-1}| - |V_c| - (|V_c| - 1 - k - t)$. By Lemma 2.2 with $X = Q_1$ and $Y = Q_1 \cup Q_2 = Q$ and because of $Q \cap V_i = \emptyset$, it follows that

$$\begin{aligned} & \frac{d(Q_1, Q_2) + d(Q_2, Q_1)}{|Q_1|} + \frac{d(Q_1, Q_2) + d(Q_2, Q_1)}{|Q_2|} \\ &= \frac{d(Q_1, Q_2)}{|Q_1|} + \frac{d(Q_1, Q_2)}{|Q_2|} \geq |Q| - |V_j| \\ &\geq |V(D)| - |V_{c-1}| - 2|V_c| + 1 + k + t. \end{aligned}$$

Thus,

$$\begin{aligned} \text{(i)} \quad & \frac{d(Q_1, Q_2)}{|Q_1|} \geq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 1 + k + t}{2} - |Y_2| + |Y_1| \text{ or} \\ \text{(ii)} \quad & \frac{d(Q_1, Q_2)}{|Q_2|} \geq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 1 + k + t}{2} + |Y_2| - |Y_1|. \end{aligned}$$

Assume that (i) holds as the case when (ii) holds can be treated similarly. Because of $R_1 = Q_1 \cup Y_1$ and $R_2 = Q_2 \cup Y_2$, Lemma 2.3 yields

$$\begin{aligned} i_g(D) \geq i_l(D) &\geq \frac{d(Q_1, V(D) - Q_1) - d(V(D) - Q_1, Q_1)}{|Q_1|} \\ &= \frac{d(Q_1, Q_2)}{|Q_1|} + \frac{d(Q_1, Y \cup Y_2) - d(Y \cup Y_2, Q_1)}{|Q_1|} \\ &\quad + \frac{d(Q_1, Z \cup Y_1) - d(Z \cup Y_1, Q_1)}{|Q_1|} - \frac{d(Q_2, Q_1)}{|Q_1|} \\ &\geq \left(\frac{|V(D)| - |V_{c-1}| - 2|V_c| + 1 + k + t}{2} - |Y_2| + |Y_1| \right) \\ &\quad + (|Y| + |Y_2|) - (|Z| + |Y_1|) \\ &= \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 1 + k + t}{2} + |Y| - |Z| \\ &\geq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 3 + 3k + t}{2}. \end{aligned}$$

This completes the proof of the theorem. \square

Theorem 3.1 with $t = 0$ leads immediately to the following result on the path covering number.

Corollary 3.2 *Let V_1, V_2, \dots, V_c be the partite sets of a semicomplete multipartite digraph D such that $|V_1| \leq |V_2| \leq \dots \leq |V_c|$. If there exists a positive integer k such that $i_g(D) \leq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 3k + 2}{2}$, then $pc(D) \leq k$.*

4 Hamiltonian paths

Another consequence of Theorem 3.1 is presented in the following result.

Theorem 4.1 *Let V_1, V_2, \dots, V_c be the partite sets of the semicomplete c -partite digraph D such that $1 \leq r = |V_1| \leq |V_2| \leq \dots \leq |V_c| \leq r + p$ for an integer $p \geq 0$. If $c \geq \max\{2, 3 + \frac{2i_g(D) - 5 + p}{r}\}$, then D contains a Hamiltonian path.*

Proof. Clearly, D contains a Hamiltonian path if and only if $pc(D) = 1$. Hence, according to Corollary 3.2 with $k = 1$, it is sufficient to show that $i_g(D) \leq \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 5}{2}$. Because of $c \geq 3 + \frac{2i_g(D) - 5 + p}{r}$, we conclude that $i_g(D) \leq \frac{(c-3)r + 5 - p}{2}$, and together with $|V_1|, |V_2|, \dots, |V_{c-2}| \geq r, |V_c| \leq r + p$ and $c \geq 2$ this implies

$$\begin{aligned} \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 5}{2} &= \frac{|V_1| + |V_2| + \dots + |V_{c-2}| - |V_c| + 5}{2} \\ &\geq \frac{(c-3)r - p + 5}{2} \geq i_g(D), \end{aligned}$$

the desired result. \square

If D is a multipartite tournament, then, according to Theorem 2.1, we can choose $p = 2i_g(D)$ in the previous theorem.

