#### MINIMAL EQUITABILITY OF HAIRY CYCLES

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ABSTRACT. Every labeling of the vertices of a graph with distinct natural numbers induces a natural labeling of its edges: the label of an edge (x,y) is the absolute value of the difference of the labels of x and y. By analogy with graceful labelings, we say that a labeling of the vertices of a graph of order n is minimally k-equitable if the vertices are labelled with  $1,2,\ldots,n$  and in the induced labeling of its edges every label either occurs exactly k times or does not occur at all. For  $m \geq 3$ , let  $C'_m$  (denoted also in the literature by  $C_m \circ K_1$  and called a corona graph) be a graph with 2m vertices such that there is a partition of them into sets U and V of cardinality m, with the property that U spans a cycle, V is independent and the edges joining U to V form a matching. Let P be the set of all pairs (m,k) of positive integers such that k is a proper divisor of 2m (i.e., a divisor different from 2m and 1) and k is odd if m is odd. We show that  $C'_m$  is minimally k-equitable if and only if  $(m,k) \in \mathcal{P}$ .

#### 1. Introduction

A labeling of a graph G is an assignment of distinct natural numbers to the vertices of G. Every labeling induces a natural labeling of the edges: the label of an edge (x,y) is the absolute value of the difference of the labels of x and y. Bloom [4] defined a labeling of a graph to be k-equitable if in the induced labeling of its edges, every label occurs exactly k times, if at all. Furthermore, a k-equitable labeling of a graph of order n is said to be minimal if the vertices are labelled with  $1, 2, \ldots, n$ . A graph is minimally k-equitable if it has a minimal k-equitable labeling.

The notion of minimally k-equitable labelings is a generalization of special labelings of trees called graceful. A labeling of a tree is graceful if and only if it is minimally 1-equitable. (Note that the induced labels of the edges must be then  $1, 2, \ldots, n-1$ .) Graceful labelings were defined by Rosa [6] in connection with a famous and very difficult conjecture of Ringel and Kotzig (see Ringel [5]) concerning decompositions of complete graphs with odd number of vertices into subgraphs isomorphic to trees. The open conjecture that every tree has a graceful labeling implies the conjecture of Ringel and Kotzig.

The notion of minimal k-equitability was first applied to cycles. Let  $C_n$  be the cycle on n vertices. Given natural numbers n and k,  $n \geq 3$ , it is easy to see that if the cycle  $C_n$  is minimally k-equitable, then k is a proper divisor of n (that is, k divides n,  $k \neq n$ , and  $k \neq 1$ ). Answering a question posed by Bloom [4], Wojciechowski [7] proved that this necessary condition is also sufficient.

Barrientos, Dejter and Hevia [3] proved a number of results concerning k-equitability of forests.

Another class of graphs to which the concept of minimal k-equitability was recently applied is the class of graphs whose vertices could be partitioned into two sets U and V such that U induces a cycle, V induces no edges, and the edges between U and V form a matching (in particular, U and V are of the same cardinality). We will call such graphs hairy cycles and denote them by  $C'_m$  (m is the cardinality of U).

There is a general construction that, given graphs G and H, produces the *corona* graph  $G \circ H$ . The hairy cycle  $C'_m$  is obtained by applying that construction to the cycle  $C_m$  and the graph  $K_1$  consisting of a single vertex. Therefore  $C'_m$  is the corona graph  $C_m \circ K_1$ .

Again it is easy to see that given natural numbers m and k,  $m \geq 3$ , if the hairy cycle  $C'_m$  is minimally k-equitable, then k is a proper divisor of 2m which is the number of vertices of  $C'_m$ . Going a little further, we get the following result.

**Proposition 1.1.** Let  $m \geq 3$  and  $k \geq 1$  be integers. If m is odd and the hairy cycle  $C'_m$  is minimally k-equitable, then k is also odd.

*Proof.* Assume that  $m \geq 3$  is arbitrary and the hairy cycle  $C'_m$  is minimally k-equitable. Let U and V be the sets that form a partition of the vertex set of  $C'_m$  such that U spans a cycle, V spans no edges, and the edges between U and V form a matching. Let E be the set of edges of the cycle spanned by U and F be the set of the edges of the matching. Let

$$g:U\cup V \to \{1,2,\ldots,2m\}$$

be any labeling of  $C'_m$ , and  $h: E \cup F \to \mathbb{N}$  be the induced labeling of the edges of  $C'_m$ . It is clear that  $\sum_{e \in E} h(e)$  is even.

We claim that

(1) 
$$\sum_{e \in F} h(e) \equiv m \mod 2.$$

If h(e) = 1 for every  $e \in F$ , then  $\sum_{e \in F} h(e) = m$  so (1) holds. Otherwise, there are vertices  $u \in U$  and  $v \in V$  such that  $|g(u) - g(v)| \ge 2$ . Let w, y be vertices of  $C'_m$  such that g(w) is between g(u) and g(v) in the standard ordering of integers, and y is joined to w by an edge of F. Let  $g': U \cup V \to \mathbb{N}$  be the labeling of  $C'_m$  obtained from g by exchanging min  $\{g(w), g(y)\}$  with

 $\max\{g(u),g(v)\}$ . For example, if

$$g(u) < g(w) < g(v) < g(y),$$

then we exchange the labels of w and v.

Let  $h': E \cup F \to \mathbb{N}$  be the labeling of the edges of  $C'_m$  induced by g'. Then

$$\sum_{e \in F} h'(e) < \sum_{e \in F} h(e)$$

and

$$\sum_{e \in F} h'(e) \equiv \sum_{e \in F} h(e) \bmod 2.$$

Therefore, induction on  $\sum_{e \in F} h(e)$  can be used to prove that (1) holds in general.

