

Ramsey $(K_{1,2}, C_4)$ -minimal Graphs

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Abstract. For any given graphs G and H , we write $F \rightarrow (G, H)$ to mean that any red-blue coloring of the edges of F contains a red copy of G or a blue copy of H . Graph F is (G, H) -minimal (Ramsey-minimal) if $F \rightarrow (G, H)$ but $F^* \not\rightarrow (G, H)$ for any proper subgraph $F^* \subset F$. The class of all (G, H) -minimal graphs is denoted by $\mathcal{R}(G, H)$. In this paper we will determine the graphs in $\mathcal{R}(K_{1,2}, C_4)$.

1 Introduction

We consider simple graphs, namely finite undirected graphs without loops and multiple edges. Let $G = (V, E)$. We say that G contains H if G contains a subgraph isomorphic to H . The subgraph of G isomorphic to C_4 is defined as a *basic cycle*.

Let G and H be graphs. We say that $F \rightarrow (G, H)$ if any red-blue coloring of the edges of F contains a red copy of G or a blue copy of H . Graph F is (G, H) -minimal (Ramsey-minimal) if $F \rightarrow (G, H)$ but $F^* \not\rightarrow (G, H)$ for any proper subgraph $F^* \subset F$. The class of all (G, H) -minimal graphs is denoted by $\mathcal{R}(G, H)$. In general we follow the terminology in [7]. Recent results on Ramsey numbers can be found in [11].

Here are some previous results dealing with the problem of finding graphs in $\mathcal{R}(G, H)$. Burr, Erdős and Lovász [6] proved that $\mathcal{R}(2K_2, 2K_2) = \{3K_2, C_5\}$ and $\mathcal{R}(K_{1,2}, K_{1,2}) = \{K_{1,3}, C_{2n+1}\}$ for $n \geq 1$. Furthermore, Burr et.al [5] determined all graphs in $\mathcal{R}(2K_2, K_3)$. Later, Burr et.al [4] showed that $\mathcal{R}(K_{1,m}, K_{1,n}) = \{K_{1,m+n-1}\}$ for odd m and n .

In [10] Mengersen and Oeckermann characterized all graphs in $\mathcal{R}(2K_2, K_{1,n})$ for $n \geq 3$. Another result is the characterization of all graphs in $\mathcal{R}(K_{1,2}, K_{1,m})$ for $m \geq 3$ by Borowiecki et.al in [2]. Then Borowiecki et.al [3] determined all graphs in $\mathcal{R}(K_{1,2}, K_3)$. Recently, Baskoro et.al [1] showed that $W_{3t+1} \in \mathcal{R}(P_3, C_3^t)$, where W_{3t+1} is a wheel with $6t+2$ edges and C_3^t is a windmill graph, i.e. a graph obtained by connecting a vertex c (called a hub) to all vertices of tK_2 .

The problem of characterizing pairs of graphs (G, H) for which the set $\mathcal{R}(G, H)$ is finite or infinite has also been investigated in numerous papers. In particular, all pairs of two forests for which the set $\mathcal{R}(G, H)$ is finite are specified in a theorem of Faudree [8]. Then Luczak [9] stated that if G is a forest other than a matching and H is a graph containing at least one cycle then $\mathcal{R}(G, H)$ is infinite. It follows that the set $\mathcal{R}(K_{1,2}, C_4)$ is infinite. In this paper we will determine the graphs in $\mathcal{R}(K_{1,2}, C_4)$.

2 Some classes of graphs

We define some classes of graphs needed to prove our main results.

Let k be positive integer, $k \geq 2$. A graph G with

$$\begin{aligned} V(G) &= \{w_i \mid 1 \leq i \leq k\} \cup \{u_i \mid 1 \leq i \leq k\} \cup \{v_j \mid 1 \leq j \leq k+1\}, \\ E(G) &= \{v_i w_i \mid 1 \leq i \leq k\} \cup \{v_i u_i \mid 1 \leq i \leq k\} \cup \{u_i v_{i+1} \mid 1 \leq i \leq k\} \\ &\quad \cup \{w_i v_{i+1} \mid 1 \leq i \leq k\} \end{aligned}$$

is called the C_4 -path. We define vertices v_1 and v_{k+1} as the *end vertices* of the C_4 -path.

A C_4 -cycle is constructed by identifying two end vertices of C_4 -path. The *length* of a C_4 -path (a C_4 -cycle) is the number of basic cycles in the C_4 -path (the C_4 -cycle).

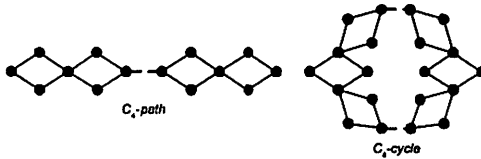


Fig. 1. C_4 -path and C_4 -cycle

We use C_4 -path and C_4 -cycle (Figure 1), graphs L_1 and L_2 (Figure 2) to define some classes of graphs below.

