

# When all minimal $k$ -vertex separators induce complete or edgeless subgraphs

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## Abstract

Define  $\mathcal{D}_k$  to be the class of graphs such that, for every independent set  $\{v_1, \dots, v_h\}$  of vertices with  $2 \leq h \leq k$ , if  $S$  is an inclusion-minimal set of vertices whose deletion would put  $v_1, \dots, v_h$  into  $h$  distinct connected components, then  $S$  induces a complete subgraph; also, let  $\mathcal{D} = \bigcap_{k \geq 2} \mathcal{D}_k$ . Similarly, define  $\mathcal{D}'_k$  and  $\mathcal{D}'$  with “complete” replaced by “edgeless,” and define  $\mathcal{D}^*_k$  and  $\mathcal{D}^*$  with “complete” replaced by “complete or edgeless.” The class  $\mathcal{D}_2$  is the class of chordal graphs, and the classes  $\mathcal{D}$ ,  $\mathcal{D}'_2$ , and  $\mathcal{D}^*_2$  have also been characterized recently. The present paper gives unified characterizations of all of the classes  $\mathcal{D}_k$ ,  $\mathcal{D}'_k$ ,  $\mathcal{D}^*_k$ ,  $\mathcal{D}$ ,  $\mathcal{D}'$ , and  $\mathcal{D}^*$ .

## 1 Introduction: minimal $k$ -vertex separators

When  $G$  is a graph with an independent set  $\mathcal{I}$  of two or more vertices, call  $S \subset V(G)$  an  $\mathcal{I}$ -separator of  $G$  if the vertices of  $\mathcal{I}$  are in  $|\mathcal{I}|$  separate connected components of the subgraph of  $G$  that is induced by  $V(G) - S$ ; a *minimal  $\mathcal{I}$ -separator* is an inclusion-minimal  $\mathcal{I}$ -separator. When  $|\mathcal{I}| = k$ , these will also be called (*minimal*)  $k$ -vertex separators. Minimal 2-vertex separators have long been studied as “minimal separators” (sometimes called “minimal vertex separators” or “minseps”); see [1, 3, 9].

The present paper will characterize the graphs for which every minimal  $k$ -vertex separator induces a complete subgraph, an edgeless subgraph, or a complete or edgeless subgraph. Each of these characterizations will imply the existence of a forbidden induced subgraph characterization of the graph class, and so the existence of a polynomial-time recognition algorithm.

Let  $v \sim w$  [and  $v \not\sim w$ ] denote that vertices  $v$  and  $w$  are adjacent [respectively, nonadjacent]. The *components* of a graph are simply its connected components. For any  $S \subset V(G)$ , let  $G - S$  denote the subgraph of

$G$  that is induced by  $V(G) - S$ . If  $S$  is a minimal  $\mathcal{I}$ -separator of  $G$ , then the minimality of  $S$  implies that each  $x \in S$  has neighbors  $x'$  and  $x''$  in, respectively, components  $G'$  and  $G''$  of  $G - S$  such that  $G'$  and  $G''$  contain different elements of  $\mathcal{I}$ . Observe that vertices in different components of  $G - S$  cannot be adjacent in  $G$ .

## 2 Complete minimal $k$ -vertex separators

Let  $\mathcal{D}_k$  denote the class of graphs such that every minimal  $\mathcal{I}$ -separator induces a complete subgraph whenever  $2 \leq |\mathcal{I}| \leq k$ , and call such graphs *complete minimal  $k$ -vertex separator graphs*. Let  $\mathcal{D} = \bigcap_{k \geq 2} \mathcal{D}_k$ , and note that  $\mathcal{D} \subseteq \dots \subseteq \mathcal{D}_4 \subseteq \mathcal{D}_3 \subseteq \mathcal{D}_2$ .

Proposition 1, from Dirac's seminal paper [2], shows that  $\mathcal{D}_2$  is the familiar class of *chordal graphs*; see [1, 9].

**Proposition 1** *A graph is in  $\mathcal{D}_2$  if and only if no induced subgraph is a cycle of length 4 or more.*

Figure 1 shows a graph that is in  $\mathcal{D}_3$  (and  $\mathcal{D}_2$ ), but is not in  $\mathcal{D}_4$ . Each minimal 3-vertex separator is a complete set (of cardinality 2), but  $\{s_1, s_2, s_3, s_4\}$  is a minimal  $\{r_1, r_2, r_3, r_4\}$ -separator that does not induce an complete subgraph.