Corollary 4.2 *Let V_1, V_2, \dots, V_c be the partite sets of a c -partite tournament D such that $1 \leq r = |V_1| \leq |V_2| \leq \dots \leq |V_c|$. If $c \geq \max\{2, \frac{4i_g(D) - 5}{r} + 3\}$, then D contains a Hamiltonian path.*

Theorem 4.3 *Let D be an almost regular c -partite tournament with the partite sets V_1, V_2, \dots, V_c such that $|V_1| \leq |V_2| \leq \dots \leq |V_c|$. Then D contains a Hamiltonian path if and only if $c \geq 3$ or $c = 2$ and $|V_2| \leq |V_1| + 1$.*

Proof. Firstly, let $c = 2$. Suppose that $|V_2| = |V_1| + 2$. Then it is obvious that D does not contain any Hamiltonian path, because the vertices of this path would alternately be part of the partite sets V_1 and V_2 . Hence let $|V_2| \leq |V_1| + 1$. In this case Theorem 2.5 shows that D contains a Hamiltonian path, since $T_{3,3}$ is Hamiltonian.

Secondly, let $c \geq 3 \geq \max\{2, 3 - \frac{1}{r}\}$ for all $r \in \mathbb{N}$. In this case, Corollary 4.2 yields the desired result. \square

The case that $i_g(D) = 2$ is more complicated as the following considerations demonstrate.

Theorem 4.4 *Let V_1, V_2, \dots, V_c be the partite sets of a c -partite tournament D such that $1 \leq r = |V_1| \leq |V_2| \leq \dots \leq |V_c|$. If $i_g(D) = 2$, $c \geq 4$ and D doesn't have one of the partition-sequences $1, 1, 2, 4$; $1, 2, 3, 5$; $1, 1, 3, 4$ and $2, 2, 4, 6$, then D contains a Hamiltonian path.*

Proof. Since $\max\{\frac{8-5}{r} + 3 | r \in \mathbb{N}\} = 6 > 3$, Corollary 4.2 yields the desired result, if $c \geq 6$.

Hence, let $c = 4$ or $c = 5$ and assume that D does not contain any Hamiltonian path. Let the sets $Y, Z, R_1, R_2, Q, Q_1, Q_2, Y_1, Y_2$ be defined as in the proof of Theorem 3.1. Since the case that $Q_2 = \emptyset$ follows analogously as the case that $Q_1 = \emptyset$, in the following we will always distinguish the two cases that $Q_1 = \emptyset$ or $Q_1, Q_2 \neq \emptyset$.

Case 1. Let $c = 5$. If $r \geq 2$, then because of $c > \max\{\frac{8-5}{r} + 3 | r \geq 2\} = \frac{9}{2} > 3$ and Corollary 4.2 D contains a Hamiltonian path and the proof is finished. If $|V_5| \leq 4$, then $\frac{|V_1|+|V_2|+|V_3|-|V_5|+5}{2} \geq 2 = i_g(D)$, a contradiction to Corollary 3.2. If $|V_5| = 5$ and $|V_1| + |V_2| + |V_3| \geq 4$, then we analogously arrive at a contradiction. Altogether, we see that there remain to consider the following five partition-sequences.

Subcase 1.1. Let $(n_i) = 1, 1, 1, 1, 5$. In this case, we see that $d^+(x) = d^-(x) = 4$, if $x \in V_1 \cup V_2 \cup V_3 \cup V_4$ and $d^+(x) = d^-(x) = 2$, if $x \in V_5$, that means $i_t(D) = 0$. Since $\frac{|V(D)|-|V_4|-2|V_5|+6}{2} = 2 > i_t(D)$, it remains to consider the case that $Q_1 = \emptyset$ in Theorem 3.1. If $Y = V_5$, then it follows that $|Z| \leq 3$ and thus $|Q_2| \geq 1$. This yields $d^-(x) \geq 5$ for all $x \in Q_2$, a contradiction to $d^-(x) \leq 4$ for all $x \in V(D)$. If $|Y| = 4$, then it follows that $|Z| \leq 2$ and thus $|Q_2| \geq 2$. This implies that there is an arc $pq \in E(D[Q_2])$. Since $d^-(q) \geq 5$, we arrive at a contradiction. If $|Y| \leq |V_5| - 2$, then $|V(D)| - 3|V_5| + 8 = 2 > i_t(D)$ contradicts Theorem 3.1 with $t = 2$.