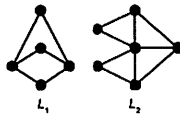


Fig. 2. L_1 with root x and L_2 with root y

Let \mathcal{A} be a family of graphs which contains

- (1) A_1, A_2 and A_3 (Figure 3),

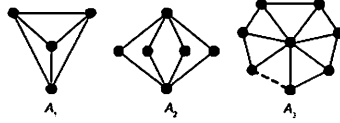


Fig. 3. A_1, A_2 and A_3

- (2) $A_4(k)$, $k \geq 0$. This graph consists of two copies of L_1 with a C_4 -path of length k joining two roots of L_1 (if $k = 0$ then we have two copies of L_1 with a common root),
- (3) $A_5(k)$, $k \geq 0$. This graph consists of a L_1 and a L_2 with a C_4 -path of length k joining the root of L_1 and the root of L_2 ,
- (4) $A_6(k)$, $k \geq 0$. This graph consists of two copies of L_2 with a C_4 -path of length k joining two roots of L_2 (for $k = 0$ we have two copies of L_2 with a common root).

Let \mathcal{B} be a family of graphs which contains

- (1) $B_1(k, t)$, $k \geq 0$, $t \geq 2$. This graph consists of a C_4 -cycle of length t and a L_1 with a C_4 -path of length k joining the root of L_1 and an arbitrary vertex of degree 4 of C_4 -cycle,
- (2) $B_2(k, t)$, $k \geq 0$, $t \geq 2$. This graph consists of a C_4 -cycle of length t and a L_2 with a C_4 -path of length k joining the root of L_2 and an arbitrary vertex of degree 4 of C_4 -cycle,
- (3) $B_3(k, t)$, $k \geq 0$, $t \geq 2$. This graph consists of a C_4 -cycle of length t and a L_1 with a C_4 -path of length k joining the root of L_1 and an arbitrary vertex of degree 2 of C_4 -cycle,
- (4) $B_4(k, t)$, $k \geq 0$, $t \geq 2$. This graph consists of a C_4 -cycle of length t and a L_2 with a C_4 -path of length k joining the root of L_2 and an arbitrary vertex of degree 2 of C_4 -cycle.

Let \mathcal{D} be a family of graphs which contains

- (1) $D_1(k, t_1, t_2)$, $k \geq 0$, $t_1 \geq 2$, $t_2 \geq 2$. This graph is constructed by joining two copies of C_4 -cycles of length t_1 and t_2 and a C_4 -path of length k . One of the end vertices of C_4 -path is identified with an arbitrary vertex of degree 4 of the first C_4 -cycle and the other end vertex is identified with an arbitrary vertex of degree 4 of the second C_4 -cycle,
- (2) $D_2(k, t_1, t_2)$, $k \geq 0$, $t_1 \geq 2$, $t_2 \geq 2$. This graph is constructed by joining two copies of C_4 -cycles of length t_1 and t_2 and a C_4 -path of length k . One of the end vertices of C_4 -path is identified with an arbitrary vertex of degree 4 of the first C_4 -cycle and the other end vertex is identified with an arbitrary vertex of degree 2 of the second C_4 -cycle,

- (3) $D_3(k, t_1, t_2)$, $k \geq 0$, $t_1 \geq 2$, $t_2 \geq 2$. This graph is constructed by joining two copies of C_4 -cycles of length t_1 and t_2 and a C_4 -path of length k . One of the end vertices of C_4 -path is identified with an arbitrary vertex of degree 2 of the first C_4 -cycle and the other end vertex is identified with an arbitrary vertex of degree 2 of the second C_4 -cycle.

Let $\mathcal{I} = \{I(t, k_1, k_2), t \geq 2, k_1 \geq 0, k_2 \geq 0\}$ where $I(t, k_1, k_2)$ is constructed by joining a C_4 -cycle of length t with two copies of C_4 -paths of length k_1 and k_2 . Two of the end vertices of the C_4 -paths are identified with two arbitrary vertices of C_4 -cycle (the C_4 -paths are attached to different vertices of C_4 -cycle).

3 Main Results

The *distance* $d(u, v)$ between two (not necessary distinct) vertices u and v in a graph G is the length of a shortest path between them. When u and v are identical, their distance is 0. When u and v are unreachable from each other, their distance is defined to be ∞ . The *diameter* of G , $diam(G)$, is the greatest distance between any two vertices in that graph. In Theorems 1, 2 and 3 we present a collection of graphs that belongs to $\mathcal{R}(K_{1,2}, C_4)$.

Theorem 1. *If $\mathcal{R}_1 = \{F \in \mathcal{R}(K_{1,2}, C_4) \mid diam(F) = 1\}$ then $\mathcal{R}_1 = \{A_1\}$.*

Proof. It can be easily seen that $A_1 \simeq K_4$. Consider any red-blue coloring of A_1 that implies an edge-decomposition $A_1 = A_{11} \oplus A_{12}$. Let $A_{11} \not\supseteq K_{1,2}$. Then A_{11} contains at most one edge. Therefore $A_{12} \supseteq C_4$. Consequently $A_1 \rightarrow (K_{1,2}, C_4)$.

To prove the minimality of A_1 , let $V(A_1) = \{v_i \mid 1 \leq i \leq 4\}$. Consider $A_1^* \simeq A_1 \setminus \{e\}$ for any fixed edge $e \in E(A_1)$. Without loss of generality, we assume that $e = v_1 v_2$. Then the edges of A_1^* can be partitioned into two classes, namely E_1 and E_2 , with $E_1 = \{v_2 v_4, v_1 v_3\}$ and $E_2 = \{v_1 v_4, v_2 v_3, v_3 v_4\}$ such that $A_1^*[E_1] \not\supseteq K_{1,2}$ and $A_1^*[E_2] \not\supseteq C_4$. Thus $A_1^* \not\rightarrow (K_{1,2}, C_4)$. Therefore $A_1 \in \mathcal{R}(K_{1,2}, C_4)$.

