

# On Reconstructing Graphs from $n - 2$ Cards

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*To Ralph G. Stanton, in memoriam.*

## Abstract

Let  $G$  and  $H$  be graphs on  $n+2$  vertices  $\{u_1, u_2, \dots, u_n, x, y\}$  such that  $G - u_i \cong H - u_i$ , for  $i = 1, 2, \dots, n$ . Recently Ramachandran, Monikandan, and Balakumar have shown in a sequence of two papers that if  $n \geq 9$ , then  $|\varepsilon(H) - \varepsilon(G)| \leq 1$ . In this paper we present a simpler proof of their theorem, using a counting lemma.

## 1 Introduction

Let  $G$  and  $H$  be graphs on  $n + 2$  vertices  $\{u_1, u_2, \dots, u_n, x, y\}$  such that  $G - u_i \cong H - u_i$ , for  $i = 1, 2, \dots, n$ . Recently Ramachandran and Monikandan [3] and Monikandan and Balakumar [1] have shown in a sequence of two papers that if  $n \geq 9$ , then  $|\varepsilon(H) - \varepsilon(G)| \leq 1$ . Their proof is based on determining the partial structure of the graphs. In this paper we present a simpler proof of their theorem, using a counting lemma that requires only the degrees of the vertices.

Let  $d(u, G)$  denote the degree of vertex  $u$  in any graph  $G$ . The following lemma is from [2].

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**Lemma 1.1** *Let  $m \geq 0$  and let  $G$  and  $H$  be as above. Suppose that  $\varepsilon(H) = \varepsilon(G) + m$ . Then  $d(u_i, H) = d(u_i, G) + m$ , for  $i = 1, 2, \dots, n$ .*

**Proof.**  $d(u_i, H) = \varepsilon(H) - \varepsilon(H - u_i) = \varepsilon(G) + m - \varepsilon(G - u_i) = d(u_i, G) + m$ .

Suppose now that  $m \geq 2$ . Choose a vertex  $u_1 \in U$ , and let  $p$  be an isomorphism mapping  $H - u_1$  to  $G - u_1$ . Then  $p$  is a permutation of  $V - u_1$ . Let  $G' = G - u_1$  and  $H' = H - u_1$ . Consider a cycle  $(v_1, v_2, \dots, v_k)$  of  $p$ . Is it possible that all  $v_i \in U$ ? Let  $\alpha = d(v_1, H')$ . In  $H'$  there are  $\alpha$  edges incident on  $v_1$ , hence in  $G'$  there are  $\alpha$  edges incident on  $p(v_1) = v_2$ , so  $d(v_2, G') = \alpha$ . But  $d(v_2, H) = d(v_2, G) + m$ , so that  $d(v_2, H') \geq \alpha + 1$ . Continuing in this way, we find that  $d(v_3, H') \geq \alpha + 2$ , etc, until we reach  $d(v_k, H') \geq \alpha + k - 1$  from which it follows that  $d(v_1, H') \geq \alpha + k$ , a contradiction. It follows that every cycle of  $p$  contains either  $x$  or  $y$ , so that there are at most two cycles.

Let  $H[u, v]$  denote the number of edges of  $H$  joining  $u$  to  $v$  (either 0 or 1). Consider any  $u \in U$ , where  $u \neq u_1$ . Since  $d(u, H) = d(u, G) + m$ , we can write:

- (a)  $d(u, H') = d(u, G') + m$  and  $H[u_1, u] = G[u_1, u]$ ; or
- (b)  $d(u, H') = d(u, G') + m - 1$  and  $H[u_1, u] = 1 + G[u_1, u]$ ; or
- (c)  $d(u, H') = d(u, G') + m + 1$  and  $H[u_1, u] + 1 = G[u_1, u]$ .

Let there be  $a$  of the first type,  $b$  of the second type, and  $c$  of the third type. Let  $h = H[u_1, x] + H[u_1, y]$  and  $g = G[u_1, x] + G[u_1, y]$

**Lemma 1.2** *Let  $m \geq 1$ . Then*

$$ma + (m - 1)b + (m + 1)c = mn - 2m + h - g$$

**Proof.** We have  $a + b + c = n - 1$ , since there are  $n - 1$  vertices in  $U - u_1$ . Then  $d(u_1, H) - d(u_1, G) = m = b - c + h - g$ , so that  $c = b - m + h - g$ , from which it follows that  $a + 2b = n + 1 - h + g$ . Hence  $ma + (m - 1)b + (m + 1)c = mn - 2m + h - g$ .

