

# Constructing $\gamma$ -Sets of Grids

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## Abstract

A lot of research has been spent determining the domination numbers,  $\gamma_{m,n}$ , of grid graphs. But relatively little effort has been given to constructing minimum dominating sets of grid graphs. In this paper, we introduce a method for constructing  $\gamma$ -sets of grid graphs  $G_{m,n}$  for all  $m \geq 16$  and  $n \geq 16$ . Further, for  $G_{m,n}$ ,  $m < 16$ ,  $m \neq 12, 13$ , we show how particular  $\gamma$ -sets can be used to construct  $\gamma$ -sets for other grid graphs.

## 1 Introduction

Let  $G = (V, E) = (V(G), E(G))$  be a graph with vertex set  $V = \{v_1, v_2, \dots, v_n\}$  and order  $n = |V|$ . The *open neighborhood* of a vertex  $v$  is the set  $N(v) = \{u | uv \in E\}$  of vertices  $u$  that are adjacent to  $v$ ; the *closed neighborhood* of  $v$  is the set  $N[v] = N(v) \cup \{v\}$ . Similarly, the *closed neighborhood of a set  $S$*  is the set  $N[S] = \bigcup_{v \in S} N[v]$ . A set  $S \subseteq V$  is a *dominating set* of  $G$  if every vertex in  $V - S$  is adjacent to at least one vertex in  $S$ , or equivalently if  $N[S] = V$ . The *domination number*  $\gamma(G)$  of  $G$  equals the minimum cardinality of a dominating set  $S$  of  $G$ ; we say that such a set  $S$  is a  $\gamma$ -set. For more on domination theory, we refer the reader to [15, 16].

An  $m \times n$  grid graph  $G_{m,n}$  has vertex set  $V = \{(i, j) | 1 \leq i \leq m, 1 \leq j \leq n\}$  with  $(i, j)$  adjacent to  $(k, l)$  if  $i = k$  and  $|j - l| = 1$  or  $j = l$  and  $|i - k| = 1$ . We define the *rectilinear distance*, *dist*, between vertices  $(i, j)$  and  $(k, l)$  as

$$\text{dist}((i, j), (k, l)) = |i - k| + |j - l|.$$

We note that  $\text{dist}$  equals the number of edges in a shortest path in  $G_{m,n}$  between  $(i, j)$  and  $(k, l)$ . Further, the set of vertices of the form  $(i, j)$ ,  $1 \leq j \leq n$ , is called the  $i$ th row of  $G_{m,n}$  and the set of vertices of the form  $(i, j)$ ,  $1 \leq i \leq m$ , is called the  $j$ th column of  $G_{m,n}$ . We will denote the domination number of  $G_{m,n}$  by  $\gamma_{m,n}$ .

Numerous papers have been published on the problem of computing the domination number of grid graphs. Beginning in 1983, Jacobson and Kinch [17] determined  $\gamma_{m,n}$  for  $1 \leq m \leq 4$  and all  $n$ . In 1993, Chang and Clark [3] extended this to  $m = 5, 6$  and all  $n$ . In 1989, Hare [4] settled specific cases for  $m = 7, 8, 9, 10, 11$ . In 1998, Fisher [8] developed a method for calculating  $\gamma_{m,n}$  that is described in Spalding's 2001 Ph.D. thesis [21], who gave the values of  $\gamma_{m,n}$  for  $m \leq 19$  and all  $n$ . Several authors [5, 6, 7, 11, 13, 14, 19, 20] have developed techniques for either computing  $\gamma_{m,n}$  exactly or establishing bounds. Alanko et al. [1] used dynamic programming to extend these results to  $m \leq 29$  and all  $n$ . Back in 1992 Chang [2] devoted his Ph.D. thesis to studying the domination numbers of grid graphs, and he conjectured that

$$\gamma_{m,n} = \left\lfloor \frac{(m+2)(n+2)}{5} \right\rfloor - 4,$$

for all values of  $m, n \geq 16$ . This conjecture has recently been proved by Goncalves et al. in 2011 [10]. The following theorem, given in [1] and [10], provides formulae for the value of  $\gamma_{m,n}$  for all  $m, n$ .

**Theorem 1.1** *Let  $G_{m,n}$ ,  $m, n \geq 1$ , be a grid graph. Then*

$$\begin{aligned} \gamma_{1,n} &= \left\lfloor \frac{n+2}{3} \right\rfloor \\ \gamma_{2,n} &= \left\lfloor \frac{n+2}{2} \right\rfloor \\ \gamma_{3,n} &= \left\lfloor \frac{3n+4}{4} \right\rfloor \\ \gamma_{4,n} &= \begin{cases} n+1 & \text{if } n = 5, 6, 9 \\ n & \text{otherwise} \end{cases} \\ \gamma_{5,n} &= \begin{cases} \left\lfloor \frac{6n+6}{5} \right\rfloor & \text{if } n = 7 \\ \left\lfloor \frac{6n+8}{5} \right\rfloor & \text{otherwise} \end{cases} \end{aligned}$$

$$\begin{aligned}
\gamma_{6,n} &= \begin{cases} \lfloor \frac{10n+10}{7} \rfloor & \text{if } n \equiv 1 \pmod{7} \\ \lfloor \frac{10n+12}{7} \rfloor & \text{otherwise} \end{cases} \\
\gamma_{7,n} &= \lfloor \frac{5n+3}{3} \rfloor \\
\gamma_{8,n} &= \lfloor \frac{15n+14}{8} \rfloor \\
\gamma_{9,n} &= \lfloor \frac{23n+20}{11} \rfloor \\
\gamma_{10,n} &= \begin{cases} \lfloor \frac{30n+37}{13} \rfloor & \text{if } n \neq 13, 16, n \equiv 0, 3 \pmod{13} \\ \lfloor \frac{30n+24}{13} \rfloor & \text{otherwise} \end{cases} \\
\gamma_{11,n} &= \begin{cases} \lfloor \frac{38n+21}{15} \rfloor & \text{if } n = 11, 18, 20, 22, 33 \\ \lfloor \frac{38n+36}{15} \rfloor & \text{otherwise} \end{cases} \\
\gamma_{12,n} &= \lfloor \frac{80n+66}{29} \rfloor \\
\gamma_{13,n} &= \begin{cases} \lfloor \frac{98n+111}{33} \rfloor & \text{if } n \not\equiv 14, 15, 17, 20 \pmod{33} \\ \lfloor \frac{98n+78}{33} \rfloor & \text{otherwise} \end{cases} \\
\gamma_{14,n} &= \begin{cases} \lfloor \frac{35n+40}{11} \rfloor & \text{if } n \equiv 18 \pmod{22} \\ \lfloor \frac{35n+29}{11} \rfloor & \text{otherwise} \end{cases} \\
\gamma_{15,n} &= \begin{cases} \lfloor \frac{44n+27}{13} \rfloor & \text{if } n \equiv 5 \pmod{26} \\ \lfloor \frac{44n+40}{13} \rfloor & \text{otherwise} \end{cases} \\
\gamma_{m,n} &= \lfloor \frac{(m+2)(n+2)}{5} \rfloor - 4, m \geq 16, n \geq 16.
\end{aligned}$$

Most of the literature on this problem has focused on developing techniques for computing the value of  $\gamma_{m,n}$ , but little focus has been given to finding a method for constructing  $\gamma$ -sets of grid graphs. Hare [14] displayed many  $\gamma$ -sets for  $m \leq 11$ , and Cockayne et al. [7] gave a method for constructing  $\gamma$ -sets for square grids  $G_{n,n}$ . In this paper, we extend the work of Cockayne et al. to construct  $\gamma$ -sets for many of the values of  $\gamma_{m,n}$  cited in Theorem 1.1.