Since F has diameter 1, F must be a complete graph. Furthermore for $n \geq 5$, K_n always contain A_1 . Thus K_n cannot be a $(K_{1,2}, C_4)$ -minimal for $n \geq 5$. \square

Theorem 2. *If $\mathcal{R}_2 = \{F \in \mathcal{R}(K_{1,2}, C_4) \mid diam(F) = 2\}$ then $\{A_2, A_3\} \subseteq \mathcal{R}_2$.*

Proof. Case 2.1 Graph A_2 .

To show that $A_2 \in \mathcal{R}_2$, we consider any red-blue coloring of A_2 that implies an edge-decomposition $A_2 = A_{21} \oplus A_{22}$. Let $A_{21} \not\supseteq K_{1,2}$. Thus A_{21} consists of at most two disjoint edges. Therefore, $A_{22} \supseteq C_4$. It follows that $A_2 \rightarrow (K_{1,2}, C_4)$.

To prove the minimality of A_2 , let $V(A_2) = \{v_i \mid 1 \leq i \leq 4\} \cup \{u_j \mid 1 \leq j \leq 2\}$ such that v_i is the vertex of degree 2 and u_j is the vertex of degree 4. Consider $A_2^* \simeq A_2 \setminus \{e\}$ for any fixed edge $e \in E(A_2)$. W.l.o.g we assume that $e = v_1 u_1$. Then the edges of A_2^* can be partitioned into two classes, namely E_1 and E_2 , with $E_1 = \{v_2 u_1, v_4 u_2\}$ and $E_2 = \{v_1 u_2, v_2 u_2, v_3 u_2, v_3 u_1, v_4 u_1\}$ such that $A_2^*[E_1] \not\supseteq K_{1,2}$ and $A_2^*[E_2] \not\supseteq C_4$. Thus $A_2^* \not\rightarrow (K_{1,2}, C_4)$. Therefore $A_2 \in \mathcal{R}_2(K_{1,2}, C_4)$.

Case 2.2 Graph A_3 .

A_3 is a wheel W_{2n+1} with odd number of spokes. We define

$$\begin{aligned} V(W_{2n+1}) &= \{c\} \cup \{v_i \mid 1 \leq i \leq 2n+1\}, \\ E(W_{2n+1}) &= E_1 \cup E_2, \text{ where} \\ E_1 &= \{cv_i \mid 1 \leq i \leq 2n+1\} \text{ and} \\ E_2 &= \{v_j v_{j+1} \mid 1 \leq j \leq 2n\} \cup \{v_{2n+1} v_1\}. \end{aligned}$$

To show that $A_3 \in \mathcal{R}_2$, we consider any red-blue coloring of A_3 that implies an edge-decomposition $A_3 = A_{31} \oplus A_{32}$. Let $A_{31} \not\supseteq K_{1,2}$. Thus A_{31} contains at most one spoke. If A_{31} contains one spoke, say cv_i for some $i, 1 \leq i \leq 2n+1$, then $\{c, v_{i-1}, v_i, v_{i+1}\}$ forms a blue C_4 . In case that A_{31} contains no spoke, the parity of $|C_{2n+1}|$ in W_{2n+1} forces two incident edges in C_{2n+1} to be in A_{32} . Therefore, $A_{32} \supseteq C_4$. It follows that $A_3 \rightarrow (K_{1,2}, C_4)$.

Let us consider the graph $A_3^* \simeq W_{2n+1} \setminus \{e\}$ for any fixed edge $e \in E(W_{2n+1})$. Consider the following two cases.

Case 2.2.1 $e \in E_1$.

W.l.o.g let $e = cv_2$. If $n = 1$ then color cv_2 by red and other edges by blue. Therefore, $W_3 \setminus \{cv_1\}$ has no red $K_{1,2}$ neither blue C_4 . If $n \geq 2$ then the edges of $A_3^*(k)$ can be partitioned into two classes, namely E_3 and E_4 , as follows

$$\begin{aligned} E_3 &= \{cv_2\} \cup \{v_{2s} v_{2s+1} \mid 2 \leq s \leq n\}, \\ E_4 &= (E_1 \setminus \{cv_2\}) \cup \{v_{2n+1} v_1, v_1 v_2, v_2 v_3, v_3 v_4\} \cup \{v_{2t-1} v_{2t} \mid 3 \leq t \leq n\}. \end{aligned}$$

Color the edges of E_3 by red and E_4 by blue. Under this coloring $A_3^*[E_3] \not\supseteq K_{1,2}$ and $A_3^*[E_4] \not\supseteq C_4$.

Case 2.2.2 $e \in E_2$.