If there is a sequence  $x, v_1, v_2, \dots, v_k$  of consecutive vertices in a cycle of  $p$ , where each  $v_i \in U$ , then  $d(v_i, H')$  increases by  $m - 1, m$ , or  $m + 1$  for each  $i$ . If  $k = n - 1$ , then starting with  $d(v_1, G') =$

$d(x, H')$ , we find that  $d(v_k, H') = d(x, H') + ma + (m - 1)b + (m + 1)c = mn - 2m + h - g + d(x, H')$ . If  $m \geq 2$ , this expression is  $\geq 2n - 4 + h - g + d(x, H') > n$ , if  $n \geq 7$ , since  $h - g \geq -2$ . But each  $d(v_i, H') \leq n$ . Hence we conclude that if  $n \geq 7$ , there are exactly two sequences of consecutive vertices of  $U$  in  $p$ , one starting from  $x$ , and one starting from  $y$ .

Let  $x, v_1, v_2, \dots, v_k$  be one sequence of consecutive vertices in a cycle of  $p$ , and let  $y, w_1, w_2, \dots, w_j$  be the other, where each  $v_i, w_i \in U$  and  $k + j = n - 1$ . It is convenient to take  $v_0 = x$  and  $w_0 = y$ . Let there be  $a_1, b_1, c_1$  vertices of type (a),(b),(c), respectively, in  $v_1, v_2, \dots, v_k$ , and  $a_2, b_2, c_2$  in  $w_1, w_2, \dots, w_j$ .

**Lemma 1.3** *Let  $0 \leq i \leq k - 1$ . Then*

$$d(v_{i+1}, H') = d(v_i, H') + m + G[u_1, v_{i+1}] - H[u_1, v_{i+1}]$$

*Similarly, if  $0 \leq i \leq j - 1$ , then*

$$d(w_{i+1}, H') = d(w_i, H') + m + G[u_1, w_{i+1}] - H[u_1, w_{i+1}]$$

**Proof.** Since  $p$  maps  $H'$  to  $G'$  we have  $d(v_{i+1}, G') = d(v_i, H')$ . We also have  $d(v_{i+1}, H) = d(v_{i+1}, G) + m$ . But  $d(v_{i+1}, H') = d(v_{i+1}, H) - H[u_1, v_{i+1}]$  and  $d(v_{i+1}, G') = d(v_{i+1}, G) - G[u_1, v_{i+1}]$ . The result follows.

It follows that  $d(v_k, H') = d(x, H') + ma_1 + (m - 1)b_1 + (m + 1)c_1$ . Similarly  $d(w_j, H') = d(y, H') + ma_2 + (m - 1)b_2 + (m + 1)c_2$ . Adding these gives  $d(v_k, H') + d(w_j, H') = ma + (m - 1)b + (m + 1)c + d(x, H') + d(y, H')$ , which reduces to  $mn - 2m + h - g + d(x, H') + d(y, H')$ .

**Lemma 1.4** *If  $n \geq 9$  then  $|\varepsilon(H) - \varepsilon(G)| \leq 2$ .*

**Proof.** We have  $d(v_k, H'), d(w_j, H') \leq n$ . If  $m > 2$ , the above formula gives  $mn - 2m + h - g + d(x, H') + d(y, H') \leq 2n$ , or  $n \leq 2m / (m - 2) - \{h - g + d(x, H') + d(y, H')\} / (m - 2)$ . Since  $h - g \geq -2$  and  $m \geq 3$ , this gives  $n \leq 8$ , a contradiction.

At this point we take  $m = 2$  and  $n \geq 9$ . The above formula becomes

$$2n - 4 + h - g + d(x, H') + d(y, H') = d(v_k, H') + d(w_j, H') \quad (*)$$

We now abbreviate the notation somewhat, in order to streamline the proof of the main theorem. We have  $H' = H - u_1$  and  $G' = G - u_1$ . Here we have chosen  $u_1$  as a vertex in  $U$  with the *largest* degree in  $H$ . Since we will be working mostly with the graphs  $H$  and  $H'$ , we write  $dx' = d(x, H')$  and  $dy' = d(y, H')$ . Similarly for  $dv'_k$  and  $dw'_j$ . Similarly, we write  $dx$  for  $d(x, H)$ ,  $du_1$  for  $d(u_1, H)$ , etc. We recall that  $h = H[u_1, x] + H[u_1, y]$  and  $g = G[u_1, x] + G[u_1, y]$ . We write  $u \rightarrow v$  to indicate that  $u$  is adjacent to  $v$ .

