# DETOUR GRAPHS

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#### **Abstract**

Let G = (V, E) be a connected simple graph. Let  $u, v \in V(G)$ . The detour distance, D(u, v), between u and v is the distance of a longest path from u to v. E.Sampathkumar defined the detour graph of G, denoted by D(G) as follows: D(G) is an edge labelled complete graph on n vertices where n = |V(G)|, the edge label for uv,  $u, v \in V(K_n)$  being D(u, v). Any edge labelled complete graph need not be the detour graph of a graph. In this paper, we characterize detour graphs of a tree. We also characterise graphs for which the detour distance sequences are given.

**Definition:-** 1 Let G = (V, E) be a simple graph. The detour graph of G denoted by D(G), is an edge labelled complete graph

 $K_{|V(G)|}$  on the same vertex set V(G) such that for any two vertices a, b in G with D(a, b) = k, the edge label of ab is k in  $K_{|V(G)|}$ .

**Result:-** 1 If a detour graph of G has all the labels 1, then  $G = K_2$ .

#### Proof:

Assume  $|V(G)| \ge 3$ . If d(G) = 1, then all the labels are (n-1) and  $(n-1) \ge 2$ . If  $d(G) = k \ge 2$ , there exists a label  $r \ge 2$ , a contradiction. Therefore  $G = K_2$ .

**Result:- 2** If a detour graph has all the edge labels 2, then the graph is  $K_3$ .

#### Proof :-

Assume  $|V(G)| \ge 4$ . Since there is no label 1,  $\delta(G) \ge 2$ . Therefore, G has a cycle. If  $d(G) \ge 3$ , then there is an edge with label greater than or equal to 3 in the detour graph, a contradiction.

If d(G) = 1, then we get a complete graph and all the labels are greater than or equal to 3, a contradiction. Therefore, d(G) = 2.

If G has a cycle  $C_4$  or  $C_5$  as a subgraph, there exists an edge which has label greater than or equal to 3, a contradiction. Therefore, the cycles in G are all triangles. Let G have two triangles  $T_1$  and  $T_2$ . If  $T_1$  and  $T_2$  have common edge, we have a  $C_4$ , a contradiction. If  $T_1$  and  $T_2$  have a common vertex, then there exists an edge with label greater than or equal to 4, a contradiction. Therefore, G cannot have two triangles. Therefore, G has a unique triangle. Since  $|V(G)| \geq 4$ , there exists an edge label greater than or equal to 3, a contradiction. Therefore,  $G = K_3$ .

**Result:- 3** If a graph is any one of the following, then the detour distance between any two points in the graph is (n-1).

(i.)  $G_1 + G_2$  where  $G_1$  and  $G_2$  are Hamiltonian graphs and  $|V(G_1)| + |V(G_2)| = n$ .

(ii.)  $G_1+K_r$  where the number of components of  $G_1$  is less than or equal to r. (that is,  $\omega(G) \leq r$ ) and  $|V(G_1)| = n - r$ .

(iii.)  $K_n$ .

#### Result:- 4

(i.) If  $G = C_{2n}$ , then  $D(C_{2n}) = nK_2$ .

(ii.) If 
$$G = C_{2n+1}$$
, then  $D(C_{2n+1}) = C_{2n+1}$ .

**Result:- 5** If  $G = P_n$ , then  $D(P_n) = P_n$  for all n.

**Result:- 6** If  $G = K_{1,r}$ , then  $D(K_{1,r}) = K_{r+1}$ .

**Result:-** 7 The maximum number of distinct integers that one can get in a detour distance sequence is n-1.

#### Proof :-

Let  $G = P_n$ . Then the detour distance sequence is  $1^{(n-1)}, 2^{(n-2)}, \cdots, (n-1)^{(1)}$  and this has maximum number of distinct integers. (Note that the number of distinct integers in the detour distance sequence of any graph on n vertices is at most n-1).

**Theorem:-** 1 A graph G on n vertices is a tree if and only if the detour graph D(G) contains (n-1) 1's.

#### Proof:-

Let the detour graph of G have (n-1) 1's. Let  $V(G) = \{u_1, u_2, u_3, \cdots, u_n\}$ . If G contains a cycle C with vertices  $u_{r_1}, u_{r_2}, u_{r_3}, \cdots, u_{r_k}$ , then all the edges of this cycle will have labels greater than 1. Then the (n-k) points, other than those in the cycle C will contribute a maximum of (n-k-1) 1's to the edges of the complete graph  $K_n$  on the n-vertices  $\{u_1, u_2, u_3, \cdots, u_n\}$ . Therefore, the number of 1's in D(G) is atmost (n-k-1) 1's, a contradiction. Therefore, G has no cycles. Since there are (n-1) 1's in D(G), there are exactly (n-1)edges in G. Therefore G is a connected graph. Therefore G is a connected acyclic graph. That is, G is a tree. The converse is obvious.

