

On the Upper Bound for the Number of Spanning Trees of a Connected Graph

Rao Li

Dept. of Mathematical Sciences
University of South Carolina Aiken
Aiken, SC 29801
Email: raol@usca.edu

Abstract

Das [4], Feng et al. [5], and Li et al. [13] obtained upper bounds for the number of spanning trees of a connected graph. Using some ideas in [4], [5], and [13] and other established results, we obtain new upper bounds for the number of spanning trees of a connected graph.

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1. Introduction

We consider only finite, undirected, connected graphs without loops or multiple edges. Notation and terminology not defined here follow those in [2]. Let G be a graph of order n . We use G^c to denote the complement of G . We assume that the vertices in G are ordered such that $d_1 \geq d_2 \geq \dots \geq d_n$, where d_i , $1 \leq i \leq n$, is the degree of vertex v_i in G . For vertex v_i , we use $N(v_i)$ to denote its neighbors. Grone and Merris in [8] introduced the definition of $d_i^*(G)$. They defined $d_i^*(G)$ as $|\{v \in V(G) : d(v) \geq i\}|$. For two distinct vertices u and v in G , we define $c_{u,v}$ as $|N(u) \cap N(v)|$; if $d(u) = d_1$ and $d(v) = d_2$, we further define

$$a_{u,v} := \begin{cases} \frac{d_2 + 2 + \sqrt{(d_2 + 2)^2 - 8d_2 + 4c_{u,v}}}{2} & \text{if } uv \in E(G), \\ \frac{d_2 + 1 + \sqrt{(d_2 + 1)^2 - 4c_{u,v}}}{2} & \text{if } uv \notin E(G). \end{cases}$$

and $a := \max\{a_{u,v} : \text{where } d(u) = d_1, d(v) = d_2, \text{ and } u \neq v\}$. Notice that $a_{u,v}$ follows from [3]. We use t to denote the number of spanning

trees of a labeled graph G . $G \vee H$ is defined as the join of the vertex-disjoint graphs of G and H . The Laplacian of a graph G is defined as $L(G) = D(G) - A(G)$, where $D(G)$ is the diagonal matrix of the degree sequence of G and $A(G)$ is the adjacency matrix of G . The eigenvalues $\lambda_1(G) \geq \lambda_2(G) \geq \dots \geq \lambda_n(G) = 0$ of $L(G)$ are called the Laplacian eigenvalues of the graph G . We use S_n to denote the star graph $K_{1, n-1}$. We also use $K_n - e$ to denote the graph obtained by removing one edge e from K_n . Finally, we use $S_n + e$ to denote the graph obtained by adding one edge e to $K_{1, n-1}$.

Using the Matrix Tree Theorem, i.e., $tn = \lambda_1 \lambda_2 \dots \lambda_{n-1}$, several authors obtained upper bounds for the number of spanning trees of a graph. In particular, Das in [4] proved the following theorem.

Theorem A Let G be a connected graph with n vertices and e edges. Then

$$t \leq \left(\frac{2e - d_1 - 1}{n - 2} \right)^{n-2}.$$

Equality holds if and only if G is K_n or S_n .

Feng et al. in [5] proved the following theorem.

Theorem B Let G be a connected graph with n vertices and e edges. Then

$$t \leq \left(\frac{d_1 + 1}{n} \right) \left(\frac{2e - d_1 - 1}{n - 2} \right)^{n-2}.$$

Equality holds if and only if G is K_n or S_n .

Li et al. in [13] proved the following theorem.

Theorem C Let G be a non-complete connected graph with $n \geq 4$ vertices and e edges. Then

$$t \leq d_n \left(\frac{2e - d_1 - 1 - d_n}{n - 3} \right)^{n-3}.$$

Equality holds if and only if G is S_n , $K_n - e$, or $K_1 \vee (K_1 \cup K_{n-2})$.

Notice that Theorem C is a corrected version of Theorem 3.1 in [13]. Using some ideas in [4], [5], and [13], we in this note will present a new upper bound for the number of spanning trees of a connected graph of order $n \geq 5$.

