

Acyclic Coloring of Central and Total Graph of Path P_n and Fan Graph $F_{m,n}$

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Abstract

A proper vertex coloring (no two adjacent vertices have the same color) of a graph G is said to be acyclic if the induced subgraph of any two color classes is acyclic. The minimum number of colors required for acyclic coloring of a graph G is said to be its acyclic chromatic number and is denoted by $a(G)$. In this paper, we find the exact value of the acyclic chromatic number of the central and total graph of path P_n and Fan graph $F_{m,n}$.

1. Introduction

All graphs considered here are simple finite and undirected. Throughout this paper, we use the term coloring for vertex coloring of graphs. A proper coloring of a graph G is a coloring of the vertices of G such that no two neighbors in G are assigned the same color.

Definition 1. A subgraph H of a graph G is an induced subgraph if it has all the edges that appear in G over the same vertex set. The subgraph induced by the vertex set $v_1, v_2, v_3, \dots, v_k$ is denoted by $\langle v_1, v_2, v_3, \dots, v_k \rangle$.

Definition 2. A vertex coloring of a graph is said to be acyclic [9] if the induced subgraph of any two color classes is acyclic. In other words, the subgraph induced by any two color classes is a forest.

Definition 3. *The minimum number of colors needed to acyclically color the vertices of a graph G is called its acyclic chromatic number and is denoted by $a(G)$.*

The concept of acyclic coloring was introduced by B.Grunbaum in 1973[9]. He has proved that any planar graph is acyclically 9-colorable and he has conjectured that any planar graph can be acyclically vertex colored with 5 colors. This conjecture was later proved by Borodin[5]. Determining the acyclic chromatic number is a hard problem from both theoretical and an algorithmic point of view. More specifically, A.V. Kostochka has proved that it is an NP - complete problem to decide for a given graph G , whether $a(G) \leq 3$ [11]. Determining $a(G)$ for the class of bipartite graphs is still an open problem. Alon et al. have given a greedy algorithm to color any graph of maximum degree Δ acyclically, using $\Delta^2 + 1$ colors [2]. Albertson et al. have improved this result and shown that $a(G) \leq \Delta(\Delta - 1) + 2$ [1]. Some of the research works have been carried out by focusing on the family of graphs with small maximum degree. The introductory work of Grunbaum started with a bound of 4 colors on acyclic vertex coloring of a graph with maximum degree 3. Burnstein has shown that any graph with maximum degree 4 can be acyclically vertex colored with 5 colors [6]. Fertin and Raspaud have given $o(n\Delta^2)$ algorithm to color any graph with $\Delta(\Delta - 1)/2$ colors[7]. They have also proved that 9 colors are sufficient to color a graph with maximum degree 4. Skulrattankulchai has given a linear time algorithm to acyclically color the vertices of a subcubic graph with 4 colors [17]. Recently, V.Satish and K.Yadav have proved that graphs with maximum degree 4,5 and 6 can be acyclically vertex colored with 5,8 and 12 colors respectively by providing a linear time algorithm that color these graphs [15, 16, 18]. They have also shown that a graph of maximum degree Δ can be acyclically vertex colored with $(3\Delta^2 + 4\Delta + 8)/8$ colors [14].

Definition 4. *A cycle in a graph G is said to be a bicolored (j, k) -cycle if all its vertices are properly colored with two colors j and k . A graph G is said to be a (j, k) -cycle free graph or C_{jk} -free graph if it does not contain any bicolored (j, k) -cycle.*

Definition 5. *Let G be a graph with vertex set $V(G)$ and edge set $E(G)$. The central graph [13] of G , denoted by $C(G)$, is obtained from G by subdividing each edge exactly once and joining all the non adjacent vertices of G .*

Definition 6. *The total graph [4] of a graph, denoted by $T(G)$, is a graph such that the vertex set of T corresponds to the vertices and edges of G and two vertices are adjacent in T if and only if their corresponding elements are either adjacent or incident in G .*

Definition 7. A Fan graph $F_{m,n}$ [8] is defined as the graph sum of null graph on m vertices with path on n vertices. Symbolically, $F_{m,n} = \overline{K}_m + P_n$.

