



# On the automorphism group of putative $(n, r)$ -arcs in $\text{PG}(2, 11)$ and $\text{PG}(2, 13)$

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## ABSTRACT

An  $(n, r)$ -arc in  $\text{PG}(2, q)$  is a set  $\mathcal{B}$  of points in  $\text{PG}(2, q)$  such that each line in  $\text{PG}(2, q)$  contains at most  $r$  elements of  $\mathcal{B}$  and such that there is at least one line containing exactly  $r$  elements of  $\mathcal{B}$ . The value  $m_r(2, q)$  denotes the maximal number  $n$  of points in the projective geometry  $\text{PG}(2, q)$  for which an  $(n, r)$ -arc exists. We show by systematically excluding possible automorphisms that putative  $(44, 5)$ -arcs,  $(90, 9)$ -arcs in  $\text{PG}(2, 11)$ , and  $(39, 4)$ -arcs in  $\text{PG}(2, 13)$ —in case of their existence—are rigid, i.e. they all would only admit the trivial automorphism group of order 1. In addition, putative  $(50, 5)$ -arcs,  $(65, 6)$ -arcs,  $(119, 10)$ -arcs,  $(133, 11)$ -arcs, and  $(146, 12)$ -arcs in  $\text{PG}(2, 13)$  would be rigid or would admit a unique automorphism group (up to conjugation) of order 2.

*Keywords:* projective geometry, arcs, linear codes

## 1. Introduction

**Definition 1.1.** An  $(n, r)$ -arc in  $\text{PG}(2, q)$  is a set  $\mathcal{B}$  of points in  $\text{PG}(2, q)$  such that each line in  $\text{PG}(2, q)$  contains at most  $r$  elements of  $\mathcal{B}$  and such that there is at least one line containing exactly  $r$  elements of  $\mathcal{B}$ .

It is well-known (e.g. see [3]) that  $(n, r)$ -arcs in  $\text{PG}(2, q)$  are closely related to error-correcting linear codes: The  $n$  points of an  $(n, r)$ -arc in  $\text{PG}(2, q)$  define the columns of a  $3 \times n$  generator matrix of linear  $[n, 3, n - r]_q$  code, which is code of length  $n$ , dimension 3, and minimum distance  $n - r$  with respect to the Hamming metric. The linear code is projective since the columns of any generator matrix are pairwise linearly independent.

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**Definition 1.2.** Let  $m_r(2, q)$  denote the maximum number  $n$  for which an  $(n, r)$ -arc in  $\text{PG}(2, q)$  exists.

A major goal in studying  $(n, r)$ -arcs in  $\text{PG}(2, q)$  is the determination of  $m_r(2, q)$ .

In general it is hard to determine the exact value of  $m_r(2, q)$  and in most cases instead of the exact value only a lower and an upper bound for  $m_r(2, q)$  are known. An explicit construction of an  $(n, r)$ -arc in  $\text{PG}(2, q)$  yields a lower bound  $m_r(2, q) \geq n$ .

The values  $m_r(2, q)$  with  $q \leq 9$  are exactly determined (see [3]). For  $m_r(2, q)$  with  $11 \leq q \leq 19$  we refer to [2] whereas a table for  $23 \leq q \leq 31$  can be compiled from several sources e.g. [4, 7, 9, 8, 12, 10, 14, 13, 11, 15, 16, 17, 20].

In this article, we show that the putative  $(n, r)$ -arcs in  $\text{PG}(2, q)$  for  $q \in \{11, 13\}$  for the open gaps between lower and upper bound on  $m_r(2, q)$ —in case of existence—are rigid or only admit a unique automorphism of order 2.

## 2. Construction of $(n, r)$ -arcs by integer linear programming

The construction of  $(n, r)$ -arcs in  $\text{PG}(2, q)$  with prescribed groups of automorphisms can be described as integer linear programming (see [6, 4]):

In the following, let

$$\left[ \begin{matrix} \text{GF}(q)^n \\ k \end{matrix} \right],$$

denote the set of  $k$ -dimensional subspaces of  $\text{GF}(q)^n$ , which is called the *Grassmannian*. Its cardinality is given by the *Gaussian number*, also called  $q$ -Binomial coefficient:

$$\left[ \begin{matrix} n \\ k \end{matrix} \right]_q = \left| \left[ \begin{matrix} \text{GF}(q)^n \\ k \end{matrix} \right] \right| = \prod_{i=0}^{k-1} \frac{q^n - q^i}{q^k - q^i}.$$

