# **Domination Polynomials of Graph Products**

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#### **Abstract**

The domination polynomials of binary graph operations, aside from union, join and corona, have not been widely studied. We compute and prove recurrence formulae and properties of the domination polynomials of families of graphs obtained by various products, including both explicit formulae and recurrences for specific families

## 1 Introduction and Definitions

This paper discusses simple undirected graphs G = (V, E). A vertex subset  $W \subseteq V$  of G is a *dominating set* in G, if for each vertex  $v \in V$  of G either v itself or an adjacent vertex is in W.

**Definition 1.1.** Let G = (V, E) be a graph. The domination polynomial D(G, x) is given by

$$D(G,x) = \sum_{i=0}^{|V|} d_i(G)x^i,$$

where  $d_i(G)$  is the number of dominating sets of size i in G. The domination number of a graph G, denoted  $\gamma(G)$ , is the smallest i such that  $d_i(G) > 0$ .

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In [20] we showed that there exist recurrence relations for the domination polynomial which allow for efficient schemes to compute the polynomial for some types of graphs. A recurrence for the domination polynomial of the *path graph* with n vertices  $(P_n)$  was shown in [3] to be

$$D(P_{n+1},x) = x(D(P_n,x) + D(P_{n-1},x) + D(P_{n-2},x))$$
(1)

where  $D(P_0, x) = 1$ ,  $D(P_1, x) = x$  and  $D(P_2, x) = x^2 + 2x$ . Note that the complete graphs  $K_j \cong P_j$  for  $0 \le j \le 2$  and that  $D(K_r, x) = (x+1)^r - 1$ .

Given any two graphs G and H we define the Cartesian product, denoted  $G \square H$ , to be the graph with vertex set  $V(G) \times V(H)$  and edges between two vertices  $(u_1, v_1)$  and  $(u_2, v_2)$  if and only if either  $u_1 = u_2$  and  $v_1v_2 \in E(H)$ , or  $u_1u_2 \in E(G)$  and  $v_1 = v_2$ . As in [14], if  $u \in G$  then the subgraph of  $G \square H$  induced by the vertices (u, v) such that  $v \in H$  will be the H-layer through u and this will be denoted by  $H^u$ . We define  $G^v$  analogously.

The Cartesian product is well known to be commutative and, if G is a disconnected graph with components  $G_1$  and  $G_2$ , then  $G \square H = (G_1 \square H) \cup (G_2 \square H)$ , so that

$$D(G \square H, x) = D(G_1 \square H, x)D(G_2 \square H, x).$$

Despite these properties, it is difficult to determine much about this product, even in such simple cases as the grid graphs  $P_n \square P_m$ , especially in the case of dominating sets. The strong product  $(G \boxtimes H)$  is the graph which is formed by taking the graph  $G \square H$  and then additionally adding edges between vertices  $(u_1, v_1)$  and  $(u_2, v_2)$  if both  $u_1u_2 \in E(G)$  and  $v_1v_2 \in E(H)$ .

The domination numbers of graph products have been extensively studied in the literature, see e.g [1, 5, 9, 10, 12, 13, 16, 19, 22, 23, 25, 26]. In particular, a large number of papers have addressed the domination number of Cartesian products, inspired by the conjecture by V. G. Vizing [27] that  $\gamma(G \square H) \ge \gamma(G) \times \gamma(H)$  (see [6] for a recent survey). In contrast, although the domination polynomial has been actively studied in recent years, almost no attention has been given to the domination polynomials of graph products.

The closed neighbourhood  $N_G[W]$  of a vertex set W in G contains W and all vertices adjacent to vertices in W. When  $W = \{v\}$  we will write  $N_G[v]$  or just N[v] if the graph we are working in is obvious. We define  $N_G(W)$  as the open neighbourhood which includes all neighbours of W that are not in W, so that  $N_G(W) := N_G[W] \setminus W$ . If S is a set of vertices from G we use G - S to mean the graph resulting from the deletion of all vertices in S from G, and let G - v be  $G - \{v\}$ . The vertex contraction G/v denotes the graph obtained from G by the removal of v and the addition of edges between any pair of non-adjacent neighbours of v.

