

# Towards the linear arboricity conjecture: graph products

Yuina Tanaka<sup>✉</sup>

## ABSTRACT

For a graph  $G$ , let  $la(G)$  denote the linear arboricity of  $G$  and  $\Delta(G)$  denote the maximum degree of  $G$ . The famous linear arboricity conjecture was made by Akiyama, Exoo, and Harary [Covering and packing in graphs. IV. Linear arboricity] in 1981. It asserts that  $la(G) \leq \left\lceil \frac{\Delta(G)+1}{2} \right\rceil$ . In this paper, we prove the linear arboricity conjecture for products of a path and a complete graph, and for products of a path and a tree.

*Keywords:* linear arboricity, linear arboricity conjecture, graph products

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## 1. Introduction

In graph theory, a *linear forest* is a forest in which every component is a path. For a graph  $G$ , the *linear arboricity* of  $G$  (defined by Harary [8] in 1970) is denoted by  $la(G)$  and is the minimum number of edge-disjoint linear forests whose union is  $E(G)$ . The maximum degree of  $G$  is denoted by  $\Delta(G)$ . We can observe

$$la(G) \geq \left\lceil \frac{\Delta(G)}{2} \right\rceil. \quad (1)$$

Also, it is well-known that, for regular graphs,

$$la(G) \geq \left\lceil \frac{\Delta(G) + 1}{2} \right\rceil. \quad (2)$$

In 1981, Akiyama, Exoo, and Harary [2] made the famous linear arboricity conjecture:

\* Corresponding author.

**Conjecture 1.1** (Akiyama, Exoo, and Harary, [2]). *For every graph  $G$ ,*

$$la(G) \leq \left\lceil \frac{\Delta(G) + 1}{2} \right\rceil.$$

This conjecture is still open, but it was proven for some graphs:

**Theorem 1.2** (Akiyama, Exoo, and Harary, [1]). *For a complete graph  $K_n$ ,*

$$la(K_n) = \left\lceil \frac{n}{2} \right\rceil.$$

**Theorem 1.3** (Akiyama, Exoo, and Harary, [1]). *For a tree  $T$ ,*

$$la(T) = \left\lceil \frac{\Delta(T)}{2} \right\rceil.$$

Also, see: [3, 5, 7, 10, 11].

In 1992, Alon and Spencer [4] proved the following upper bound:

**Theorem 1.4** (Alon and Spencer, [4]). *For every graph  $G$ ,*

$$la(G) \leq \frac{\Delta(G)}{2} + O\left(\Delta(G)^{\frac{2}{3}} \log(\Delta(G))^{\frac{1}{3}}\right).$$

In 2020, Ferber, Fox, and Jain [6] proved the following improved upper bound:

**Theorem 1.5** (Ferber, Fox, and Jain, [6]). *For every graph  $G$ ,*

$$la(G) \leq \frac{\Delta(G)}{2} + C\Delta(G)^{\frac{2}{3}-\alpha},$$

for constants  $C$  and  $\alpha$ .

In this paper, we prove the linear arboricity conjecture for products of a path and a complete graph, and for products of a path and a tree.

For two graphs  $G_1$  and  $G_2$ , the *Cartesian product* of them is denoted by  $G_1 \square G_2$ , whose vertex set is

$$V(G_1 \square G_2) = \{(v_1, v_2) \mid v_1 \in V(G_1), v_2 \in V(G_2)\},$$

and edge set is

$$E(G_1 \square G_2) = \{(v_1, v_2)(v'_1, v'_2) \mid v_1 = v'_1 \text{ and } v_2 v'_2 \in E(G_2), \text{ or } v_1 v'_1 \in E(G_1) \text{ and } v_2 = v'_2\}.$$

For two graphs  $G_1$  and  $G_2$ , the *direct product* of them is denoted by  $G_1 \times G_2$ , whose vertex set is

$$V(G_1 \times G_2) = \{(v_1, v_2) \mid v_1 \in V(G_1), v_2 \in V(G_2)\},$$

and edge set is

$$E(G_1 \times G_2) = \{(v_1, v_2)(v'_1, v'_2) \mid v_1 v'_1 \in E(G_1) \text{ and } v_2 v'_2 \in E(G_2)\}.$$

For two graphs  $G_1$  and  $G_2$ , the *strong product* of them is denoted by  $G_1 \boxtimes G_2$ , whose vertex set is

$$V(G_1 \boxtimes G_2) = \{(v_1, v_2) \mid v_1 \in V(G_1), v_2 \in V(G_2)\},$$

and edge set is

$$E(G_1 \boxtimes G_2) = E(G_1 \square G_2) \cup E(G_1 \times G_2).$$

In the following theorems, we assume  $n, m \geq 2$ , because if  $n = 1$ , then the linear arboricity conjecture is proven by Theorem 1.2. If  $m = 1$ , then  $P_n \square K_1$  is only one path and  $P_n \times K_1$  does not have an edge.