In Section 2 we extend the work in [7] to give a method for constructing  $\gamma$ -sets of grid graphs for all  $m, n \geq 16$ . In Section 3 we show how this method can be used to construct  $\gamma$ -sets for some values of  $m, n < 16$ ,  $m, n \neq 12, 13$ , and how some of these  $\gamma$ -sets for a fixed  $m$  can be used to construct  $\gamma$ -sets for other values of  $n$ . We also show how these  $\gamma$ -sets define the rate of growth given in Theorem 1.1 for values of  $m < 16$ .

## 2 Constructing $\gamma$ -Sets when $m, n \geq 16$

In this section we show how to construct a  $\gamma$ -set for  $G_{m,n}$  of size  $\lfloor \frac{(m+2)(n+2)}{5} \rfloor - 4$  for  $m, n \geq 16$ . These same techniques can be applied in some cases to construct  $\gamma$ -sets for  $G_{m,n}$  when  $m, n < 16$ . Assume that  $m, n \geq 4$ . Further, let  $G^*$  be an  $(m+2) \times (n+2)$  grid graph where  $(i, j) \in V(G^*)$  if  $0 \leq i \leq m+1$  and  $0 \leq j \leq n+1$ , and  $G_{m,n}$  is a subgraph of  $G^*$ , such that the  $(i, j)$  entry of  $G_{m,n}$  is the  $(i+1, j+1)$  entry of  $G^*$ . We call any vertex in  $V(G^*) - V(G_{m,n})$  a *boundary vertex*. Further, we define a  $(2,1)$ -slant grid  $S(2,1)$  as an infinite graph with vertex set  $V(S(2,1)) = \{(i, j) : i \in \mathcal{Z}, j \in \mathcal{Z} \text{ with the property that if } (i, j) \in V(S(2,1)) \text{ then so are } (i+2, j+1), (i+1, j-2), (i-2, j-1), \text{ and } (i-1, j+2)\}$ . The edges that connect vertices  $(i+2, j+1)$ ,  $(i+1, j-2)$ ,  $(i-2, j-1)$ , and  $(i-1, j+2)$  to  $(i, j)$  are in the edge set  $E(S(2,1))$ . We note that there are 5 isomorphic graphs of  $S(2,1)$  depending on the smallest value of  $k$ ,  $0 \leq k \leq 4$  for which the vertex  $(0, k)$  of  $G^*$  is a vertex of  $S(2,1)$ . We denote these five slant grids as  $S^k(2,1)$ ,  $0 \leq k \leq 4$ . Figure 1 shows a partial slant grid with  $k = 4$ .

Given  $G^*$ , we can overlay  $S^k(2,1)$  on top of  $G^*$  in 5 ways depending on the 'starting' vertex of  $(0, k)$ ,  $0 \leq k \leq 4$ . Let  $V^k = V(S^k(2,1)) \cap G^*$ . In Figure 1, we show this overlay for  $m = 10, n = 10, k = 4$ . To highlight  $G_{m,n}$  as a subgraph of  $G^*$ , we have removed the edges connecting boundary vertices of  $G^*$  to vertices in  $G_{m,n}$ .

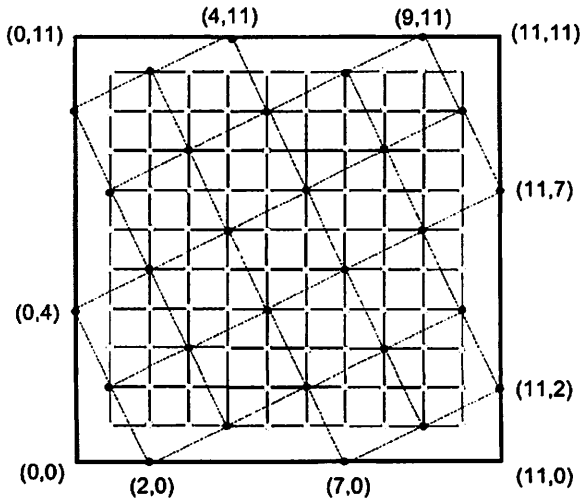


Figure 1:  $G^*$  and  $S^4(2,1)$  when  $m = 10, n = 10, k = 4$

**Lemma 2.1** For all  $0 \leq k \leq 4$  and all  $m, n \geq 1$ ,  $V^k = V(S^k(2, 1)) \cap G^*$  is a dominating set of  $V(G_{m,n})$ .

**Proof.** Let  $(x, y) \in V(G_{m,n})$ , and let  $k \in \{0, 1, 2, 3, 4\}$ . Let  $(i, j) \in V^k$  so that  $(i, j)$  minimizes the rectilinear distance  $dist$  to  $(x, y)$ . If  $dist = 0$ , then  $(x, y) = (i, j) \in V^k$ . Otherwise, we wish to show that  $dist = 1$ . Assume that  $dist \geq 2$ . By definition of  $S^k(2, 1)$ , if  $(i, j) \in V(S^k(2, 1))$  then so are  $(i + 2, j + 1)$ ,  $(i + 1, j - 2)$ ,  $(i - 2, j - 1)$ , and  $(i - 1, j + 2)$ . Note that each of these 4 vertices are rectilinear distance 3 from  $(i, j)$ . Further, any vertex of rectilinear distance 2 to  $(i, j)$  is adjacent to one of these four vertices. Hence, if  $dist \geq 2$ ,  $(x, y)$  is closer to one of  $(i + 2, j + 1)$ ,  $(i + 1, j - 2)$ ,  $(i - 2, j - 1)$ , and  $(i - 1, j + 2)$  than it is to  $(i, j)$ , a contradiction. Hence,  $dist \leq 1$  and  $(x, y)$  is dominated by  $(i, j) \in V^k$ .  $\square$

We now show that we can move the boundary vertices in  $V^k$  to an adjacent vertex in  $G_{m,n}$  to construct a dominating set for  $G_{m,n}$  that is composed entirely of vertices in  $G_{m,n}$ . Let  $A = (0, k_1)$ ,  $B = (k_2, m + 1)$ ,  $C = (n + 1, m + 1 - k_3)$ , and  $D = (n + 1 - k_4, 0)$ , be the closest boundary vertices in  $V^k$ , moving counterclockwise, to the four corners  $((0, 0)$ ,  $(0, m + 1)$ ,  $(n + 1, m + 1)$ , and  $(n + 1, 0)$  of  $G^*$ . We note that  $0 \leq k_i \leq 4$  for  $i \in \{1, 2, 3, 4\}$ . In Figure 1, each of  $k_1, k_2, k_3$ , and  $k_4$  is 4 and  $A = (0, 4)$ ,  $B = (4, 11)$ ,  $C = (11, 7)$ , and  $D = (7, 0)$ .