The general reduction formula for any  $u \in V(G)$  given in [20] is the following:

$$D(G,x) = xD(G/u,x) + D(G-u,x) + xD(G-N[u],x) - (1+x)p_u(G,x),$$
 (2)

where  $p_u(G,x)$  is the polynomial which counts the dominating sets of G-u which do not include any vertex from  $N_G(u)$ . Note that if the vertices of N(u) induce a complete

graph then  $G/u \cong G - u$  and so

$$D(G,x) = (x+1)D(G-u,x) + xD(G-N[u],x) - (1+x)p_u(G,x).$$
 (3)

An outline of the paper is as follows. In section 2 we give decomposition formulae for the domination polynomials of the Cartesian product of an arbitrary graph G with  $K_2$  and of the strong product of G with  $K_r$ , then generalise these results. Section 3 gives a recurrence relation for the domination polynomial of any graph which contains  $P_n \square K_2$  that uses only six smaller graphs. A generalisation of the result in section 3 is given in section 4, where we give a recurrence for  $P_n \square K_r$ . In section 5 we give the polynomial for a family of graphs which generalise path graphs.

# 2 Domination Polynomials of Products with Complete Graphs

Let us suppose that  $V(K_2) := \{u, v\}$  in the product  $G \square K_2$  and let G be any non-null graph. We will concentrate first on the vertices in  $G^u$ : every vertex subset W of  $G^u$  can be a subset of some dominating set S in  $G \square K_2$  if some vertices in  $G^v$  are also in S. Let  $W \subseteq V(G)$ , so, by definition, all vertices (y, u) are dominated for  $y \in N_G[W]$  as well as the vertices (w, v) for  $w \in W$ . If S is a dominating set for  $G \square K_2$  such that  $S \cap V(G^u) = W$ , all vertices (y, u) such that  $y \in V(G) \setminus N[W]$  must then be dominated by (y, v), their only neighbour outside of  $G^u$ .

**Theorem 2.1.** Let  $J_W$  be formed from the subgraph of G induced by  $N_G[W]$  by adding a new vertex z joined to the union of W and  $N(V(G) - N_G[W])$ . The domination polynomial for  $D(G \square K_2, x)$  is then equal to

$$\frac{x^{|V(G)|}}{x+1} \times \sum_{W \subset V(G)} \frac{(D(J_W/z, x) + D(J_W - N_{J_W}[z], x) + D(J_W, x) - D(J_W - z, x))}{x^{|N_G(W)|}}.$$

*Proof.* Suppose that  $W \subseteq V(G)$ , so that we know that, in any dominating set for  $G \square K_2$ , if the only vertices from  $G^u$  are W then we must also include all vertices  $(y, v) \in G^v$  where  $y \notin N_G[W]$ . In this way all vertices in  $G^u$  are dominated by

$$|W| + |V(G) \setminus N_G[W]| = |W| + |V(G)| - |N_G[W]| = |V(G)| - |N_G(W)|$$

vertices, giving the powers of x as in the theorem.

It now remains to ensure that all of the vertices in  $G^{\nu}$  are dominated. Using the vertices forced to dominate  $G^{\mu}$  we see that, in  $G^{\nu}$ , every vertex in either W or in N[V(G) - N[W]] is dominated. The only vertices not dominated are therefore those which are in N(W) but have no neighbours outside of N[W]. Let us call this set  $T_W$ .

We now introduce the graph  $J_W$  which is formed by taking the subgraph of G induced by N[W] and adding a new vertex z which is adjacent to every vertex either in W or N(V(G) - N[W]). The vertices which z is joined to are exactly those *not* in  $T_W$ . Thus we want to count all sets of vertices in  $J_W \setminus \{z\}$  such that  $T_W$  is dominated.

As defined in Equation (2),  $p_z(J_W, x)$  generates the dominating sets for  $J_W - N[z]$  which additionally dominate the vertices of N(z). Each of these sets when combined with z is a dominating set for  $J_W$  in which  $T_W$  is dominated and z is only dominated by itself. All other sets which dominate  $T_W$  must then include a vertex from N(z) and hence they will be a dominating set for both  $J_W$  and  $J_W - z$ . The difference of domination polynomials  $D(J_W, x) - D(J_W - z, x)$  generates all such sets which include z and so  $p_z(J_W, x) + D(J_W, x) - D(J_W - z, x)$  generates all sets of vertices in  $J_W$  that dominate  $T_W$  and include z.

Since z is not adjacent to any vertex of  $T_W$  the generating function counting all sets of vertices in  $J_W \setminus \{z\}$  such that  $T_W$  is dominated satisfies the following relation, using Equation (2) for the expansion of  $p_z(J_W, x)$ :

$$\frac{p_{z}(J_{W},x) + D(J_{W},x) - D(J_{W} - z,x)}{x}$$

$$= \frac{xD(J_{W}/z,x) + xD(J_{W} - N_{J_{w}}[z],x) + D(J_{W} - z,x) - D(J_{W},x)}{x(x+1)}$$

$$+ \frac{D(J_{W},x) - D(J_{W} - z,x)}{x}$$

$$= \frac{D(J_{W}/z,x) + D(J_{W} - N_{J_{W}}[z],x)}{x+1} + \frac{D(J_{W},x) - D(J_{W} - z,x)}{x+1} .$$

Putting this together with our first observation finishes the proof.