This work is an extended version of [3]. The purpose of this paper is to find the exact value of the acyclic chromatic number of central and total graph of path P_n and Fan graph $F_{m,n}$.

2. Acyclic Coloring of $C(P_n)$

Let P_n be a path on n vertices v_1, v_2, \dots, v_n . Let $x_k (1 \leq k \leq n-1)$ be the newly introduced vertex on the edge joining v_k and v_{k+1} . In this section, we first present a coloring algorithm of $C(P_n)$ and then we prove the coloring is acyclic in the immediate following theorem.

2.1. Structural Properties of $C(P_n)$

1. The vertex v_1 is adjacent with v_3, v_4, \dots, v_n whereas the vertex v_n is adjacent with v_1, v_2, \dots, v_{n-2} .
2. For each $k = 2$ to $n-1$, v_k is adjacent to v_j , for all $j = 1, 2, \dots, (k-2), (k+2), \dots, n$.

2.2. Coloring Algorithm

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Input :  $C(P_n)$ 
 $V \leftarrow \{v_1, v_2, \dots, v_n, x_1, x_2, \dots, x_{n-1}\}$ 
 $v_1, v_2 \leftarrow 1; v_3, v_4 \leftarrow 2;$ 
for  $k = 5$  to  $n$ 
    {
         $v_k \leftarrow k - 2;$ 
    }
end for
for  $k = 1$  to  $n - 1$ 
    {
         $x_k \leftarrow n - 1;$ 
    }
end for
end procedure
Output: vertex colored  $C(P_n)$ 

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Theorem 8. For any path on n vertices, the acyclic chromatic number is

$$a[C(P_n)] = n - 1, \quad \text{for all } n \geq 4.$$

Proof. Color the vertices of $C(P_n)$ as given in the algorithm 2.2. The color class of 1 is $\{v_1, v_2\}$ and that of 2 is $\{v_3, v_4\}$. The color class of j is

$\{v_{j+2}, 5 \leq j \leq n-2\}$ whereas the color class of $n-1$ is $\{x_k; 1 \leq k \leq n-1\}$.
Case(i): Consider the colors 1 and 2. The induced subgraph of these color classes contain a bicolored $(1-2)$ -path, $v_2v_4v_1v_3$. Since v_2 and v_3 are non-adjacent, the induced subgraph is a C_{12} -free graph.

Case(ii): Consider the colors 1 and j , $3 \leq j \leq n-2$. The induced subgraph of these color classes contains a bicolored path $v_2v_{j+2}v_1$. As v_1 and v_2 are non-adjacent, the induced subgraph is C_{1j} -free graph.

Case(iii): Consider the colors 2 and k , $3 \leq k \leq n-2$. By the same argument as in case (ii), the induced subgraph is a C_{2k} -free graph.

Case(iv): Consider i and $n-1$, $i = 1, 2$. The induced subgraph of color classes of i and $n-1$ contains only the bicolored $[i-(n-1)]$ path, $v_i x_i v_{i+1} x_{i+1}$, but not bicolored cycle. Hence, the induced subgraph is a $C_{i(n-1)}$ -free graph.

Case(v): Consider j and $n-1$, $3 \leq j \leq n-2$. The induced subgraph of these color classes contains only the bicolored $[j-(n-1)]$ -path $x_{j+1} v_{j+2} x_{j+2}$. So, the induced subgraph is a $C_{j(n-1)}$ -free graph. Thus, the induced subgraph of any two color classes is acyclic and therefore, the coloring given in the algorithm 2.2 is an acyclic coloring. As minimum $(n-1)$ colors are required for acyclic coloring, we have $a[C(P_n)] = n-1$, for all $n \geq 4$.