Subcase 1.2. Let $(n_i) = 1, 1, 1, 2, 5$. Since $i_g(D) = 2$, this is impossible.

Subcase 1.3. Let $(n_i) = 1, 1, 1, 3, 5$. This yields $d^+(x) = d^-(x) = 5$, if $x \in V_1 \cup V_2 \cup V_3$, $d^+(x) = d^-(x) = 4$ or $\{d^+(x), d^-(x)\} = \{3, 5\}$, if $x \in V_4$ and $d^+(x) = d^-(x) = 3$, if $x \in V_5$.

Firstly, we assume that $Q_1 = \emptyset$. If $Y = V_5$, then we conclude that $|Z| \leq 3$ and thus $|Q_2| \geq 3$. Since $d^+(x) = d^-(x) = 3$ for all $x \in V_5$, it follows that $|Z| = |Q_2| = 3$ and $Z \rightarrow Y$. It is obvious that there are

either in Q_2 or in Z two vertices of different partite sets. Hence, there is an arc $p \rightarrow q$ that is either in $E(D[Q_2])$ or in $E(D[Z])$. If $pq \in E(D[Q_2])$, then $d^-(q) \geq 6$, and if $pq \in E(D[Z])$, then $d^+(p) \geq 6$, in both cases a contradiction. If $|Y| = 4$, then we see that $|Z| \leq 2$ and thus $|Q_2| \geq 4$, a contradiction to $d^+(x) = 3$ for all $x \in V_5$. Hence, let $|Y| \leq |V_5| - 2$. But now, $|V(D)| - 3|V_5| + 8 = 4 > i_t(D)$ contradicts Theorem 3.1.

Consequently, it remains to consider the case that $Q_1 \neq \emptyset$ and $Q_2 \neq \emptyset$. Note that $\frac{|V(D)| - |V_4| - 2|V_5| + 6}{2} = 2$. According to the proof of Theorem 3.1 the inequalities in the last inequality-chain of Case 3 have to be equalities, that means especially $t = 0$ and $|Y| = |Z| + 2$. This implies $Y = V_5$ and $|Z| = |Y| - 2 = 3$. Hence, we conclude that $|Q| = 3$ and, without loss of generality, let $|Q_1| = 1$ and $|Q_2| = 2$. If there is an arc leading from Q_1 to Q_2 and $q_1 \in Q_1$, then it follows that $d^+(q_1) \geq |Y| + 1 = 6$, a contradiction. Since $Q_1 \rightsquigarrow Q_2$, we obtain $Q = V_4$ and thus $Z = V_1 \cup V_2 \cup V_3$. Let $D' = D[V_1 \cup V_2 \cup V_3]$. Because of $Q_2 \rightarrow Z \rightarrow Q_1$, we have on the one hand that

$$\begin{aligned} 15 &= \sum_{x \in V(D')} d^+(x) = d(V(D'), Q_1) + \sum_{x \in V(D')} d_{D'}^+(x) + d(V(D'), Y) \\ &= 3 + 3 + d(V(D'), Y), \end{aligned}$$

that means $d(V(D'), Y) = 9$. Since $Q_1 \rightarrow Y \rightarrow Q_2$, we observe on the other hand that

$$15 = \sum_{y \in Y} d^-(y) = d(Q_1, Y) + d(V(D'), Y) = 5 + d(V(D'), Y),$$

that means $d(V(D'), Y) = 10$, a contradiction.

Subcase 1.4. Let $(n_i) = 1, 1, 1, 4, 5$. Since $i_g(D) = 2$, this is impossible.