**Theorem:- 2** A graph G is a path if and only if the detour distance sequence is  $1^{(n-1)}, 2^{(n-2)}, \dots, (n-1)^{(1)}$ 

#### Proof:-

Suppose the detour distance sequence of a graph G is  $1^{(n-1)}$ ,  $2^{(n-2)}, \dots, (n-1)^{(1)}$ . Then G is a tree. Since the label (n-1) appears exactly once, there exists exactly one longest path(say) between u and v of length (n-1). Let  $\{u=u_1, u_2, u_3, \dots, u_n=v\}$ , be the longest path. If u or v has degree greater than 1, then there will be a longest path of length greater than (n-1). But there exists no label in the detour distance sequence greater than (n-1). Therefore, u and v are pendent vertices. Further, all the n-points are in this longest path. Therefore, G is a path. The converse is obvious.

**Remark:-** 1 In the above proof, we use only two facts. In the detour distance sequence, (i) there are exactly (n-1) 1's and (ii) one (n-1). These two facts are enough to determine that the graph is a path. That is, if the detour distance sequence of G contains  $1^{(n-1)}$  and  $(n-1)^{(1)}$ , then 2 exactly appear (n-2) times, 3 appears exactly (n-3) times,  $\cdots (n-2)$  appears exactly 2 times.

**Theorem:- 3** A graph G is a star if and only if the detour distance sequence of G is  $1^{n-1}$ ,  $2^{(n_{c_2}-(n-1))}$ 

### Proof:-

Since the detour distance sequence of G contains  $1^{(n-1)}$ , G is a tree.

Claim: There exists a point of degree (n-1).

Let u be a point of G of degree greater than 1. Let  $v, w \in N(u)$ . If |V(G)| = 3, then G is a star with center u. Suppose  $x \in V(G) - \{u, v, w\}$ . Suppose x is not adjacent to x. Then the detour distance between x and x is greater than or equal to x, a contradiction (since the detour distance sequence contains no x). Therefore, x is adjacent to every vertex of x. Since x is a tree, all the vertices other than x are pendent vertices. Therefore, x is a star. The converse is obvious.

**Theorem:-** 4 A graph G is a cycle of length 2n if and only if the detour distance sequence of G is  $n^{(n)}$ ,  $(n+1)^{(2n)}$ ,  $\cdots$ ,  $(2n-1)^{(2n)}$ .

#### Proof :-

Suppose G is acyclic. The absence of  $\infty$  in the detour distance sequence ensures that G is a tree. 1 must occur (n-1) times in the detour distance sequence. But 1 is not present in the detour distance sequence. Therefore, G contains a cycle.

G has no pendent vertex since 1 is absent in the detour distance sequence of G. Therefore,  $\delta(G) \geq 2$ .

Since  $(2n-1)^{(2n)}$  occurs in the detour distance sequence of G, there are exactly 2n longest paths of length (2n-1). Let u and v be two points, the length of the longest path between them being (2n-1). From the sequence it is clear that G has order 2n. Let  $V(G) = \{u_1, u_2, u_3, \cdots, u_{2n}\}$ . Let, without loss of generality, the longest path between  $u_1$  and  $u_{2n}$  be  $u_1, u_2, u_3, \cdots, u_{2n}$ . Since there are 2n longest paths of length (2n-1), there exist  $u_i$  and  $u_j$  ( $1 \le i < j \le 2n$ ) such that the longest path between  $u_i$  and  $u_j$  is (2n-1). This shows that  $u_i$  and  $u_j$  are adjacent and  $u_1$  and  $u_{2n}$  are adjacent. Hence G is a cycle of length 2n. The converse is obvious.

Remark:- 2 A graph G is a cycle of length (2n+1) if and only if the detour distance sequence of G is  $(n+1)^{(2n+1)}$ ,  $(n+2)^{(2n+1)}$ ,  $(n+3)^{(2n+1)}$ ,  $\cdots$ ,  $(2n)^{(2n+1)}$ .

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