If m is odd and g is a minimal k-equitable labeling, then  $\sum_{e \in E \cup F} h(e)$  is odd and divisible by k implying that k is odd.

Let  $\mathcal{P}$  be the set of all pairs (m,k) of positive integers such that  $m \geq 3$ , k is a proper divisor of 2m and k is odd if m is odd. Then for the graph  $C'_m$  to be minimally k-equitable it is necessary that  $(m,k) \in \mathcal{P}$ . Acharya and Bhat-Nayak [1] [2] proved that the condition  $(m,k) \in \mathcal{P}$  is sufficient for  $C'_m$  to be minimally k-equitable when  $k \in \{3,4\}$ . We will prove that this condition is sufficient in general.

**Theorem 1.2.** Let  $m \geq 3$  and  $k \geq 2$  be integers. If  $(m, k) \in \mathcal{P}$ , then the graph  $C'_m$  is minimally k-equitable.

We shall call a graph G an integer graph if its vertex set is a finite subset of  $\mathbb{N}$ . If e = (u, v) is an edge of G, then we will say that e has length |u - v|. Observe that a finite graph H has a k-equitable labeling if and only if it is isomorphic to an integer graph G with either 0 or k edges of any length. We will call such G a k-equitable representation of the graph H. Note also that a finite graph H is minimally k-equitable if it has a k-equitable representation G whose vertices are consecutive integers. We will call such G a minimal k-equitable representation of H.

In the following proofs, to show that  $C'_m$  is k-equitable, we will construct an integer graph  $G^k_m$  that will be a minimal k-equitable representation of  $C'_m$ . In the included figures of the graphs  $G^k_m$  the vertices will be placed on the real line, the edges that are part of the cycle (cycle edges) will be marked by thick lines below the real line and the edges of the matching (matching edges) above the real line.

The proof of Theorem 1.2 will be split into several lemmas. The proofs of these lemmas will be given in the remaining sections. Assume that  $(m,k) \in \mathcal{P}$ . We have then two possibilities: either k divides m or it does not divide m. In section 2 we will present the starting point of our construction in the case when k divides m, namely, we will show that  $C'_m$  is minimally

k-equitable when m=k or m=2k. The proof that  $C'_{\ell k}$  is minimally k-equitable for  $\ell \geq 3$  will be given in section 4 for even k and in section 5 for odd k. If k does not divide m, then the definition of  $\mathcal P$  implies that  $k\equiv 0 \mod 4$  and  $m=\ell k/2$  for some odd integer  $\ell \geq 3$ . Let  $k\equiv 0 \mod 4$ . We will prove in section 3 that the graph  $C'_{3k/2}$  is minimally k-equitable. The proof that  $C'_{\ell k/2}$  is minimally k-equitable for odd  $\ell \geq 5$  will be given in section 4.

### 2. Minimal k-equitability of $C'_k$ and $C'_{2k}$

**Lemma 2.1.** Let  $k \geq 3$  be an integer. The graph  $C'_k$  is minimally k-equitable.

*Proof.* Let  $G_k^k$  be the integer graph with the vertex set  $\{1, 2, ..., 2k\}$  and the edge set consisting of all the edges listed below:

$$\bullet$$
 (1,2), (2,3), ...,  $(k, k+1)$ 

— the above edges are k edges of length 1, the first of which is a matching edge, and all the others are cycle edges;

• 
$$(2, k+1), (3, k+2), \ldots, (k+1, 2k)$$

— the above edges are k edges of length k-1, the first of which is a cycle edge, and all the others are matching edges.

Figure 1 shows the graph  $G_3^3$  and Figure 2 shows the graph  $G_k^k$  for arbitrary  $k \geq 3$ .

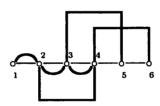


FIGURE 1. The graph  $G_3^3$ 

It is clear that  $G_k^k$  is a minimal k-equitable representation of  $C_k'$  implying that  $C_k'$  is minimally k-equitable.

**Lemma 2.2.** Let  $k \geq 3$  be an odd integer. The graph  $C'_{2k}$  is minimally k-equitable.

*Proof.* Let  $G_{2k}^k$  be the integer graph with the vertex set  $\{1, 2, ..., 4k\}$  and the edge set consisting of all the edges listed below:

• 
$$(1, \frac{k+3}{2}), (2, \frac{k+5}{2}), \ldots, (\frac{k+1}{2}, k+1)$$

— the edges listed above are  $\frac{k+1}{2}$  edges of length  $\frac{k+1}{2}$ , all of which are matching edges;

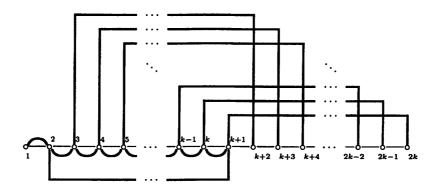


FIGURE 2. The graph  $G_k^k$ 

• 
$$(k+2,\frac{3k+5}{2}), (k+3,\frac{3k+7}{2}), \ldots, (\frac{3k+1}{2},2k+1)$$

— the edges listed above are  $\frac{k-1}{2}$  edges of length  $\frac{k+1}{2}$ , the last of which is a cycle edge, and all the others are matching edges;

• 
$$(\frac{k+3}{2}, \frac{k+5}{2}), (\frac{k+7}{2}, \frac{k+9}{2}), \dots, (\frac{3k+1}{2}, \frac{3k+3}{2})$$

— these are  $\frac{k+1}{2}$  edges of length 1, the last of which is a matching edge, and all the others are cycle edges;

• 
$$(2k+2,2k+3)$$
,  $(2k+4,2k+5)$ , ..., $(3k-1,3k)$ 

— these are  $\frac{k-1}{2}$  edges of length 1, all of which are cycle edges;

• 
$$(\frac{k+3}{2}, 2k+1), (\frac{k+5}{2}, 2k+2), \ldots, (\frac{3k+1}{2}, 3k)$$

— these are k edges of length  $\frac{3k-1}{2}$ , all of which are cycle edges;

• 
$$(2k+1,3k+1), (2k+2,3k+2), \ldots, (3k,4k)$$

— these are k edges of length k, all of which are matching edges.