Subcase 1.5. Let $(n_i) = 1, 1, 1, 5, 5$. This yields that $d^+(x) = d^-(x) = 6$, if $x \in V_1 \cup V_2 \cup V_3$ and $d^+(x) = d^-(x) = 4$, if $x \in V_4 \cup V_5$, that means $i_t(D) = 0$. Since $|V(D)| - 3|V_5| + 3 = 1 > i_t(D)$ and $\frac{|V(D)| - |V_4| - 2|V_5| + 5}{2} = \frac{3}{2} > i_t(D)$, this contradicts Theorem 1.3.

Case 2. Let $c = 4$. If $r \geq 3$, then $4 \geq \max\{2, \frac{3}{r} + 3\}$ and Corollary 4.2 yields that D contains a Hamiltonian path, a contradiction. If $r = 2$ and $|V_4| \leq 5$ or $r = 2$, $|V_4| = 6$ and $|V_1| + |V_2| \geq 5$, then $\frac{|V(D)| - |V_3| - 2|V_4| + 5}{2} \geq 2 = i_g(D)$ leads to a contradiction to Corollary 3.2. If $r = 1$ and $|V_4| \leq 3$ or $r = 1$, $|V_4| = 4$ and $|V_1| + |V_2| \geq 3$ or $r = 1$, $|V_4| = 5$ and $|V_1| + |V_2| \geq 4$, then $\frac{|V_1| + |V_2| - |V_4| + 5}{2} \geq 2 = i_g(D)$, a contradiction to Corollary 3.2.

Summarizing our results, we see that, according to the assertion of this theorem, there remain to treat 14 partition-sequences.

Subcase 2.1. Since $i_g(D) = 2$, the partition-sequences 1, 1, 1, 5; 1, 1, 3, 5; 1, 1, 5, 5; 1, 2, 2, 5; 1, 2, 4, 5; 2, 2, 3, 6 and 2, 2, 5, 6 are impossible.

Subcase 2.2. Let $(n_i) = 1, 1, 2, 5$. In this case, we obtain that $d^+(x) = d^-(x) = 4$, if $x \in V_1 \cup V_2$, $\{d^+(x), d^-(x)\} = \{3, 4\}$, if $x \in V_3$, and $d^+(x) = d^-(x) = 2$, if $x \in V_5$, that means $i_i(D) = 1$. Since $\frac{|V_1|+|V_2|-|V_4|+6}{2} = \frac{3}{2} > i_i(D)$, it remains to consider the case that $Q_1 = \emptyset$ in Theorem 3.1. If $Y = V_4$, then we conclude that $|Z| \leq 3$ and thus $|Q_2| \geq 1$, a contradiction to $d^-(x) \leq 4$ for all $x \in V(D)$. If $|Y| = 4$, then it follows that $|Z| \leq 2$ and $|Q_2| \geq 2$. This implies that there is an arc $p \rightarrow q$ that is either in $E(D[Q_2])$ or in $E(D[Z])$. Since $d^-(x) = 2$ for all $x \in Y$, we conclude that $|Z| = 2$ and $Z \rightarrow Y$. If $pq \in E(D[Q_2])$, then $d^-(q) \geq 5$ and if $pq \in E(D[Z])$, then $d^+(p) \geq 5$, in both cases a contradiction. Hence, let $|Y| \leq 3 = |V_4| - 2$. Then because of $|V(D)| - 3|V_4| + 8 = 2 > i_i(D)$ we arrive at a contradiction to Theorem 3.1.

Subcase 2.3. Let $(n_i) = 1, 1, 4, 5$. This yields that $d^+(x) = d^-(x) = 5$, if $x \in V_1 \cup V_2$, $\{d^+(x), d^-(x)\} = \{3, 4\}$, if $x \in V_3$, and $d^+(x) = d^-(x) = 3$, if $x \in V_5$, that means $i_i(D) = 1$. Since $\frac{|V_1|+|V_2|-|V_4|+6}{2} = \frac{3}{2} > i_i(D)$, it remains to consider the case that $Q_1 = \emptyset$ in Theorem 3.1. If $Y = V_4$, then it follows that $|Z| \leq 3$ and thus $|Q_2| \geq 3$. This implies that there is a vertex $q_2 \in Q_2 \cap V_3$, since $Y \rightarrow Q_2$ a contradiction to $d^-(x) \leq 4$ for all $x \in V_3$. Consequently, let $|Y| \leq 4$. In this case, the fact that $|V(D)| - 3|V_4| + 6 = 2 > i_i(D)$ contradicts Theorem 3.1.