Theorem 1 Let G be a non-complete connected graph with $n \geq 5$ vertices and e edges. If $G \neq S_n$, then

$$t \leq (n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4},$$

with equality if and only if G is $K_n - e$, $K_2 \vee K_{n-2}^c$, $K_2 \vee S_{n-2}^c$, $K_1 \vee S_{n-1}^c$, $S_n + e$, or $K_1 \vee (K_1 \cup S_{n-2})$.

2. Proofs of Theorem 1

In order to prove Theorem 1, we need several established results. Theorem 2 below is Lemma 13.1.3 on Page 280 in [7].

Theorem 2 Let G be a graph of order n . Then $\lambda_i(G) = n - \lambda_{n-i}(G^c)$ for each i with $1 \leq i \leq n - 1$. In particular, $\lambda_1(G) \leq n$.

The following Theorem 3 was proven by Fiedler in [6].

Theorem 3 Let G be a non-complete graph of order n . Then $\lambda_{n-1}(G) \leq \kappa(G) \leq \kappa'(G) \leq d_n$, where $\kappa(G)$ and $\kappa'(G)$ are the connectivity and edge-connectivity of G , respectively.

Kirkland in [10] and Li and Fan in [11] characterized the graphs with $\lambda_{n-1}(G) = \kappa(G)$. Their results are stated in Theorem 4 below.

Theorem 4 Let G be a graph of order n with $1 \leq \kappa(G) \leq n - 2$. Then $\lambda_{n-1}(G) = \kappa(G) = k$ if and only if there exists a vertex subset $S \subset V(G)$ with $|S| = k$, such that $G = G[S] \vee (G_1 \cup G_2 \cup \dots \cup G_m)$, $m \geq 2$, and $\kappa(G[S]) \geq 2k - n$ if $\lfloor \frac{n}{2} \rfloor < k \leq n - 2$, where G_1, G_2, \dots, G_m are the components of $G[V(G) - S]$.

Theorem 5 below was proven by Grone and Merris in [8].

Theorem 5 Let G be a graph containing at least one edge. Then $\lambda_1(G) \geq d_1(G) + 1$. Moreover, if G is connected on $n > 1$ vertices, the equality holds if and only if $d_1(G) = n - 1$.

The statement in Theorem 6 was conjectured by Grone and Merris in [8]. Bai in [1] proved that the conjecture is true.

Theorem 6 Let G be a graph of order n . Then $\sum_{i=1}^k \lambda_i(G) \leq \sum_{i=1}^k d_i^*(G)$ for each k with $1 \leq k \leq n$.

Both Theorem 7 and Theorem 8 below were proven by Das in [4].

Theorem 7 Let G be a simple connected graph of order n . Then $\lambda_2(G) = \lambda_3(G) = \dots = \lambda_{n-1}(G)$ if and only if G is K_n or S_n or K_{d_1, d_1} .

Theorem 8 Let G be a simple connected graph of order n . Then $\lambda_1(G) = \lambda_2(G) = \dots = \lambda_{n-2}(G)$ if and only if G is K_n or $K_n - e$.

Theorem 9 below follows from results obtained by Das in [3].

Theorem 9 Let G be a simple connected graph of order $n > 2$. If G is not S_n , then $\lambda_2 \geq a$.

Now we will prove Theorem 1.

Proof of Theorem 1. Let G be a non-complete graph connected graph with $n \geq 5$ vertices and e edges. Suppose that $G \neq S_n$. We assume that x and y are two distinct vertices in G such that $a_{x,y} = a$ with $d(x) = d_1$ and $d(y) = d_2$. Notice first that $0 \leq c_{x,y} \leq d_2 - 1$ if $xy \in E$, and $0 \leq c_{x,y} \leq d_2$ if $xy \notin E$. By the definition of a , we can easily verify that $d_2 \leq a \leq d_2 + 1$.

From Theorem 6, we have that $\lambda_1 + \lambda_2 \leq d_1^* + d_2^* = n + d_2^*$. Theorem 5 then implies that $\lambda_2 \leq n + d_2^* - d_1 - 1$.