Depending on the value of  $k$ , we make one of five adjustments to  $V^k$  so that these four boundary vertices in  $V^k$  are moved to vertices in  $G_{m,n}$  to form a dominating set of  $G_{m,n}$ . Some of these adjustments are given in [7] and are referred to as *redomination*. Since these redomination techniques are not exhaustive, we extend them here. To start, all boundary vertices in  $V^k$  that are not involved in one of these five adjustments are replaced by the vertices in  $G_{m,n}$  that are closest in rectilinear distance. Let  $S^k$  equal the set  $V^k$  where the boundary vertices that are not involved in an adjustment have been replaced by their nearest rectilinear distance neighbor in  $G_{m,n}$ . We describe the adjustments to the four boundary vertices closest to the four corners of  $G^*$  in terms of the boundary vertex  $B = (k_2, m + 1)$ . However, these adjustments are equivalent for all four corners.

We first describe the adjustments if  $k_2$  is even. When  $k_2 = 0$  and  $B = (0, m + 1)$ , we simply delete the boundary vertex  $B$  from  $S^k$ , and let  $S^k = S^k - B$ . When  $k_2 = 2$  and  $B = (2, m + 1)$ , we let  $S^k = S^k - \{B, (0, m)\} \cup \{(1, m)\}$ . When  $k = 4$  and  $B = (4, m + 1)$ , let  $S^k = S^k - \{B, (2, m), (0, m - 1)\} \cup \{(3, m), (1, m - 1)\}$ . Note that in each of these cases we remove boundary vertices from the dominating set and add vertices in  $G_{m,n}$  so that all of the vertices of  $G_{m,n}$  remain dominated by  $S$ . Figure 2 shows these adjustments when  $k$  is even.

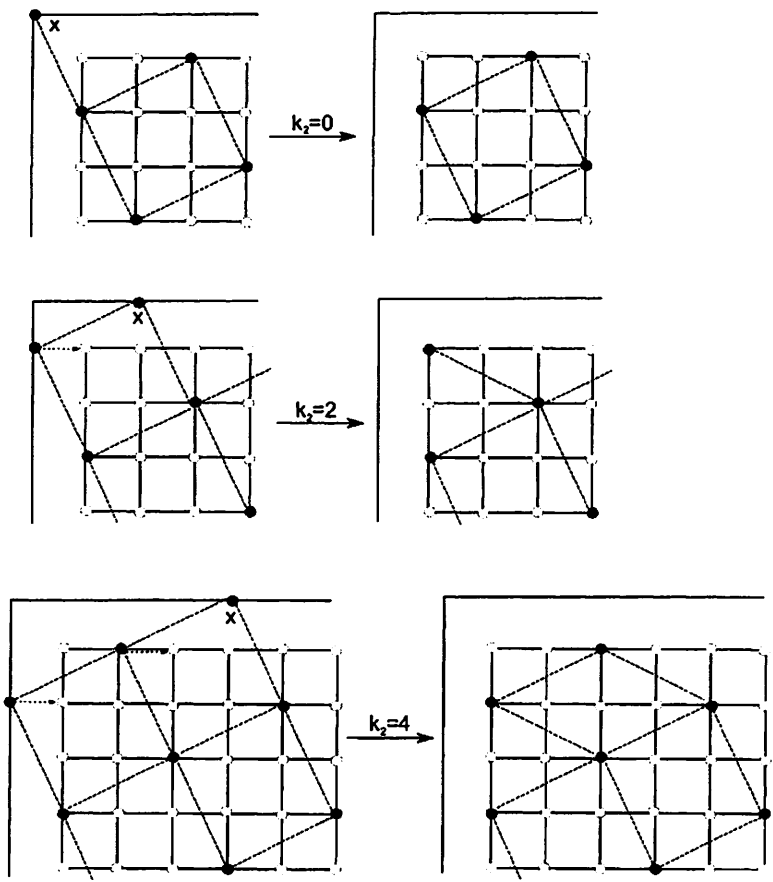


Figure 2: Adjustments for  $k_2 = 0, 2, 4$

When  $k_2$  is odd, there are two types of adjustments: a full adjustment and a partial adjustment. When  $k_2 = 1$ , we first describe the partial adjustment. Here we let  $S^k = S^k - \{B, (0, m-2)\} \cup \{(1, m-1)\}$ . For the full adjustment, we let  $S^k = S^k - \{B, (0, m-2), (2, m-1), (4, m), (6, m+1)\} \cup \{(1, m), (1, m-2), (3, m-1), (5, m)\}$ . These adjustments are shown in Figure 3. We note in Figure 3, the arrows and the  $X$  that have subscripts with a  $p$  denote movement that occurs only in the partial adjustment. The arrows without the  $p$  subscript denote movement that occurs only in the full adjustment.

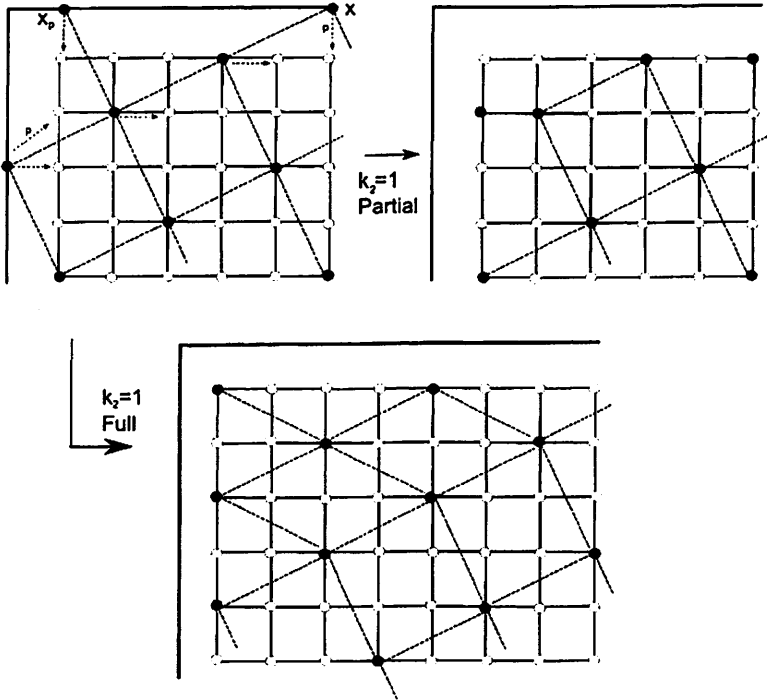


Figure 3: Full and Partial Adjustments for  $k_2 = 1$

When  $k_2 = 3$  and  $B = (3, m+1)$ , we first describe the partial adjustment. Here we let  $S^k = S^k - \{B, (1, m), (0, m-3)\} \cup \{(2, m), (1, m-2)\}$ . For the full adjustment, we let  $S^k = S^k - \{B, (1, m), (0, m-3), (2, m-2), (4, m-1), (6, m), (8, m+1)\} \cup \{(1, m-1), (3, m), (1, m-3), (3, m-2), (5, m-1), (7, m)\}$ . Figures 4 and 5 show the full and partial adjustments when  $k_2 = 3$ .

Figure 6 shows the full adjustments at each of the four corners of  $G_{10,10}$  in Figure 1.