Since the graphs involved in the summation have at most around half the number of vertices of the product it is significantly faster to use Theorem 2.1 to calculate the domination polynomial even with the summation over all subsets. Additionally, it can be used to get a closed form solution for some highly symmetric graphs as we show in Corollary 2.2.

**Corollary 2.2.** For 
$$r \ge 1$$
,  $D(K_r \square K_2, x) = ((x+1)^r - 1)^2 + 2x^r$ .

*Proof.* When  $G = K_r$  we have  $J_W/z = J_W - z$  for all W since all vertices in G are joined to all others. Unless  $W = \emptyset$  or W = V(G) the sum is therefore  $(x+1)((x+1)^r - 1)$  since  $J_W$  is then  $K_r$  with z joined to the vertices in W; if we combine any non-empty subset of W with or without z we get a dominating set for exactly one of  $J_W$  or  $J_W - N_{J_W}[z]$ . By Theorem 2.1 we then have

$$D(K_r \square K_2, x) = \frac{x^r}{x+1} \left( x+1 + (x+1)^{r+1} + \sum_{j=1}^{r-1} {r \choose j} \frac{(x+1)((x+1)^r - 1)}{x^{r-j}} \right)$$

$$= x^r \left( 1 + (x+1)^r + ((x+1)^r - 1) \sum_{j=1}^{r-1} {r \choose j} \frac{1}{x^{r-j}} \right)$$

$$= x^r + x^r (x+1)^r + ((x+1)^r - 1) \left( \left( \sum_{j=0}^r {r \choose j} x^j \right) - (1+x^r) \right)$$

$$= x^r + ((x+1)^r - 1)^2 + x^r.$$

The following result was also proven independently in [7] as their Lemma 3:

**Theorem 2.3.** For any graph G

$$D(G \boxtimes K_r, x) = D(G, (x+1)^r - 1).$$

*Proof.* Let u be a vertex of G and  $v \in V(K_r)$ ; the closed neighbourhood of the vertex (u,v) is  $(N_G[u],K_r)$ . For any  $X \subseteq V(G)$ , let  $\{A_x \mid x \in X\}$  be a family of arbitrary nonempty subsets of  $V(K_r)$ . We then have that such a set X is a dominating set of G if and only if

$$\bigcup_{x\in X}\{(x,\nu)\mid \nu\in A_x\}$$

is a dominating set of  $G \boxtimes K_r$ . Consequently, each vertex u of a dominating set of G corresponds to all non-empty subsets of the  $K_r$  through u in  $G \boxtimes K_r$ , which are counted by the generating function  $(x+1)^r - 1$ .

Theorem 2.3 can be used to generalise recurrence relations for the domination polynomial of any families of graphs, such as for  $H_{n,r} := P_n \boxtimes K_r$  as follows:

Corollary 2.4. For any integers  $n \ge 3$  and  $r \ge 1$ ,

$$D(H_{n+1,r},x) = ((x+1)^r - 1) \left( D(H_{n,r},x) + D(H_{n-1,r},x) + D(H_{n-2,r},x) \right).$$

Proof. From Equation (1) and using Theorem 2.3 we have

$$D(H_{n+1,r},x) = D(P_{n+1} \boxtimes K_r,x)$$

$$= D(P_{n+1},(x+1)^r - 1)$$

$$= ((x+1)^r - 1)(D(P_n,(x+1)^r - 1) + D(P_{n-1},(x+1)^r - 1) + D(P_{n-2},(x+1)^r - 1))$$

$$= ((x+1)^r - 1)(D(H_{n,r},x) + D(H_{n-1,r},x) + D(H_{n-2,r},x))$$

as required.

Note that, as shown in [3], the same recurrence as Equation (1) holds for the cycle graphs  $C_n$  hence there is an identical generalisation for the domination polynomial of  $C_n \boxtimes K_r$ .