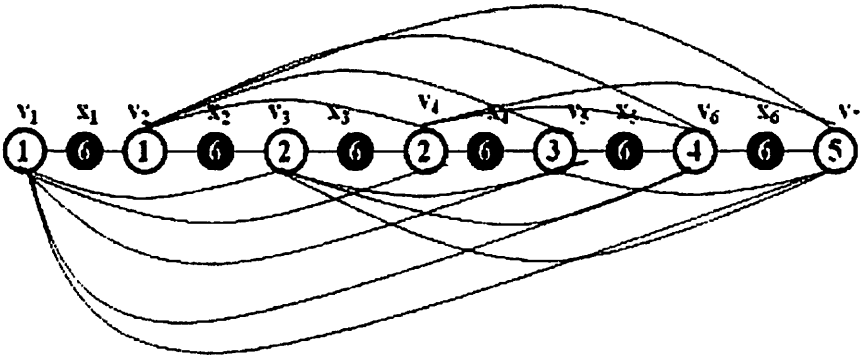


Fig.1. $a[C(P_7)] = 6$

3. Acyclic Coloring of $C[F_{m,n}]$

In Fan graph $F_{m,n}$, let v_1, v_2, \dots, v_m be the vertices of null graph on m vertices and w_1, w_2, \dots, w_n be the vertices of path P_n . Let us denote the newly added vertex on the edge joining v_i and w_j by f_{ij} ($1 \leq i \leq m$, $1 \leq j \leq n$). Let us denote the newly added vertex on the edge joining w_k and w_{k+1} by g_k ($1 \leq k \leq n-1$).

3.1. Structural properties of $C[F_{m,n}]$

1. $\langle v_k, k = 1 \text{ to } m \rangle$ form a clique of order $m + 1$
2. For each $1 \leq i \leq m, 1 \leq j \leq n$, the neighbors of f_{ij} are $\{v_i, w_j\}$.
3. For each $k = 1 \text{ to } n - 1$, the neighbors of g_k are $\{v_k, v_{k+1}\}$.
4. For each $k = 2 \text{ to } n - 1$, the neighbors of v_k are $\{g_{k-1}, g_k, v_1, v_2, \dots, v_{k-2}, v_{k+2}, \dots, v_n\}$ and the neighbors of v_1 are $\{g_1, v_3, \dots, v_n\}$ whereas the neighbors of v_n are $\{g_{n-1}, v_1, v_2, \dots, v_{n-2}\}$.
5. $\langle v_i, x_j ; i = 1 \text{ to } n \text{ and } j = 1 \text{ to } n - 1 \rangle$ is a $C(P_n)$.

3.2 Coloring Algorithm

Consider the acyclic coloring of $C[F_{m,n}]$ as follows. Assign i to $v_i, 1 \leq i \leq m$. Assign 1 to w_1, w_2 and 2 to w_3, w_4 and j to $w_{j+2}, 3 \leq j \leq n - 2, (n - 1)$ to $g_i, 1 \leq i \leq n - 1$. Now, assign colors to $f_{ij}, 1 \leq i \leq m$ and $1 \leq j \leq n$ as follows.

Case (i): Suppose v_i and w_j (adjacent vertices of f_{ij}) have different colors [(i-e) i and $j - 2$ are distinct]. Then, assign k to f_{ij} where $k \neq i, j - 2$ and $1 \leq k \leq \max(m, n - 1)$.

Case (ii): Suppose v_i and w_j have the same color. We have the following cases. (a) Both v_i and w_j have the same color 1. (w_1 and v_1 as well as v_1 and v_2). So, consider f_{11} and f_{12} . Assign 2 to f_{11} , but not to f_{12} . If we assign 2 to both f_{11} and f_{12} , then $v_1 f_{11} w_1 w_4 w_2 f_{12} v_2 v_1$ will be a bicolored (1 - 2) cycle in $C[F_{m,n}]$. Hence, assign $n - 1$ to f_{12} .

(b). Consider f_{23} and f_{24} . Assign 3 to f_{23} and $n - 1$ to f_{24} . (c). Both v_i and w_j have the same color, say k , other than 1 and 2. Then, assign r to f_{ij} only if k is not assigned to f_{ij}' whose adjacent vertices have color r . (i-e) suppose f_{ij} is the newly added vertex on edge joining v_i and w_j which have same color k and f_{ij}' is the newly added vertex on the edge joining v_i' and w_j' with same color r . Then, assign r to f_{ij} only if k is not assigned to f_{ij}' . Otherwise, $v_i f_{ij} w_j w_j' f_{ij}' v_i' v_i$ will become a bicolored ($k - r$) cycle (Refer fig.2).