**Theorem 1.5** *Let  $G$  and  $H$  be graphs on  $n+2$  vertices, where  $n \geq 9$ , such that  $H - u_i \cong G - u_i$  for  $i = 1, \dots, n$ . Then  $|\varepsilon(H) - \varepsilon(G)| \leq 1$ .*

The proof requires only the degrees of the vertices. We assume that  $\varepsilon(H) = \varepsilon(G) + 2$  and obtain a contradiction. We use the inequality (\*) to limit the values of the parameters  $dx', dy', dv'_k, dw'_j$ , and compare the smallest degrees of  $H - v_1, H - w_1, G - v_1$ , and  $G - w_1$ .

**Case 1.**  $dx' = dy' = 0$ .

This implies that  $dv'_k, dw'_j \leq n - 2$ . It then follows from (\*) that  $2n - 4 + h - g \leq 2n - 4$ , so that  $h \leq g$ . We observe first that if  $h = 0$ , then  $H$  and  $H''$  both have at least two vertices of degree 0 (namely  $x$  and  $y$ ), whereas  $G''$  has at most one vertex of degree 0 (namely  $v_1$  or  $w_1$ ). Therefore  $H'' \not\cong G''$ . Hence, we can assume that  $h \geq 1$ .

If  $h = 1$ , then  $H - v_1$  and  $H - w_1$  both have exactly one vertex of degree 0. Therefore  $G - v_1$  and  $G - w_1$  must each have one vertex of degree 0. It follows that  $u_1 \not\rightarrow v_1, w_1$  in  $G$ . But then  $G - v_2$  has two vertices of degree 0, whereas  $H - v_2$  has just one. Hence, we can assume that  $h = 2$ .

We now find that  $H - v_1$  and  $H - w_1$  both have exactly two vertices of degree 1. Therefore  $G - v_1$  and  $G - w_1$  also have exactly two vertices of degree 1. Therefore  $u_1 \rightarrow v_1, w_1$  in  $G$ , so that  $d(v_1, G) = d(w_1, G) = 1$ , and one of  $v_2, w_2$  has degree 1 in  $G$ . Consequently  $dv_1 = dw_1 = 3$ , so that  $dv'_1, dw'_1 \geq 2$ . But this implies that  $d(v_2, G'), d(w_2, G') \geq 2$ , a contradiction.

**Case 2.**  $dx' \neq 0$  and  $dy' = 0$ .

This implies that  $dv'_k, dw'_j \leq n - 1$ . It then follows from (\*) that  $2n - 4 + h - g + dx' \leq 2n - 2$ . If  $u_1 \not\rightarrow y$  in  $H$ , then  $H - w_1$  will have one vertex of degree 0, whereas  $G - w_1$  will have no vertices of degree 0. Hence we must have  $u_1 \rightarrow y$  in  $H$ , so that  $h \geq 1$ .

We next observe that in  $H - w_1$ , vertex  $y$  has degree 1. Therefore  $G - w_1$  must have a vertex of degree 1, which can only be  $v_1$ . It follows that  $u_1 \not\rightarrow v_1$  in  $G$ , and that  $1 = d(v_1, G') = dx'$ . Therefore  $v_1$  is adjacent to exactly one vertex  $z$  in  $G'$ , and in  $H'$ ,  $x$  is adjacent to only  $p^{-1}(z)$ . Then  $G - z$  has a vertex of degree 0, but  $H - z$  does not, a contradiction.

**Case 3.**  $dx' = 0$  and  $dy' \neq 0$ .

This is identical to Case 2, interchanging  $x$  and  $y$ , and  $j$  and  $k$ .

**Case 4.**  $dx' \neq 0$  and  $dy' \neq 0$ .

We have  $dx', dy' \geq 1$ . Let  $\delta = dx' + dy' - 2$ . Then  $\delta \geq 0$ . Without loss of generality, we take  $n \geq dv'_k \geq dw'_j$ .

**4.1**  $dv'_k = n$ . Then  $du_1 \geq dv_k \geq n$ . If  $u_1 \rightarrow v_k$  in  $H$ , then  $du_1 = dv_k = n + 1$ , so that  $u_1 \rightarrow x, y$  in  $H$ , which implies that  $h = 2$ . If  $u_1 \not\rightarrow v_k$  in  $H$ , then since  $du_1 \geq n$ , we again have  $u_1 \rightarrow x, y$  in  $H$  and  $h = 2$ . By (\*),  $2n - 4 + 2 - g + 2 + \delta = n + dw'_j$ , which reduces to  $g = n - dw'_j + \delta$ . Now if  $dw'_j = n$ , then  $w_j \rightarrow x, y$  in  $H$ , so that  $\delta \geq 2$ , which implies that  $g = \delta = 2$ . If  $dw'_j = n - 1$ , then  $w_j$  is adjacent to at least one of  $x, y$  in  $H$ , so that  $\delta \geq 1$ , which gives  $g = 2$  and  $\delta = 1$ . If  $dw'_j = n - 2$ , this gives  $g = 2$  and  $\delta = 0$ . So  $g$  always equals 2, and  $dx', dy'$  are forced. They are either  $(2, 2), (2, 1)$  or  $(1, 1)$  according as  $dw'_j$  is  $n, n - 1$ , or  $n - 2$ . Note that  $dw'_j \not\leq n - 3$ , since this would give  $g \geq 3$ .