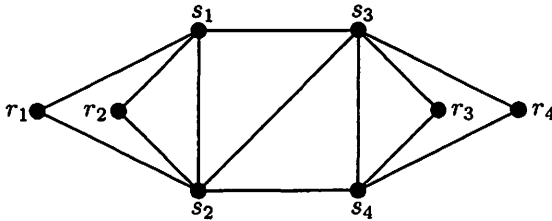


Figure 1: A graph in  $\mathcal{D}_3$  but not  $\mathcal{D}_4$ .

Proposition 2, from [7], characterizes  $\mathcal{D}$ , where  $C_n$  [and  $P_n$ ] denote the cycle [respectively, path] on  $n$  vertices (so  $C_n$  has length  $n$  and  $P_n$  has length  $n - 1$ ), and  $2P_3$  denotes the graph that consists of two components, each a copy of  $P_3$ .

**Proposition 2** *A graph is in  $\mathcal{D}$  if and only if no induced subgraph is isomorphic to  $C_4$ ,  $C_5$ ,  $P_5$ , or  $2P_3$ .*

The graph  $G$  in Figure 1 is not in  $\mathcal{D}$ , since deleting the vertices  $s_2$  and  $s_3$  would leave an induced subgraph isomorphic to  $2P_3$ . Inserting an additional edge  $s_1s_4$  into  $G$  would produce a graph that is in  $\mathcal{D}$ .

Define a *chord of a set*  $\{\pi_1, \dots, \pi_s\}$  of paths (or a *chord of a path*  $\pi_1$  if  $s = 1$ ) to be an edge  $xy$  where  $x, y \in V(\pi_1) \cup \dots \cup V(\pi_s)$  but  $xy \notin E(\pi_1) \cup \dots \cup E(\pi_s)$ . Define a  $\pi_i$ -to- $\pi_j$  *chord* to be a chord  $x_i x_j$  of  $\{\pi_i, \pi_j\}$  where  $x_i \in V(\pi_i) - V(\pi_j)$  and  $x_j \in V(\pi_j) - V(\pi_i)$ . Let  $\pi_i^\circ$  denote the set of vertices of the path  $\pi_i$  that are not endpoints of  $\pi_i$ , and call paths  $\pi_i$  and  $\pi_j$  *internally disjoint* when  $\pi_i^\circ \cap \pi_j^\circ = \emptyset$ . Motivated by the definition of "chordless cycle" in [11], define a *chordless path* to be a path of length at least 2 that has no chord—thus, chordless paths are induced paths that are long enough for a chord to have been possible.

Theorem 3 will generalize (and prove again) both Propositions 1 and 2.

**Theorem 3** *A graph is in  $\mathcal{D}_k$  if and only if no two internally disjoint chordless paths  $\pi_1$  and  $\pi_2$  have a total of at most  $\min\{k, 4\}$  distinct endpoints with no  $\pi_1$ -to- $\pi_2$  chord. Consequently,  $\mathcal{D}_k = \mathcal{D}_4$  whenever  $k \geq 4$  and  $\mathcal{D} = \mathcal{D}_4$ .*

**Proof.** First suppose two internally disjoint chordless paths  $\pi_1$  and  $\pi_2$  of  $G$  have a total of at most  $\min\{k, 4\}$  distinct endpoints with no  $\pi_1$ -to- $\pi_2$  chord (arguing by contraposition). Let  $\mathcal{I}$  be the set of endpoints of  $\pi_1$  and  $\pi_2$  (so  $\mathcal{I}$  is independent with  $2 \leq |\mathcal{I}| \leq 4$ ), and let  $S = \{x_1, x_2\}$  where  $x_1 \not\sim x_2$  and each  $x_i \in \pi_i^\circ$ . But then  $S$  is a minimal  $\mathcal{I}$ -separator of the subgraph of  $G$  induced by  $V(\pi_1) \cup V(\pi_2)$ , and so  $S$  is contained in a minimal  $\mathcal{I}$ -separator of  $G$  where the subgraph induced by  $S$  is not complete (since  $x_1 \not\sim x_2$ ) and  $2 \leq |\mathcal{I}| \leq \min\{k, 4\} \leq k$ . Therefore,  $G \notin \mathcal{D}_k$ .