Figures 3, 4, and 5 show the graphs  $G_6^3$ ,  $G_{10}^5$ , and  $G_{14}^7$  respectively; figures 6 and 7 show the general graph  $G_{2k}^k$  for arbitrary  $k \equiv 1 \mod 4$ , and  $k \equiv 3 \mod 4$  respectively.

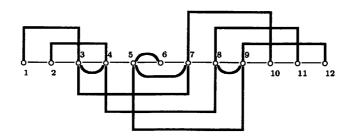


FIGURE 3. The graph  $G_6^3$ 

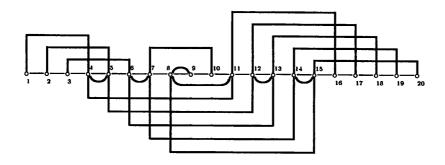


FIGURE 4. The graph  $G_{10}^5$ 

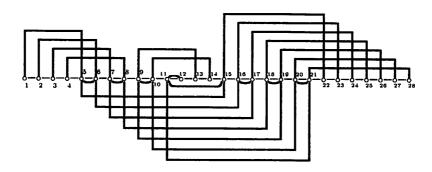


FIGURE 5. The graph  $G_{14}^7$ 

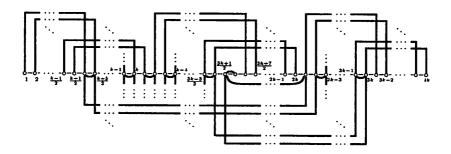


FIGURE 6. The graph  $G_{2k}^k$  when  $k \equiv 1 \mod 4$ 

It is clear that  $G_{2k}^k$  is a minimal k-equitable representation of  $C_{2k}'$  implying that  $C_{2k}'$  is minimally k-equitable.

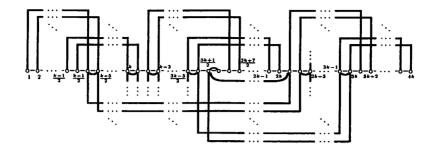


FIGURE 7. The graph  $G_{2k}^k$  when  $k \equiv 3 \mod 4$ 

**Lemma 2.3.** Let  $k \geq 2$  be an even integer. The graph  $C'_{2k}$  is minimally k-equitable.

*Proof.* Let  $G_4^2$  be the integer graph with the vertex set  $\{1, 2, ..., 8\}$  and the edge set consisting of the following edges (see Figure 8):

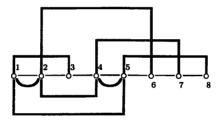


FIGURE 8. The graph  $G_4^2$ 

- (1,2), (4,5) 2 edges of length 1 both of them cycle edges;
- (1,3), (2,4) 2 edges of length 2 the first a matching edge, the second a cycle edge;
- (4,7), (5,8) 2 edges of length 3 both of which being matching edges;
- (1,5), (2,6) 2 edges of length 4 the first a cycle edge, the second a matching edge.

For  $k \geq 4$  let  $G_{2k}^k$  be the integer graph with the vertex set  $\{1, 2, \ldots, 4k\}$  and the edge set consisting of all the edges listed below:

- $(1, \frac{k+2}{2}), (2, \frac{k+4}{2}), \ldots, (\frac{k}{2}, k)$
- these are  $\frac{k}{2}$  edges of length  $\frac{k}{2}$ , all of which are matching edges;
  - $(k+1,\frac{3k+2}{2}), (k+2,\frac{3k+4}{2}), \ldots, (\frac{3k}{2},2k)$
- these are  $\frac{k}{2}$  edges of length  $\frac{k}{2}$ , the last of which is a cycle edge, and all the others are matching edges;

• 
$$(\frac{k+2}{2}, \frac{k+4}{2}), (\frac{k+6}{2}, \frac{k+8}{2}), \ldots, (\frac{3k-2}{2}, \frac{3k}{2})$$

— these are  $\frac{k}{2}$  edges of length 1, all of which are cycle edges;

• 
$$(2k, 2k+1), (2k+2, 2k+3), \ldots, (3k-2, 3k-1)$$

— these are  $\frac{k}{2}$  edges of length 1, all of which are cycle edges;

• 
$$(\frac{k+2}{2}, 2k+1), (\frac{k+4}{2}, 2k+2), \ldots, (\frac{3k}{2}, 3k)$$

— these are k edges of length  $\frac{3k}{2}$ , the last of which is a matching edge, and all the others are cycle edges;

• 
$$(2k, 3k+1), (2k+1, 3k+2), \ldots, (3k-1, 4k)$$

— these are k edges of length k+1, all of which are matching edges.

Figures 9 and 10 show the graphs  $G_8^4$  and  $G_{12}^6$  respectively; Figures 11 and 12 show the general graph  $G_{2k}^k$  for arbitrary  $k \equiv 0 \mod 4$ , and  $k \equiv 2 \mod 4$  respectively.

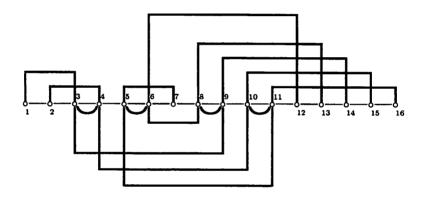


FIGURE 9. The graph  $G_8^4$ 

It is clear that, for every even  $k \geq 2$ , the graph  $G_{2k}^k$  is a minimal k-equitable representation of  $C_{2k}'$ , implying that  $C_{2k}'$  is minimally k-equitable.