We now find that the two smallest degrees of  $H - v_1$  and  $H - w_1$  are  $dx' + 1, dy' + 1$ . These must also be the smallest degrees of  $G - v_1$  and  $G - w_1$ . It follows that  $u_1 \rightarrow v_1, w_1$  in  $G$ , so that  $d(v_1, G) = dx' + 1$  and  $d(w_1, G) = dy' + 1$ . Therefore  $dv_1 = dx' + 3$  and  $dw_1 = dy' + 3$ , so that  $dv'_1 \geq dx' + 2$  and  $dw'_j \geq dy' + 2$ . But then the smallest degrees of  $G - v_1$  and  $G - w_1$  cannot be  $dx' + 1, dy' + 1$ , a contradiction.

**4.2**  $dv'_k = n - 1$ . By (\*), we have  $2n - 4 + h - g + 2 + \delta = n - 1 + dw'_j$ , which reduces to  $g = n - 1 + h + \delta - dw'_j$ . If  $u_1 \not\rightarrow v_k, w_j$  in  $H$ , then

since  $du_1 \geq n - 1$ , we must have  $u_1 \rightarrow x, y$  in  $H$ , so that  $h = 2$ . If  $u_1$  is adjacent to  $v_k$  but not to  $w_j$ , then  $du_1 \geq n$ , so that we again have  $h = 2$ . Then  $g = n + 1 + \delta - dw'_j$ . We must have  $dw'_j = n - 1$ ,  $g = 2$ , and  $\delta = 0$ . Therefore  $dx' = dy' = 1$ .

If  $u_1$  is adjacent to  $w_j$  but not to  $v_k$ , then  $h \geq 1$ . If  $u_1 \rightarrow v_k, w_j$  in  $H$ , then  $du_1 \geq n$  and so  $h \geq 1$ . Then  $g \geq n + \delta - dw'_j$ . We either have  $dw'_j = n - 1$  (which forces  $h = 2$ ),  $g = 2$ , and  $\delta = 0$  as in the previous paragraph, or else  $dw'_j = n - 2$ ,  $h = 1$ ,  $g = 2$ , and  $\delta = 0$ . Therefore  $dx' = dy' = 1$ .

In the situation when  $h = 2$ , we find that  $H - v_1$  and  $H - w_1$  both have two vertices of degree 2 as the smallest degrees, namely  $x$  and  $y$ . Therefore  $G - v_1$  and  $G - w_1$  must also have two vertices of degree two. Hence  $u_1 \rightarrow v_1, w_1$  in  $G$ , so that  $d(v_1, G) = d(w_1, G) = 2$ , from which  $dv_1 = dw_1 = 4$  and  $dv'_1, dw'_1 \geq 3$ . But then  $G - v_1$  and  $G - w_1$  do not have two vertices of degree two, a contradiction.

In the situation when  $h = 1$ ,  $H - v_1$  and  $H - w_1$  will have vertices  $x$  and  $y$  with degrees 1 and 2. Therefore  $G - v_1$  and  $G - w_1$  also have smallest degrees 1 and 2. It follows that  $u_1 \not\rightarrow v_1, w_1$  in  $G$ . Therefore  $dv_1 = dw_1 = 3$ , and  $u_1 \rightarrow v_1, w_1$  in  $H$ . Then at least one of  $H - v_1$  and  $H - w_1$  will have vertices  $v_k, w_j$  with largest degrees  $n - 1, n - 2$ . But since  $g = 2$ , the two largest degrees of  $G - v_1$  and  $G - w_1$  will both be  $n - 1$ , a contradiction.

**4.3**  $dv'_k = n - 2$ . By (\*), we have  $2n - 4 + h - g + 2 + \delta \leq 2n - 4$ , which gives  $g \geq 2 + h + \delta$ , which requires  $g = 2$ ,  $h = 0$ , and  $\delta = 0$ . But  $du_1 \geq n - 2$ , so that  $u_1$  must be adjacent to at least one of  $v_k, w_j$ . This forces  $du_1 \geq n$ , so that  $h \geq 1$ , a contradiction.

This completes the proof of the theorem.

## References

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