In terms of vector spaces, an  $(n, r)$ -arc in  $\text{PG}(2, q)$  corresponds to a set  $\mathcal{B} \subseteq \left[ \begin{matrix} \text{GF}(q)^3 \\ 1 \end{matrix} \right]$  such that for all  $H \in \left[ \begin{matrix} \text{GF}(q)^3 \\ 2 \end{matrix} \right]$  holds:

$$|\{P \in \mathcal{B} \mid H \supseteq P\}| \leq r.$$

If  $\left[ \begin{matrix} \text{GF}(q)^3 \\ 1 \end{matrix} \right] = \{P_1, \dots, P_{q^2+q+1}\}$  and  $\left[ \begin{matrix} \text{GF}(q)^3 \\ 2 \end{matrix} \right] = \{H_1, \dots, H_{q^2+q+1}\}$ , where  $\left[ \begin{matrix} 3 \\ 1 \end{matrix} \right]_q = \left[ \begin{matrix} 3 \\ 2 \end{matrix} \right]_q = q^2 + q + 1$ , we define the  $(q^2 + q + 1) \times (q^2 + q + 1)$  incidence matrix

$$A(q) = (a_{ij}),$$

with entries

$$a_{ij} := \begin{cases} 1 & \text{if } H_i \supseteq P_j, \\ 0 & \text{otherwise.} \end{cases}$$

**Lemma 2.1.** If  $u = (1, \dots, 1)^T$  denotes the all-one vector any binary column vector  $x$  satisfying

$$A(q) \cdot x \leq r \cdot u,$$

is equivalent to a  $(u^T \cdot x, r)$ -arc in  $\text{PG}(2, q)$ .

**Corollary 2.2.** The determination of  $m_r(2, q)$  corresponds to the following integer linear programming problem:

$$m_r(2, q) = \max_{x \in \{0,1\}^{q^2+q+1}} \{u^T \cdot x \mid A(q) \cdot x \leq r \cdot u\}.$$

The incidence preserving bijections (automorphisms) of our ambient space for  $(n, r)$ -arcs—the projective geometry  $\text{PG}(2, q)$ —are defined by the projective semi-linear group  $\text{P}\Gamma\text{L}(3, q)$  (see [1]). It acts transitively on the Grassmannian  $\left[ \begin{smallmatrix} \text{GF}(q)^3 \\ k \end{smallmatrix} \right]$ .

Hence, any subgroup  $G \leq \text{P}\Gamma\text{L}(3, q)$  partitions the Grassmannian into  $G$ -orbits. If  $\alpha \in \text{P}\Gamma\text{L}(3, q)$  and  $S \in \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ k \end{smallmatrix} \right]$  we denote by

$$\alpha S := \{\alpha x \mid x \in S\},$$

the transformed subspace and by

$$G(S) := \{\alpha S \mid \alpha \in G\},$$

the  $G$ -orbit of  $S$ . The set of all  $G$ -orbits will be written as

$$G \backslash \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ k \end{smallmatrix} \right].$$

**Definition 2.3.** An  $(n, r)$ -arc  $\mathcal{B}$  in  $\text{PG}(2, q)$  admits a subgroup  $G \leq \text{P}\Gamma\text{L}(3, q)$  as a *group of automorphisms* if and only if  $\mathcal{B}$  consists of  $G$ -orbits on  $\left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 1 \end{smallmatrix} \right]$ . The maximal group of automorphisms of  $\mathcal{B}$  is called the *automorphism group* of  $\mathcal{B}$  and abbreviated by

$$\text{Aut}(\mathcal{B}).$$

**Definition 2.4.** Let  $m_r^G(2, q)$  denote the maximal size  $n$  of an  $(n, r)$ -arc in  $\text{PG}(2, q)$  admitting  $G \leq \text{P}\Gamma\text{L}(3, q)$  as a group of automorphisms.

**Corollary 2.5.** For any  $G \leq \text{P}\Gamma\text{L}(3, q)$  we get a lower bound

$$m_r^G(2, q) \leq m_r(2, q).$$

In particular, for the trivial group  $G = \{1\}$  we have

$$m_r^{\{1\}}(2, q) = m_r(2, q).$$

If  $\{P_1, \dots, P_\ell\}$  denotes a set of representatives of the orbits  $G \backslash \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 1 \end{smallmatrix} \right]$  and  $\{H_1, \dots, H_\ell\}$  a transversal of the orbits  $G \backslash \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 2 \end{smallmatrix} \right]$  for any  $G \leq \text{P}\Gamma\text{L}(3, q)$  we define the  $G$ -incidence matrix  $A(G) = (a_{ij})$  with

$$a_{ij} := |\{P \in G(P_j) \mid H_i \supseteq P\}|.$$

Furthermore, by  $w(G) = (w_1, \dots, w_\ell)^T$  we denote the vector of the lengths of  $G$ -orbits on  $\left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 1 \end{smallmatrix} \right]$ , i.e.