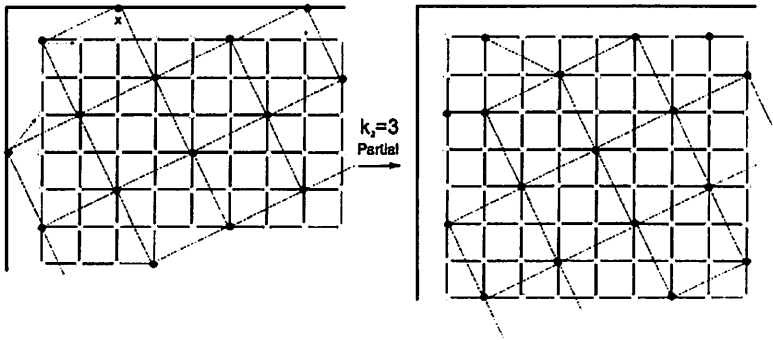


Figure 4: Partial Adjustment for  $k_2 = 3$

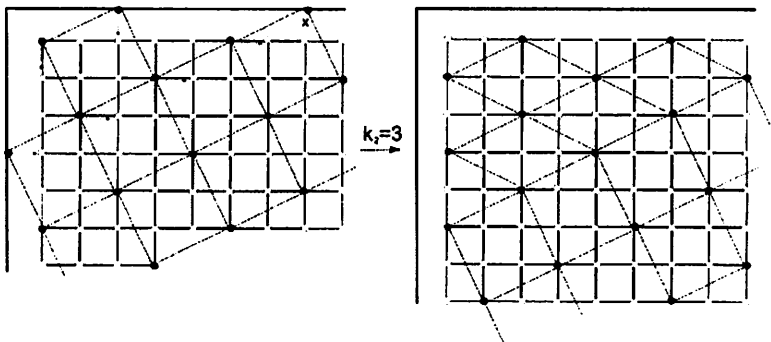


Figure 5: Full Adjustment for  $k_2 = 3$



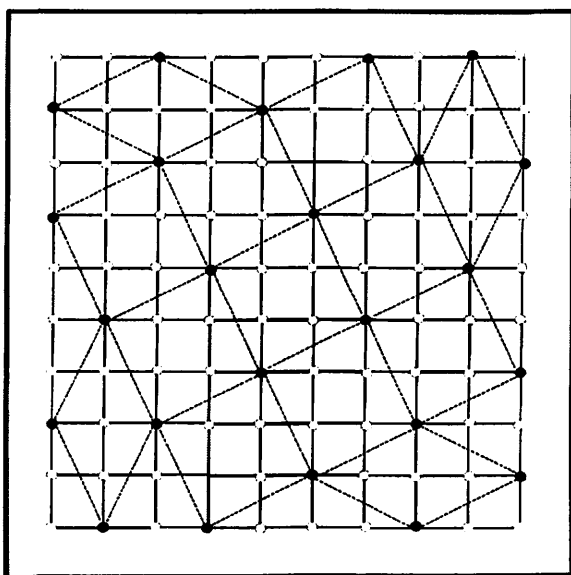


Figure 6: A  $\gamma$ -set for  $G_{10,10}$

We note that not all adjustments at a given corner are possible in conjunction with adjustments at other corners for some small values of  $m$  and  $n$ . For instance, consider the example in Figure 7, where  $m = 5$ ,  $n = 7$ , and  $k = 4$ .

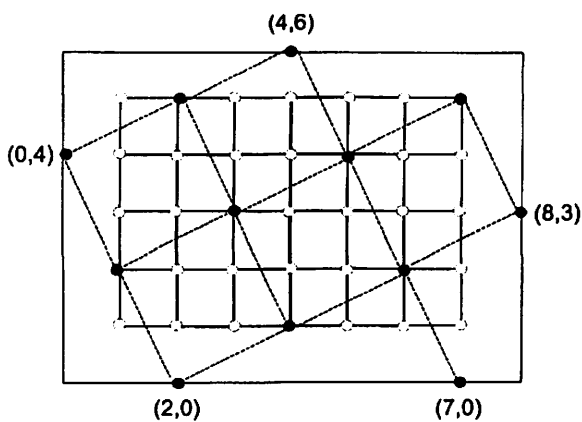


Figure 7:  $S^4(2,1)$  overlaid on  $G_{5,7}$

Note that  $k_1 = 4$  and  $k_2 = 4$ , but the adjustment for  $k_1$  cannot be completed in conjunction with the adjustment for  $k_2$  since the  $k_1$  adjustment requires the deletion of vertex  $(0, 4)$  from  $S^k$  without replacement whereas the  $k_2$  adjustment requires that  $(0, 4)$  be replaced by  $(1, 4)$ . If these two adjustments are made together, we do not end up making full adjustments around each corner. For grids  $G_{m,n}$  where  $m$  and  $n$  take on values that do not allow for full adjustments, we simply make as many adjustments as possible to form a dominating set for  $G_{m,n}$ . Table 1 shows the lower bounds that  $m$  and  $n$  need to be if  $k_i = a$  and  $k_{i+1} = b$  in order to ensure full adjustments. Partial adjustments can be made in the cases where  $k = 1$  and  $k = 3$ . Each of these partial adjustments require  $m, n \geq 7$ .

		b				
		0	1	2	3	4
a	0	4	10	6	12	8
	1	7	13	9	15	11
	2	5	11	7	13	9
	3	8	14	10	16	12
	4	6	12	8	14	10

Table 1: Lower Bounds on  $m$  and  $n$  Ensuring Full Adjustments

Knowing  $k$ , it is easy to compute the four values of  $k_1, k_2, k_3$ , and  $k_4$  for each value of  $m$  and  $n$  modulo 5. Given  $k = k_1$ , we note that the vertex  $(0, (-3k \bmod 5))$  is in  $S^k(2, 1)$  but not  $G^*$ . However, each vertex of the form  $(0, (5q - 3k) \bmod 5) \in G^*$  for all  $1 \leq q \leq \lfloor \frac{n+1-(3k \bmod 5)}{5} \rfloor$ . Hence,  $k_4 = (n + 1 - 3k) \bmod 5$ . We can apply the same reasoning to see that  $k_3 = (m + 1 - k - 3(m + 1)) \bmod 5$  and  $k_2 = (3k + 2(m + 1)) \bmod 5$ . Table 5, in the Appendix, shows the values of  $k_1, k_2, k_3$ , and  $k_4$  for each value of  $m$  and  $n$  modulo 5. In the table,  $k_1$  is listed in the lower left-hand corner,  $k_2$  the upper left-hand corner,  $k_3$  the upper right-hand corner, and  $k_4$  the lower right-hand corner for each  $m$  and  $n$ .

So under the conditions on  $m$  and  $n$  as defined by Tables 1 and 5 that allow for full adjustments or partial adjustments in the case of  $k = 1$  or  $k = 3$ , adjusting  $V^k$  to form set  $S^k$ , we see that  $S^k$  contains one fewer vertex per corner adjustment than  $V^k$ , so that the final set  $S^k$  has four fewer vertices than  $V^k$ . Since in each case we replace a boundary vertex with one from  $G_{m,n}$ , we have shown the following lemma.