**Corollary 2.5.** For any integers n > 3 and  $r \ge 1$ ,  $D(C_{n+1} \boxtimes K_r, x) =$ 

$$((x+1)^r-1)(D(C_n\boxtimes K_r,x)+D(C_{n-1}\boxtimes K_r,x)+D(C_{n-2}\boxtimes K_r,x)).$$

Corollary 2.2 can be generalised in the following way:

**Theorem 2.6.** The domination polynomial for  $K_r \square K_s$  is, for  $r \ge 2$  and  $s \ge 2$ ,

$$D(K_r \square K_s, x) = ((x+1)^r - 1)^s - \sum_{k=1}^{s-1} \binom{s}{k} (-1)^k \left( (x+1)^{s-k} - 1 \right)^r.$$

*Proof.* We can imagine the vertices of  $K_r \square K_s$  as elements of an  $r \times s$  matrix; for a dominating set in this graph we need to have at least one element in every row and column. The simplest way this can be achieved is to have at least one vertex in every column and the ordinary generating function that generates such sets is  $((x+1)^r - 1)^s$ . However, it is also possible to have empty sets in some columns, so long as each row contains at least one element:

There are s choices for the case of one empty column and, given that choice, the generating function counting non-empty rows of s-1 elements is  $((x+1)^{s-1}-1)^r$ . However, some of the sets counted in this way will have more than one empty column; by the principle of inclusion-exclusion, we now need to subtract the  $\binom{s}{2}$  ways to choose a pair of columns to be empty.

The polynomial counting dominating sets with at least two columns empty is

$$((x+1)^{s-2}-1)^r$$

but this then includes sets with more than two empty columns and so the inclusion-exclusion process will continue. The final case will be when we have all but one column empty, in which case the only possible dominating set contains all r vertices from one column. The term counting all such sets will be  $sx^r = \binom{s}{s-1}((s+1)-1)^r$ , which matches the term in the sum in the theorem when k = s - 1. Combining all of these cases together completes the proof.

**Corollary 2.7.** The domination polynomial for  $K_r \square K_3$  is, for r > 1,

$$((x+1)^r-1)^3+3x^r((x+2)^r-1).$$

*Proof.* Substituting s = 3 into Theorem 2.6 we get

$$D(K_r \square K_3, x) = ((x+1)^r - 1)^3 - \sum_{k=1}^2 {3 \choose k} (-1)^k ((x+1)^{3-k} - 1)^r$$

$$= ((x+1)^r - 1)^3 + 3 (((x+1)^2 - 1)^r - ((x+1) - 1)^r)$$

$$= ((x+1)^r - 1)^3 + 3(((x(x+2))^r - x^r))$$

$$= ((x+1)^r - 1)^3 + 3x^r ((x+2)^r - 1)$$

# **3** The Domination Polynomial for $P_n \square K_2$

Let  $L_n$  be the graph  $P_n \square K_2$  and label the vertices of the two copies of  $P_n$  as  $u_1, \ldots, u_n$  and  $v_1, \ldots, v_n$  where  $u_i$  and  $v_i$  are adjacent,  $i = 1, \ldots, n$ . Note that the graph  $L_{n-1}$  is formed from  $L_n$  by deletion of  $u_n$  and  $v_n$ . The domination polynomials of the first six graphs in the family are given in Table 1.

We first prove a small result which will be used in the main theorem of this section.

Table 1: The domination polynomials for the graphs  $P_n \square K_2$ 

n	$D(P_n \square K_2, x)$
1	$x^2 + 2x$
2	$x^4 + 4x^3 + 6x^2$
3	$x^6 + 6x^5 + 15x^4 + 16x^3 + 3x^2$
4	$x^8 + 8x^7 + 28x^6 + 52x^5 + 48x^4 + 12x^3$
5	$x^{10} + 10x^9 + 45x^8 + 116x^7 + 178x^6 + 148x^5 + 47x^4 + 2x^3$
6	$x^{12} + 12x^{11} + 66x^{10} + 216x^9 + 453x^8 + 604x^7 + 470x^6 + 168x^5 + 17x^4$

**Lemma 3.1.** The polynomial  $A_n(x)$  counting the dominating sets of  $L_n$  such that both  $u_n$  and  $v_n$  are included is

$$A_n(x) := x^2 (D(L_{n-1}, x) + D(L_{n-2}, x) - A_{n-2}(x)).$$

*Proof.* Every dominating set for either  $L_{n-1}$  or  $L_{n-2}$  will be a dominating set for  $L_n$  when combined with  $u_n$  or  $v_n$  since these two vertices dominate themselves and their neighbours. Any set S which is a dominating set in both  $L_{n-1}$  and  $L_{n-2}$  cannot contain either  $u_{n-1}$  or  $v_{n-1}$  since they are not in  $L_{n-2}$  and hence S must contain both  $u_{n-2}$  and  $v_{n-2}$  in order for the former pair of vertices to be dominated. Thus exactly  $x^2A_{n-2}(x)$  sets are counted twice and this is subtracted to give our result.