From Theorem 2, Theorem 9, and the AM-GM inequality, we have that

$$\begin{aligned} t &= \frac{1}{n} \prod_{i=1}^{n-1} \lambda_i = \frac{\lambda_1 \lambda_2 \lambda_{n-1}}{n} \prod_{i=3}^{n-2} \lambda_i \\ &\leq (n + d_2^* - d_1 - 1) \lambda_{n-1} \left(\frac{\sum_{i=3}^{n-2} \lambda_i}{n-4} \right)^{n-4} \\ &\leq (n + d_2^* - d_1 - 1) \lambda_{n-1} \left(\frac{2e - \lambda_1 - \lambda_2 - \lambda_{n-1}}{n-4} \right)^{n-4} \\ &\leq (n + d_2^* - d_1 - 1) \lambda_{n-1} \left(\frac{2e - d_1 - 1 - a - \lambda_{n-1}}{n-4} \right)^{n-4}. \end{aligned}$$

Now consider the function

$$f(z) = z \left(\frac{2e - d_1 - 1 - a - z}{n-4} \right)^{n-4},$$

where $0 < z \leq d_n$. It is easy to see that

$$f'(z) = \frac{(2e - d_1 - 1 - a - z)^{n-5} [2e - d_1 - 1 - a - (n-3)z]}{(n-4)^{n-4}}.$$

Notice first that $2e - d_1 - 1 - a - z \geq d_1 + d_2 + d_3 + d_4 + \dots + d_{n-1} + d_n - d_1 - 1 - d_2 - 1 - d_n \geq 0$.

If $d_3 \geq 2$, then $2e - d_1 - 1 - a - (n-3)z \geq d_1 + d_2 + d_3 + d_4 + \dots + d_{n-1} + d_n - d_1 - 1 - d_2 - 1 - (n-3)d_n \geq d_3 - 2 \geq 0$.

If $d_3 = 1$, then $d_3 = d_4 = \dots = d_n = 1$. Obviously, $d_2 \neq 1$. Otherwise by the assumption that G is connected we have that $v_i v_j \notin E$, where $2 \leq i \neq j \leq n$, which implies that G is S_n , a contradiction. Now we assume that $d_2 \geq 2$. Again since G is connected, we have that $v_i v_j \notin E$, where $3 \leq i \neq j \leq n$. Hence, for each i with $3 \leq i \leq n$, v_i is adjacent to exactly one of v_1 and v_2 . Since $d(v_1) = d_1 \geq d(v_2) = d_2 = 2 > d_3 = d_4 = \dots = d_n = 1$, we have that $\{x, y\} = \{v_1, v_2\}$. Since G is connected, $xy \in E$. Clearly, $c_{x,y} = 0$. Hence, by the definition of a , we have that $a = d_2$. So

$$\begin{aligned} & 2e - d_1 - 1 - a - (n-3)z \\ & \geq d_1 + d_2 + d_3 + d_4 + \dots + d_{n-1} + d_n - d_1 - 1 - d_2 - (n-3)d_n \\ & \geq d_3 - 1 \geq 0. \end{aligned}$$

Thus $f(z)$ is increasing when $0 < z \leq d_n$. So

$$\begin{aligned} t & \leq (n + d_2^* - d_1 - 1)f(\lambda_{n-1}) \leq (n + d_2^* - d_1 - 1)f(d_n) \\ & = (n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n-4} \right)^{n-4}. \end{aligned}$$

Suppose that

$$t = (n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n-4} \right)^{n-4}.$$

In review of the proof above, we have that

$$\lambda_1 = n = d_1 + 1, \lambda_2 = n + d_2^* - d_1 - 1 = d_2^* = a, \lambda_3 = \dots = \lambda_{n-2}, \lambda_{n-1} = d_n.$$

Since $\lambda_1 \geq \lambda_2$, $\lambda_2 \geq \lambda_3$, and $\lambda_{n-2} \geq \lambda_{n-1}$, we have the following eight cases.

Case 1. $\lambda_1 = \lambda_2 = \lambda_3$, and $\lambda_{n-2} = \lambda_{n-1}$.

This implies $d_1 + 1 = \lambda_1 = \lambda_2 = \dots = \lambda_{n-1} = d_n$. This is contrary to $d_1 \geq d_n$. Hence this case is impossible.

Case 2. $\lambda_1 = \lambda_2 = \lambda_3$, and $\lambda_{n-2} > \lambda_{n-1}$.

By Theorem 8 and the assumption that $G \neq K_n$, we have that G is $K_n - e$ with $n \geq 5$.