**Lemma 2.2** *The set  $S^k \subset V(G_{m,n})$  is a dominating set for  $G_{m,n}$  of size  $|S^k| = |V^k| - 4$ , provided  $m, n$  are of sufficient size, as defined in Table 1, that allow for 4 full or partial corner adjustments to be performed.*

Since we know that  $S^k$  is a dominating set for  $G_{m,n}$  for certain values of  $m$  and  $n$ , we turn our attention to the size of  $V^k$ . We will show that either  $|V^k| = \lfloor \frac{(m+2)(n+2)}{5} \rfloor$  or  $|V^k| = \lfloor \frac{(m+2)(n+2)}{5} \rfloor + 1$ , by examining how many vertices in  $V^k$  appear in each column of  $G^*$ . For each  $1 \leq i \leq n+2$ , let  $\bar{j}_i$  be the smallest value of  $j$  such that  $(i, j) \in V^k$ . Note that  $\bar{j}_0 = k$  and for each  $1 \leq i \leq n+2$ ,  $\bar{j}_i = (\bar{j}_{i-1} + 2) \bmod 5$ . Further,  $(i, \bar{j}_i + 5q)$ ,  $1 \leq q \leq \lfloor \frac{m+1-\bar{j}_i}{5} \rfloor$ , will also be in  $V^k$ . Hence, given  $\bar{j}_i$ , there will be  $\lfloor \frac{m+1-\bar{j}_i}{5} \rfloor + 1 = \lfloor \frac{m+6-\bar{j}_i}{5} \rfloor$  vertices in column  $i$ . We note that  $j$  ranges from 0 to 4 and depends on the value of  $k$ . Note that any span of 5 columns of  $G^*$  contains

$$\sum_{q=0}^4 \lfloor \frac{m+6-q}{5} \rfloor = \sum_{q=0}^4 \lfloor \frac{m+2+q}{5} \rfloor = m+2 \text{ vertices of } V^k. \quad (1)$$

Thus, if  $n+2 = 5s+r$ ,

$$\begin{aligned} |V^k| &= (m+2) \lfloor \frac{n+2}{5} \rfloor + \sum_{q=0}^{r-1} \lfloor \frac{m+2+(4-\bar{j}_q)}{5} \rfloor, \\ &= \frac{(m+2)(n+2-r)}{5} + \sum_{q=0}^{r-1} \lfloor \frac{m+2+(4-\bar{j}_q)}{5} \rfloor. \end{aligned}$$

If  $m+2+(4-\bar{j}_q) = 5t_q + d_q$ , for each  $0 \leq q \leq r-1$ , we have

$$\begin{aligned} |V^k| &= \frac{(m+2)(n+2)}{5} - \frac{r(m+2)}{5} + \sum_{q=0}^{r-1} \frac{m+2+(4-\bar{j}_q) - d_q}{5}, \\ &= \frac{(m+2)(n+2)}{5} - \frac{r(m+2)}{5} + \frac{r(m+2)}{5} + \sum_{q=0}^{r-1} \frac{(4-\bar{j}_q - d_q)}{5}, \\ &= \frac{(m+2)(n+2)}{5} + \sum_{q=0}^{r-1} \frac{q^*}{5}, \end{aligned}$$

where  $q^* = 4 - \bar{j}_q - d_q$ . Let's examine the term  $q^*$  more closely. This term represents the amount we round up or round down from  $m+2$  depending on the value of  $\bar{j}_q$ . So, if  $4 - \bar{j}_q \geq 5 - (m+2)$ , then  $q^* = 5 - (m+2)$  otherwise  $q^* = -(m+2)$ . Given (1) above, we know that

$$\sum_{q=0}^4 \frac{(4 - \bar{j}_q - d)}{5} = 0.$$

However, for values of  $r < 5$ , this sum could vary away from 0. Since  $\bar{j}_i = (\bar{j}_{i-1} + 2) \pmod{5}$ ,

$$-\frac{4}{5} \leq \sum_{q=0}^4 \frac{(4 - \bar{j}_q - d)}{5} \leq \frac{4}{5}.$$

Hence,

$$\left\lfloor \frac{(m+2)(n+2)}{5} \right\rfloor \leq |V^k| \leq \left\lfloor \frac{(m+2)(n+2)}{5} \right\rfloor + 1.$$

Table 2 shows the cases for each value of  $m$ ,  $n$  and  $k$  for which  $V^k$  either is exactly  $\lfloor \frac{(m+2)(n+2)}{5} \rfloor$  or is one more than this value.

n (mod 5)	m = 0 (mod 5)					m = 1 (mod 5)				
	k = 0	k = 1	k = 2	k = 3	k = 4	k = 0	k = 1	k = 2	k = 3	k = 4
0	1	1	1	1	0	0	0	1	0	0
1	1	0	0	0	0	1	1	1	0	1
2	1	1	1	0	0	0	1	1	0	0
3	0	0	0	0	0	0	0	0	0	0
4	1	1	0	0	0	1	1	1	0	0
n (mod 5)	m = 2 (mod 5)					m = 3 (mod 5)				
	k = 0	k = 1	k = 2	k = 3	k = 4	k = 0	k = 1	k = 2	k = 3	k = 4
0	1	0	1	1	0	0	0	0	0	0
1	1	0	1	0	0	0	0	0	0	0
2	0	0	1	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0	0	0	0
4	1	1	1	1	0	0	0	0	0	0
n (mod 5)	m = 4 (mod 5)									
	k = 0	k = 1	k = 2	k = 3	k = 4					
0	1	0	1	0	0					
1	1	0	1	0	1					
2	1	1	0	0	1					
3	0	0	0	0	0					
4	1	0	0	0	0					

Table 2:  $|V^k| - (\lfloor \frac{(m+2)(n+2)}{5} \rfloor)$

It is easy to see from Table 2 that for each value of  $m$  and  $n$ , there is at least one value of  $k$  in which we can construct a set  $V^k$  of cardinality  $\lfloor \frac{(m+2)(n+2)}{5} \rfloor$  and hence a dominating set  $S^k$  of cardinality  $\lfloor \frac{(m+2)(n+2)}{5} \rfloor - 4$ . From [10], we know that for each  $m, n \geq 16$ ,  $\gamma_{m,n} = \lfloor \frac{(m+2)(n+2)}{5} \rfloor - 4$ . Hence, we can use the above construction methods to create a  $\gamma$ -set for  $G_{m,n}$ , for  $m, n \geq 16$ . Thus, we have proven the following.

**Theorem 2.3** For all  $m, n \geq 16$ , there exists a  $k$ ,  $0 \leq k \leq 4$ , such that  $S^k$  is a  $\gamma$ -set of  $G_{m,n}$ .

### 3 Constructing $\gamma$ -sets when $m, n \leq 15$ , $m \neq 12, 13$

As outlined in [10] and [1], when  $m, n \leq 15$ ,  $\gamma_{m,n}$  may or may not equal  $\left\lfloor \frac{(m+2)(n+2)}{5} \right\rfloor - 4$ . In this section we will show that in some of the cases where  $m, n \leq 15$ ,  $\gamma_{m,n} = \left\lfloor \frac{(m+2)(n+2)}{5} \right\rfloor - 4$ , and our methods can be used to construct a  $\gamma$ -set for  $G_{m,n}$ . In other cases, we will show how a repeatable  $\gamma$ -pattern can be used to construct  $\gamma$ -sets.