W.l.o.g let $e = v_1 v_2$. The edges of $A_3^*(k)$ can be partitioned into two classes, namely E_5 and E_6 , as follows

$$\begin{aligned} E_5 &= \{cv_1\} \cup \{v_{2p} v_{2p+1} \mid 1 \leq p \leq n\}, \\ E_6 &= (E_1 \setminus \{cv_1\}) \cup \{v_{2n+1} v_1\} \cup \{v_{2q-1} v_{2q} \mid 2 \leq q \leq n\}. \end{aligned}$$

Color the edges of E_5 by red and E_6 by blue. Consequently, $A_3^*[E_5] \not\supseteq K_{1,2}$ and $A_3^*[E_6] \not\supseteq C_4$. Thus we have that $A_3^* \rightarrow (K_{1,2}, C_4)$. Therefore $A_3 \in \mathcal{R}_2(K_{1,2}, C_4)$. \square

Theorem 3. *If $\mathcal{R}_3 = \{F \in \mathcal{R}(K_{1,2}, C_4) \mid \text{diam}(F) \geq 4\}$ then*

$$\mathcal{R}_3 \supseteq \{A_4(k), A_5(k), A_6(k), B_1(k, t), B_2(k, t), D_1(k, t_1, t_2)\}.$$

Proof. Case 3.1 Graph $A_4(k)$.

Consider any red-blue coloring of $A_4(k)$ that implies an edge-decomposition $A_4(k) = A_{41}(k) \oplus A_{42}(k)$. Let $A_{41}(k) \not\supseteq K_{1,2}$. Observe that if there exists a blue C_4 in the C_4 -path that belongs to $A_{42}(k)$ then the proof is complete. So now

we assume that in each basic cycle in C_4 -path there is an edge belongs to $A_{41}(k)$. We call such an edge as a *red edge*.

Define the *end cycles* as the basic cycles that are attached to L_1 . Thus the end vertices of the C_4 -path belong to the end cycles. We claim that one of the end vertices, say v_1 , is incident to a red edge. The claim is easy to justify by noting that if the end vertex is not incident to that red edge then the condition $A_{41}(k) \not\supseteq K_{1,2}$ forces the other end vertex, say v_{k+1} , to be incident to a red edge (See Figure 4). Then there is a blue C_4 in one of the L_1 s. It follows that $A_{42}(k) \supseteq C_4$. Thus $A_4(k) \rightarrow (K_{1,2}, C_4)$.

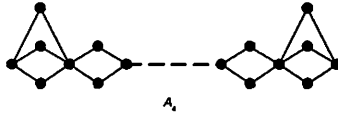


Fig. 4. $A_4(k)$

We denote the set of vertices of the L_1 s by $\{x_i \mid 1 \leq i \leq 4\} \cup \{y_j \mid 1 \leq j \leq 6\}$ where x_i is the vertex of degree 3 and y_j is the vertex of degree 2. The root of the first L_1 , x_2 , is identified with v_1 and the root of the second L_1 , x_3 , is identified with v_{k+1} .

To prove the minimality of $A_4(k)$, let us consider the graph $A_4^* \simeq A_4 \setminus \{e\}$ for any fixed edge $e \in E(A_4)$. Consider the following two cases.

Case 3.1.1 e is an edge of one of the L_1 s.

W.l.o.g assume that e is in the first L_1 s, say $e = x_1y_1$. The edges of $A_4^*(k)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{x_1y_2, x_3y_5, x_4y_4\} \cup \{v_iw_i \mid 1 \leq i \leq k\}, \\ E_2 &= \{x_1y_3, x_2y_1, x_2y_2, x_2y_3, x_3y_4, x_3y_6, x_4y_5, x_4y_6\} \\ &\quad \cup \{v_iu_i \mid 1 \leq i \leq k\} \cup \{u_iv_{i+1} \mid 1 \leq i \leq k\} \cup \{w_iv_{i+1} \mid 1 \leq i \leq k\}. \end{aligned}$$

Color the edges of E_1 by red and E_2 by blue. Under this coloring $A_4^*[E_1] \not\supseteq K_{1,2}$ and $A_4^*[E_2] \not\supseteq C_4$.

Case 3.1.2 e is an edge of the C_4 -path.

W.l.o.g assume that $e = v_1w_1$. Then the edges of $A_4^*(k)$ can be partitioned into two classes, namely E_3 and E_4 , as follows

$$\begin{aligned} E_3 &= \{x_1y_2, x_2y_1, x_3y_5, x_4y_4\} \cup \{v_jw_j \mid 2 \leq j \leq k\}, \\ E_4 &= \{x_1y_1, x_1y_3, x_2y_2, x_2y_3, x_3y_4, x_3y_6, x_4y_5, x_4y_6\} \\ &\quad \cup \{v_iu_i \mid 1 \leq i \leq k\} \cup \{u_iv_{i+1} \mid 1 \leq i \leq k\} \cup \{w_iv_{i+1} \mid 1 \leq i \leq k\}. \end{aligned}$$

Color the edges of E_3 by red and E_4 by blue. Under this coloring $A_4^*[E_3] \not\supseteq K_{1,2}$ and $A_4^*[E_4] \not\supseteq C_4$. Thus $A_4^* \rightarrow (K_{1,2}, C_4)$. Therefore $A_4 \in \mathcal{R}_2(K_{1,2}, C_4)$.

Case 3.2 Graph $A_5(k)$.

The proof is similar with the above case. Consider any red-blue coloring of $A_5(k)$ that implies an edge-decomposition $A_5(k) = A_{51}(k) \oplus A_{52}(k)$. Let $A_{51}(k) \not\supseteq K_{1,2}$. If there is no blue C_4 in the C_4 -path that belongs to $A_{52}(k)$ then we claim that a red edge is incident to one of the end vertices of the C_4 -path (in this case we define the end cycles as the basic cycles that are attached to L_1 and L_2). Then there is a blue C_4 in L_1 or L_2 (depends on whether the red edge is incident to the root of L_1 or the root of L_2). It follows that $A_{52}(k) \supseteq C_4$. Thus $A_5(k) \rightarrow (K_{1,2}, C_4)$.