Conversely, suppose  $G \notin \mathcal{D}_k$ , say with  $S$  a minimal  $\mathcal{I}$ -separator of  $G$  that does not induce a complete subgraph and has  $2 \leq |\mathcal{I}| \leq k$ . To be specific, suppose  $x_1, x_2 \in S$  where  $x_1 \not\sim x_2$ . By the minimality of  $S$ , each  $x_i$  has neighbors  $x'_i$  and  $x''_i$  in, respectively, components  $G'_i$  and  $G''_i$  of  $G - S$ .

*Case 1'*:  $G'_1 \neq G'_2$ . Define paths  $\tau'_1$  and  $\tau'_2$  to consist of, respectively, the single edge  $x'_1 x_1$  and the single edge  $x'_2 x_2$ .

*Case 2'*:  $G'_1 = G'_2$ . Let  $\sigma'$  denote a chordless  $x_1$ -to- $x_2$  path in  $G$  with  $\sigma' \subset V(G'_1)$ , and choose  $v' \in \sigma'^\circ$ . Define paths  $\tau'_1$  and  $\tau'_2$  to consist of, respectively, the  $v'$ -to- $x_1$  and  $v'$ -to- $x_2$  subpaths of  $\sigma'$ .

Now consider similar *cases 1'' and 2''* for the components  $G''_1$  and  $G''_2$  of  $G - S$  by replacing each superscript  $'$  with  $''$  in the corresponding case above so as to define paths  $\tau''_1$  and  $\tau''_2$ . Finally, for each  $i \in \{1, 2\}$ , define  $\pi_i$  to be the chordless path  $\tau'_i \cup \tau''_i$  in  $G$  (noting that each  $x_i \in \pi_i^\circ$ ). In every combination of cases 1', 2' and 1'', 2'', the two paths  $\pi_1, \pi_2$  will be internally disjoint and chordless with a total of at most  $\min\{k, 4\}$  distinct endpoints, with no  $\pi_1$ -to- $\pi_2$  chord.

Consequently, if  $k \geq 4$ , then  $\min\{k, 4\} = 4$  and the preceding two paragraphs show  $\mathcal{D}_k = \mathcal{D}_4$ , and so  $\mathcal{D} = \mathcal{D}_4$ .  $\square$

To illustrate Theorem 3 in Figure 1 with  $k = 4$ ,  $\mathcal{I} = \{r_1, r_2, r_3, r_4\}$ , and  $S = \{s_1, s_2, s_3, s_4\}$ , observe that the paths  $\pi_1 = r_1, s_1, r_2$  and  $\pi_2 = r_3, s_4, r_4$  have four distinct endpoints with no  $\pi_1$ -to- $\pi_2$  chord.

Theorem 3 has Propositions 1 and 2 as consequences; in the latter's characterization of  $\mathcal{D} = \mathcal{D}_4$ , the subgraphs  $C_4$  and  $C_5$  arise when  $\pi_1$  and  $\pi_2$  share the same two endpoints, and the subgraphs  $P_5$  and  $2P_3$  arise when  $\pi_1$  and  $\pi_2$  have, respectively, a total of three or four endpoints. Similarly,  $\mathcal{D}_3$  can be characterized by no induced subgraph being isomorphic to  $C_4$ ,  $C_5$ , or  $P_5$ .

### 3 Edgeless minimal $k$ -vertex separators

Let  $\mathcal{D}'_k$  denote the class of graphs such that every minimal  $\mathcal{I}$ -separator induces an edgeless subgraph whenever  $2 \leq |\mathcal{I}| \leq k$ , and call such graphs *edgeless minimal  $k$ -vertex separator graphs*. Let  $\mathcal{D}' = \bigcap_{k \geq 2} \mathcal{D}'_k$ , and note that  $\mathcal{D}' \subseteq \dots \subseteq \mathcal{D}'_4 \subseteq \mathcal{D}'_3 \subseteq \mathcal{D}'_2$ .

Proposition 4, from [6], shows that  $\mathcal{D}'_2$  is also the graph class characterized in [10], whose members are called *unchord-free graphs* in [4, 5].

**Proposition 4** *A graph is in  $\mathcal{D}'_2$  if and only if no cycle has a unique chord.*

Figure 2 shows a graph that is in  $\mathcal{D}'_3$  (and  $\mathcal{D}'_2$ ), but is not in  $\mathcal{D}'_4$ . Each minimal 3-vertex separator is an independent set (of cardinality 3 or 4), but  $\{s_1, s_2, s_3, s_4\}$  is a minimal  $\{r_1, r_2, r_3, r_4\}$ -separator that does not induce an edgeless subgraph.