Case 3. $\lambda_1 = \lambda_2 > \lambda_3$, and $\lambda_{n-2} = \lambda_{n-1}$.

From $d_2 + 1 \geq a = \lambda_2 = \lambda_1 = d_1 + 1$, we have that $(a - 1) = d_1 = d_2 = (n - 1)$. Thus $d(x) = d(y) = (n - 1)$. Hence $N(x) = V - \{x\}$ and $N(y) = V - \{y\}$. In particular, $xy \in E$.

Since $G \neq K_n$ and $\lambda_{n-1} = d_n$, Theorem 4 implies that G is the same as the graph $G[S] \vee (G_1 \cup G_2 \cup \dots \cup G_m)$, where $|S| = d_n$ and $m \geq 2$. In fact, from the proof of Theorem 1 in [11], we can see that S is a cut-set of G and G_1, G_2, \dots, G_m are the components of $G[V(G) - S]$. Since $d(x) = d(y) = (n - 1)$, both x and y are in S .

We claim that $|S| = 2$. Suppose, to the contrary, that $|S| \geq 3$. Since $N(x) = V - \{x\}$ and $N(y) = V - \{y\}$, G^c has at least four components. Notice that the Laplacian eigenvalues of G are

$$(n, n, \lambda_3, \dots, \lambda_{n-1}, 0).$$

Thus the Laplacian eigenvalues of G^c are

$$((n - \lambda_{n-1}), (n - \lambda_{n-2}), \dots, (n - \lambda_3), 0, 0, 0)$$

and G^c has three components, a contradiction. Hence $|S| = 2$.

Now we have that $\lambda_3 = \dots = \lambda_{n-1} = d_n = |S| = 2$. Therefore $G^c - \{x, y\}$ has Laplacian eigenvalues

$$((n - 2), (n - 2), \dots, (n - 2), 0).$$

By Theorem 8, we have that $G^c - \{x, y\}$ is K_{n-2} or $K_{n-2} - e$. Since $G^c - \{x, y\}$ and $K_{n-2} - e$ have different sets of Laplacian eigenvalues, $G^c - \{x, y\} \neq K_{n-2} - e$. Thus $G^c - \{x, y\} = K_{n-2}$. Hence G is $K_2 \vee K_{n-2}^c$ with $n \geq 5$.

Case 4. $\lambda_1 = \lambda_2 > \lambda_3$, and $\lambda_{n-2} > \lambda_{n-1}$.

Using arguments similar to those in Case 3, we have that $(a - 1) = d_1 = d(x) = d_2 = d(y) = (n - 1)$, $N(x) = V - \{x\}$, $N(y) = V - \{y\}$, and $xy \in E$. Moreover, $S = \{x, y\}$ is a cutset of G and G is the same as the graph $G[S] \vee (G_1 \cup G_2 \cup \dots \cup G_m)$, where $|S| = d_n = \lambda_{n-1} = 2$ and $m \geq 2$.

Therefore $G^c - \{x, y\}$ has Laplacian eigenvalues

$$((n - \lambda_{n-1}), (n - \lambda_{n-2}), \dots, (n - \lambda_3), 0).$$

By Theorem 7, we have that $G^c - \{x, y\}$ is K_{n-2} , S_{n-2} , or $K_{\frac{n-2}{2}, \frac{n-2}{2}}$. Since $(n - \lambda_{n-1}) > (n - \lambda_{n-2})$, $G^c - \{x, y\}$ cannot be K_{n-2} . If $G^c - \{x, y\}$ is $K_{\frac{n-2}{2}, \frac{n-2}{2}}$, the minimum degree of G is $\frac{n}{2}$, which is not equal to $d_n = 2$ when $n \geq 5$. Thus $G^c - \{x, y\}$ cannot be $K_{\frac{n-2}{2}, \frac{n-2}{2}}$. Therefore $G^c - \{x, y\} = S_{n-2}$. Hence G is $K_2 \vee S_{n-2}^c$ with $n \geq 5$.

Case 5. $\lambda_1 > \lambda_2 = \lambda_3$, and $\lambda_{n-2} = \lambda_{n-1}$.