Let $V(L_1) = \{x_i \mid 1 \leq i \leq 2\} \cup \{y_j \mid 1 \leq j \leq 3\}$, where x_i is the vertex of degree 3 and y_j is the vertex of degree 2. The root of L_1 , x_2 , is identified with v_1 . Let $V(L_2) = \{p_j \mid 1 \leq j \leq 3\} \cup \{q_i \mid 1 \leq i \leq 2\} \cup \{r_1\}$ where p_j is the vertex of degree 3, q_j is the vertex of degree 2 and r_1 is the vertex of degree 4. The root of L_2 , p_1 , is identified with v_{k+1} .

It can be proved similarly that $A_5(k) \setminus \{e\} \rightarrow (K_{1,2}, C_4)$ for any fixed edge $e \in E(A_5(k))$. Therefore $A_5(k) \in \mathcal{R}_2(K_{1,2}, C_4)$.

Case 3.3 Graph $A_6(k)$.

The proof is similar with two cases above. Consider any red-blue coloring of $A_6(k)$ that implies an edge-decomposition $A_6(k) = A_{61}(k) \oplus A_{62}(k)$. Let $A_{61}(k) \not\supseteq K_{1,2}$. If there is no blue C_4 in the C_4 -path that belongs to $A_{62}(k)$ then we claim that a red edge is incident to one of the end vertices of the C_4 -path (in this case we define the end cycles as the basic cycles that are attached to L_2 s). Then there is a blue C_4 in one of the L_2 s. It follows that $A_{62}(k)[E_2] \supseteq C_4$. Thus $A_6(k) \rightarrow (K_{1,2}, C_4)$.

We denote the set of vertices of the L_2 s by $\{p_j \mid 1 \leq j \leq 6\} \cup \{q_i \mid 1 \leq i \leq 4\} \cup \{r_t \mid 1 \leq t \leq 2\}$ where p_j is the vertex of degree 3, q_i is the vertex of degree 2 and r_t is the vertex of degree 4. The root of the first L_2 , p_1 , is identified with v_1 and the root of the second L_2 , p_4 , is identified with v_{k+1} .

It can be proved similarly that $A_6(k) \setminus \{e\} \rightarrow (K_{1,2}, C_4)$ for any fixed edge $e \in E(A_6(k))$. Therefore $A_6(k) \in \mathcal{R}_2(K_{1,2}, C_4)$.

Case 3.4 Graph $B_1(k, t)$.

Consider any edge-decomposition $B_1(k, t) = B_{11}(k, t) \oplus B_{12}(k, t)$ by a red-blue coloring. Let $B_{11}(k, t) \not\supseteq K_{1,2}$. If there is no blue C_4 in the C_4 -path that belongs to $B_{12}(k, t)$ then we claim that a red edge is incident to one of the end vertices of the C_4 -path. If the red edge is incident to v_1 then we have a blue C_4 in one of the basic cycles of C_4 -cycle. If the red edge is incident to v_{k+1} then there is a blue C_4 in L_1 . It follows that $B_{12}(k, t) \supseteq C_4$. Thus $B_1(k, t) \rightarrow (K_{1,2}, C_4)$.

The vertices of C_4 -cycle of length t are denoted by $\{z_s \mid 1 \leq s \leq t\} \cup \{a_s \mid 1 \leq s \leq t\} \cup \{b_s \mid 1 \leq s \leq t\}$, where z_s is the vertex of degree 4, a_s and b_s are the vertices of degrees 2. For L_1 , the notation of their vertices is similar with the L_1 in $A_4(k)$. W.l.o.g, we assume that the first end vertex of C_4 -path, v_1 , is identified with z_s for an arbitrary s , $1 \leq s \leq t$ and the second end vertex, v_{k+1} , is identified with x_2 , the root of L_1 .

To prove the minimality of $B_1(k, t)$, we consider $B_1^*(k, t) = B_1(k, t) \setminus \{e\}$ for any fixed edge $e \in E(B_1(k, t))$. Consider the following cases.

Case 3.4.1 e is an edge in C_4 -cycle.

W.l.o.g we assume that $e = z_t a_t$. Then the edges of $B_1^*(k, t)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{a_p z_{p+1} \mid 1 \leq p \leq t-1\} \cup \{v_i w_i \mid 1 \leq i \leq k\} \cup \{x_1 y_2, x_2 y_1\}, \\ E_2 &= E(B_1^*(k, t)) \setminus E_1. \end{aligned}$$

Color the edges of E_1 by red and E_2 by blue. Under this coloring $B_1^*[E_1] \not\cong K_{1,2}$ and $B_1^*[E_2] \not\cong C_4$.

Case 3.4.2 e is an edge in C_4 -path.

W.l.o.g we assume that $e = v_1 w_1$. Then the edges of $B_1^*(k, t)$ can be partitioned into two classes, namely E_3 and E_4 , as follows

$$\begin{aligned} E_3 &= \{z_s a_s \mid 1 \leq s \leq t\} \cup \{v_j w_j \mid 2 \leq j \leq k\} \cup \{x_1 y_1, x_2 y_2\}, \\ E_4 &= E(B_1^*(k, t)) \setminus E_3. \end{aligned}$$

Color the edges of E_3 by red and E_4 by blue. Under this coloring $B_1^*[E_3] \not\cong K_{1,2}$ and $B_1^*[E_4] \not\cong C_4$.