We begin with an example. Consider  $G_{5,6}$ . As in the previous section, we form  $G^*$  and overlay a slant grid. From Table 5 we see that when  $m \equiv 0 \pmod{5}$  and  $n \equiv 1 \pmod{5}$  there is a value of  $k$ , namely  $k = 0$ , that produces values  $k_1 = 0, k_2 = 2, k_3 = 0$  and  $k_4 = 2$  that allow for full adjustments to be made at each corner. Given  $m, n$ , we say that any value of  $k$  that allows for four full or partial adjustments to be made is a *compatible* value of  $k$ . Also, note for this value of  $k$ ,  $k_1 = k_3$  and  $k_2 = k_4$ . In these cases, we say that this value of  $k$  produces *symmetric* adjustments. In the previous example given in Figures 1 and 6, the compatible value of  $k = 4$  produced symmetric adjustments. In Figure 8, we show how performing these adjustments produces a dominating set of size 8 for  $G_{5,6}$  which, from Theorem 1.1, is optimum.

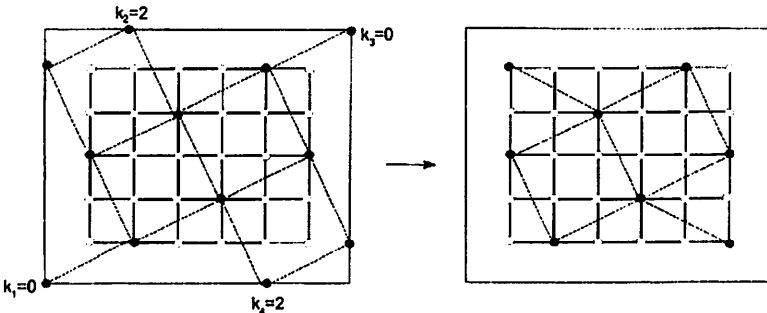


Figure 8: Constructing a  $\gamma$ -set of  $G_{5,6}$

Further, note that when symmetric adjustments are made, in some cases of  $m$  and  $n$ , the resulting  $\gamma$ -set can be combined with itself to produce a

$\gamma$ -sets of larger grids. For example, the  $\gamma$ -set we constructed for  $G_{5,6}$  can be flipped over a line of symmetry drawn on column 6 (or column 1) to produce a dominating set of  $G_{5,11}$  of size  $2 * \gamma_{5,6} - \gamma(n)$ , where  $\gamma(n)$  equals the number of  $\gamma$ -vertices that lie in column  $n$  of  $G_{m,n}$ , in this case  $n = 6$ . We note that flipping  $\gamma_{5,6}$  over a line of symmetry drawn on row 1 or row  $m$  also produces a  $\gamma$ -set for  $\gamma_{9,6}$ . Figure 9 shows how this flipping can create a  $\gamma$ -set for  $G_{5,16}$ .

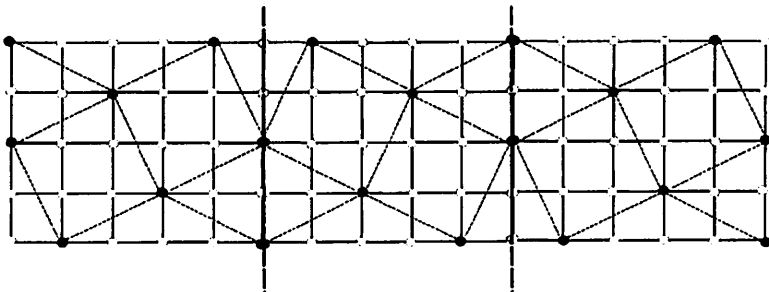


Figure 9: Constructing a  $\gamma$ -set of  $G_{5,16}$

In general, for  $m, n \leq 15$ , a compatible value for  $k$  that allows symmetric adjustments to produce a  $\gamma$ -set for  $G_{m,n}$  is called *repeatable*, and this  $\gamma$ -set can be flipped over a line of symmetry drawn at column  $n$  (or column 1) of  $G_{m,n}$ . We say that this repeatable  $\gamma$ -set has a *rate of growth*,  $R$ , where

$$R = \frac{\gamma_{m,n} - \gamma(n)}{n - 1}.$$

For our example highlighted in Figure 8, note that the rate of growth is  $R = \frac{8-2}{5} = \frac{6}{5}$  which matches the rate of growth highlighted in Theorem 1.1 for  $\gamma_{5,n}$ . If  $R < \frac{(m+2)(n+2)}{5}$ , then a repeatable  $\gamma$ -set has a rate of growth slower than that of our construction method to produce set  $S$ , and thus can be used to produce a dominating set of size smaller than  $\lfloor \frac{(m+2)(n+2)}{5} \rfloor - 4$ . In Figures 10 - 14, we show the repeatable  $\gamma$ -sets matching the growth rates indicated in Theorem 1.1 for  $2 \leq m \leq 11$  and  $14 \leq m \leq 15$ .

The key to achieving a rate of growth slower than  $\frac{(m+2)(n+2)}{5}$  lies in maximizing  $\gamma(n)$  or  $\gamma(1)$ . As an example, note that when  $m = 8$ ,  $k = k_1 = 2$  implies  $k_2 = 4$ , and there are 3  $\gamma$ -vertices on the first column border. This value of  $k$  maximizes the number of  $\gamma$ -vertices that exist on the first column border for  $m = 8$ . From Table 2, we note that  $k = 2$  allows for the creation of a  $\gamma$ -set of size  $\lfloor \frac{(m+2)(n+2)}{5} \rfloor - 4$  for all values of  $n \geq 8$ . In Figures 15 and

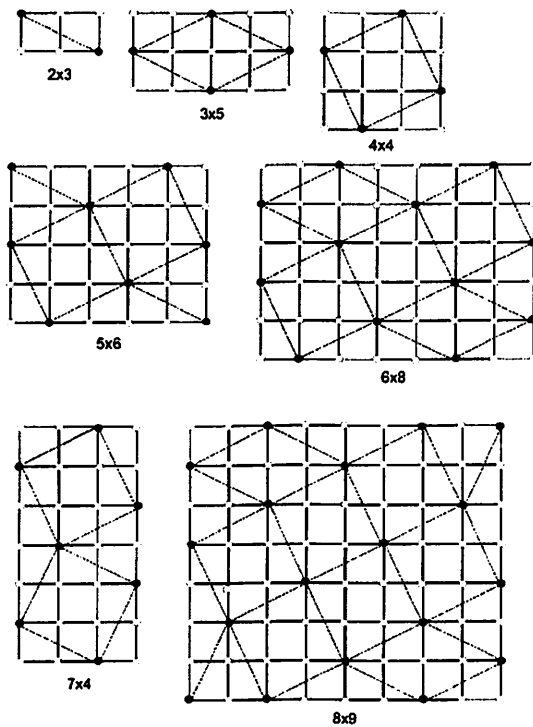
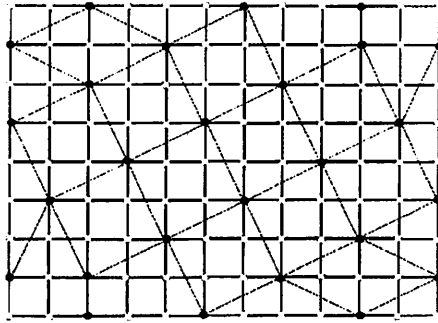
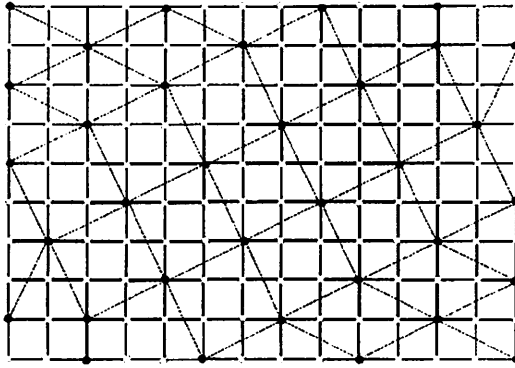