By Theorem 7 and the assumptions that $G \neq K_n$ and $G \neq S_n$, we have that G is K_{d_1, d_1} . Since $d_1 = (n - 1)$, G must be K_2 , contradicting the assumption that $n \geq 5$. Hence this case is impossible.

Case 6. $\lambda_1 > \lambda_2 = \lambda_3$, and $\lambda_{n-2} > \lambda_{n-1}$.

From $d_1 = (n - 1)$, we have that $d(x) = (n - 1)$ and $N(x) = V - \{x\}$. Since G is not K_n and $\lambda_{n-1} = d_n$, Theorem 4 implies that G is the same as the graph $G[S] \vee (G_1 \cup G_2 \cup \dots \cup G_m)$, where $|S| = d_n$ and $m \geq 2$. Clearly, x is in S .

We claim that $|S| = 1$. Suppose, to the contrary, that $|S| \geq 2$. Since $N(x) = V - \{x\}$, G^c has at least three components. Notice that the Laplacian eigenvalues of G are

$$(n, \lambda_2, \lambda_3, \dots, \lambda_{n-1}, 0).$$

Thus the Laplacian eigenvalues of G^c are

$$((n - \lambda_{n-1}), (n - \lambda_{n-2}), \dots, (n - \lambda_2), 0, 0)$$

and therefore G^c has two components, a contradiction. Hence $|S| = 1$.

Now we have that $\lambda_{n-1} = d_n = |S| = 1$. Therefore $G^c - \{x\}$ has Laplacian eigenvalues

$$((n - 1), (n - \lambda_{n-2}), \dots, (n - \lambda_2), 0).$$

By Theorem 7, we have that $G^c - \{x\}$ is K_{n-1} , S_{n-1} , or $K_{\frac{n-1}{2}, \frac{n-1}{2}}$. Since $(n-1) = (n-\lambda_{n-1}) > (n-\lambda_{n-2})$, $G^c - \{x\}$ cannot be K_{n-1} . If $G^c - \{x\}$ is $K_{\frac{n-1}{2}, \frac{n-1}{2}}$, the minimum degree of G is $\frac{n-1}{2}$, which is not equal to $d_n = 1$ when $n \geq 5$. Thus $G^c - \{x\}$ cannot be $K_{\frac{n-1}{2}, \frac{n-1}{2}}$. Therefore $G^c - \{x\} = S_{n-1}$. Hence G is $K_1 \vee S_{n-1}^c$ with $n \geq 5$.

Case 7. $\lambda_1 > \lambda_2 > \lambda_3$, and $\lambda_{n-2} = \lambda_{n-1}$.

Again, since $d_1 = (n-1)$, we have that $d(x) = (n-1)$ and $N(x) = V - \{x\}$. Since G is not K_n and $\lambda_{n-1} = d_n$, Theorem 4 implies that G is the same as the graph $G[S] \vee (G_1 \cup G_2 \cup \dots \cup G_m)$, where $|S| = d_n$ and $m \geq 2$. Clearly, x is in S .

We claim again that $|S| = 1$. Suppose, to the contrary, that $|S| \geq 2$. Since $N(x) = V - \{x\}$, G^c has at least three components. Since the Laplacian eigenvalues of G are

$$(n, \lambda_2, \lambda_3, \dots, \lambda_{n-1}, 0),$$

the Laplacian eigenvalues of G^c are

$$((n-d_n), (n-d_n), \dots, (n-d_n), (n-\lambda_2), 0, 0).$$

Since $n-\lambda_2 = \lambda_1 - \lambda_2 > 0$, G^c has two components, a contradiction. Hence $|S| = 1$.

Now we have that $\lambda_{n-1} = d_n = |S| = 1$. Therefore G^c has Laplacian eigenvalues

$$((n-1), (n-1), \dots, (n-1), (n-\lambda_2), 0, 0),$$

and $G^c - \{x\}$ has Laplacian eigenvalues

$$((n-1), (n-1), \dots, (n-1), (n-\lambda_2), 0).$$

By Theorem 8, we have that $G^c - \{x\}$ is K_{n-1} or $K_{n-1} - e$. Since $(n-1) = (n-d_n) > (n-\lambda_2)$, $G^c - \{x\}$ cannot be K_{n-1} . Therefore $G^c - \{x\} = K_{n-1} - e$. Hence G is $S_n + e$ with $n \geq 5$.