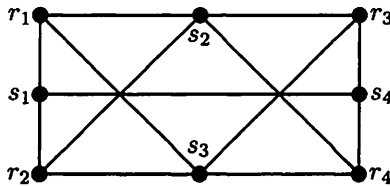


Figure 2: A graph in  $\mathcal{D}'_3$  but not  $\mathcal{D}'_4$ .

Theorem 5 will generalize (and reprove) Proposition 4.

**Theorem 5** *A graph is in  $\mathcal{D}'_k$  if and only if no two internally disjoint chordless paths  $\pi_1$  and  $\pi_2$  have a total of at most  $\min\{k, 4\}$  distinct endpoints with a unique  $\pi_1$ -to- $\pi_2$  chord  $x_1x_2$  and each  $x_i \in \pi_i^?$ . Consequently,  $\mathcal{D}'_k = \mathcal{D}'_4$  whenever  $k \geq 4$  and  $\mathcal{D}' = \mathcal{D}'_4$ .*

**Proof.** The proof given above for Theorem 3 will also prove Theorem 5 after making the following four substitutions: replace each  $\mathcal{D}$  (subscripted

or not) with  $\mathcal{D}'$ ; replace “no  $\pi_1$ -to- $\pi_2$  chord” with “a unique  $\pi_1$ -to- $\pi_2$  chord  $x_1x_2$  and each  $x_i \in \pi_i^o$ ”; replace “ $x_1 \not\sim x_2$ ” with “ $x_1 \sim x_2$ ”; and replace “complete” with “edgeless.”  $\square$

To illustrate Theorem 5 in Figure 2 with  $k = 4$ ,  $\mathcal{I} = \{r_1, r_2, r_3, r_4\}$ , and  $S = \{s_1, s_2, s_3, s_4\}$ , observe that the paths  $\pi_1 = r_1, s_1, r_2$  and  $\pi_2 = r_3, s_4, r_4$  have four distinct endpoints with the unique  $\pi_1$ -to- $\pi_2$  chord  $s_1s_4$  and  $s_1 \in \pi_1^o$  and  $s_4 \in \pi_2^o$ .

Theorem 5 has Proposition 4 as a consequence. Moreover, a graph is in  $\mathcal{D}' = \mathcal{D}'_4$  if and only if it contains none of the graphs shown in Figure 3 as an induced subgraph; in each of these six graphs, the “hollow vertices”

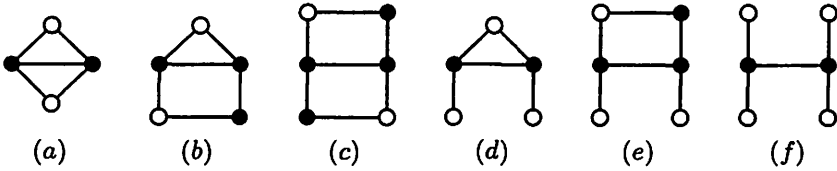


Figure 3: The six forbidden induced subgraphs for  $\mathcal{D}' = \mathcal{D}'_4$ .

show possible endpoints for the two internally disjoint paths  $\pi_1$  and  $\pi_2$ , with the “horizontal edge” halfway up being their unique  $\pi_1$ -to- $\pi_2$  chord. In subgraphs (a), (b), and (c), the two paths share the same two endpoints; in (d) and (e), they share a total of three endpoints; in (f), they have a total of four endpoints. (Similarly,  $\mathcal{D}'_3$  can be characterized by no induced subgraph being a cycle with pendant edges attached at two consecutive vertices or being one of the subgraphs (a), (b), or (c) in Figure 3.)

The graph  $G$  in Figure 2 is not in  $\mathcal{D}'$ , since deleting vertices  $s_2$  and  $s_3$  would leave an induced subgraph isomorphic to subgraph (f) in Figure 3. Deleting the edge  $s_1s_4$  from  $G$  would produce a graph that is in  $\mathcal{D}'$ .