$$w_j := |G(P_j)|.$$

Note that the number of orbits of  $G$  on the set of points and hyperplanes is equal

$$\ell = \left| G \backslash \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 1 \end{smallmatrix} \right] \right| = \left| G \backslash \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 2 \end{smallmatrix} \right] \right| \leq q^2 + q + 1.$$

**Theorem 2.6.** Any binary vector  $x$  of length  $\ell = |G \backslash \left[ \begin{smallmatrix} \text{GF}(q)^3 \\ 1 \end{smallmatrix} \right]|$  with

$$A(G) \cdot x \leq r \cdot u,$$

corresponds to a  $(w(G)^T \cdot x, r)$ -arc in  $\text{PG}(2, q)$  admitting  $G \leq \text{P}\Gamma\text{L}(3, q)$  as a group of automorphisms. In addition, we obtain the following integer linear programming:

$$m_r^G(2, q) = \max_{x \in \{0,1\}^\ell} \{w(G)^T \cdot x \mid A(G) \cdot x \leq r \cdot u\}.$$

### 3. Excluding automorphisms

An open gap is an entry in the tables of  $m_r(2, q)$  for which upper and lower bound differ:

$$\ell \leq m_r(2, q) \leq u \quad \text{where} \quad \ell < u.$$

In that case the question is whether an  $(\ell + 1, r)$ -arc in  $\text{PG}(2, q)$  exists or not. We call such an arc a *putative* arc in  $\text{PG}(2, q)$ . In [5] it was shown that for the gap  $100 \leq m_{10}(2, 11) \leq 101$  a putative  $(101, 10)$ -arc in  $\text{PG}(2, 11)$  admits—in case of its existence—only the trivial automorphism group of order 1.

In this paper we consider the remaining gaps for  $q = 11$  and  $q = 13$  which are given in Table 1.

**Table 1.** Open cases of  $m_r(2, q)$  for  $q = 11, 13$

$m_5(2, 11)$	$\in$	$\{43, 44, 45\}$
$m_9(2, 11)$	$\in$	$\{89, 90\}$
$m_4(2, 13)$	$\in$	$\{38, 39, 40\}$
$m_5(2, 13)$	$\in$	$\{49, 50, 51, 52, 53\}$
$m_6(2, 13)$	$\in$	$\{64, 65, 66\}$
$m_{10}(2, 13)$	$\in$	$\{118, 119\}$
$m_{11}(2, 13)$	$\in$	$\{132, 133\}$
$m_{12}(2, 13)$	$\in$	$\{145, 146, 147\}$

**Definition 3.1.** Let  $\mathcal{B}, \mathcal{B}'$  be an  $(n, r)$ -arcs in  $\text{PG}(2, q)$ . The two sets  $\mathcal{B}$  and  $\mathcal{B}'$  are defined to be isomorphic if and only if there exists  $\alpha \in \text{P}\Gamma\text{L}(3, q)$  such that

$$\alpha\mathcal{B} := \{\alpha P \mid P \in \mathcal{B}\} = \mathcal{B}'.$$

The set of all arcs that are *isomorphic* to  $\mathcal{B}$  is denoted by

$$\text{P}\Gamma\text{L}(3, q)(\mathcal{B}) := \{\alpha\mathcal{B} \mid \alpha \in \text{P}\Gamma\text{L}(3, q)\}.$$

Note that due to the incidence preserving property of  $\text{P}\Gamma\text{L}(3, q)$  isomorphic arcs have the same parameters.

The following lemma is well-known from the theory of group actions (cf. [21]) and states that the automorphism groups of isomorphic objects are conjugated.

**Lemma 3.2.** *Let  $\mathcal{B}$  be an  $(n, r)$ -arc in  $\text{PG}(2, q)$  and let  $\alpha \in \text{P}\Gamma\text{L}(3, q)$ . Then we obtain:*

$$\text{Aut}(\alpha\mathcal{B}) = \alpha \text{Aut}(\mathcal{B})\alpha^{-1} = \{\alpha\beta\alpha^{-1} \mid \beta \in \text{Aut}(\mathcal{B})\}.$$

If  $\mathcal{B}$  in an  $(n, r)$ -arc in  $\text{PG}(2, q)$  with  $G \leq \text{Aut}(\mathcal{B})$  then any isomorphic arc  $\mathcal{B}' = \alpha\mathcal{B}$  for  $\alpha \in \text{P}\Gamma\text{L}(3, q)$  admits the conjugated group  $G' = \alpha G\alpha^{-1}$  satisfies

$$G' = \alpha G\alpha^{-1} \leq \alpha \text{Aut}(G)\alpha^{-1} = \text{Aut}(\alpha\mathcal{B}) = \text{Aut}(\mathcal{B}'),$$

which means that the conjugated group  $G'$  also occurs as a group of automorphisms of  $\mathcal{B}'$ .