Case 3.4.3 $e \in L_1$.

W.l.o.g we assume that $e = x_1 y_1$. Then the edges of $B_1^*(k, t)$ can be partitioned into two classes, namely E_5 and E_6 , as follows

$$\begin{aligned} E_5 &= \{z_s a_s \mid 1 \leq s \leq t\} \cup \{w_i v_{i+1} \mid 1 \leq i \leq k\} \cup \{x_1 y_2\}, \\ E_6 &= E(B_1^*(k, t)) \setminus E_5. \end{aligned}$$

Color the edges of E_5 by red and E_6 by blue. Under this coloring $B_1^*[E_5] \not\cong K_{1,2}$ and $B_1^*[E_6] \not\cong C_4$. Thus $B_1^*(k, t) \rightarrow (K_{1,2}, C_4)$. Therefore $B_1(k, t) \in \mathcal{R}(K_{1,2}, C_4)$.

Case 3.5 Graph $B_2(k, t)$.

The notation for C_4 -cycle is similar with the above case, while the notation for L_2 is similar with the notation of L_2 in $A_5(k)$. W.l.o.g we assume that the first end vertex of C_4 -path, v_1 , is identified with z_s for an arbitrary s , $1 \leq s \leq t$ and the second end vertex, v_{k+1} , is identified with p_1 , the root of L_2 .

The proof is similar with Case 3.4 above. We consider any edge-decomposition $B_2(k, t) = B_{21}(k, t) \oplus B_{22}(k, t)$ by a red-blue coloring. Let $B_{21}(k, t) \not\cong K_{1,2}$. If there is no blue C_4 in the C_4 -path that belongs to $B_{22}(k, t)$ then we claim that a red edge is incident to one of the end vertices of the C_4 -path. If the red edge is incident to v_1 then we have a blue C_4 in one of the basic cycles of C_4 -cycle. If the red edge is incident to v_{k+1} then there is a blue C_4 in L_2 . It follows that $B_{22}(k, t) \supseteq C_4$. Thus $B_2(k, t) \rightarrow (K_{1,2}, C_4)$.

It can be proved similarly that $B_2(k, t) \setminus \{e\} \rightarrow (K_{1,2}, C_4)$ for any fixed edge $e \in E(B_2(k, t))$. Therefore $B_2(k, t) \in \mathcal{R}(K_{1,2}, C_4)$.

Case 3.6 Graph $D_1(k, t_1, t_2)$.

We denote the set of vertices of C_4 -cycle of length t_1 by

$$\{z_{mt_1} \mid 1 \leq m \leq t_1\} \cup \{a_{mt_1} \mid 1 \leq m \leq t_1\} \cup \{b_{mt_1} \mid 1 \leq m \leq t_1\}$$

and the set of vertices of the C_4 -cycle of length t_2 by

$$\{z_{nt_2} \mid 1 \leq n \leq t_2\} \cup \{a_{nt_2} \mid 1 \leq n \leq t_2\} \cup \{b_{nt_2} \mid 1 \leq n \leq t_2\}.$$

The degrees of z_{mt_1} and z_{nt_2} are 4, while the vertices a_{mt_1}, b_{mt_1} and a_{nt_2}, b_{nt_2} are of degrees 2. W.l.o.g we identify one of the end vertex of C_4 -path, v_1 , with z_{1t_1} and the other end vertex, v_{k+1} , is identified with z_{1t_2} .

Consider any edge-decomposition $D_1(k, t_1, t_2) = D_{11}(k, t_1, t_2) \oplus D_{12}(k, t_1, t_2)$ by a red-blue coloring. Let $D_{11}(k, t_1, t_2) \not\subseteq K_{1,2}$. If there is no blue C_4 in the C_4 -path that belongs to $D_{12}(k, t_1, t_2)$ then we claim that a red edge is incident to one of the end vertices of the C_4 -path. If the red edge is incident to v_1 then we have a blue C_4 in one of the basic cycles of the first C_4 -cycle. If the red edge is incident to v_{k+1} then there is a blue C_4 in the second C_4 -cycle. It follows that $D_{12}(k, t_1, t_2) \supseteq C_4$. Thus $D_3(k, t_1, t_2) \rightarrow (K_{1,2}, C_4)$.

To prove the minimality of $D_1(k, t_1, t_2)$, we consider $D_1^*(k, t_1, t_2) = D_1(k, t_1, t_2) \setminus \{e\}$ for any fixed edge $e \in E(D_1(k, t_1, t_2))$. Consider the following cases.

Case 3.6.1 e is an edge of one of the C_4 -cycles.

W.l.o.g we assume that $e = z_{1t_1}a_{1t_1}$. Then the edges of $D_1^*(k, t_1, t_2)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{z_{pt_1}a_{pt_1} \mid 2 \leq p \leq t_1\} \cup \{v_i w_i \mid 1 \leq i \leq k\} \\ &\quad \cup \{a_{st_2}z_{(s+1)t_2} \mid 1 \leq s \leq t_2 - 1\} \cup \{a_{t_2t_2}z_{1t_2}\}, \\ E_2 &= E(D_1^*(k, t_1, t_2)) \setminus E_1. \end{aligned}$$

Color the edges of E_1 by red and E_2 by blue. Under this coloring $D_1^*(k, t_1, t_2)^* [E_1] \not\subseteq K_{1,2}$ and $D_1^*(k, t_1, t_2) [E_2] \not\subseteq C_4$.