## 4 Extreme minimal $k$ -vertex separators

Let  $\mathcal{D}_k^*$  denote the class of graphs such that every minimal  $\mathcal{I}$ -separator induces a complete or edgeless subgraph whenever  $2 \leq |\mathcal{I}| \leq k$ , and (motivated by [8]) call such graphs *extreme minimal  $k$ -vertex separator graphs*. Let  $\mathcal{D}^* = \bigcap_{k \geq 2} \mathcal{D}_k^*$ , and note that  $\mathcal{D}^* \subseteq \dots \subseteq \mathcal{D}_4^* \subseteq \mathcal{D}_3^* \subseteq \mathcal{D}_2^*$ . Of course,  $\mathcal{D}_k \cup \mathcal{D}'_k \subseteq \mathcal{D}_k^*$  and  $\mathcal{D} \cup \mathcal{D}' \subseteq \mathcal{D}^*$ .

Proposition 6, from [8], characterizes  $\mathcal{D}_2^*$ , where a *block* is an inclusion-maximal 2-connected subgraph and a graph  $G$  is the *edge sum* of graphs  $G_1$  and  $G_2$  if  $V(G) = V(G_1) \cup V(G_2)$  and  $E(G) = E(G_1) \cup E(G_2)$  where  $V(G_1) \cap V(G_2) = \{x, y\}$  and  $E(G_1) \cap E(G_2) = \{xy\}$ .

**Proposition 6** *A graph is in  $\mathcal{D}_2^*$  if and only if every block is the edge sum of graphs in  $\mathcal{D}_2 \cup \mathcal{D}_2'$ .*

Both Figure 1 and 2 show graphs that are in  $\mathcal{D}_3^*$  (and  $\mathcal{D}_2^*$ ), but are not in  $\mathcal{D}_4^*$ . Each minimal 3-vertex separator is an independent set (of cardinality 3 or 4), but  $\{s_1, s_2, s_3, s_4\}$  is a minimal  $\{r_1, r_2, r_3, r_4\}$ -separator that induces neither a complete nor an edgeless subgraph.

Although there does not seem to be a direct generalization of Proposition 6 to  $\mathcal{D}_k^*$  when  $k \geq 3$ , Theorem 7 will characterize each class  $\mathcal{D}_k^*$  (along with  $\mathcal{D}^*$ ) while generalizing the results and proofs of sections 1 and 2.

**Theorem 7** *A graph is in  $\mathcal{D}_k^*$  if and only if no three internally disjoint chordless paths  $\pi_1, \pi_2$ , and  $\pi_3$  have a total of at most  $\min\{k, 6\}$  distinct endpoints with no  $\pi_1$ -to- $\pi_2$  chord and a unique  $\pi_2$ -to- $\pi_3$  chord  $x_2x_3$  and  $x_2 \in \pi_2^\circ$  and  $x_3 \in \pi_3^\circ$ . Consequently,  $\mathcal{D}_k^* = \mathcal{D}_6^*$  whenever  $k \geq 6$  and  $\mathcal{D}^* = \mathcal{D}_6^*$ .*

**Proof.** First suppose three internally disjoint chordless paths  $\pi_1, \pi_2$ , and  $\pi_3$  of  $G$  have a total of at most  $\min\{k, 6\}$  distinct nonadjacent endpoints with no  $\pi_1$ -to- $\pi_2$  chord and a unique  $\pi_2$ -to- $\pi_3$  chord  $x_2x_3$  and  $x_2 \in \pi_2^\circ$  and  $x_3 \in \pi_3^\circ$  (arguing by contraposition). Let  $\mathcal{I}$  be the set of endpoints of  $\pi_1, \pi_2$ , and  $\pi_3$  (so  $\mathcal{I}$  is independent with  $2 \leq |\mathcal{I}| \leq 6$ ), and let  $S = \{x_1, x_2, x_3\}$  where  $x_1 \not\sim x_2 \sim x_3$  and each  $x_i \in \pi_i^\circ$ . But then  $S$  is a minimal  $\mathcal{I}$ -separator of the subgraph of  $G$  induced by  $V(\pi_1) \cup V(\pi_2) \cup V(\pi_3)$ , and so  $S$  is contained in a minimal  $\mathcal{I}$ -separator of  $G$  where the subgraph induced by  $S$  is neither complete (because  $x_1 \not\sim x_2$ ) nor edgeless (since  $x_2 \sim x_3$ ) and  $2 \leq |\mathcal{I}| \leq \min\{k, 6\} \leq k$ . Therefore,  $G \notin \mathcal{D}_k^*$ .