## 3. Minimal k-equitability of $C_{3k/2}'$ when $k \equiv 0 \mod 4$

**Lemma 3.1.** Let  $k \geq 4$  be an integer such that  $k \equiv 0 \mod 4$ . The graph  $C'_{3k/2}$  is minimally k-equitable.

*Proof.* Let  $G_{3k/2}^k$  be the integer graph with the vertex set  $\{1, 2, ..., 3k\}$  and the edge set consisting of all the edges listed below:

• 
$$(1, \frac{k+2}{2}), (2, \frac{k+4}{2}), \ldots, (\frac{k-2}{2}, k-1)$$

— these are  $\frac{k-2}{2}$  edges of length  $\frac{k}{2}$ , all of which are matching edges;

• 
$$(\frac{3k+4}{2}, 2k+2), (\frac{3k+6}{2}, 2k+3), \dots, (2k-1, \frac{5k-2}{2})$$

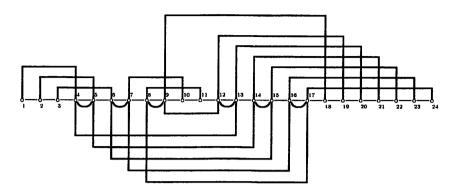


FIGURE 10. The graph  $G_{12}^6$ 

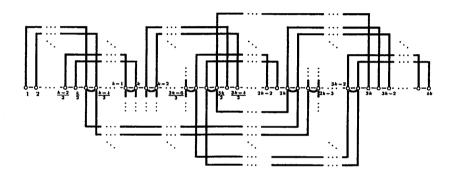


FIGURE 11. The graph  $G_{2k}^k$  when  $k \equiv 0 \mod 4$ 

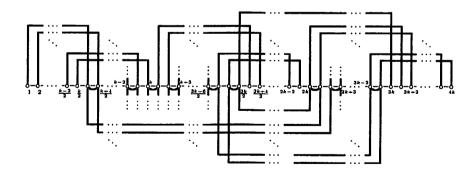


FIGURE 12. The graph  $G_{2k}^k$  when  $k \equiv 2 \mod 4$ 

— these are  $\frac{k-4}{2}$  edges of length  $\frac{k}{2}$ , all of which are matching edges, (they appear only when  $k \geq 8$ );

• 
$$(k+1,\frac{3k+2}{2}), (\frac{3k}{2},2k), (2k+1,\frac{5k+2}{2})$$

— these are 3 edges of length  $\frac{k}{2}$ , the first and the third of which are matching edges, and the second is a cycle edge;

• 
$$(\frac{k+2}{2}, \frac{k+4}{2}), (\frac{k+6}{2}, \frac{k+8}{2}), \ldots, (k-1,k)$$

— these are  $\frac{k}{4}$  edges of length 1, all of which are cycle edges;

• 
$$(k, k+1), (k+1, k+2), \ldots, (\frac{3k-2}{2}, \frac{3k}{2})$$

— these are  $\frac{k}{2}$  edges of length 1, all of which are cycle edges;

• 
$$(2k, 2k+1), (2k+2, 2k+3), \ldots, (\frac{5k-4}{2}, \frac{5k-2}{2})$$

— these are  $\frac{k}{4}$  edges of length 1, all of which are cycle edges;

• 
$$\left(\frac{k}{2},2k\right),\left(\frac{k+2}{2},2k+1\right),\ldots,\left(k,\frac{5k}{2}\right)$$

— these are  $\frac{k+2}{2}$  edges of length  $\frac{3k}{2}$ , the first and the last of which are matching edges, and all the others are cycle edges;

• 
$$(k+2,\frac{5k+4}{2}), (k+3,\frac{5k+6}{2}), \ldots, (\frac{3k}{2},3k)$$

— these are  $\frac{k-2}{2}$  edges of length  $\frac{3k}{2}$ , all of which are matching edges.

Figures 13, 14 and 15 show the graphs  $G_6^4$ ,  $G_{12}^8$ , and  $G_{18}^{12}$  respectively; figure 16 shows the general graph  $G_{3k/2}^k$  for arbitrary  $k \equiv 0 \mod 4$ .

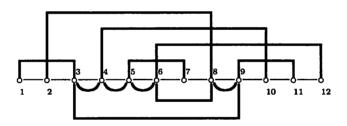


FIGURE 13. The graph  $G_6^4$ 

It is clear that, for every  $k \equiv 0 \mod 4$ , the graph  $G^k_{3k/2}$  is a minimal k-equitable representation of  $C'_{3k/2}$ , implying that  $C'_{3k/2}$  is minimally k-equitable.  $\blacksquare$ 

# 4. Minimal k-equitability of $C_m'$ when k is even

Assume that  $k \geq 2$  is even. We are going to show that  $C'_m$  is minimally k-equitable for every m such that  $(m, k) \in \mathcal{P}$ .