Figure 10: Repeating Patterns for  $n = 2$  through  $n = 8$

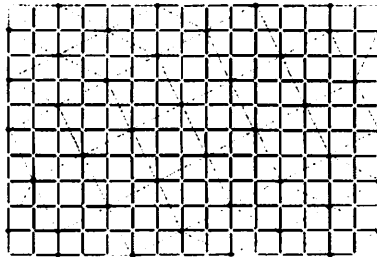


9x12



10x14

Figure 11: Repeatable Patterns for  $n = 9$  and  $n = 10$



11x16

Figure 12: Repeatable Pattern for  $n = 11$



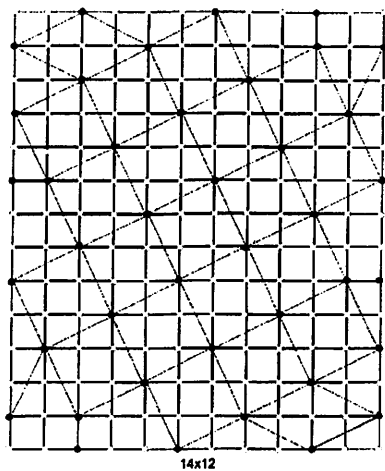


Figure 13: Repeatable Pattern for  $n = 14$

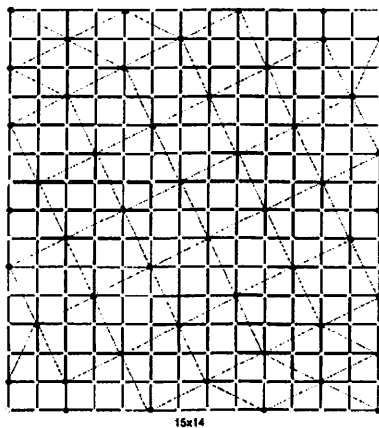


Figure 14: Repeatable Pattern for  $n = 15$

16 we show the optimal  $\gamma$ -sets for  $G_{8,n}$  for  $8 \leq n \leq 14$ .

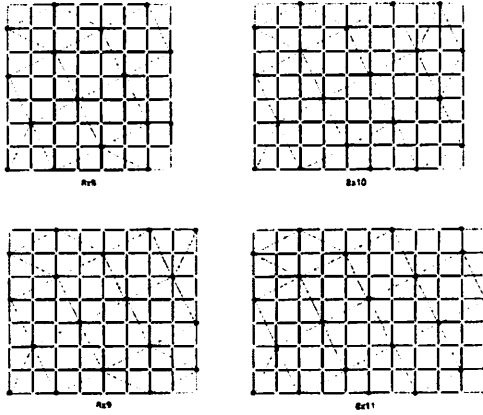


Figure 15:  $k = 2$   $\gamma$ -sets for  $n = 8$

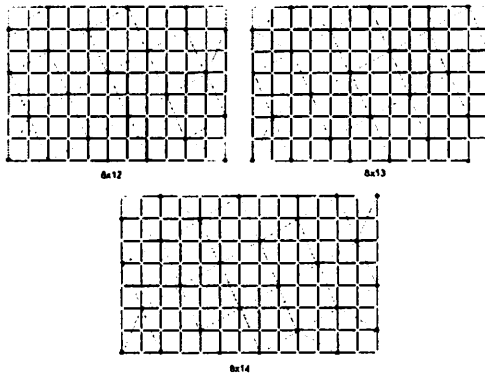


Figure 16:  $k = 2$   $\gamma$ -sets for  $n = 8$

Although the  $\gamma$ -set for  $G_{8,9}$  is used to define the rate of growth of  $\gamma$ -sets for  $G_{8,n}$  for all  $n \geq 8$ , the  $\gamma$ -set in Figure 15 for  $G_{8,8}$  can be repeated one time (since it is not symmetric) to create a  $\gamma$ -set for  $G_{8,15}$  as seen in the Figure 17. It can then be combined with the symmetric  $\gamma$ -set for  $G_{8,9}$  to create a  $\gamma$ -set for  $G_{8,23}$ . Continuing this process for other values of  $n$  shows that we can create a  $\gamma$ -set for  $G_{8,n}$  for any value of  $n \geq 8$ .

Tables 3 and 4 shows how to construct a  $\gamma$ -set for  $G_{m,n}$  by indicating the value of  $k$  to use or by showing how to use a smaller, repeatable  $\gamma$ -

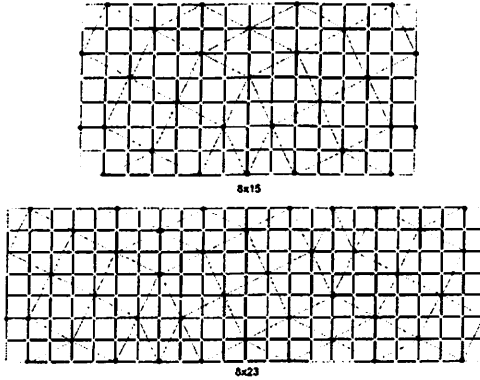


Figure 17:  $\gamma$ -sets for  $G_{8,15}$  and  $G_{8,23}$

set to construct the  $\gamma$ -set. The notation  $a|a$  indicates that the  $\gamma$ -set for  $G_{m,2a-1}$  is constructed by flipping the repeatable  $\gamma$ -set for  $G_{m,a}$  once. The notation  $a|a|a|a$  indicates that the  $\gamma$ -set for  $G_{m,4a-3}$  is formed by flipping the symmetric, repeatable  $\gamma$ -set for  $G_{m,a}$  four times. The notation  $a - a$  means that the  $\gamma$ -set for  $G_{m,2a}$  is formed from two copies of  $G_{m,a}$  that share no vertices. This occurs for some special cases when  $m \leq 4$  and  $m = 7$ . When  $m = 10$ , the  $\gamma$ -set for  $G_{10,4}$  is the same as the  $\gamma$ -set indicated in the chart of  $G_{4,10}$ . It is denoted as  $4^*$ . This is a repeatable  $\gamma$ -set and can be used to generate  $\gamma$ -sets of  $G_{10,3j+1}$  for  $1 \leq j \leq 16$ .