Case 3.6.2 e is an edge in C_4 -path.

W.l.o.g we assume that $e = v_1 w_1$. Then the edges of $D_1^*(k, t_1, t_2)$ can be partitioned into two classes, namely E_3 and E_4 , as follows

$$\begin{aligned} E_3 &= \{a_{rt_1}z_{(r+1)t_1} \mid 1 \leq r \leq t_1 - 1\} \cup \{v_j w_j \mid 2 \leq j \leq k\} \\ &\quad \cup \{a_{st_2}z_{(s+1)t_2} \mid 1 \leq s \leq t_2 - 1\} \cup \{a_{t_2t_2}z_{1t_2}\}, \\ E_4 &= E(D_1^*(k, t_1, t_2)) \setminus E_3. \end{aligned}$$

Color the edges of E_3 by red and E_4 by blue. Under this coloring $D_1^*(k, t_1, t_2) [E_3] \not\subseteq K_{1,2}$ and $D_1^*(k, t_1, t_2) [E_4] \not\subseteq C_4$. Therefore $D_1(k, t_1, t_2) \in \mathcal{R}(K_{1,2}, C_4)$. \square

Theorem 4. $B_3(k, t), B_4(k, t) \notin \mathcal{R}(K_{1,2}, C_4)$.

Proof. **Case 4.1** Graph $B_3(k, t)$.

The notation of $V(B_3(k, t))$ is similar with notation of $V(B_1(k, t))$. W.l.o.g we assume that the first end vertex of C_4 -path, v_1 , is identified with a_1 and the second end vertex, v_{k+1} , is identified with x_2 , the root of L_1 .

The edges of $B_3(k, t)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{z_s b_s \mid 1 \leq s \leq t\} \cup \{v_i w_i \mid 1 \leq i \leq k\} \cup \{x_1 y_1, x_2 y_2\}, \\ E_2 &= E(B_3(k, t)) \setminus E_1. \end{aligned}$$

Under this coloring $B_3(k, t)[E_1] \not\cong K_{1,2}$ and $B_3(k, t)[E_2] \not\cong C_4$.

Thus

$B_3(k, t) \rightarrow (K_{1,2}, C_4)$. Therefore $B_3(k, t) \notin \mathcal{R}(K_{1,2}, C_4)$.

Case 4.2 Graph $B_4(k, t)$.

The notation of $V(B_4(k, t))$ is similar with notation of $V(B_2(k, t))$. W.l.o.g we assume that the first end vertex of C_4 -path, v_1 , is identified with a_1 and the second end vertex, v_{k+1} , is identified with p_1 , the root of L_2 .

Similar with Case 4.1 above, the edges of $B_4(k, t)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{z_s b_s \mid 1 \leq s \leq t\} \cup \{v_i w_i \mid 1 \leq i \leq k\} \cup \{p_1 p_3, p_2 q_1, q_2 r_1\}, \\ E_2 &= E(B_4(k, t)) \setminus E_1. \end{aligned}$$

Under this coloring $B_4(k, t)[E_1] \not\cong K_{1,2}$ and $B_4(k, t)[E_2] \not\cong C_4$. Thus $B_4(k, t) \rightarrow (K_{1,2}, C_4)$. Therefore $B_4(k, t) \notin \mathcal{R}(K_{1,2}, C_4)$. \square

Theorem 5. $D_2(k, t_1, t_2), D_3(k, t_1, t_2) \notin \mathcal{R}(K_{1,2}, C_4)$.

Proof. **Case 5.1 Graph $D_2(k, t_1, t_2)$.**

The notation of $V(D_2(k, t_1, t_2))$ is similar with $V(D_1(k, t_1, t_2))$. W.l.o.g we assume that v_1 is identified with z_{1t_1} and v_{k+1} is identified with a_{1t_2} .

The edges of $D_2(k, t_1, t_2)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{a_{rt_1} z_{(r+1)t_1} \mid 1 \leq r \leq t_1 - 1\} \cup \{a_{t_1 t_1} z_{1t_1}\} \\ &\quad \cup \{w_i v_{i+1} \mid 1 \leq i \leq k\} \cup \{z_{nt_2} b_{nt_2} \mid 1 \leq n \leq t_2\}, \\ E_2 &= E(D_2(k, t_1, t_2)) \setminus E_1. \end{aligned}$$

Under this coloring $D_2(k, t_1, t_2)[E_1] \not\cong K_{1,2}$ and $D_2(k, t_1, t_2)[E_2] \not\cong C_4$. Thus $D_2(k, t_1, t_2) \rightarrow (K_{1,2}, C_4)$. Therefore $D_2(k, t_1, t_2) \notin \mathcal{R}(K_{1,2}, C_4)$.

Case 5.2 Graph $D_3(k, t_1, t_2)$.

The notation of $V(D_3(k, t_1, t_2))$ is similar with $V(D_1(k, t_1, t_2))$. W.l.o.g we assume that v_1 is identified with a_{1t_1} and v_{k+1} is identified with a_{1t_2} .