Given a minimal k-equitable representation G of  $C'_m$ , we say that a set S of edges of G is a k-socket if the following conditions are satisfied:

(1) S consists of  $\frac{k}{2}$  cycle edges of length 1 whose endpoints form a set of consecutive k integers  $s, s+1, \ldots, s+k-1$ ;

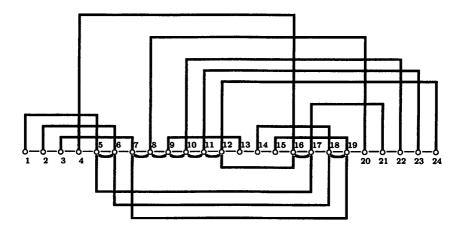


FIGURE 14. The graph  $G_{12}^8$ 

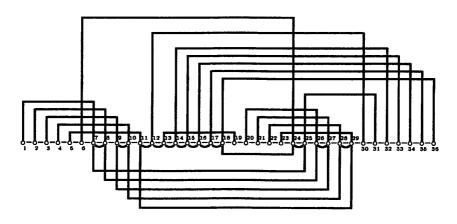


FIGURE 15. The graph  $G_{18}^{12}$ 

(2) if a is the smallest integer in the vertex set of G and b is the largest integer in the vertex set of G, then either b-s+1 or s+k-a is not the length of any edge of G.

Note that if the set of endpoints of the edges of S consists of either the largest k vertices of G or the smallest k vertices of G, then the second condition above is satisfied.

A minimal k-equitable representation of  $C'_m$  with a k-socket will be called a k-proper representation of  $C'_m$ .

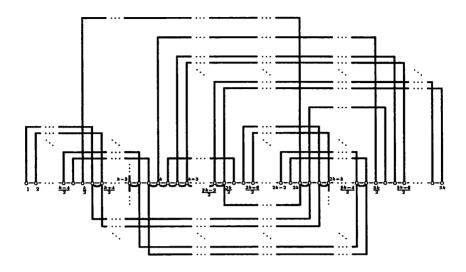


FIGURE 16. The graph  $G_{3k/2}^k$  when  $k \equiv 0 \mod 4$ 

**Lemma 4.1.** The graph  $G_{2k}^k$  is a k-proper representation of  $C'_{2k}$  for every even  $k \geq 2$ .

*Proof.* Let  $k \geq 2$  be an even integer. It follows from the proof of Lemma 2.3 that  $G_{2k}^k$  is a minimally k-equitable representation of  $C_{2k}'$ . If k=2, then let  $\mathcal{S}=\{(4,5)\}$ . (See Figure 8.) With b=8 and s=4 the integer b-s+1=5 is not a length of any edge of  $G_4^2$ . Thus  $\mathcal{S}$  is a 2-socket in  $G_4^2$  implying that  $G_4^2$  is a 2-proper representation of  $C_4'$ .

If  $k \geq 4$ , then let (see Figures 11 and 12)

$$S = \left\{ \left( \frac{k+2}{2}, \frac{k+4}{2} \right), \left( \frac{k+6}{2}, \frac{k+8}{2} \right), \dots, \left( \frac{3k-2}{2}, \frac{3k}{2} \right) \right\}.$$

With b=4k and  $s=\frac{k+2}{2}$ , the integer  $b-s+1=\frac{7k}{2}$  is not a length of any edge of  $G_{2k}^k$ , implying that  $G_{2k}^k$  is a k-proper representation of  $G_{2k}'$ .

**Lemma 4.2.** The graph  $G_{3k/2}^k$  is a k-proper representation of  $C_{3k/2}'$  for every integer  $k \geq 4$  such that  $k \equiv 0 \mod 4$ .

*Proof.* Let  $k \ge 4$  be an integer with  $k \equiv 0 \mod 4$ . It follows from the proof of Lemma 3.1 that  $G_{3k/2}^k$  is a minimally k-equitable representation of  $C_{3k/2}'$ . Let (see Figure 16)

$$S = \left\{ \left( \frac{k+2}{2}, \frac{k+4}{2} \right), \left( \frac{k+6}{2}, \frac{k+8}{2} \right), \dots, \left( \frac{3k-6}{2}, \frac{3k-4}{2} \right), \left( \frac{3k-2}{2}, \frac{3k}{2} \right) \right\}$$

With b=3k and  $s=\frac{k+2}{2}$ , the integer  $b-s+1=\frac{5k}{2}$  is not a length of any edge of  $G_{3k/2}^k$ , implying that  $G_{3k/2}^k$  is a k-proper representation of  $G_{3k/2}^k$ .

**Lemma 4.3.** If k is even and there is a k-proper representation of  $C'_m$ , then there is a k-proper representation of  $C'_{m+k}$ .

**Proof.** Let G be a k-proper representation of  $C'_m$  with a k-socket S and let s be the smallest integer in the set of endpoints of the edges in S. Let a be the smallest integer which is a vertex of G and let b be the largest integer which is a vertex of G. Let H be the graph obtained from G by performing the following operations:

- remove the edges of S;
- add 2 sets of k vertices each at both ends of the graph G, namely add the vertices a-k, a-k+1, ..., a-1 and b+1, b+2, ..., b+k;
- add k edges of length b-a+k+1 matching the new vertices, namely add the edges (a-k,b+1), (a-k+1,b+2), ..., (a-1,b+k);

Case 1: if G has no edges of length b-s+1, then

- add k edges of length b-s+1 joining the endpoints of the edges in S to the new vertices whose value is larger than b, namely add the edges  $(s, b+1), (s+1, b+2), \ldots, (s+k-1, b+k)$ ;
- add  $\frac{k}{2}$  edges of length 1 that form a matching of the set of all the new vertices whose value is larger than b, namely add the edges  $(b+1,b+2), (b+3,b+4), \ldots, (b+k-1,b+k)$ .

Figure 17 shows the new edges of the graph H in case 1.