Subcase 2.4. Let $(n_i) = 1, 2, 5, 5$. This yields that $d^+(x) = d^-(x) = 6$, if $x \in V_1$, $\{d^+(x), d^-(x)\} = \{5, 6\}$, if $x \in V_2$, and $d^+(x) = d^-(x) = 4$, if $x \in V_3 \cup V_4$, that means $i_i(D) = 1$. Since $\frac{|V_1|+|V_2|-|V_4|+5}{2} = \frac{3}{2} > i_i(D)$ and $|V(D)| - 3|V_4| + 3 = 1 = i_i(D)$, we have a contradiction to Theorem 1.3.

Subcase 2.5. Let $(n_i) = 1, 1, 1, 4$. In this case, we observe that $d^+(x) = d^-(x) = 3$, if $x \in V_1 \cup V_2 \cup V_3$, and $\{d^+(x), d^-(x)\} = \{1, 2\}$, if $x \in V_4$, that means $i_i(D) = 1$. Because of $\frac{|V_1|+|V_2|-|V_4|+6}{2} = 2 > i_i(D)$, it remains to consider the case that $Q_1 = \emptyset$ in Theorem 3.1. If $Y = V_4$, then we conclude that $|Z| \leq 2$ and thus $|Q_2| \geq 1$, a contradiction to $d^-(x) \leq 3$ for all $x \in V(D)$. If $|Y| = 3$, then $|Z| \leq 1$ and $|Q_2| \geq 2$. Hence, there exists an arc $p \rightarrow q$ with $p, q \in Q_2$. Since $d^-(q) \geq 4$, we arrive at a contradiction. Consequently, let $|Y| \leq 2$. In this case, $|V(D)| - 3|V_4| + 8 = 3 > i_i(D)$ contradicts Theorem 3.1.

Subcase 2.6. Let $(n_i) = 1, 1, 4, 4$. This yields $\{d^+(x), d^-(x)\} = \{4, 5\}$, if $x \in V_1 \cup V_2$ and $d^+(x) = d^-(x) = 3$, if $x \in V_3 \cup V_4$, that means $i_i(D) = 1$. Since $\frac{|V_1|+|V_2|-|V_4|+5}{2} = \frac{3}{2} > i_i(D)$ and $|V(D)| - 3|V_4| + 3 = 1 = i_i(D)$, we arrive at a contradiction to Theorem 1.3.

Subcase 2.7. Let $(n_i) = 2, 2, 2, 6$. In this case, we observe that $d^+(x) = d^-(x) = 5$, if $x \in V_1 \cup V_2 \cup V_3$, and $d^+(x) = d^-(x) = 3$, if $x \in V_4$, that means $i_i(D) = 0$. Because of $\frac{|V_1|+|V_2|-|V_4|+6}{2} = 2 > i_i(D)$, it remains to consider the case that $Q_1 = \emptyset$ in Theorem 3.1. If $Y = V_4$, then we conclude that $|Z| \leq 4$ and thus $|Q_2| \geq 2$. Since $Y \rightarrow Q_2$, we have a contradiction to $d^-(x) \leq 5$ for all $x \in V(D)$. If $|Y| = 5$, then it follows $|Z| \leq 3$ and $|Q_2| \geq 3$. This yields the existence of an arc $p \rightarrow q$ with $p, q \in Q_2$. Hence, $d^-(q) \geq 6$, a contradiction. Consequently, let $|Y| \leq 4$. But now $|V(D)| - 3|V_4| + 8 = 2 > i_i(D)$ contradicts Theorem 3.1.