The edges of $D_3(k, t_1, t_2)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{z_{mt_1} a_{mt_1} \mid 1 \leq m \leq t_1\} \cup \{w_i v_{i+1} \mid 1 \leq i \leq k\} \\ &\quad \cup \{b_{st_2} z_{(s+1)t_2} \mid 1 \leq s \leq t_2 - 1\} \cup \{b_{t_2 t_2} z_{1t_2}\}, \\ E_2 &= E(D_3(k, t_1, t_2)) \setminus E_1. \end{aligned}$$

Under this coloring $D_3(k, t_1, t_2)[E_1] \not\cong K_{1,2}$ and $D_3(k, t_1, t_2)[E_2] \not\cong C_4$. Thus $D_3(k, t_1, t_2) \rightarrow (K_{1,2}, C_4)$. Therefore $D_3(k, t_1, t_2) \notin \mathcal{R}(K_{1,2}, C_4)$. \square

Theorem 6. $I(t, k_1, k_2) \notin \mathcal{R}(K_{1,2}, C_4)$.

Proof. We denote the set of vertices of C_4 -path of length k_1 by

$$\{v_{pk_1} \mid 1 \leq p \leq k_1 + 1\} \cup \{w_{qk_1} \mid 1 \leq q \leq k_1\} \cup \{u_{qk_1} \mid 1 \leq q \leq k_1\}$$

and the set of vertices of C_4 -path of length k_2 by

$$\{v_{rk_2} \mid 1 \leq r \leq k_2 + 1\} \cup \{w_{sk_2} \mid 1 \leq s \leq k_2\} \cup \{u_{sk_2} \mid 1 \leq s \leq k_2\}.$$

The notation of the vertices of C_4 -cycle of length t are similar with its notation in Case 3.4. Consider these following cases.

Case 6.1 $v_{(k_1+1)k_1}$ and v_{1k_2} are identified with two vertices of degrees 4 of the C_4 -cycle.

W.l.o.g we assume that $v_{(k_1+1)k_1}$ is identified with z_1 and v_{1k_2} is identified with z_t . Then the edges of $I(t, k_1, k_2)$ can be partitioned into two classes, namely E_1 and E_2 , as follows

$$\begin{aligned} E_1 &= \{v_{ik_1}w_{ik_1} \mid 1 \leq i \leq k_1\} \cup \{w_{jk_2}v_{(j+1)k_2} \mid 1 \leq j \leq k_2\} \\ &\quad \cup \{a_s z_{s+1} \mid 1 \leq s \leq t-1\} \cup \{a_t z_1\}, \\ E_2 &= E(I(t, k_1, k_2)) \setminus E_1. \end{aligned}$$

Under this coloring $I(t, k_1, k_2)[E_1] \not\cong K_{1,2}$ and $I(t, k_1, k_2)[E_2] \not\cong C_4$.

Case 6.2 $v_{(k_1+1)k_1}$ is identified with a vertex of degrees 4 and v_{1k_2} is identified with a vertex of degree 2 of the C_4 -cycle.

W.l.o.g we assume that $v_{(k_1+1)k_1}$ is identified with z_1 and v_{1k_2} is identified with a_t . Then the edges of $I(t, k_1, k_2)$ can be partitioned into two classes, namely E_3 and E_4 , as follows

$$\begin{aligned} E_3 &= \{v_{ik_1}w_{ik_1} \mid 1 \leq i \leq k_1\} \cup \{w_{jk_2}v_{(j+1)k_2} \mid 1 \leq j \leq k_2\} \\ &\quad \cup \{a_s z_{s+1} \mid 1 \leq s \leq t-1\} \cup \{a_t z_1\}, \\ E_4 &= E(I(t, k_1, k_2)) \setminus E_3. \end{aligned}$$

Under this coloring $I(t, k_1, k_2)[E_3] \not\cong K_{1,2}$ and $I(t, k_1, k_2)[E_4] \not\cong C_4$.

Case 6.3 $v_{(k_1+1)k_1}$ and v_{1k_2} are identified with two vertices of degrees 2 of the C_4 -cycle.

W.l.o.g we assume that $v_{(k_1+1)k_1}$ is identified with a_1 and v_{1k_2} is identified with a_t . Then the edges of $I(t, k_1, k_2)$ can be partitioned into two classes, namely E_5 and E_6 , as follows

$$\begin{aligned} E_5 &= \{v_{ik_1}w_{ik_1} \mid 1 \leq i \leq k_1\} \cup \{w_{jk_2}v_{(j+1)k_2} \mid 1 \leq j \leq k_2\} \\ &\quad \cup \{a_s z_{s+1} \mid 1 \leq s \leq t-1\} \cup \{a_t z_1\}, \\ E_6 &= E(I(t, k_1, k_2)) \setminus E_5. \end{aligned}$$

Under this coloring $I(t, k_1, k_2)[E_5] \not\cong K_{1,2}$ and $I(t, k_1, k_2)[E_6] \not\cong C_4$. Thus $I(t, k_1, k_2) \rightarrow (K_{1,2}, C_4)$. Therefore $I(t, k_1, k_2) \notin \mathcal{R}(K_{1,2}, C_4)$. \square

As a final remark, we present some problems that are raised from this paper.

- (1) Characterize all graphs $F \in \mathcal{R}(K_{1,2}, C_4)$ with $\text{diam}(F) = 2$,
- (2) Does there exist a graph F such that $F \in \mathcal{R}(K_{1,2}, C_4)$ with $\text{diam}(F) = 3$?
- (3) Characterize all graphs $F \in \mathcal{R}(K_{1,2}, C_4)$ with $\text{diam}(F) \geq 4$.

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