In 2013, Tao and Lin [9] proved the linear arboricity conjecture for Cartesian product of a path and a complete graph.

**Theorem 1.6** (Tao and Lin, [9]). *For Cartesian product of a path and a complete graph,*

$$la(P_n \square K_m) = \left\lceil \frac{m+1}{2} \right\rceil \leq \left\lceil \frac{\Delta(P_n \square K_m) + 1}{2} \right\rceil,$$

for  $n, m \geq 2$ .

In section 2, we first prove the following theorem, which proves the linear arboricity conjecture for direct product of a path and a complete graph.

**Theorem 1.7.** *For direct product of a path and a complete graph,*

$$la(P_n \times K_m) = m - 1 = \left\lceil \frac{\Delta(P_n \times K_m)}{2} \right\rceil,$$

for  $n \geq 3$  and  $m \geq 2$ .

$$la(P_2 \times K_m) = \left\lceil \frac{m}{2} \right\rceil = \left\lceil \frac{\Delta(P_2 \times K_m) + 1}{2} \right\rceil,$$

for  $m \geq 2$ .

By combining Theorems 1.6 and 1.7, we further prove the following theorem for  $n \geq 3$ , which proves the linear arboricity conjecture for strong product of a path and a complete graph. The case  $n = 2$  is an open problem.

**Theorem 1.8.** *For strong product of a path and a complete graph,*

$$la(P_n \boxtimes K_m) = \left\lceil \frac{3m-1}{2} \right\rceil = \left\lceil \frac{\Delta(P_n \boxtimes K_m)}{2} \right\rceil,$$

for  $n \geq 3$  and  $m \geq 2$ .

In Section 3, we study Cartesian, direct, and strong products of a path  $P_n$  with  $n \geq 2$  and a tree  $T$ . If  $n = 1$ , then the linear arboricity conjecture is proven by Theorem 1.3. In the following theorems, we assume  $\Delta(T)$  is even. The case  $\Delta(T)$  is odd is an open problem.

**Theorem 1.9.** *For a path and a tree,*

(a)

$$la(P_n \square T) = \left\lceil \frac{\Delta(P_n \square T)}{2} \right\rceil,$$

for  $n \geq 2$  and even  $\Delta(T)$ .

(b)

$$la(P_n \times T) = \left\lceil \frac{\Delta(P_n \times T)}{2} \right\rceil,$$

for  $n \geq 2$  and even  $\Delta(T)$ .

(c)

$$la(P_n \boxtimes T) = \left\lceil \frac{\Delta(P_n \boxtimes T)}{2} \right\rceil,$$

for  $n \geq 2$  and even  $\Delta(T)$ .

## 2. Products of a path and a complete graph

We first prove Theorem 1.7:

**Proof of Theorem 1.7.** Let

$$V(P_n) = \{u_1, u_2, \dots, u_n\},$$

and

$$E(P_n) = \{u_1u_2, u_2u_3, \dots, u_{n-1}u_n\}.$$

Let

$$V(K_m) = \{v_1, v_2, \dots, v_m\}.$$

For convenience, we will also use notations  $v_{m+1}, v_{m+2}, \dots, v_{2m-1}$ , which are defined by  $v_{m+1} = v_1, v_{m+2} = v_2, \dots, v_{2m-1} = v_{m-1}$ .

We first assume  $n \geq 3$ . So, for  $2 \leq k \leq n-1$  and  $1 \leq l \leq m$ , the degree of  $(u_k, v_l)$  is  $2m-2$ . For  $k=1$  or  $k=n$  and  $1 \leq l \leq m$ , the degree of  $(u_k, v_l)$  is  $m-1$ . So, in  $P_n \times K_m$ , the maximum degree is  $2m-2$  and  $\left\lceil \frac{\Delta(P_n \times K_m)}{2} \right\rceil = m-1$ .