Conversely, suppose  $G \notin \mathcal{D}_k^*$ , say with  $S$  a minimal  $\mathcal{I}$ -separator of  $G$  that does not induce a complete or edgeless subgraph and has  $2 \leq |\mathcal{I}| \leq k$ . To be specific, suppose  $x_1, x_2, x_3 \in S$  where  $x_1 \not\sim x_2 \sim x_3$  ( $x_1$  might or might not be adjacent to  $x_3$ ). By the minimality of  $S$ , each  $x_i$  has neighbors  $x_i'$  and  $x_i''$  in, respectively, components  $G_i'$  and  $G_i''$  of  $G - S$ . Assume in addition that  $S$  and  $x_1, x_2, x_3$  have been chosen to make  $|V(G'')|$  a minimum.

*Case 1'*: no two of  $G_1', G_2', G_3'$  are equal. For each  $i \in \{1, 2, 3\}$ , define a path  $\tau_i'$  to consist of the single edge  $x_i'x_i$ .

*Case 2'*: exactly two of  $G_1', G_2', G_3'$  are equal—say  $G_i' = G_j' \neq G_h'$  where  $\{h, i, j\} = \{1, 2, 3\}$ . Let  $\sigma'$  denote a chordless  $x_i$ -to- $x_j$  path in  $G$  with  $\sigma'^\circ \subset V(G_i')$ , and choose  $v' \in \sigma'^\circ$ . Define paths  $\tau_i'$  and  $\tau_j'$  to consist of, respectively, the  $v'$ -to- $x_i$  and  $v'$ -to- $x_j$  subpaths of  $\sigma'$ . Define  $\tau_h'$  to consist of the single edge  $x_h'x_h$ .

*Case 3'*:  $G_1' = G_2' = G_3'$ . Let  $H'$  be the subgraph of  $G$  induced by  $V(G_1') \cup \{x_1, x_2, x_3\}$  except without the edge  $x_2x_3$  (and without  $x_1x_3$  if  $x_1 \sim x_3$ ). Let  $T'$  be a subtree of  $H'$  with leaves  $x_1, x_2, x_3$  and with  $|E(T')|$  a minimum; specifically, say  $T'$  consists of a vertex  $v'$  and three  $v'$ -to- $x_i$

paths (each either chordless or the single edge  $v'x_i$ ). Let  $G'$  be the subgraph of  $H'$  induced by  $V(T')$ , and note that  $\{x_1, x_2, x_3\}$  is independent in  $G'$ . Each edge  $E(G') - E(T')$  must join two neighbors of  $v'$  in  $T'$  (otherwise, that edge could be inserted into  $T'$  and two or more edges deleted so as to contradict the minimality of  $|E(T')|$ ). Similarly, there can be at most two edges in  $E(G') - E(T')$  that join neighbors of  $v'$  in  $T'$  (if there were three, then two of them could be inserted into  $T'$  and the three edges incident with  $v'$  deleted so as to contradict the minimality of  $|E(T')|$ ). Thus, one of the following three subcases occurs:

- $3'_0$ : The subgraph  $G' = T'$  has no triangles and so consists of the three internally disjoint  $v'$ -to- $x_i$  paths; take  $a' = v'$ .
- $3'_1$ : The subgraph of  $G'$  induced by  $V(T')$  consists of one triangle  $a'b'c'$  and three vertex disjoint paths from  $\{a', b', c'\}$  to  $\{x_1, x_2, x_3\}$  where at most one of  $a', b', c'$  is in the independent set  $\{x_1, x_2, x_3\}$ ; without loss of generality, suppose  $a' \notin \{x_1, x_2, x_3\}$ .
- $3'_2$ : The subgraph of  $G'$  induced by  $V(T')$  consists of two triangles with vertices in  $\{a', b', c', d'\}$  and three vertex disjoint paths from  $\{a', b', c', d'\}$  to  $\{x_1, x_2, x_3\}$  where either  $\{a', b', c', d'\} \cap \{x_1, x_2, x_3\} = \emptyset$  or exactly one of the two vertices of  $\{a', b', c', d'\}$  that are in both triangles (and neither of the other two vertices) is in the independent set  $\{x_1, x_2, x_3\}$ ; without loss of generality, suppose  $a' \notin \{x_1, x_2, x_3\}$ .

In each subcase, define internally disjoint paths  $\tau'_1, \tau'_2, \tau'_3$  such that each  $\tau_i$  is the  $a'$ -to- $x_i$  path in  $G'$  (and so is chordless or a single edge).