**Theorem 3.2.** The domination polynomial for  $L_n$  satisfies this recurrence for  $n \ge 6$ :

$$D(L_n,x) = x(x+2)D(L_{n-1},x) + x(x+1)D(L_{n-2},x) + x^2(x+1)D(L_{n-3},x) - x^3D(L_{n-4},x) - x^3D(L_{n-5},x)$$

*Proof.* Let T be a dominating set for  $L_n$  and set  $T_1 := T \setminus \{u_n, v_n\}$ . If  $T_1 = T$  then (in order to have  $u_n$  and  $v_n$  dominated) we can conclude that  $|T \cap \{u_{n-1}, v_{n-1}\}| = 2$  and the polynomial counting such sets will be  $A_{n-1}(x)$  as in Lemma 3.1. This gives us the contribution  $x^2(D(L_{n-2}, x) + D(L_{n-3}, x) - A_{n-3}(x))$  for our summation.

Now suppose  $|T \cap \{u_n, v_n\}| \ge 1$ ; if  $T_1$  is a dominating set for  $L_{n-1}$  then T will be a dominating set for  $L_n$ . Thus we get the term  $x(x+2)D(L_{n-1},x)$ , the x(x+2) coming from that we can use  $u_n$  and/or  $v_n$  with  $T_1$  to form a dominating set.

However, there are circumstances under which  $T_1$  does not have to be a dominating set for  $L_{n-1}$ , since  $u_{n-1}$  and  $v_{n-1}$  in  $L_{n-1}$  might be only dominated by  $u_n$  or  $v_n$  in T. Let us now consider the ways that exist such that  $u_{n-1}$  and  $v_{n-1}$  are not dominated in  $T_1$  but dominated in T.

If both  $u_{n-1}$  and  $v_{n-1}$  are undominated by  $T_1$  then we must have  $|T \cap \{u_n, v_n\}| = 2$  to dominate those vertices and also  $|T \cap \{u_{n-3}, v_{n-3}\}| = 2$  to dominate  $u_{n-2}$  and  $v_{n-2}$ , (and neither  $u_{n-1}$  nor  $v_{n-1}$ ) giving the term  $x^2A_{n-3}(x)$  which will cancel that term introduced at the start of the proof.

We are now left to count just the dominating sets for  $L_{n-2}$  which include only one of  $u_{n-2}$  and  $v_{n-2}$ . These sets will make a previously uncounted dominating set

for  $L_n$  when combined with  $v_n$  and/or  $u_n$  respectively. These are the four different possibilities, defining  $S := T \cap \{u_n, v_n, u_{n-1}, v_{n-1}, u_{n-2}, v_{n-2}\}$ :

- (i)  $S = \{u_n, v_n, v_{n-2}\}$
- (ii)  $S = \{u_n, v_{n-2}\}$
- (iii)  $S = \{v_n, u_{n-2}\}$
- (iv)  $S = \{u_n, v_n, u_{n-2}\}$

To count these possibilities we can now consider the different ways that exactly one of  $u_n$  or  $v_n$  can be combined with a dominating set for  $L_{n-2}$  which will lead to the contribution of the term  $xD(L_{n-2},x)$  to our sum. Suppose Q is a dominating set for  $L_{n-2}$ ; we will split into subcases depending on  $r := |Q \cap \{u_{n-2}, v_{n-2}\}|$  as follows:

Every set Q satisfying r=2 can be converted into a set of the type of possibility (i) (by adding  $u_n$  and switching  $v_n$  for  $u_{n-2}$ ), but this new set will not be a dominating set for  $L_n$  when  $u_{n-3}$  is solely dominated by  $u_{n-2}$  in Q; that is when  $Q \cap \{u_{n-3}, v_{n-3}, u_{n-4}\} = \emptyset$ . Let the sets of this form which have  $v_{n-4} \in Q$  be counted by the polynomial J(x) and such sets which also do not include  $v_{n-4}$  are necessarily  $x^2 A_{n-5}(x)$  as in Lemma 3.1.

When r=1 we can add  $u_n$  or  $v_n$  as appropriate and have possibilities (ii) and (iii) for S. In the case when r=0, Q must include both  $u_{n-3}$  and  $v_{n-3}$  to be dominating. No such set can be combined with just one more vertex to make a dominating set for  $L_n$ , and we can count the sets with r=0 (and one additional unspecified vertex) using the polynomial  $xA_{n-3}(x)$ . Putting these terms together, we see that possibilities (i),(ii) and (iii) are counted by

$$x(D(L_{n-2},x)-J(x)-x^2A_{n-5}-A_{n-3}(x)).$$

Finally, we can count the dominating sets for  $L_n$  with S as in possibility (iv) by using  $x^3D(L_{n-3},x)+xJ(x)$ . We make a slight adjustment in the same way as in the subcase when r=0 since a set in which only  $u_{n-3}$  is not dominated in  $L_{n-3}$  will still be a dominating set in  $L_n$  when combined with this S, and the polynomial counting such sets exactly matches the definition of xJ(x).