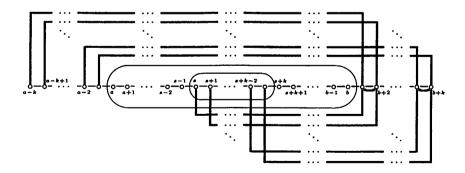


FIGURE 17. The new edges of the graph H in case 1

Case 2: if G has some edge of length b-s+1, (so that it has no edges of length s+k-a), then

- add k edges of length s+k-a joining the endpoints of the edges in S to the new vertices whose value is smaller than a, namely add the edges (a-k,s), (a-k+1,s+1), ..., (a-1,s+k-1);
- add  $\frac{k}{2}$  edges of length 1 that form a matching of the set of the new vertices whose value is smaller than a, namely add the edges  $(a-k, a-k+1), (a-k+2, a-k+3), \ldots, (a-2, a-1)$ .

Figure 18 shows the new edges of the graph H in case 2. Figures 19 and 20 show the results of applying the above construction to the graphs  $G_4^2$  and  $G_6^4$  respectively.

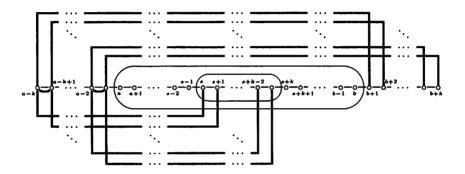


FIGURE 18. The new edges of the graph H in case 2

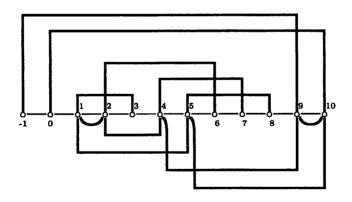


FIGURE 19. A 2-proper representation of  $C_6^\prime$  obtained from  $G_4^2$ 

Since each edge of S in G is replaced by a path of length 3 in H, the cycle of G gives rise to a cycle of length m+k in H. Moreover, the new k edges of length b-a+k+1 in H are matching edges implying that H is isomorphic

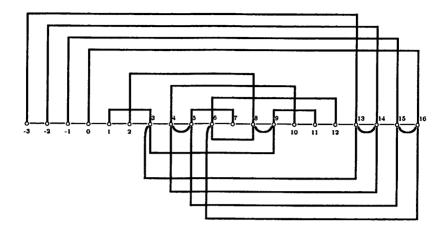


FIGURE 20. A 2-proper representation of  $C'_{10}$  obtained from  $G_6^4$ 

to  $C'_{m+k}$ . It is clear that H is a minimally k-equitable representation of  $C'_{m+k}$ . The set of the new edges of length 1 in H is a k-socket since the set of endpoints of these edges consists of either the k vertices of H having the smallest label or the k vertices of H having the largest label. Thus H is a k-proper representation of  $C'_{m+k}$ .

**Proposition 4.4.** Let  $(m, k) \in \mathcal{P}$  with k being even. Then  $C'_m$  is minimally k-equitable.

*Proof.* Assume that  $m = \ell k$  for some integer  $\ell$ . The minimal k-equitability of  $C'_m$  follows from Lemma 2.1 when  $\ell = 1$ , and from Lemma 2.3 when  $\ell = 2$ . If  $\ell \geq 3$ , then the minimal k-equitability of  $C'_m$  follows by induction on  $\ell$  using Lemmas 4.1 and 4.3.

If k is not a divisor of m, then  $k \equiv 0 \mod 4$  and  $m = \ell k/2$  for some odd integer  $\ell \geq 3$ . The minimal k-equitability of  $C'_m$  follows from Lemma 3.1 when  $\ell = 3$ , and follows by induction using Lemmas 4.2 and 4.3 when  $\ell \geq 5$ .

### 5. Minimal k-equitability of $C_m'$ when k is odd

Assume that  $k \geq 3$  is odd. We are going to show that  $C'_m$  is minimally k-equitable for every m such that  $(m, k) \in \mathcal{P}$ .

If e = (u, v) is an edge of an integer graph and u < v, then we say that u is the *left endpoint* of e and v is the *right endpoint* of e. Given a minimal k-equitable representation G of  $C'_m$ , we say that a pair  $\mathbb{T} = (T, \mathcal{C})$  is a k-thread in G if the following conditions are satisfied:

- (1) C consists of k edges that are cycle edges and have the same length;
- (2) T is either the set of all left endpoints or the set of all right endpoints of the edges in C and it consists of k consecutive integers t, t+1, ..., t+k-1;
- (3) if a is the vertex of G with the smallest label, b is the vertex of G with the largest label, s is the other endpoint of the edge in C that has t as one of its endpoints,  $A_G$  is the set of lengths of the edges of G, and  $R_{G,\mathbb{T}}$ ,  $L_{G,\mathbb{T}}$ ,  $W_G$  are infinite sets of integers defined as follows:

$$\begin{split} R_{G,\mathbb{T}} &= \{b-s+1\} \cup \{b+2ik-t+1: i \geq 0\} \\ &\quad \cup \{t-a+2ik: i \geq 1\}, \\ L_{G,\mathbb{T}} &= \{s-a+k\} \cup \{b+(2i+1)k-t+1: i \geq 0\} \\ &\quad \cup \{t-a+(2i+1)k: i \geq 0\}, \\ W_G &= \{b-a+ik+1: i \geq 1\}, \end{split}$$

then either  $L_{G,\mathbb{T}}$  or  $R_{G,\mathbb{T}}$  is disjoint with  $A_G \cup W_G$ .

The following three figures illustrate the definition of a k-thread in a graph G when t < s; the pictures require obvious modifications when s < t. Figure 21 shows the edges of the graph G that belong to the set C. Figure 22 shows the relationship between the sets  $R_{G,\mathbb{T}}$ ,  $W_G$  and the vertices of the graph G, where the integers in  $R_{G,\mathbb{T}}$  and  $W_G$  are represented by edges of the corresponding lengths. The set  $R_{G,\mathbb{T}}$  is represented above the line containing the vertices of G and the set  $W_G$  below it. Figure 23 shows analogous relationship between the sets  $L_{G,\mathbb{T}}$ ,  $W_G$  and the graph G.