Now consider similar cases  $1''$ ,  $2''$ , and  $3''$  (with three subcases  $3''_i$ ) for the components  $G''_1, G''_2, G''_3$  of  $G - S$  by replacing each superscript  $'$  with  $''$  in the corresponding case above so as to define paths  $\tau''_1, \tau''_2$ , and  $\tau''_3$ . Subcases  $3''_1$  and  $3''_2$  cannot occur (otherwise, one or two of  $x_1, x_2, x_2$  could have been replaced by one or two vertices of  $G''$ , contradicting the assumed minimality of  $|V(G'')|$ ). Thus, case  $3''$  will always have  $G'' = T''$ . Finally, for each  $i \in \{1, 2, 3\}$ , define  $\pi_i$  to be the chordless path  $\tau'_i \cup \tau''_i$  in  $G$  (noting that each  $x_i \in \pi_i^\circ$ ).

If neither subcase  $3'_1$  nor  $3'_2$  occurs, then no chord will exist between  $\tau'_i$  and  $\tau'_j$  in  $G'$  and the paths  $\pi_1, \pi_2, \pi_3$  will be as described in the theorem.

On the other hand, if subcase  $3'_1$  or  $3'_2$  does occur, then there were at least two choices for  $a'$ , and so  $a'$  can be chosen so as either to have  $\{i, j\} = \{1, 3\}$  for every chord between  $\tau'_i$  and  $\tau'_j$  in  $G'$  or to have  $\{i, j\} = \{1, 2\}$  for every such chord. (Either way, there will only be one or two such chords.) If  $\{i, j\} = \{1, 3\}$  always holds, then the three paths  $\pi_1, \pi_2, \pi_3$  will be as described in the theorem. If  $\{i, j\} = \{1, 2\}$  always holds, then interchanging  $\pi_1$  and  $\pi_2$  (do this if  $x_1 \sim x_3$ ) or interchanging  $\pi_2$  and  $\pi_3$  (do this if  $x_1 \not\sim x_3$ ) will make the edges of  $E(G') - E(T')$  all into  $\pi_1$ -to- $\pi_3$  chords, and the new  $\pi_1, \pi_2, \pi_3$  will be as described in the theorem.

Therefore, in every combination of cases  $1', 2', 3'$  and  $1'', 2'', 3''$ , the three paths  $\pi_1, \pi_2, \pi_3$  will end up being internally disjoint and chordless with a total of at most  $\min\{k, 6\}$  distinct endpoints, with no  $\pi_1$ -to- $\pi_2$  chord and with a unique  $\pi_2$ -to- $\pi_3$  chord  $x_2x_3$  and  $x_2 \in \pi_2^\circ$  and  $x_3 \in \pi_3^\circ$ .

Consequently, if  $k \geq 6$ , then  $\min\{k, 6\} = 6$  and the preceding paragraphs show  $\mathcal{D}_k^* = \mathcal{D}_6^*$ , and so  $\mathcal{D}^* = \mathcal{D}_6^*$ .  $\square$

To illustrate Theorem 7 in Figure 1 with  $k = 4$ ,  $\mathcal{I} = \{r_1, r_2, r_3, r_4\}$ , and  $S = \{s_1, s_2, s_3, s_4\}$ , observe that the paths  $\pi_1 = r_1, s_1, r_2$  and  $\pi_2 = r_3, s_4, r_4$  and  $\pi_3 = r_1, s_2, r_2$  share four distinct endpoints with no  $\pi_1^\circ$ -to- $\pi_2^\circ$  chord and the unique  $\pi_2^\circ$ -to- $\pi_3^\circ$  chord  $s_2s_4$ . Similarly in Figure 2, the paths  $\pi_1 = r_1, s_2, r_2$  and  $\pi_2 = r_1, s_1, r_2$  and  $\pi_3 = r_3, s_4, r_4$  share four distinct endpoints with no  $\pi_1^\circ$ -to- $\pi_2^\circ$  chord and the unique  $\pi_2^\circ$ -to- $\pi_3^\circ$  chord  $s_1s_4$ .

Theorem 7 would also lead to forbidden induced subgraph characterizations, much as Theorems 3 and 5 did; for instance, characterizing  $\mathcal{D}^* = \mathcal{D}_6^*$  would require 20 forbidden subgraphs. By Proposition 2 and Theorems 5 and 7, the “house graph”—shown as (b) in Figure 3—is the smallest graph that is in  $\mathcal{D}^*$  but is not in  $\mathcal{D}$  or  $\mathcal{D}'$ .

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