Using Lemma 3.1 again, we get that

$$x^{3}A_{n-5}(x) + xA_{n-3}(x) = x^{3} (D(L_{n-4}, x) + D(L_{n-5}, x))$$

and so, summing all of our terms together, we can count all possible dominating sets T for  $L_n$  by using the polynomial in the statement in the theorem.

Note that at no point did we either concern ourselves with the structure beyond  $u_{n-5}$  and  $v_{n-5}$  or utilise the symmetry of  $P_n \square K_2$ , and hence this same recurrence also holds for any family of graphs with  $P_6 \square K_2$  as a pendant subgraph.

We can again use Theorem 2.3 as in Corollary 2.4 to find the domination polynomial for the strong product  $Z_{n,r} := L_n \boxtimes K_r$ :

Corollary 3.3. For any integers  $n \ge 6$  and  $r \ge 1$ ,

$$D(Z_{n,r},x) = ((x+1)^{2r}-1)D(Z_{n-1,r},x) +((x+1)^r-1)(x+1)^rD(Z_{n-2,r},x) +((x+1)^r-1)^2(x+1)^rD(Z_{n-3,r},x) -((x+1)^r-1)^3(D(Z_{n-4,r},x)+D(Z_{n-5,r},x)).$$

*Proof.* Let us substitute  $y := (x+1)^r - 1$  to simplify calculations.

$$D(Z_{n,r},x) = D(L_n,(x+1)^r - 1)$$

$$= D(L_n,y)$$

$$= y(y+2)D(L_{n-1},y) + y(y+1)D(L_{n-2},y)$$

$$+y^2(y+1)D(L_{n-3},y) - y^3D(L_{n-4},y) - y^3D(L_{n-5},y)$$

$$= y(y+2)D(Z_{n-1,r},x) + y(y+1)D(Z_{n-2,r},x)$$

$$+y^2(y+1)D(Z_{n-3,r},x) - y^3D(Z_{n-4,r},x) - y^3D(Z_{n-5,r},x).$$

Utilising now that  $y + 1 := (x + 1)^r$  we get the desired result.

# 4 The Domination Polynomial for $P_n \square K_r$

We denote by  $M_{n,r} := P_n \square K_r$  the Cartesian product of the path  $P_n$  and the complete graph  $K_r$ , where n and r are non-negative integers. We will utilise the linear structure of  $P_n$  and refer to the copy of  $K_r$  corresponding to one of the vertices of degree one in  $P_n$  as at the first  $K_r$  layer and the copy of  $K_r$  adjacent to it as the second  $K_r$  layer. Let  $m_{n,r}^t(x)$  be the polynomial counting the vertex subsets of  $M_{n,r}$  such that all vertices outside of the first  $K_r$  layer are dominated and a particular subset of t of the t vertices of the first  $K_r$  layer is not necessarily dominated.

Let  $\delta_{t,r} := [t = r]$  denote the Kronecker delta function. The graph  $M_{0,r}$  is the null graph and  $M_{1,r} = K_r$  and so only the case of the empty dominating set needs to be considered carefully. For n = 2 the case t = 0 and t > 0 corresponds to Corollary 2.2 and the proof of Theorem 2.6 can be generalised to give the result here.

$$m_{0,r}^{t}(x) = 1$$

$$m_{1,r}^{t}(x) = (x+1)^{r} - 1 + \delta_{t,r}$$

$$m_{2,r}^{t}(x) = (x+1)^{2r} - 2(x+1)^{r} + x^{r} + 1 + x^{(r-t)}(x+1)^{t} - \delta_{t,r}$$
(4)

From these equations we can establish the following recurrence relations for  $m_{n,r}^t$  in general and  $D(M_{n,r},x)=m_{n,r}^0(x)$  in particular.

**Theorem 4.1.** The domination polynomial for  $P_n \square K_r$  (with  $n \ge 3$  and  $r \ge 3$ ) satisfies

$$D(M_{n,r},x) = \sum_{p=1}^{r} {r \choose p} x^{p} m_{n-1,r}^{p}(x) + \sum_{i=0}^{t} {t \choose i} x^{r-i} m_{n-2,r}^{r-i}(x) + \delta_{r,t}(m_{n-1,r}^{0} - m_{n-1,r}^{r})$$

where the  $m_{i,r}^{t}(x)$  terms can be evaluated recursively.