$m : n$	2	3	4	5	6	7	8	9	10	11	12	13	14	15
2	$k=0$	$k=2$	$k=2$	3 3	3-3	3 3 3	3 3-3	3 3 3 3	3 3 3-3	3 3 3 3 3	...	...	...	...
3		$k=0$	$k=3$	$k=2^*$	$k=2$	$k=2$	$k=2$	5 5	5-5	$k=2$	$k=2$	5 5 5	5 5-5	5-5-5
4			$k=0$	$k=0$	$k=0$	4 4	4-4	$k=0$	4 4 4	4 4-4	4-4-4	4 4 4 4	4 4 4-4	4 4-4-4
5				$k=4$	$k=0$	$k=1$	$k=0$	$k=0$	$k=0$	6 6	6 6-6	6 6	6 6	6 10
6					$k=3$	$k=0$	$k=0$	$k=0$	$k=0$	$k=0$	$k=0$	$k=0$	8 7	8 8
7			4	$k=0$	$k=3$	4 4	4-4	$k=2$	4 4 4	4 4-4	4-4-4	4 4 4 4	4 4 4-4	4 4-4-4
8							$k=2$	$k=2$	$k=2$	$k=2$	$k=2$	$k=2$	$k=2$	8 8
9							$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	8 8
10			4*			4 4	$k=2$	$k=3$	4 4 4	$k=3$	$k=3$	4 4 4 4	$k=3$	$k=4$
11							$k=4$	$k=4$	$k=4$	$k=3$	$k=4$	$k=4$	$k=3$	8 8
14						$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	$k=3$	8 8
15							$k=3$	$k=3$	$k=4$	$k=3$	$k=3$	$k=3$	$k=3$	$k=4$

Table 3: Generation of  $\gamma$ -sets of  $G_{m,n}$  for  $m \leq 15$  and  $n \leq 15$ ,  $m \neq 12, 13$

There are a couple of special cases in the table above. The  $\gamma$ -sets for  $G_{11,18}$ ,  $G_{11,20}$ , and  $G_{11,22}$  do not follow the construction methods from Section 1 or the repeatable pattern of this section. Figure 18 shows the  $\gamma$ -set for  $G_{11,18}$ . The  $\gamma$ -sets for  $G_{11,20}$  and  $G_{11,22}$  have a similar construction to  $\gamma_{11,18}$ . The construction method for these  $\gamma$ -sets are more complicated and resemble the construction of  $\gamma$ -sets when  $m = 12, 13$ . These constructions do not use

$m : n$	16	17	18	19	20	21	22	23	24	25	26	27
8	8 9	9 9	9 10	9 11	9 12	9 13	9 14	8 8 9	8 9 9	9 9 9	9 9 10	...
9	8 9	9 9	9 10	8 12	9 12	10 12	11 12	12 12	12 13	12 14	8 12 8	8 12 9
10	4 4 4 4 4	4 4 4 4 - 4	15 4	4 4 4 4 4	$k = 4$	15 4 4	4 4 4 4 4 4	20 4 4	15 4 4 4	4 4 4 4 4 4 4	20 4 4	14 14
11	$k = 3$	8 10	**	8 12	**	12 10	**	12 12	9 16	12 14	11 16	14 14
14	8 9	8 10	8 11	8 12	9 12	10 12	11 12	12 12	12 13	12 14	8 12 8	8 12 9
15	$k = 3$	$k = 3$	$k = 3$	$k = 3$	$k = 4$	11 11	11 12	12 12	11 14	12 14	13 14	14 14

Table 4: Generation of  $\gamma$ -sets of  $G_{m,n}$ ,  $m \leq 15$ ,  $m \neq 12, 13$ ,  $16 \leq n \leq 27$

a repeatable pattern to create larger  $\gamma$ -sets and are left for future work.

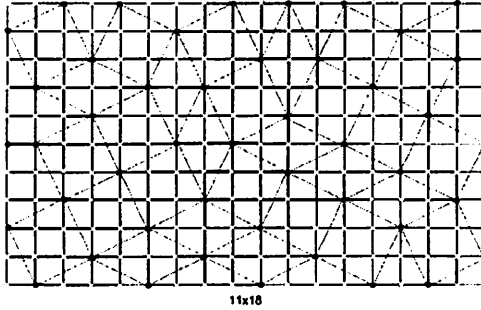


Figure 18:  $\gamma$ -set for  $G_{11,18}$

## 4 Conclusion and Future Work

In this paper, we present a method for constructing  $\gamma$ -sets for  $G_{m,n}$ , for all  $m \neq 12, 13$  and all  $n$ , that exhibits the rate of growth noted in Theorem 1.1. However, there is still work to be done in this area. For  $m < 16$ , our method constructs one  $\gamma$ -set which can be used to create other  $\gamma$ -sets through a repeatable transformation. In some cases, when  $m = 12, 13$  for instance, known  $\gamma$ -sets of  $G_{m,n}$  do not follow these repeatable transformations but are much more complicated. From Theorem 1.1, the growth rate  $R$  when  $m = 12$  is  $\frac{80}{29}$  and when  $m = 13$  is  $\frac{98}{33}$ . If our construction approach works in this case,  $G_{12,30}$  should have a  $\gamma$ -set of size 85 with 5  $\gamma$ -vertices on columns 1 and 30. Further,  $G_{13,34}$  should have a  $\gamma$ -set of size 103 with 5  $\gamma$ -vertices on columns 1 and 34. However, the authors have not been able to construct such  $\gamma$ -sets using various construction techniques. In fact, known  $\gamma$ -sets for these cases indicate more complicated transformations that do not conform with the ones presented here.

Although we have given a method for constructing a  $\gamma$ -set of a grid graph

$G_{m,n}$ , there are other  $\gamma$ -sets of  $G_{n,m}$  that are not constructible using our method, and that are not just flips and rotations of the  $\gamma$ -sets that we construct. We do not know of techniques or transformations to create these other  $\gamma$ -sets.

It seems reasonable that the techniques discussed here could be used to construct  $\gamma$ -sets of related graphs. For instance, the methods presented here could be extended to construct  $\gamma$ -sets of  $P_m \square C_n$  and  $C_m \square C_n$ . Further, these techniques could construct sets that are optimal for other types of domination. The  $\gamma$ -sets constructed here for  $m, n \geq 16$  are independent sets, proving that  $\gamma_{m,n} = i_{m,n}$  for  $m, n \geq 16$ , where  $i_{m,n}$  denotes the independent domination number of  $G_{n,m}$ . However, if  $m, n \leq 15$ , many of the  $\gamma$ -sets created have adjacent vertices (c.f. Figures 11-14). For these cases, we do not know if it is possible to transform the vertices so that there is an independent  $\gamma$ -set of the same size. Identifying the values of  $m$  and  $n$  where  $\gamma_{m,n} < i_{m,n}$  for  $m, n < 16$  is an open question for future work.

Finally, the computation of the total domination numbers of grid graphs has not been studied as much as the domination numbers of grid graphs. Gravier [9] has determined the total domination numbers of  $G_{m,n}$  for  $m \leq 4$  and all  $n$ , while Klobucar [18] has determined these numbers for  $m = 5, 6$  and all  $n$ . However, it appears that an analysis similar to that given here applies for the construction of  $\gamma_t$ -sets of grids. Figure 1 in [9] suggests that a transformation of the slant grid overlay given here could be applied to produce  $\gamma_t$ -sets of grid graphs.

## 5 Appendix

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