Subcase 2.8. Let $(n_i) = 2, 2, 6, 6$. This implies that $d^+(x) = d^-(x) = 7$ for all $x \in V_1 \cup V_2$ and $d^+(x) = d^-(x) = 5$ for all $x \in V_3 \cup V_4$, that means $i_i(D) = 0$. Since $\frac{|V_1|+|V_2|-|V_4|+6}{2} = \frac{3}{2} > i_i(D)$ and $|V(D)| - 3|V_4| + 3 = 1 > i_i(D)$, this is a contradiction to Theorem 1.3. \square

For the partition-sequences 1, 1, 2, 4; 2, 2, 4, 6; 1, 2, 3, 5 and 1, 1, 3, 4 there are multipartite tournaments that do not have any Hamiltonian path as the following four examples demonstrate.

Example 4.5 Let D be a 4-partite tournament with the partite sets $V_1 = \{u\}$, $V_2 = \{v\}$, $V_3 = \{x_1, x_2\}$ and $V_4 = \{y_1, y_2, y_3, y_4\}$ such that $(V_1 \cup V_2) \rightarrow x_1 \rightarrow V_4 \rightarrow x_2 \rightarrow (V_1 \cup V_2)$ and $u \rightarrow v \rightarrow y_4 \rightarrow u \rightarrow y_1 \rightarrow v \rightarrow y_3 \rightarrow u \rightarrow y_2 \rightarrow v$ (see Figure 2). Then $i_g(D) = 2$ and D has the partition-sequence 1, 1, 2, 4 but no Hamiltonian path.

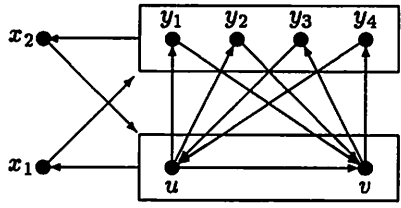


Figure 2: A 4-partite tournament D with $i_g(D) = 2$ and the partition-sequence 1, 1, 2, 4 that does not contain a Hamiltonian path

Example 4.6 Let D be a 4-partite tournament with the partite sets $V_1 = \{u_1, u_2\}$, $V_2 = \{v_1, v_2\}$, $V_3 = \{x_1, x_2, x_3, x_4\}$ and $V_4 = \{y_1, y_2, y_3, y_4, y_5, y_6\}$ such that $V_4 \rightarrow \{x_3, x_4\} \rightarrow (V_1 \cup V_2) \rightarrow \{x_1, x_2\} \rightarrow V_4$, $u_1 \rightarrow v_1 \rightarrow u_2 \rightarrow v_2 \rightarrow u_1$ and $\{y_1, y_2, y_3\} \rightarrow \{u_2, v_2\} \rightarrow \{y_4, y_5, y_6\} \rightarrow \{u_1, v_1\} \rightarrow \{y_1, y_2, y_3\}$ (see Figure 3). Then D is a 4-partite tournament with $i_g(D) =$

2 and the partition-sequence 2, 2, 4, 6 that does not contain any Hamiltonian path.

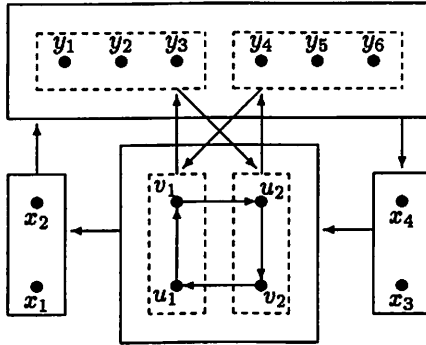


Figure 3: A 4-partite tournament D with $i_g(D) = 2$ and the partition-sequence 2, 2, 4, 6 that does not contain a Hamiltonian path

Example 4.7 Let D be a 4-partite tournament with the partite sets $V_1 = \{u\}$, $V_2 = \{v_1, v_2\}$, $V_3 = \{x_1, x_2, x_3\}$ and $V_4 = \{y_1, y_2, y_3, y_4, y_5\}$ such that $V_4 \rightarrow \{x_1, x_2\} \rightarrow (V_1 \cup V_2) \rightarrow x_3 \rightarrow V_4$, $V_2 \rightarrow V_1$, $\{v_2, u\} \rightarrow \{y_1, y_2\} \rightarrow v_1$, $\{v_1, u\} \rightarrow \{y_3, y_4\} \rightarrow v_2$ and $V_2 \rightarrow y_5 \rightarrow V_1$ (see Figure 4). Then D is a 4-partite tournament with $i_g(D) = 2$ and the partition-sequence 1, 2, 3, 5 that does not contain a Hamiltonian path.