*Proof.* We consider the graph  $M_{n,r}$  for  $n \ge 3$  and note that, on deletion of its first  $K_r$  layer, we get a copy of the graph  $M_{n-1,r}$ . For  $m_{n,r}^r(x)$  we are looking for sets of vertices in which all of the vertices of the first  $K_r$  layer of  $M_{n,r}$  are not necessarily dominated but all other vertices are. We can combine a set of q vertices in the first  $K_r$  layer with a set in  $M_{n-1,r}$  counted by  $m_{n-1,r}^q(x)$  and form a subset of the vertices in  $M_{n,r}$  such that none of the vertices in the first  $K_r$  layer of  $M_{n,r}$  are necessarily dominated but all other vertices are and so

$$m_{n,r}^{r}(x) = \sum_{q=0}^{r} {r \choose q} x^{q} m_{n-1,r}^{q}(x).$$
 (5)

Now suppose that  $0 \le t < r$  and note that this implies that, in every subset S of  $M_{n,r}$  which is counted by  $m_{n,r}^t(x)$ , there must be at least one vertex in the second  $K_r$  layer and therefore all vertices outside of the first  $K_r$  layer are automatically dominated. We need to consider two cases as to whether or not there is a vertex in S from the first  $K_r$  layer.

1. If not then it will be possible to have a set in  $M_{n,r}$  in which the t particular vertices are not necessarily dominated by adding a set including all of the other r-t vertices (and perhaps some others) from the second  $K_r$  layer to a subset of  $M_{n-2,r}$  in which the corresponding vertices in the first  $K_r$  layer of  $M_{n-2,r}$  are not necessarily dominated; this will contribute a term equal to

$$x^{r-t} \sum_{i=0}^{t} {t \choose j} x^{j} m_{n-2,r}^{r-t+j}(x) = \sum_{i=0}^{t} {t \choose i} x^{r-i} m_{n-2,r}^{r-i}(x).$$

2. Now we can assume there is at least one vertex in the first  $K_r$  layer of S and so all vertices will be dominated and the value of t becomes immaterial. We can take any p > 0 vertices in the first  $K_r$  layer and combine them with a set counted by  $m_{n-1,r}^p(x)$  to count such sets. Putting these cases together we get that, for  $0 \le t < r$ ,

$$m_{n,r}^{l}(x) = \sum_{p=1}^{r} {r \choose p} x^{p} m_{n-1,r}^{p}(x) + \sum_{i=0}^{l} {t \choose i} x^{r-i} m_{n-2,r}^{r-i}(x).$$
 (6)

It is possible to combine these two results as follows, after noting that, when t = r we can apply Equation (5) as follows:

$$\sum_{i=0}^{t} {t \choose i} x^{r-i} m_{n-2,r}^{r-i}(x) = \sum_{i=0}^{r} {r \choose i} x^{r-i} m_{n-2,r}^{r-i}(x)$$

$$= \sum_{j=0}^{r} {r \choose j} x^{j} m_{n-2,r}^{j}(x)$$

$$= m_{n-1,r}^{r}.$$

Thus, combining Equations (5) and (6) using the Kronecker delta and noting that the first summations in the two equations differ only by one term,

$$m_{n,r}^{t}(x) = \sum_{p=1}^{r} {r \choose p} x^{p} m_{n-1,r}^{p}(x) + \sum_{i=0}^{t} {t \choose i} x^{r-i} m_{n-2,r}^{r-i}(x) + \delta_{r,t}(m_{n-1,r}^{0} - m_{n-1,r}^{r})$$
 (7)

Equations (4) and (7) can produce all polynomials necessary for this result.  $\Box$ 

## 5 Domination Polynomials of k-path graphs

A different way of combining complete graphs and paths was introduced by Beineke and Pippert in [4]. The k-path graph of length  $n \ge k$  is defined as follows;  $P_n^k$  is an *n*-vertex graph with vertices  $v_1$  to  $v_k$  all being joined to each other and for j > k add edges from vertex  $v_j$  to all vertices from  $v_{j-1}$  to  $v_{j-k}$ .

In [17] the domination polynomial was given for  $P_n^k$  with  $n \le 2k + 6$ , but we can simplify and extend the results given as follows:

**Theorem 5.1.** For  $k \ge 2$  and  $k < n \le 2k + 1$  there is the following recursion:

$$D(P_{n-1}^k, x) = D(P_{n-1}^{k-1}, x) + x(1+x)^{n-1}.$$

*Proof.* Since  $k+1 \le n \le 2k+1$ ,  $\nu_{k+1}$  is adjacent to every other vertex and its removal leaves the (k-1)-path graph with n-1 vertices. Any other vertices when combined with  $\nu_{k+1}$  give a dominating set, leading to the term  $x(1+x)^{n-1}$  and  $\nu_{k+1}$  will be dominated by any other vertex in a dominating set without it, giving us  $D(P_{n-1}^{k-1}, x)$ .