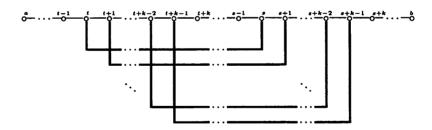


FIGURE 21. Edges of the graph G that belong to C

A minimal k-equitable representation of  $C'_m$  with a k-thread will be called a k-proper representation of  $C'_m$ .

**Lemma 5.1.** The graph  $G_{2k}^k$  is a k-proper representation of  $C'_{2k}$  for every odd  $k \geq 3$ .

*Proof.* Let  $k \geq 3$  be an odd integer. It follows from the proof of Lemma 2.2 that  $G_{2k}^k$  is a minimally k-equitable representation of  $C_{2k}'$ .

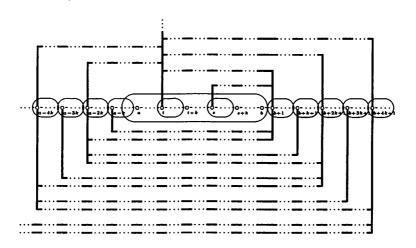


Figure 22. Edges whose length is in  $R_{G,\mathbb{T}}$  or  $W_G$ 

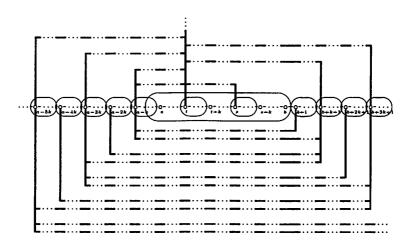


Figure 23. Edges whose length is in  $L_{G,\mathbb{T}}$  or  $W_G$ 

Let (see Figures 6 and 7)

$$\mathcal{C} = \left\{ \left(\frac{k+3}{2}, 2k+1\right), \left(\frac{k+5}{2}, 2k+2\right), \dots, \left(\frac{3k+1}{2}, 3k\right) \right\},$$

and

$$T = \left\{ \frac{k+3}{2}, \frac{k+5}{2}, \dots, \frac{3k+1}{2} \right\}.$$

With  $\mathbb{T}=(T,\mathcal{C}), a=1, b=4k, t=\frac{k+3}{2}$  and s=2k+1, we have

$$R_{G,\mathbb{T}} = \left\{2k\right\} \cup \left\{\frac{\left(4i-1\right)k-1}{2}: i \geq 2\right\} \cup \left\{\frac{\left(4i+1\right)k+1}{2}: i \geq 1\right\}.$$

Since

$$W_G = \{ik : i \geq 5\},\,$$

and

$$A_G = \left\{1, \frac{k+1}{2}, k, \frac{3k-1}{2}\right\},\,$$

it is clear that

$$R_{G,\mathbf{T}}\cap (A_G\cup W_G)=\emptyset.$$

It follows that  $\mathbb{T}$  is a k-thread in  $G_{2k}^k$ , implying that  $G_{2k}^k$  is a k-proper representation of  $G_{2k}^{\prime}$ .

**Lemma 5.2.** If k is odd and there is a k-proper representation of  $C'_m$ , then there is a k-proper representation of  $C'_{m+k}$ .

*Proof.* Let G be a k-proper representation of  $C'_m$  with a k-thread  $\mathbb{T} = (T, \mathcal{C})$ , let t be the smallest label of T, let s be the other endpoint of the edge in  $\mathcal{C}$  that has t as one of its endpoints, and let

$$S = \{s, s+1, \ldots, s+k-1\}.$$

Let a be the smallest label of the vertices of G and let b be the largest label of the vertices of G. Let H be the graph obtained from G by performing the following operations:

- remove the edges of C;
- add the vertices a-k, a-k+1, ..., a-1 and b+1, b+2, ..., b+k;
- add k edges of length b-a+k+1 matching the new vertices, namely add the edges (a-k,b+1), (a-k+1,b+2), ..., (a-1,b+k);

Case 1: if  $R_{G,\mathbb{T}} \cap (A_G \cup W_G) = \emptyset$ , then

- add k edges of length b-s+1 joining the vertices in S to the new vertices whose label is larger than b, namely add the edges  $(s,b+1), (s+1,b+2), \ldots, (s+k-1,b+k)$ ;
- add k edges of length b-t+1 joining the vertices in T to the new vertices whose label is larger than b, namely add the edges (t, b+1),  $(t+1, b+2), \ldots, (t+k-1, b+k)$ .

Figure 24 shows the new edges of the graph H in case 1.

Case 2: if  $R_{G,\mathbb{T}} \cap (A_G \cup W_G) \neq \emptyset$ , (so that  $L_{G,\mathbb{T}} \cap (A_G \cup W_G) = \emptyset$ ) then

• add k edges of length s + k - a joining the vertices in S to the new vertices whose label is smaller than a, namely add the edges  $(a - k, s), (a - k + 1, s + 1), \ldots, (a - 1, s + k - 1);$ 

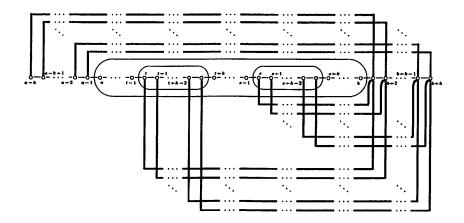


FIGURE 24. The new edges of the graph H in case 1

• add k edges of length t+k-a joining the vertices in T to the new vertices whose label is smaller than a, namely add the edges  $(a-k,t), (a-k+1,t+1), \ldots, (a-1,t+k-1)$ .