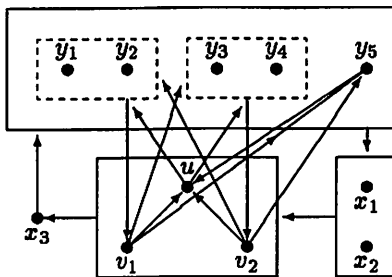


Figure 4: A 4-partite tournament D with $i_g(D) = 2$ and the partition-sequence 1, 2, 3, 5 that does not contain a Hamiltonian path

Example 4.8 Let D be a 4-partite tournament with the partite sets $V_1 = \{u\}$, $V_2 = \{v\}$, $V_3 = \{x_1, x_2, x_3\}$ and $V_4 = \{y_1, y_2, y_3, y_4\}$ such that $V_4 \rightarrow \{x_1, x_2\} \rightarrow (V_1 \cup V_2) \rightarrow x_3 \rightarrow V_4$, $u \rightarrow v \rightarrow \{y_3, y_4\} \rightarrow u$, $u \rightarrow \{y_1, y_2\}$ and $y_1 \rightarrow v \rightarrow y_2$ (see Figure 5). Then D is a 4-partite tournament with $i_g(D) = 2$ and the partition-sequence 1, 1, 3, 4 that does not contain any Hamiltonian path.

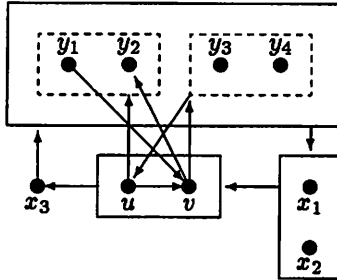


Figure 5: A 4-partite tournament D with $i_g(D) = 2$ and the partition-sequence 1, 1, 3, 4 that does not contain a Hamiltonian path

In the case that $c = 2$ or $c = 3$ and $i_g(D) \geq 2$, there are infinitely many digraphs D that do not contain any Hamiltonian path as we can see in the following example.

Example 4.9 Let D be a bipartite tournament with $i_g(D) \geq 2$. If V_1, V_2 are the partite sets of D such that $|V_1| + 2 \leq |V_2| \leq |V_1| + 2i_g(D)$, then clearly, D does not contain a Hamiltonian path.

Now, let D be a 3-partite tournament with the partite sets V_1, V_2, V_3 . If $|V_1| = |V_2| = r$ and $|V_3| = r + i_g(D)$ with $i_g(D) \geq 2$ such that $V_1 \rightarrow V_2 \rightarrow V_3 \rightarrow V_1$, then D does not contain a Hamiltonian path.

All the presented examples show that Theorem 4.4 is best possible. Combining Corollary 4.2 with Example 4.9, we observe the following.

Corollary 4.10 Let $i \geq 2$ be an arbitrary integer. Then all, except a finite number, of c -partite tournaments with $i_g \leq i$ and $c \geq 4$ have a Hamiltonian path. Furthermore the bound $c \geq 4$ is best possible.

Proof. If $r \geq 4i - 5$, then Corollary 4.2 yields that a c -partite tournament D with $c \geq 4$ and at least r vertices in each partite set contains a Hamiltonian path. Because of Lemma 2.1 there are only finitely many c -partite tournaments with $i_g \leq i$ and at most $4i - 6$ vertices in one partite set. Thus, the first part of this corollary is proved.

Example 4.9 demonstrates that there are infinitely many 3-partite tournaments with $i_g \geq 2$ that do not contain any Hamiltonian path and the proof of this corollary is complete. \square

An interesting extension to our results would be the solution of the following open problem.

Problem 4.11 *For all i find the smallest value, $g(i)$, with the property that all c -partite tournaments with $i_g \leq i$ and $c \geq g(i)$ have a Hamiltonian path.*

According to the Theorems 1.1, 4.3 and 4.4, it is already shown that $g(0) = 1$, $g(1) = 3$ and $g(2) = 5$.

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