Case 8. $\lambda_1 > \lambda_2 > \lambda_3$, and $\lambda_{n-2} > \lambda_{n-1}$.

Using arguments similar to those in Case 7, we have that $d(x) = (n-1)$, $N(x) = V - \{x\}$, G is the same as the graph $G[S] \vee (G_1 \cup G_2 \cup \dots \cup G_m)$ with $|S| = d_n = \lambda_{n-1} = 1$ and $m \geq 2$, and x is in S .

Since $d_n = 1$, v_n is only adjacent to x in G and v_n is adjacent to each vertex of $V - \{x\}$ in G^c . Since the Laplacian eigenvalues of G are

$$(n, \lambda_2, \dots, \lambda_{n-2}, \lambda_{n-1}, 0),$$

the Laplacian eigenvalues of G^c are

$$((n-1), (n-\lambda_{n-2}), \dots, (n-\lambda_3), (n-\lambda_2), 0, 0).$$

This implies that the Laplacian eigenvalues of $G^c - \{x\}$ are

$$((n-1), (n-\lambda_{n-2}), \dots, (n-\lambda_3), (n-\lambda_2), 0).$$

Therefore the Laplacian eigenvalues of $(G^c - \{x\})^c$ are

$$((n-1) - (n-\lambda_2), (n-1) - (n-\lambda_3), \dots, (n-1) - (n-\lambda_{n-2}), 0, 0),$$

which further implies that the Laplacian eigenvalues of $(G^c - \{x\})^c - \{v_n\}$, i.e., $G - \{x, v_n\}$, are

$$((\lambda_2 - 1), (\lambda_3 - 1), \dots, (\lambda_{n-2} - 1), 0).$$

Since $(\lambda_3 - 1) = \dots = (\lambda_{n-2} - 1)$, we, by Theorem 7, have that $(G^c - \{x\})^c - \{v_n\} = G - \{x, v_n\}$ is K_{n-2} , S_{n-2} , or $K_{\frac{n-2}{2}, \frac{n-2}{2}}$. Since $\lambda_2 - 1 > \lambda_3 - 1$, $G - \{x, v_n\}$ cannot be K_{n-2} . If $G - \{x, v_n\}$ is $K_{\frac{n-2}{2}, \frac{n-2}{2}}$, then $\lambda_2 = d_2^* = (n-1) \neq \frac{n}{2} + 1 = a$ when $n \geq 5$. Thus $G - \{x, v_n\}$ cannot be $K_{\frac{n-2}{2}, \frac{n-2}{2}}$. Therefore $G - \{x, v_n\} = S_{n-2}$. Hence G is $K_1 \vee (K_1 \cup S_{n-2})$ with $n \geq 5$.

Next we show that if G is $K_n - e$, $K_2 \vee K_{n-2}^c$, $K_2 \vee S_{n-2}^c$, $K_1 \vee S_{n-1}^c$, $S_n + e$, or $K_1 \vee (K_1 \cup S_{n-2})$, then

$$t = (n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4}.$$

If G is $K_n - e$ with $n \geq 5$, then the Laplacian eigenvalues of G^c are

$$(2, 0, \dots, 0, 0, 0).$$

Thus the Laplacian eigenvalues of G are

$$(n, n, \dots, n, (n-2), 0).$$

Hence the number of spanning trees of G is $t = \prod_{i=1}^{n-1} \lambda_i / n = (n-2)n^{n-3}$. Now

$$(n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4} = (n-2)n^{n-3},$$

which is the number of spanning trees of $K_n - e$ with $n \geq 5$.

If G is $K_2 \vee K_{n-2}^c$ with $n \geq 5$, then the Laplacian eigenvalues of G^c are

$$((n-2), (n-2), \dots, (n-2), 0, 0, 0).$$

Thus the Laplacian eigenvalues of G are

$$(n, n, 2, 2, \dots, 2, 0).$$

Hence the number of spanning trees of G is $t = \prod_{i=1}^{n-1} \lambda_i/n = n2^{n-3}$. Now

$$(n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4} = n2^{n-3},$$

which is the number of spanning trees of $K_2 \vee K_{n-2}^c$ with $n \geq 5$.