Figure 25 shows the new edges of the graph H in case 2. Figure 26 shows the results of applying the above construction to the graph  $G_6^3$ .

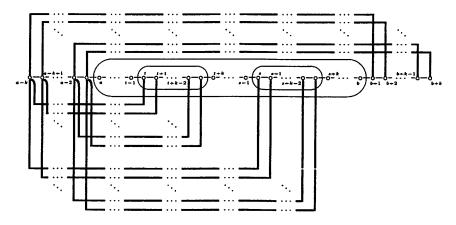


FIGURE 25. The new edges of the graph H in case 2

Since each edge of  $\mathcal C$  in G is replaced by a path of length 2 in H, the cycle of G gives rise to a cycle of length m+k in H. Moreover, the new k edges of length b-a+k+1 in H are matching edges, implying that H is isomorphic to  $C'_{m+k}$ . Since

$$b-a+k+1>b-a,$$

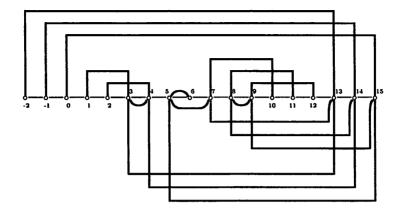


FIGURE 26. A 3-proper representation of  $C_9'$  obtained from  $G_8^3$ 

there are no edges in G of length b-a+k+1. Moreover the lengths of the new cycle edges of H are in  $L_{G,\mathbb{T}}$  in case 1 and in  $R_{G,\mathbb{T}}$  in case 2, implying that they are not in  $A_G$ . Therefore, the graph H is a minimally k-equitable representation of  $C'_{m+k}$ . To prove that H is a k-proper representation of  $C'_{m+k}$  it remains to show that there is a k-thread in H.

In case 1, let C' be the set of the new edges of H of length b-t+1 that join the vertices in T to the new vertices whose label is larger than b. We claim that the pair  $\mathbb{T}' = (T, C')$  is a k-thread in the graph H. Let a' = a - k, b' = b + k, and let s' = b + 1 be the other endpoint of the edge of C' that has t as one of its endpoints. We have then

$$\begin{array}{rcl} L_{H,\mathbb{T}'} &=& \{s'-a'+k\} \cup \{b'+(2i+1)\,k-t+1: i \geq 0\} \\ && \cup \{t-a'+(2i+1)\,k: i \geq 0\} \\ &=& \{b-a+2k+1\} \cup \{b+2ik-t+1: i \geq 1\} \\ && \cup \{t-a+2ik: i \geq 1\} \\ &=& R_{G,\mathbb{T}} \cup \{b-a+2k+1\} \setminus \{b-s+1, b-t+1\} \,. \end{array}$$

Moreover

$$A_H = A_G \cup \{b-a+k+1, b-s+1, b-t+1\} \setminus \{|t-s|\},\,$$

and

$$W_H = \{b' - a' + ik + 1 : i \ge 1\}$$

$$= \{b - a + ik + 1 : i \ge 3\}$$

$$= W_G \setminus \{b - a + k + 1, b - a + 2k + 1\},$$

SO

$$A_H \cup W_H = (A_G \cup W_G) \cup \{b-s+1, b-t+1\} \setminus \{b-a+2k+1, |t-s|\}.$$
 Since  $R_{G,\mathbb{T}} \cap (A_G \cup W_G) = \emptyset$ , it follows that

$$L_{H,\mathbf{T}'}\cap (A_H\cup W_H)=\emptyset,$$

so the pair  $\mathbb{T}' = (T, \mathcal{C}')$  is a k-thread in the graph H.

In case 2, let C' be the set of the new edges of H of length t+k-a that join the vertices in T to the new vertices whose label is smaller than a. We claim that the pair  $\mathbb{T}' = (T, C')$  is a k-thread in the graph H. Let a' = a - k, b' = b + k, and let s' = a - k be the other endpoint of the edge of C' that has t as one of its endpoints. Then

$$\begin{array}{lll} R_{H,\mathbb{T}'} & = & \{b'-s'+1\} \cup \{b'+2ik-t+1: i \geq 0\} \cup \{t-a'+2ik: i \geq 1\} \\ \\ & = & \{b-a+2k+1\} \cup \{b+(2i+1)k-t+1: i \geq 0\} \\ \\ & \cup \{t-a+(2i+1)k: i \geq 1\} \\ \\ & = & L_{G,\mathbb{T}} \cup \{b-a+2k+1\} \setminus \{s+k-a,t+k-a\} \,. \end{array}$$

Moreover

$$A_H = A_G \cup \{b-a+k+1, s+k-a, t+k-a\} \setminus \{|t-s|\},$$

and

$$W_H = W_G \setminus \{b-a+k+1, b-a+2k+1\}$$

so

$$A_H \cup W_H = (A_G \cup W_G) \cup \{s+k-a,t+k-a\} \setminus \{b-a+2k+1,|t-s|\}.$$
 Since  $L_{G,\mathbb{T}} \cap (A_G \cup W_G) = \emptyset$ , it follows that

$$R_{H,\mathbb{T}'}\cap (A_H\cup W_H)=\emptyset,$$

so the pair  $\mathbb{T}' = (T, \mathcal{C}')$  is a k-thread in the graph H.

Thus H is a k-proper representation of  $C'_{m+k}$ .

**Proof of Theorem 1.2.** Assume the  $(m,k) \in \mathcal{P}$ . If k is even then it follows from Proposition 4.4 that  $C'_m$  is k-equitable. If k is odd, then it is a divisor of m. The k-equitability of  $C'_m$  follows from Lemma 2.1 if k=m, and it follows by induction using Lemmas 5.1 and 5.2 if k is a proper divisor of m.

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