As a consequence, when aiming for  $(n, r)$ -arcs in  $\text{PG}(2, q)$  with prescribed groups of automorphisms it is sufficient to consider representatives of conjugacy classes of subgroups of  $\text{PGL}(3, q)$  as possible candidates for potential groups to be prescribed.

Furthermore, any  $(n, r)$ -arc  $\mathcal{B}$  in  $\text{PG}(2, q)$  with  $\{1\} < G \leq \text{Aut}(\mathcal{B})$  also admits all cyclic subgroups  $C \leq G$  as groups of automorphisms.

**Corollary 3.3.** *If we can show for all representatives  $C$  of conjugacy classes of nontrivial cyclic subgroups of  $\text{PGL}(3, q)$  that no  $(n, r)$ -arc in  $\text{PG}(2, q)$  exists with  $C$  as group as automorphisms, either the automorphism group of such arcs are trivial or arcs with that set of parameters do not exist.*

In case of a prime field  $\text{GF}(q)$  the projective semi-linear group is exactly the projective linear group

$$\text{PFL}(3, q) = \text{PGL}(3, q).$$

In the following, a transversal of conjugacy classes of cyclic subgroups of  $\text{PGL}(3, q)$  will be abbreviated by

$$\text{Conj}(q).$$

Its cardinality is given by (see [22]):

$$|\text{Conj}(q)| = \begin{cases} q^2 + q + 2 & \text{if 3 divides } q - 1, \\ q^2 + q & \text{otherwise.} \end{cases}$$

**Lemma 3.4.** *Let  $q$  be a prime. If*

$$m_r^C(2, q) < n \quad \forall C \in \text{Conj}(q) \setminus \{\{1\}\},$$

*one of the following conditions holds:*

1.  $m_r(2, q) < n$ .
2.  $(n, r)$ -arcs  $\mathcal{B}$  in  $\text{PG}(2, q)$  exist where  $\text{Aut}(\mathcal{B}) = \{1\}$ .

We now apply this corollary to the parameters  $(q, n, r) = (11, 44, 5)$ ,  $(q, n, r) = (11, 90, 9)$ , and  $(q, n, r) = (13, 50, 5)$ . There are  $|\text{Conj}(11)| = 132$  conjugacy classes of cyclic subgroups of  $\text{PGL}(3, 11)$  and  $|\text{Conj}(13)| = 184$  classes in  $\text{PGL}(3, 13)$ . We compute the representatives using GAP [18]. By solving the integer linear programming  $m_r^C(2, q)$  according to Theorem 2.6 for all  $C \in \text{Conj}(q) \setminus \{1\}$  using Gurobi [19] we obtain with a runtime less than 3 hours on a 1.2 GHz Intel Core m3 processor the following result:

**Theorem 3.5.** *In case of their existence the automorphism groups of  $(44, 5)$ -arcs,  $(90, 9)$ -arcs in  $\text{PG}(2, 11)$ , and  $(50, 5)$ -arcs in  $\text{PG}(2, 13)$  would be trivial of order 1.*

For the remaining parameters  $r \in \{5, 6, 10, 11, 12\}$  for  $q = 13$  we apply a slightly adapted version of Lemma 3.4 since for these open cases exactly one cyclic subgroup

$$C_0 := \left\langle \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 12 \end{pmatrix} \right\rangle,$$

of order 2 could not directly be excluded to be a group of automorphisms of putative arcs since we cancelled the ILP solver for  $m_r^{C_0}(2, 13)$  after 5000 seconds (for each value  $r$ ).

But it is obvious to guess that this group can probably also be excluded if we spend more running time on the ILP solver.

**Lemma 3.6.** *Let  $q$  be a prime. Let  $C_0 \in \text{Conj}(q)$ . If*

$$m_r^C(2, q) < n \quad \forall C \in \text{Conj}(q) \setminus \{\{1\}, C_0\},$$

*one of the following conditions holds:*

1.  $m_r(2, q) < n$ .
2.  $(n, r)$ -arcs  $\mathcal{B}$  in  $\text{PG}(2, q)$  exist where either  $\text{Aut}(\mathcal{B}) = \{1\}$  or  $\text{Aut}(\mathcal{B})$  is conjugated to  $C_0$ .

Finally, we get

**Theorem 3.7.** *In case of their existence the automorphism groups of  $(50, 5)$ -arcs,  $(65, 6)$ -arcs,  $(119, 10)$ -arcs,  $(133, 11)$ -arcs, and  $(146, 12)$ -arcs in  $\text{PG}(2, 13)$  would either be trivial of order 1 or would be the following cyclic subgroup of order 2 (up to conjugation):*

$$C_0 := \left\langle \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 12 \end{pmatrix} \right\rangle.$$

## Conflicts of interests

The author declares that he has no competing interests.

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