We construct  $m-1$  linear forests by defining their edge sets as follows:

For  $1 \leq t \leq m-1$ , let

$$E_t = \{(u_i, v_j)(u_{i+1}, v_{j+t}) \mid 1 \leq i \leq n-1, i \text{ is odd}, 1 \leq j \leq m\}$$

$$\cup \{(u_k, v_{l+t})(u_{k+1}, v_l) \mid 1 \leq k \leq n-1, k \text{ is even}, 1 \leq l \leq m\}.$$

Then, the edges in  $E_t$  form a linear forest, which has  $m$  disjoint paths. If  $n$  is even, then these  $m$  disjoint paths are

$$\begin{aligned} & \{(u_1, v_1), (u_2, v_{1+t}), (u_3, v_1), (u_4, v_{1+t}), \dots, (u_n, v_{1+t})\}, \\ & \{(u_1, v_2), (u_2, v_{2+t}), (u_3, v_2), (u_4, v_{2+t}), \dots, (u_n, v_{2+t})\}, \\ & \dots \\ & \{(u_1, v_m), (u_2, v_{m+t}), (u_3, v_m), (u_4, v_{m+t}), \dots, (u_n, v_{m+t})\}. \end{aligned}$$

If  $n$  is odd, then these  $m$  disjoint paths are

$$\begin{aligned} & \{(u_1, v_1), (u_2, v_{1+t}), (u_3, v_1), (u_4, v_{1+t}), \dots, (u_n, v_1)\}, \\ & \{(u_1, v_2), (u_2, v_{2+t}), (u_3, v_2), (u_4, v_{2+t}), \dots, (u_n, v_2)\}, \\ & \dots \\ & \{(u_1, v_m), (u_2, v_{m+t}), (u_3, v_m), (u_4, v_{m+t}), \dots, (u_n, v_m)\}. \end{aligned}$$

The  $m-1$  linear forests cover all  $m(m-1)(n-1)$  edges in  $P_n \times K_m$ , and each edge belongs to exactly one linear forest. So,  $la(P_n \times K_m) \leq m-1 = \left\lceil \frac{\Delta(P_n \times K_m)}{2} \right\rceil$ . By combining this result with (1), we get  $la(P_n \times K_m) = m-1 = \left\lceil \frac{\Delta(P_n \times K_m)}{2} \right\rceil$ .

We then assume  $n=2$ . So,  $\Delta(P_2 \times K_m) = m-1 = \Delta(K_m)$ . We need to prove  $la(P_2 \times K_m) = \left\lceil \frac{m}{2} \right\rceil = \left\lceil \frac{\Delta(P_2 \times K_m) + 1}{2} \right\rceil$ .

For  $K_m$ , assume the  $la(K_m) = \left\lceil \frac{m}{2} \right\rceil$  linear forests are  $F_1, F_2, \dots, F_{la(K_m)}$ . For  $P_2 \times K_m$ , we construct  $la(K_m)$  linear forests by defining their edge sets as follows:

For  $1 \leq t \leq la(K_m)$ , let

$$E_t = \{(u_i, v_j)(u_k, v_l) \mid u_i u_k \in E(P_2), v_j v_l \in E(F_t)\}.$$

Then, the edges in  $E_t$  form a linear forest, which has twice as many paths as  $F_t$ . So,  $la(P_2 \times K_m) \leq la(K_m) = \left\lceil \frac{m}{2} \right\rceil$ .  $P_2 \times K_m$  is a regular graph. By combining this result with (2), we get  $la(P_2 \times K_m) = la(K_m) = \left\lceil \frac{m}{2} \right\rceil = \left\lceil \frac{\Delta(P_2 \times K_m) + 1}{2} \right\rceil$ .  $\square$

We combine Theorems 1.6 and 1.7 to prove Theorem 1.8:

**Proof of Theorem 1.8.** We assumed  $n \geq 3$ . So,  $\Delta(P_n \boxtimes K_m) = 3m-1$ . By  $E(P_n \boxtimes K_m) = E(P_n \square K_m) \cup E(P_n \times K_m)$ , we have

$$\begin{aligned} la(P_n \boxtimes K_m) & \leq la(P_n \square K_m) + la(P_n \times K_m) \\ & = \left\lceil \frac{m+1}{2} \right\rceil + m-1 \\ & = \left\lceil \frac{3m-1}{2} \right\rceil \\ & = \left\lceil \frac{\Delta(P_n \boxtimes K_m)}{2} \right\rceil. \end{aligned}$$

By combining this result with (1), we get  $la(P_n \boxtimes K_m) = \left\lceil \frac{\Delta(P_n \boxtimes K_m)}{2} \right\rceil$ .  $\square$

### 3. Products of a path and a tree

In this section, we prove Theorem 1.9:

#### Proof of Theorem 1.9.

Let

$$V(P_n) = \{u_1, u_2, \dots, u_n\},$$

and

$$E(P_n) = \{u_1u_2, u_2u_3, \dots, u_{n-1}u_n\}.$$

Let

$$V(T) = \{v_1, v_2, \dots, v_m\}.$$

(a) We first prove

$$la(P_n \square T) = \left\lceil \frac{\Delta(P_n \square T)}{2} \right\rceil.$$