If G is $K_2 \vee S_{n-2}^c$ with $n \geq 5$, then the Laplacian eigenvalues of G^c are

$$((n-2), 1, \dots, 1, 0, 0, 0).$$

Thus the Laplacian eigenvalues of G are

$$(n, n, (n-1), \dots, (n-1), 2, 0).$$

Hence the number of spanning trees of G is $t = \prod_{i=1}^{n-1} \lambda_i/n = 2n(n-1)^{n-4}$. Now

$$(n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4} = 2n(n-1)^{n-4},$$

which is the number of spanning trees of $K_2 \vee S_{n-2}^c$ with $n \geq 5$.

If G is $K_1 \vee S_{n-1}^c$ with $n \geq 5$, then the Laplacian eigenvalues of G^c are

$$((n-1), 1, \dots, 1, 0, 0).$$

Thus the Laplacian eigenvalues of G are

$$(n, (n-1), \dots, (n-1), 1, 0).$$

Hence the number of spanning trees of G is $t = \prod_{i=1}^{n-1} \lambda_i/n = (n-1)^{n-3}$. Now

$$(n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4} = (n-1)^{n-3},$$

which is the number of spanning trees of $K_1 \vee S_{n-1}^c$ with $n \geq 5$.

If G is $S_n + e$ with $n \geq 5$, then

$$(n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4} = 3,$$

which is the number of spanning trees of $S_n + e$ with $n \geq 5$.

If G is $K_1 \vee (K_1 \cup S_{n-2})$ with $n \geq 5$, from the above proof of Case 8, we have that the Laplacian eigenvalues of $G - \{x, v_n\}$ are

$$((\lambda_2 - 1), (\lambda_3 - 1), \dots, (\lambda_{n-2} - 1), 0).$$

Since $G - \{x, v_n\}$ is S_{n-2} , $\lambda_2 - 1 = n - 2$ and $\lambda_3 - 1 = \dots = \lambda_{n-2} - 1 = 1$. Thus $\lambda_1 = n$, $\lambda_2 = n - 1$, $\lambda_3 = \dots = \lambda_{n-2} = 2$, and $\lambda_{n-1} = 1$. Hence the number of spanning trees of G is $t = \prod_{i=1}^{n-1} \lambda_i / n = (n - 1)2^{n-4}$. Now

$$(n + d_2^* - d_1 - 1)d_n \left(\frac{2e - d_1 - 1 - a - d_n}{n - 4} \right)^{n-4} = (n - 1)2^{n-4},$$

which is the number of spanning trees of $K_1 \vee (K_1 \cup S_{n-2})$ with $n \geq 5$.

This completes the proof of Theorem 1.

Obviously, if G is $K_2 \vee K_{n-2}^c$, $K_2 \vee S_{n-2}^c$, $S_n + e$, or $K_1 \vee (K_1 \cup S_{n-2})$ with $n \geq 5$, then the upper bound in Theorem 1 is less than the upper bounds in Theorems A, B, or C.

Recall that Li and Pan in [12] proved that $\lambda_2(G) \geq d_2(G)$ for a connected graph of order $n \geq 3$ and Guo in [9] proved that $\lambda_3(G) \geq d_3(G) - 1$ for a connected graph of order $n \geq 4$. Using those results, we can present another upper bound for the number of spanning trees in a connected graph of order $n \geq 6$.

Theorem 10 Let G be a non-complete connected graph with $n \geq 6$ vertices and e edges. Then

$$t \leq (n + d_2^* - d_1 - 1)(n + d_2^* + d_3^* - d_1 - d_2 - 1)d_n \left(\frac{2e - d_1 - d_2 - d_3 - d_n}{n - 5} \right)^{n-5}.$$

The proof of Theorem 10 is similar to the proof of Theorem 1. It can be fulfilled as long as we realize that $\lambda_3 = \lambda_1 + \lambda_2 + \lambda_3 - \lambda_1 - \lambda_2 \leq d_1^* + d_2^* + d_3^* - d_1 - 1 - d_2 = n + d_2^* + d_3^* - d_1 - d_2 - 1$. The complete proof

Theorem 10 is skipped here. It is interesting to decide if the upper bound in Theorem 10 is attainable for some graphs.

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