It is possible to find a recursive formula for the polynomial for large n compared to k:

**Theorem 5.2.** For  $n \ge 3k + 2$  we have

$$D(P_n^k, x) = (x+1)D(P_{n-1}^k, x) - xD(P_{n-2(k+1)}^k, x).$$

*Proof.* We use Equation (3) with vertex  $u := v_n$ , after noting that u is adjacent to all vertices from  $v_{n-k}$  to  $v_{n-1}$  we thus have  $P_n^k - N[v_n] = P_{n-k-1}^k$  and hence

$$D(P_n^k, x) = (x+1)D(P_{n-1}^k, x) + xD(P_{n-k-1}^k, x) - (1+x)p_u(P_n^k, x).$$
 (8)

By the definition of  $P_n^k$ , the polynomial  $p_u(P_n^k, x)$  counts the dominating sets for  $P_{n-1}^k$  which do not include any of  $v_{n-k}$  to  $v_{n-1}$ . Aside from these vertices  $v_{n-1}$  is only adjacent to  $v_{n-k-1}$  and so  $p_u(P_n^k, x)$  actually counts the dominating sets for  $P_{n-k-1}^k$  which include  $v_{n-k-1}$ .

Let S be such a set; either  $S - v_{n-k-1}$  is a dominating set for  $P_{n-k-1}^k$  or it does not contain any of the vertices  $v_{n-2k-1}$  to  $v_{n-k-2}$  and thus is a dominating set for  $P_{n-2k-2}^k$ . Therefore we can say that, by counting both the sets S and  $S - v_{n-k-1}$ ,

$$\frac{(1+x)p_{\nu_n}(P_n^k,x)}{x} = D(P_{n-k-1}^k,x) + D(P_{n-2k-2}^k,x),$$

and so we have the relation in the theorem by substituting this into Equation (8).

For n = 2k + 2 the expression given is correct:

$$D(P_{2k+2}^k, x) = (x+1)D(P_{2k+1}^k, x) + xD(P_{k+1}^k, x) - x(1+x)^k.$$

**Theorem 5.3.** The expression given in [17] for the range  $2k + 3 \le n \le 2k + 6$  which simplifies to

$$D(P_{n-k-1}^k, x) = (x+1)D(P_{n-k-1}^k, x) + xD(P_{n-k-1}^k, x) - x(1+x)^{n-2k-2}((1+x)^{k+1} - 1)$$

is actually true for  $2k+3 \le n \le 3k+3$ .

Proof. After using Equation (8) in this case,

$$p_{\nu_n}(P_n^k, x) = x(1+x)^{n-2k-3}((1+x)^{k+1}-1).$$

This is because, as in Theorem 5.2, vertex  $v_{n-k-1}$  has to be in our dominating set S for  $P_{n-k-1}^k$ , and the vertices  $v_{n-2k-2}$  to  $v_{n-k-2}$  are therefore dominated by  $v_{n-k-1}$ . However, since  $n \ge 2k+3$ , vertex  $v_1$  is not dominated by this vertex.

The vertices  $v_1$  to  $v_{k+1}$  form a clique and so will be dominated so long as there is at least one of these vertices in S, giving the factor of  $(1+x)^{k+1}-1$  in our expression. As  $n \le 3k+3$ , we have  $n-2k-2 \le k+1$  and hence all of the other (n-k-1)-1-(k+1)=n-2k-3 vertices are dominated, and so any combination of them can be in S, giving the term  $(1+x)^{n-2k-3}$ .

This completes the calculation of  $D(P_n^k, x)$  for all n and k.

## 6 Future Work

In this paper we investigated the domination polynomials of families of graphs given by products. In a future paper we will be outlining why such recurrence relations can be deduced to exist for many graph products and show implications of their existence to properties of sequences of coefficients of the domination polynomial. Additionally the computional complexity of domination polynomial can be studied and, intriguingly, the evaluation of D(G, -2) can be shown to be potentially significant.

While our results cover some important families of graphs obtained by products, there remain some open problems which we believe deserve attention.

### Problem 6.1.

- 1. How can Theorem 2.1 be extended to deal with basic Cartesian product families such as  $G \square K_{\Sigma}$ ,  $G \square P_{\Sigma}$ ,  $G \square P_{\Sigma}$ , etc.?
- 2. Can analogues of Theorem 2.3 be found for  $G \boxtimes P_s$ ,  $G \boxtimes C_s$ , etc.?
- 3. What other families of graphs obtained using graph products have simple explicit formulae in the spirit of Theorem 2.6?

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