If  $n \geq 3$ , then  $\Delta(P_n \square T) = \Delta(T) + 2$ . We need to prove  $la(P_n \square T) = \lceil \frac{\Delta(T)+2}{2} \rceil = \frac{\Delta(T)}{2} + 1$ . (We assumed  $\Delta(T)$  is even.) If  $n = 2$ , then  $\Delta(P_n \square T) = \Delta(T) + 1$ . We still need to prove  $la(P_n \square T) = \lceil \frac{\Delta(T)+1}{2} \rceil = \frac{\Delta(T)}{2} + 1$ .

In Theorem 1.3, we have  $la(T) = \lceil \frac{\Delta(T)}{2} \rceil = \frac{\Delta(T)}{2}$ . Assume the  $la(T)$  linear forests are  $F_1, F_2, \dots, F_{la(T)}$ . For  $P_n \square T$ , we construct  $\frac{\Delta(T)}{2} + 1$  linear forests by defining their edge sets as follows:

For  $1 \leq t \leq la(T) = \frac{\Delta(T)}{2}$ , let

$$E_t = \{(u_i, v_j)(u_i, v_k) \mid 1 \leq i \leq n, 1 \leq j, k \leq m, v_jv_k \in E(F_t)\}.$$

Let

$$E_{\frac{\Delta(T)}{2}+1} = \{(u_i, v_j)(u_{i+1}v_j) \mid 1 \leq i \leq n-1, 1 \leq j \leq m\}.$$

Our idea is: For  $1 \leq t \leq la(T) = \frac{\Delta(T)}{2}$ , we let  $E_t$  be the set of edges in the  $n$  copies of  $F_t$ . Then, we let  $E_{\frac{\Delta(T)}{2}+1}$  be the set of edges in the  $m$  copies of  $P_n$ . So, the edges in each of  $E_1, E_2, \dots, E_{\frac{\Delta(T)}{2}+1}$  form a linear forest.

So,  $la(P_n \square T) \leq \frac{\Delta(T)}{2} + 1 = \lceil \frac{\Delta(P_n \square T)}{2} \rceil$ . By combining this result with (1), we get  $la(P_n \square T) = \lceil \frac{\Delta(P_n \square T)}{2} \rceil$ .

(b) We prove

$$la(P_n \times T) = \left\lceil \frac{\Delta(P_n \times T)}{2} \right\rceil.$$

We first assume  $n \geq 3$ . So,  $\Delta(P_n \times T) = 2\Delta(T)$ . We need to prove  $la(P_n \times T) = \Delta(T)$ .

In Theorem 1.3, we have  $la(T) = \lceil \frac{\Delta(T)}{2} \rceil = \frac{\Delta(T)}{2}$ . (We assumed  $\Delta(T)$  is even.) Assume the  $la(T)$  linear forests are  $F_1, F_2, \dots, F_{la(T)}$ . We divide the vertices in  $T$  into two parts  $V^+$  and  $V^-$ :

We choose a vertex  $v$  and put it in  $V^+$ . Then, every vertex with even distance from  $v$  is put in  $V^+$  and every vertex with odd distance from  $v$  is put in  $V^-$ . Because  $T$  is a tree, every vertex in  $V^+$  is only adjacent to vertices in  $V^-$  and every vertex in  $V^-$  is only adjacent to vertices in  $V^+$ .

For  $P_n \times T$ , we construct  $2la(T) = \Delta(T)$  linear forests by defining their edge sets as follows:

For  $1 \leq t \leq la(T) = \frac{\Delta(T)}{2}$ , let

$$E_t^1 = \{(u_i, v_j)(u_{i+1}, v_k) \mid 1 \leq i \leq n-1, v_j v_k \in E(F_t), v_j \in V^+, v_k \in V^-\},$$

and

$$E_t^2 = \{(u_i, v_j)(u_{i+1}, v_k) \mid 1 \leq i \leq n-1, v_j v_k \in E(F_t), v_j \in V^-, v_k \in V^+\}.$$

The edges in  $E_t^1$  form a linear forest, because, for every  $1 \leq i \leq n-1$ ,  $\{(u_i, v_j)(u_{i+1}, v_k) \mid v_j v_k \in E(F_t), v_j \in V^-, v_k \in V^+\}$  form a linear forest which is isomorphic to  $F_t$ . If we take different  $i' \neq i$ , then  $\{(u_{i'}, v_j)(u_{i'+1}, v_k) \mid v_j v_k \in E(F_t), v_j \in V^-, v_k \in V^+\}$  is disjoint with  $\{(u_i, v_j)(u_{i+1}, v_k) \mid v_j v_k \in E(F_t), v_j \in V^-, v_k \in V^+\}$ . So,  $E_t^1$  is a linear forest, because it is a disjoint union of  $n-1$  linear forests.