Theorem 9. For any Fan graph $F_{m,n}$, the acyclic chromatic number is

$$\begin{aligned} a[C(F_{m,n})] &= m, & \text{if } m \geq n \text{ and } m \geq 3 \\ &= n - 1, & \text{if } m < n \text{ and } n \geq 5. \end{aligned}$$

Proof. Color the vertices of $C(F_{m,n})$ as given in the algorithm 3.2.

Case (i): Consider the subgraph induced by the vertices $\{v_1, v_2, \dots, v_m\}$. As it is a clique, it contains no bicolored cycle.

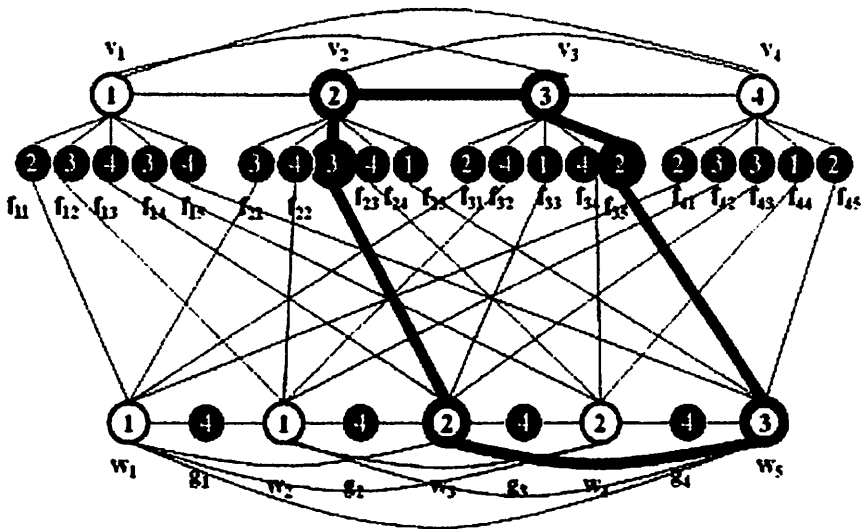


Fig.2. Existence of bicolored cycle

Case (ii): Consider the subgraph induced by $\{w_1, w_2, \dots, w_n, g_1, g_2, \dots, g_{n-1}\}$. It is a $C(P_n)$. As the vertices are colored as given in the algorithm 2.2, it does not contain any bicolored cycle by theorem 2.3.

Case (iii): Any bicolored cycle in $C(F_{m,n})$, if such a cycle exists, should contain the path (of length 2) connecting the vertices v_i and w_j with the same color. But, there exists no such a bicolored cycle in the above coloring. Hence, the coloring given in the algorithm 3.2 is acyclic. To determine the acyclic chromatic number of $C(F_{m,n})$, we have the following cases.

Case (a): $m \geq n$ and $n > 24$. As $\langle v_1, v_2, \dots, v_m \rangle$ form a clique of order m , minimum m colors are required for coloring this induced subgraph and $(n - 1)$ colors are required for $C(P_n)$. Since $m > (n - 1)$, we have $a[C(F_{m,n})] = m, n > 2$. Case (b): $m < n$ and $n \geq 5$. By the same argument as in case (a), we need minimum $(n - 1)$ colors for acyclic coloring of $C(F_{m,n})$ and hence, $a[C(F_{m,n})] = n - 1, n \geq 5$.

4. Acyclic Coloring of $T(P_n)$

4.1. Structural Properties of $T(P_n)$

1. For each $k, k = 2$ to $n - 1$, the neighbors of v_k are $\{x_{k-1}, x_k, v_{k-1}, v_{k+1}\}$ whereas the neighbors of v_1 and v_n are $\{x_1, v_2\}$ and $\{x_{n-1}, v_{n-1}\}$ respectively.

2. For each k , $k = 2$ to $n - 2$, the neighbors of x_k are $\{x_{k-1}, x_{k+1}, v_k, v_{k+1}\}$ whereas the neighbors of x_1 and x_{n-1} are $\{v_1, v_2, x_2\}$ and $\{v_{n-1}, v_n, x_{n-2}\}$ respectively.

4.2. Coloring Algorithm

Input : $T(P_n)$
 $V \leftarrow \{v_1, v_2, v_3, \dots, v_n, x_1, x_2, \dots, x_{n-1}\}$
 for $k = 1$ to $\lfloor (n+2)/3 \rfloor$
 {
 $v_{3k-2} \leftarrow 1$; $x_{3k-2} \leftarrow 2$;
 }
 end for
 for $k = 1$ to $\lfloor (n+1)/3 \rfloor$
 {
 $v_{3k-1} \leftarrow 3$; $x_{3k-1} \leftarrow 1$;
 }
 end for
 for $k = 1$ to $\lfloor n/3 \rfloor$
 {
 $v_{3k} \leftarrow 2$; $x_{3k} \leftarrow 3$;
 }
 end for
 end procedure.
 Output: vertex colored $T(P_n)$

Theorem 10. *The acyclic chromatic number of $T(P_n)$ is*

$$a[T(P_n)] = 3, \quad n \geq 2.$$

Proof. Color the vertices of $T(P_n)$ as given in the algorithm 4.2. The color class of 1 is $\{v_{3k-2}, x_{3k'-1}; k = 1$ to $\lfloor (n+2)/3 \rfloor$ and $k' = 1$ to $\lfloor (n+1)/3 \rfloor\}$ and the color class of 2 is $\{v_{3k}, x_{3k'-2}; k = 1$ to $\lfloor n/3 \rfloor$ and $k' = 1$ to $\lfloor (n+2)/3 \rfloor\}$. The color class of 3 is $\{v_{3k-1}, x_{3k'}; k = 1$ to $\lfloor (n+1)/3 \rfloor$ and $k' = 1$ to $\lfloor n/3 \rfloor\}$.

Case (i): The induced subgraph of the color classes of 1 and 2 is a bicolored path,

- (a). $v_1x_1x_2v_3 \dots x_{n-1}v_n$ when $n = 0(\text{mod}3)$
- (b). $v_1x_1x_2v_3 \dots v_{n-1}v_n$ when $n = 1(\text{mod}3)$
- (c). $v_1x_1x_2v_3 \dots v_{n-1}x_{n-1}$ when $n = 2(\text{mod}3)$

In all the cases, the induced subgraph of 1 and 2 is a C_{12} -free graph.

Case (ii): The induced subgraph of the color classes of colors 1 and 3 is a bicolored path

(a). $v_1v_2x_2x_3 \dots v_{n-1}x_{n-1}$ when $n = 0(mod3)$

(b). $v_1v_2x_2x_3 \dots x_{n-1}v_n$ when $n = 1(mod3)$

(c). $v_1v_2x_2x_3 \dots v_{n-1}v_n$ when $n = 2(mod3)$

In all the cases, the induced subgraph of 1 and 3 is a C_{13} -free graph.

Case (iii): The induced subgraph of the color classes of colors 2 and 3 is a bicolored path

(a). $x_1v_2v_3x_3 \dots v_{n-1}v_n$ when $n = 0(mod3)$

(b). $x_1v_2v_3x_3 \dots v_{n-1}x_{n-1}$ when $n = 1(mod3)$

(c). $x_1v_2v_3x_3 \dots x_{n-1}v_n$ when $n = 2(mod3)$

In all the cases, the induced subgraph of 2 and 3 is a C_{23} -free graph. Thus, the induced subgraph of any two color classes is acyclic and therefore the coloring given in the algorithm 4.2 is an acyclic coloring of $T(P_n)$. As $T(P_n)$ has paths $v_1v_2 \dots v_n$ and $x_1x_2 \dots x_{n-1}$, minimum 3 colors are required for the acyclic coloring. Therefore, we have $a[T(P_n)] = 3$, $n \geq 2$.

5. Conclusion

In this paper, we obtain the following results.

1. $a[C(P_n)] = n - 1$, $n \geq 4$.
2. $a[C(F_{m,n})] = m$, if $m \geq n$ and $m \geq 3$.
 $= n - 1$, if $m < n$ and $n \geq 5$.
3. $a[T(P_n)] = 3$, $n \geq 2$.

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