Similarly, the edges in  $E_t^2$  also form a linear forest. So,  $la(P_n \times T) \leq 2la(T) = \Delta(T)$ . By combining this result with (1), we get  $la(P_n \times T) = \Delta(T) = \lceil \frac{\Delta(P_n \times T)}{2} \rceil$ .

We then assume  $n = 2$ . So,  $\Delta(P_n \times T) = \Delta(T)$ . We need to prove  $la(P_n \times T) = \frac{\Delta(T)}{2}$ , which can be proven by a similar construction: We construct  $E_t^1$  and  $E_t^2$  in the same way, and we can combine  $E_t^1$  and  $E_t^2$ , because  $\{(u_1, v_j)(u_2, v_k) \mid v_j v_k \in E(F_t), v_j \in V^+, v_k \in V^-\}$  and  $\{(u_1, v_j)(u_2, v_k) \mid v_j v_k \in E(F_t), v_j \in V^-, v_k \in V^+\}$  are disjoint. So,  $E_t^1 \cup E_t^2$  form a linear forest when  $n = 2$ .

So, we get  $\frac{\Delta(T)}{2}$  linear forests and  $la(P_n \times T) \leq \frac{\Delta(T)}{2}$ . By combining this result with (1), we get  $la(P_n \times T) = \frac{\Delta(T)}{2}$ .

(c) We prove

$$la(P_n \boxtimes T) = \left\lceil \frac{\Delta(P_n \boxtimes T)}{2} \right\rceil.$$

We first assume  $n \geq 3$ . So,  $\Delta(P_n \boxtimes T) = 3\Delta(T) + 2$ . By  $E(P_n \boxtimes T) = E(P_n \square T) \cup E(P_n \times T)$ , we have

$$\begin{aligned} la(P_n \boxtimes T) &\leq la(P_n \square T) + la(P_n \times T) \\ &= \frac{\Delta(T)}{2} + 1 + \Delta(T) \\ &= \frac{3\Delta(T) + 2}{2} \\ &= \left\lceil \frac{\Delta(P_n \boxtimes T)}{2} \right\rceil. \end{aligned}$$

By combining this result with (1), we get  $la(P_n \boxtimes T) = \left\lceil \frac{\Delta(P_n \boxtimes T)}{2} \right\rceil$ .

We then assume  $n = 2$ . So,  $\Delta(P_n \boxtimes T) = 2\Delta(T) + 1$ . By  $E(P_n \boxtimes T) = E(P_n \square T) \cup E(P_n \times T)$ , we have

$$\begin{aligned} la(P_n \boxtimes T) &\leq la(P_n \square T) + la(P_n \times T) \\ &= \frac{\Delta(T)}{2} + 1 + \frac{\Delta(T)}{2} \\ &= \frac{2\Delta(T) + 2}{2} \\ &= \left\lceil \frac{2\Delta(T) + 1}{2} \right\rceil \\ &= \left\lceil \frac{\Delta(P_n \boxtimes T)}{2} \right\rceil. \end{aligned}$$

By combining this result with (1), we get  $la(P_n \boxtimes T) = \left\lceil \frac{\Delta(P_n \boxtimes T)}{2} \right\rceil$ . □

## 4. Open problem

In this paper, we proved the linear arboricity conjecture for products of a path and a complete graph, and for products of a path and a tree. However, in Theorem 1.8, we assumed  $n \geq 3$ . In Theorem 1.9, we assumed  $\Delta(T)$  is even. So, there are two open problems.

**Problem 4.1** (about Theorem 1.8). *Prove the linear arboricity conjecture for  $P_n \boxtimes K_m$  when  $n = 2$ .*

**Problem 4.2** (about Theorem 1.9). *Prove the linear arboricity conjecture for  $P_n \square T$ ,  $P_n \times T$ , and  $P_n \boxtimes T$  when  $\Delta(T)$  is odd.*

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Yuina Tanaka

Hosei University, 3-7-2 Kajino-cho, Koganei-shi, Tokyo, 184-8584, Japan

Liberty University, Lynchburg, VA 24515, USA

E-mail [tanakayuina@outlook.com](mailto:tanakayuina@outlook.com)